

Lecture 12:

Testing Distributions

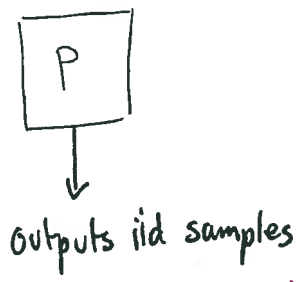
- Uniformity

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Turning to a new model:

prob dists

Probability distributions - get samples of distribution



Domain  $D$ ,  $|D|=n$  ← known  
 $p_i = \Pr[p \text{ outputs } i]$  ← unknown

← this is all we can learn from

Examples:

Lottery data

Shopping choices

experimental outcomes

⋮

What do we want to know?

is it uniform? eg. lottery

is it high entropy?

large support? (many distinct elements have  $>0$  probability)

is it monotone increasing, k-modal, monotone hazard rate...?

how can we do it?

$\chi^2$  test

plug in estimate

learn distribution, Maximum likelihood estimates

Goal: sample complexity **SUBLINEAR** in  $n$

## Testing Uniformity

The goal:

Uniform dist on  $D$

- if  $P \equiv U_D$  then tester outputs PASS  $\leftarrow$  with prob  $\geq 3/4$
- if  $\underbrace{\text{dist}(P, U_D)} > \epsilon$  then tester outputs FAIL

which measure of distance?

$l_1, l_2, \text{KL-divergence, Earth mover, Jensen-Shannon}$

$\uparrow \uparrow$   
today's focus

### Distances

$l_1$ -distance :  $\|p-q\|_1 = \sum_{i \in D} |p_i - q_i|$

$l_2$ -distance :  $\|p-q\|_2 = \sqrt{\sum_{i \in D} (p_i - q_i)^2}$

$\|p-q\|_2 \leq \|p-q\|_1 \leq n^{1/2} \|p-q\|_2$

examples:

①  $p = (1, 0, 0, \dots, 0)$

$q = (\frac{1}{n}, \frac{1}{n}, \dots, \frac{1}{n})$



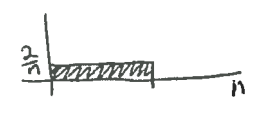
$l_1$  distance:  
 $\|p-q\|_1 = (\frac{n-1}{n}) + (n-1) \cdot \frac{1}{n} \approx 2$

$l_2$  distance:  
 $\|p-q\|_2^2 = (1 - \frac{1}{n})^2 + (n-1)(\frac{1}{n})^2 \approx 1$

②

$p = (\frac{2}{n}, \frac{2}{n}, \dots, \frac{2}{n}, 0, 0, \dots, 0)$

$q = (0, 0, \dots, 0, \frac{2}{n}, \frac{2}{n}, \dots, \frac{2}{n})$



$l_1$  distance:  
 $\|p-q\|_1 = n \cdot (\frac{2}{n}) = 2$

$l_2$  distance:  $\|p-q\|_2^2 = n \cdot (\frac{2}{n})^2 = \frac{4}{n}$

$\|p-q\|_2 = \frac{2}{\sqrt{n}}$

# "Plug-in" Estimate:

Algorithm:

- take  $m$  samples from  $p$

- estimate  $p(x) \forall x$  via

$$\hat{p}(x) = \frac{\# \text{ times } x \text{ occurs in sample}}{m}$$

- if  $\sum_x |\hat{p}(x) - \frac{1}{n}| > \epsilon$  reject  
else accept.

Analysis: (better analyses exist)

$$\forall x, |\hat{p}(x) - p(x)| < \frac{\epsilon}{n} \Rightarrow \|\hat{p} - p\|_1 < \epsilon$$

so, if  $p = U_n$   
then  $p$  passes

pick  $m$  st.  $\forall x, |\hat{p}(x) - p(x)| < \frac{\epsilon}{n}$   
by  $\Delta \neq$ , if  $\|p - \hat{p}\|_1 < \epsilon + \|\hat{p} - U_n\|_1 < \epsilon$   
then  $\|p - U_n\|_1 < 2\epsilon$ .

so, if  $\|p - U_n\|_1 > 2\epsilon$   
this test is likely to Fail

how many samples?  $\Omega(\frac{n}{\epsilon})$  maybe even worse ...

$\Theta(n)$ ? Can we do better?

for each  $x$ , need to see it at least once in order to give non zero estimate.

Better analysis:

Claim  $E[\|\hat{p}-p\|_1] \leq \sqrt{\frac{n}{m}}$

Pf

$$E[\|\hat{p}-p\|_1] = \sum_x E[|\hat{p}(x)-p(x)|] \leftarrow \text{note: } E[\hat{p}(x)] = \frac{1}{m} E\left[\sum_{i=1}^m \mathbb{1}_{i^{\text{th}} \text{ sample is } x}\right]$$

$$= \frac{1}{m} \sum_{i=1}^m E[\mathbb{1}_{i^{\text{th}} \text{ sample is } x}] = \frac{m \cdot p(x)}{m} = p(x)$$

$$\leq \sum_x \sqrt{E[(\hat{p}(x)-p(x))^2]} \leftarrow \text{Jensen's } \neq$$

$$= \sum_x \sqrt{\text{Var}(\hat{p}(x))}$$

$$\leq \sum_x \sqrt{\frac{p(x)}{m}}$$

$$\leq \frac{1}{\sqrt{m}} \cdot \sqrt{n} \leftarrow \text{since } \max_{p \in \text{prob dist over domain of size } n} \sum \sqrt{p(x)} \text{ is } \sqrt{n}$$

$$\text{Var}(\hat{p}(x)) = \frac{1}{m^2} m p(x)(1-p(x)) \leq \frac{p(x)}{m}$$

So picking  $m = \Omega\left(\frac{n}{\epsilon^2}\right)$  gives

$$E[\|\hat{p}-p\|_1] \leq \frac{\epsilon}{2}$$

by Markov's  $\neq$ : with prob  $1-\frac{1}{2}$ ,  $\|\hat{p}-p\|_1 \leq \epsilon$

Note, this says can "learn" (approximate) any dist wr.t.  $L_1$  distance in  $\Theta(n/\epsilon^2)$  samples

### L<sub>2</sub> - Distance (squared):

$$\begin{aligned} \|p - u\|_2^2 &= \sum_{i \in [n]} (p_i - \frac{1}{n})^2 \\ &= \sum p_i^2 - \underbrace{\frac{2}{n} \sum p_i}_{=1} + \underbrace{\sum (\frac{1}{n})^2}_{=\frac{1}{n}} \\ &= \sum p_i^2 - \frac{1}{n} \end{aligned}$$

Collision probability of  $p$ :

$$\|p\|_2^2 \equiv \Pr_{s, t \in p} [s = t] = \sum p_i^2$$

for  $p = u$ ,  $\|p\|_2^2 = \frac{1}{n}$

for  $p \neq u$ ,  $\|p\|_2^2 > \frac{1}{n}$

$$= \|p\|_2^2 - \|u\|_2^2$$

we can estimate this

we know this since we know  $n$

### Algorithm

1. take  $s$  samples from  $p$  ① how many samples?
2. let  $\hat{c} \leftarrow$  estimate of  $\|p\|_2^2$  from sample ② how?
3. if  $\hat{c} < \frac{1}{n} + \delta$  pass ③ what should  $\delta$  be?  
     else fail

First:  
How to estimate  $\|p\|_2^2$ ?

Naive idea:

take two raw samples:

$$X_i \leftarrow \begin{cases} 1 & \text{if samples are equal} \\ 0 & \text{o.w} \end{cases}$$

" gives  $\theta(k)$  samples of collision probability from  $k$  samples of  $p$ "

Better idea: recycle - use all pairs in sample

" gives  $\theta(k^2)$  samples of collision probability from  $k$  samples of  $p$ "

Estimate by recycling:

• Take  $s$  samples from  $p$ :  $X_1, \dots, X_s$

• for each  $1 \leq i < j \leq s$

$$b_{ij} \leftarrow \begin{cases} 1 & \text{if } X_i = X_j \\ 0 & \text{if } X_i \neq X_j \end{cases}$$

}  $b_{ij}$ 's not independent so can't use Chernoff

• Output  $\hat{c} \leftarrow \frac{\sum_{i < j} b_{ij}}{\binom{s}{2}}$

Analysis:  $E[\hat{c}] = \frac{1}{\binom{s}{2}} \cdot \binom{s}{2} \cdot E[b_{ij}] = \|p\|_2^2$



How well do we need to estimate  $\|p\|_2^2$ ?

Assumption  $\star$ :  $|\hat{C} - \|p\|_2^2| < \Delta$   
 will take enough samples so that this holds with prob  $\geq 3/4$   
 this is our parameter that determines whether our approximation is good. Spoiler: will set  $\Delta = \frac{\epsilon^2}{2}$

What happens if  $\star$  holds with  $\Delta = \frac{\epsilon^2}{2}$ ?

Correct behavior!

• if  $p = U_{[n]}$  then  $\hat{C} \leq \|U_{[n]}\|_2^2 + \Delta = \frac{1}{n} + \frac{\epsilon^2}{2}$

so test will PASS

• if  $\|p - U_{[n]}\|_2 > \epsilon$  then  $\|p - U_{[n]}\|_2^2 > \epsilon^2$

but 
$$\|p\|_2^2 = \|p - U_{[n]}\|_2^2 + \frac{1}{n}$$

$$> \epsilon^2 + \frac{1}{n}$$

← see p.6

+ 
$$\hat{C} > \|p\|_2^2 - \Delta$$

$$\geq \epsilon^2 + \frac{1}{n} - \Delta = \epsilon^2 + \frac{1}{n} - \frac{\epsilon^2}{2} = \frac{\epsilon^2}{2} + \frac{1}{n}$$

←  $\star$

so test will FAIL

Remaining Question:

How many samples do we need to estimate  $\hat{C}$  to within  $\Delta$ ?

Analysis

$$E [b_{ij}] = \Pr [b_{ij} = 1] \\ = \|p\|_2^2$$

$$E [\hat{c}] = \frac{1}{\binom{s}{2}} \binom{s}{2} E [b_{ij}] = \|p\|_2^2$$

$$\Pr [ |\hat{c} - \|p\|_2^2| > \rho ] \leq \frac{\text{Var} [\hat{c}]}{\rho^2}$$

Chebyshev  $\neq$

Fact  $\text{Var} [aX] = a^2 \text{Var} [X]$

$$\text{So } \text{Var} [\hat{c}] = \text{Var} \left[ \frac{1}{\binom{s}{2}} \cdot \sum_{i < j} b_{ij} \right] \\ = \frac{1}{\binom{s}{2}^2} \text{Var} \left[ \sum_{i < j} b_{ij} \right]$$

Lemma  $\text{Var} [\sum b_{ij}] \leq 4 \left( \binom{s}{2} \|p\|_2^2 \right)^{3/2}$

Fact  $\Rightarrow$   
 $\text{Var} [\hat{c}] \leq \frac{4 \cdot \left( \binom{s}{2} \|p\|_2^2 \right)^{3/2}}{\binom{s}{2}^2} \leq \theta \left( \|p\|_2^3 / s \right)$

Why? (proof...)

def.  $\bar{b}_{ij} = b_{ij} - E [b_{ij}]$

← trick - will rewrite variance as  $E [\bar{b}_{ij}^2]$ .

so  $E [\bar{b}_{ij}] = 0$

Also  $\because E [\bar{b}_{ij} \bar{b}_{kl}] \leq E [b_{ij} b_{kl}]$

verify at home? (or trust...)

- $\left( \sum p(x)^3 \right)^{1/3} \leq \left( \sum p(x)^2 \right)^{1/2}$
- $s^2 \leq 3 \binom{s}{2}$
- $\binom{s}{3} \leq s^3 / 6$

e.g.  $(a^3 + b^3)^2 \leq (a^2 + b^2)^3$   
 $a^6 + 2a^3b^3 + b^6 \leq a^6 + b^6 + 3a^4b^2 + 3a^2b^4$

So

$$\text{Var} \left[ \sum_{i < j} \bar{\delta}_{ij} \right] = E \left[ \left( \sum_{i < j} \bar{\delta}_{ij} - E \left[ \sum_{i < j} \bar{\delta}_{ij} \right] \right)^2 \right]$$

$$= E \left[ \left( \sum_{i < j} \bar{\delta}_{ij} \right)^2 \right]$$

$$\begin{aligned} &+ \sum_{i < j} \bar{\delta}_{ij} \bar{\delta}_{il} \quad (5) \\ &+ \sum_{i < j} \bar{\delta}_{ij} \bar{\delta}_{ki} \quad (6) \end{aligned}$$

$$= E \left[ \underbrace{\sum_{i < j} \bar{\delta}_{ij}^2}_{(1)} + \underbrace{\sum_{\substack{i < j \\ k < l \\ i, j, k, l \text{ distinct}}} \bar{\delta}_{ij} \bar{\delta}_{kl}}_{(2)} + \underbrace{\sum_{\substack{i < j \\ k < l \\ i, j, l \text{ distinct}}} \bar{\delta}_{ij} \bar{\delta}_{kl}}_{(3)} + \underbrace{\sum_{\substack{i < j \\ k < l \\ i, j, k \text{ distinct}}} \bar{\delta}_{ij} \bar{\delta}_{kl}}_{(4)} \right]$$

(1)  $E \left[ \sum_{i < j} \bar{\delta}_{ij}^2 \right] \leq E \left[ \sum_{i < j} \delta_{ij}^2 \right] = \binom{s}{2} \|p\|_2^2$

$E[\delta_{ij}] = E[\delta_{ij}^2]$  since  $\delta_{ij}$  is indicator var

(2) independent

(2)  $E \left[ \sum_{\substack{i < j \\ k < l \\ \text{all 4 distinct}}} \bar{\delta}_{ij} \bar{\delta}_{kl} \right] \leq \sum E[\bar{\delta}_{ij}] E[\bar{\delta}_{kl}] = 0$

(3)  $E \left[ \sum_{\substack{i, j, l \\ \text{distinct}}} \bar{\delta}_{ij} \bar{\delta}_{il} \right] \leq E \left[ \sum_{\substack{i, j, l \\ \text{distinct}}} \delta_{ij} \cdot \delta_{il} \right] = \sum_{\substack{i, j, l \\ \text{distinct}}} \text{pr}[X_i = X_j = X_l]$

$\leq \binom{s}{3} \sum_x p(x)^3$  expected # 3-way collisions

$$\frac{1}{6} \binom{s}{3}^{3/2} < \frac{\left( 3 \binom{s}{2} \right)^{3/2}}{6} = \frac{\sqrt{3}}{2} \binom{s}{2}^{3/2}$$

$\leq \frac{s^3}{6} \left( \sum_x p(x)^2 \right)^{3/2}$   
 $\leq \frac{\sqrt{3}}{2} \binom{s}{2}^{3/2} (\|p\|_2^2)^{3/2}$  by the facts

⑪  
p.d.

④ same as 3  
⑤  
⑥

In total:

$$\begin{aligned}\text{Var} \left[ \sum_{i < j} \delta_{ij} \right] &\leq \text{Var} \left[ \sum_{i < j} \bar{\delta}_{ij} \right] \\ &\leq \binom{s}{2} \|p\|_2^2 + 0 + 4 \cdot \frac{\sqrt{3}}{2} \left( \binom{s}{2} \|p\|_2^2 \right)^{3/2} \\ &\leq 4 \left[ \binom{s}{2} \|p\|_2^2 \right]^{3/2}\end{aligned}$$



Putting lemma into Chebyshev:

(12).  
p.d

use  $p = \frac{\epsilon^2}{2}$

$$\Pr[|\hat{c} - \|p\|_2^2| > \frac{\epsilon^2}{2}] \leq \frac{\text{Var}[\hat{c}]}{\epsilon^4} \cdot 4$$

Recall this  
comes from  
cons. in proof

$$\leq \frac{4 \left[ \binom{s}{2} \|p\|_2^2 \right]^{3/2}}{\binom{s}{2}^2 \epsilon^4} \cdot 4 \leq \frac{32}{\epsilon^4} \cdot \frac{1}{s} \cdot \|p\|_2^3$$

note  $\frac{1}{\binom{s}{2}^2} \leq \frac{1}{\sqrt{\frac{s^2}{2}}} \leq \frac{2}{s}$

So pick  $s \geq \left( \frac{1}{\epsilon^4} \right)$

also want  
this to  
be  $\leq 1$

Note: Can get better bound

1) Testing closeness to any known distribution — reduce to uniform case!

2) lower bound

How to estimate  $\|p-u\|_1$ ?

1)  $\|p-u\|_1 = 0 \Leftrightarrow \|p-u\|_2^2 = 0 \Leftrightarrow \|p\|_2^2 = \frac{1}{n}$

2) if  $\|p-u\|_1 > \epsilon \Rightarrow \|p-u\|_2 > \frac{\epsilon}{\sqrt{n}}$

$\Rightarrow \|p-u\|_2^2 > \frac{\epsilon^2}{n}$

$\Rightarrow \|p\|_2^2 \geq \frac{1}{n} + \frac{\epsilon^2}{n}$

either additive estimate with error  $\leq \frac{\epsilon^2}{2n}$

or mult error  $\leq (1 \pm \frac{\epsilon^2}{3})$

suffices

would have this if have additive error  $\leq \frac{\epsilon^2}{3n} \cdot \|p\|_2^2$

to get additive error  $\leq \frac{\epsilon^2}{3n} \|p\|_2^2$

suffices to have

$s \geq \frac{\text{const} \cdot \sqrt{n}}{\epsilon^2}$

samples

since  $\Pr[|\hat{c} - \|p\|_2^2| \geq \gamma \|p\|_2^2] \leq \frac{k \cdot \|p\|_2^3}{s \cdot \gamma^2 (\|p\|_2^2)^2} \leq \frac{k}{s \cdot \gamma^2 \|p\|_2}$

[note  $\|p\|_2^2 > \frac{1}{n}$  so  $\|p\|_2 > \frac{1}{\sqrt{n}}$  so  $\frac{1}{\|p\|_2} < \sqrt{n}$ ]

$\leq \frac{k \cdot \sqrt{n}}{s \cdot \gamma^2}$

[note: we need  $\gamma \approx \frac{\epsilon^2}{3}$ ]

so picking  $s \gg \frac{\sqrt{n}}{\epsilon^4}$  will give small probability of error  $\Rightarrow$

$\approx \frac{k \cdot \sqrt{n}}{s} \cdot \frac{1}{\epsilon^4}$