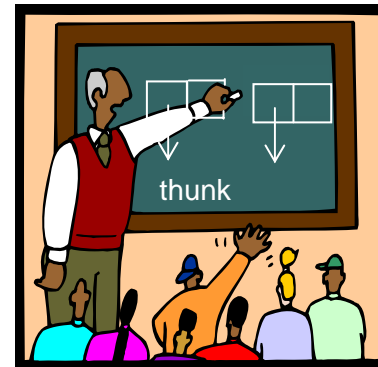


6.001 recitation 20

5/2/07

- lazy eval
- streams



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extending our evaluator: lazy evaluation

Key ideas:

- procedure args are not evaluated until needed
- represent delayed args as objects
- lazy eval can be added easily
- add new syntax

example a: applicative order

```
(define (foo x)
```

```
  (display 'foo)
```

```
  (+ x x))
```

```
(foo (bar 2))
```

```
(define (bar x)
```

```
  (display 'arg)
```

```
  (display x)
```

```
  x)
```

What is printed out? (via display and as a final return value)

a. *applicative order*:

(foo (bar 2)) *output*

printout

example b: normal (lazy) order

```
(define (foo x)
```

```
  (display 'foo)
```

```
  (+ x x))
```

```
(foo (bar 2))
```

```
(define (bar x)
```

```
  (display 'arg)
```

```
  (display x)
```

```
  x)
```

What is printed out? (via display and as a final return value)

^{ab} b. *normal order (foo's parameter x is delayed):*

(foo (bar 2)) *output*

printout

example c: normal with memoization

```
(define (foo x)
  (display 'foo)
  (+ x x))

(define (bar x)
  (display 'arg)
  (display x)
  x)

(foo (bar 2))
```

What is printed out? (via display and as a final return value)

c. *normal order with memoization (foo's parameter is delayed and stored)*

(foo (bar 2)) *output*

printout

problem 1a: applicative order

```
1. (define y 5)
   (define (foo x)
     (display 'foo)
     (+ x x))

   (define (baz x)
     (display 'arg)
     (set! y (+ y x))
     (display y)
     y)

   (foo (baz 2))
```

What is printed out? (via display and as a final return value)

a. *applicative order*

printout

problem 1b: normal (lazy) order

```
1. (define y 5)
   (define (foo x)
     (display 'foo)
     (+ x x))

   (define (baz x)
     (display 'arg)
     (set! y (+ y x))
     (display y)
     y)

   (foo (baz 2))
```

What is printed out? (via display and as a final return value)

b. *normal order (foo's parameter x)*

printout

problem 1c: normal order with memoization

```
1. (define y 5)
   (define (foo x)
     (display 'foo)
     (+ x x))

   (define (baz x)
     (display 'arg)
     (set! y (+ y x))
     (display y)
     y)

   (foo (baz 2))
```

What is printed out? (via display and as a final return value)

c. normal order with memoization (foo's parameter x)

printout

problem 2a: applicative order

```
2. (define (initialized-list f n)
    (define (helper n lst)
      (if (= n 0) lst
          (helper (- n 1) (cons (f n) lst))))
    (helper n '()))
```

; example output:

```
(initialized-list (lambda(x) (* x x)) 5)
```

```
; value (1 4 9 16 25)
```

```
(define (accum)
  (let ((count 0))
    (lambda (x)
      (set! count (+ x count))
      count)))
```

What is the value of the statement `(initialized-list (accum) 5)`

a. *applicative order*

printout

problem 2b: normal (lazy) order

```
2. (define (initialized-list f n)
    (define (helper n lst)
      (if (= n 0) lst
          (helper (- n 1) (cons (f n) lst))))
    (helper n '()))
```

; example output:

```
(initialized-list (lambda(x) (* x x)) 5)
```

```
; value (1 4 9 16 25)
```

```
(define (accum)
  (let ((count 0))
    (lambda (x)
      (set! count (+ x count))
      count)))
```

What is the value of the statement `(initialized-list (accum) 5)`

b. *normal order (initialized-list's parameter f)*

printout

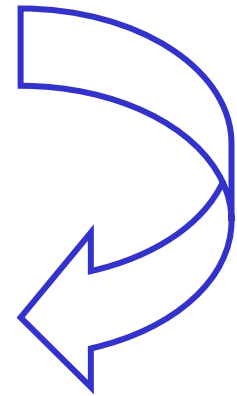
representing delayed objects: thanks

thunks: delay-it, force-it (without memoization)

```
(define (delay-it exp env) (list 'thunk exp env))
(define (thunk? obj) (tagged-list? obj 'thunk))
(define (thunk-exp thunk) (cadr thunk))
(define (thunk-env thunk) (caddr thunk))
```

```
(define (force-it obj)
  (cond ((thunk? obj)
        (actual-value (thunk-exp obj)
                       (thunk-env obj)))
        (else obj)))
```

```
(define (actual-value exp env)
  (force-it (l-eval exp env)))
```



thunks: memoizing implementation

```
(define (evaluated-thunk? obj)
  (tagged-list? obj 'evaluated-thunk))

(define (thunk-value evaluated-thunk)
  (cadr evaluated-thunk))

(define (force-it obj)
  (cond ((thunk? obj)
        (let ((result (actual-value (thunk-exp obj)
                                     (thunk-env obj))))
          (set-car! obj 'evaluated-thunk)
          (set-car! (cdr obj) result)
          (set-cdr! (cdr obj) '())
          result))
        ((evaluated-thunk? obj) (thunk-value obj))
        (else obj))))
```

controlling argument evaluation: new syntax

(lambda (a (b lazy) (c lazy-memo))

↓
eval before
proc applic

↓
delayed;
re-evaluated
each time
needed

thunk

↓
delayed;
evaluated first time needed,
value saved

thunk-memo