Lecture 9:

Lower bounds for testing

Δ- freeness:

superpoly dependence on ε

is required!

A lower bound for testing Δ-freeness

In a previous lecture:

- saw property test for Δ-freeness
  - const time in terms of n
  - $\underbrace{\text{dependence on } \varepsilon \text{ horrible}}_{\text{is this required?}}$ - tower of 2's

Today:

- answer this question partially (for 1-sided testers)
- When testing H-freeness property,

interesting characterization of bipartiteness
  - if H bipartite, poly(1/ε) is enough
  - if H not bipartite no poly(1/ε) suffices

(We'll actually prove special case of H = Δ only)

Thm (adj matrix model)

∃ const c s.t. any 1-sided tester for whether graph G is Δ-free needs $\geq \left(\frac{c}{\varepsilon}\right)^{c \log 1/\varepsilon}$ queries.

Main Tools:

(1) Goldreich-Trevisan Thm: (homework)

Adj matrix model
Property P
Tester T with $q(n,\varepsilon)$ queries

$\Rightarrow$ Tester T': "Natural Tester"

pick $q(n,\varepsilon)$ nodes
query submatrix
decide

$\Big\}\ O(q^2)$ queries

Consequences:

- l.b. for natural tester of $\Omega(q')$
  $\Rightarrow$ l.b. for any tester of $\Omega(\sqrt{q'})$
- note, reduction preserves 1-sidedness,
  so l.b. implication does too.

Main tools (cont.):

(2) Additive Number theory lemma

#TheoryLemma $\forall\ m,\quad \exists\ X \subset M = \{1,2,\dots,m\}$

of size $\geq \frac{m}{e^{10\sqrt{\log m}}}$

with no non trivial soln to $x_1 + x_2 = 2x_3$

i.e. $x_1 = x_2 = x_3$ is the trivial soln.

Will use to construct graphs st.

- far from Δ-free
- natural algorithm needs $\Omega\left(\left(\frac{c}{\varepsilon}\right)^{c \log 1/\varepsilon}\right)$ queries

examples

Bad X : $\{1,2,3\}$

$\{5,9,13\}$

Good X ? $\{1,2,4,5,\cancel{7},\cancel{8},\cancel{9},\cancel{11},10,\dots\}$ ← how big??

$\{1,2,4,8,16,32,\dots\}$ ← only size log m

Proof of lemma

- let d be integer (later, set to $e^{10\sqrt{\log m}}$)

$k \leftarrow \left\lfloor \frac{\log m}{\log d} \right\rfloor - 1$ (so $k \approx \frac{\log m}{10\sqrt{\log m}} \approx \frac{\sqrt{\log m}}{10}$)

Proof of lemma (cont.)

the two constraints will be used later in a nice way

define $X_B = \{ \sum_{i=0}^{k} x_i d^i \mid \underbrace{x_i < \frac{d}{2}}_{①} \text{ for } 0 \le i \le k$

$\& \ \underbrace{\sum_{i=0}^{k} x_i^2 = B}_{②} \}$

view each $x \in M$ as represented in base d

where $x = (\underbrace{x_0 \dots x_k}_{\text{"digits" of } x})$

Claim $X_B \subseteq M$

Why? largest number in X_B

$\le d^{k+1} \le d^{(\lfloor \frac{\log m}{\log d} \rfloor - 1) + 1} \le d^{\log_d m} = m^{\log_d d} = m$

What is B? Pick s.t. $|X_B|$ maximized

Why the constraints?

① x_i's $< \frac{d}{2}$ $\Rightarrow$ summing pairs of elements in X_B doesn't generate a carry in <u>any</u> location!

we'll see why this is useful soon

(2) will use (along with (1)) to show that X_B is "sum-free"

Claim X_B is "sum free" ie. $\nexists\, x,y,z \in X_B$ s.t.
$x+y=2z$

Pf of claim

for $x,y,z \in X_B$

$$x+y=2z \iff \sum_{i=0}^{k} x_i d^i + \sum_{i=0}^{k} y_i d^i = 2\sum_{i=0}^{k} z_i d^i$$

$$\iff \left.\begin{array}{l} x_0+y_0=2z_0 \\ x_1+y_1=2z_1 \\ \vdots \\ x_k+y_k=2z_k \end{array}\right\} \text{since no carries}$$

Note $\forall i \;\; x_i+y_i=2z_i \Rightarrow \forall i \;\; x_i^2+y_i^2 \geq 2z_i^2$
with equality only if $x_i=y_i=z_i$

Why? $f(a)=a^2$ is convex

use Jensens ≠: $\dfrac{\sum_{i=1}^{n} f(a_i)}{n} \geq f\left(\dfrac{\sum_{i=1}^{n} a_i}{n}\right)$ with equality only if a_i's are all =

$$\Rightarrow \frac{x_i^2+y_i^2}{2} \geq \left(\frac{2z_i}{2}\right)^2 = z_i^2 \text{ + equal only if}$$

$x_i=y_i=2z_i$

■ (proof of note)

finishing proof of claim:

if x, y, z s.t. not$(x=y=z)$

then $\exists\ i$ s.t. not$(x_i = y_i = z_i)$

then note $\Rightarrow$ $x_i^2 + y_i^2 > 2z_i^2$

+ for all other j, $x_j^2 + y_j^2 \geq 2z_j^2$

but then:

$$\underbrace{\sum x_i^2}_{=B} + \underbrace{\sum y_i^2}_{=B} > \sum 2z_i^2 = \underbrace{2\underbrace{\sum z_i^2}_{=B}}_{2B}$$

$\rightarrow\!\leftarrow$

but how do we know that X_B is big?

- $B \leq (k+1)\left(\frac{d}{2}\right)^2 < kd^2$

  ↑ bound on digits of B

- $$\left|\bigcup_B X_B\right| \geq \left(\frac{d}{2}\right)^{k+1} > \left(\frac{d}{2}\right)^k$$

  $$\| \\ \sum_B |X_B|$$

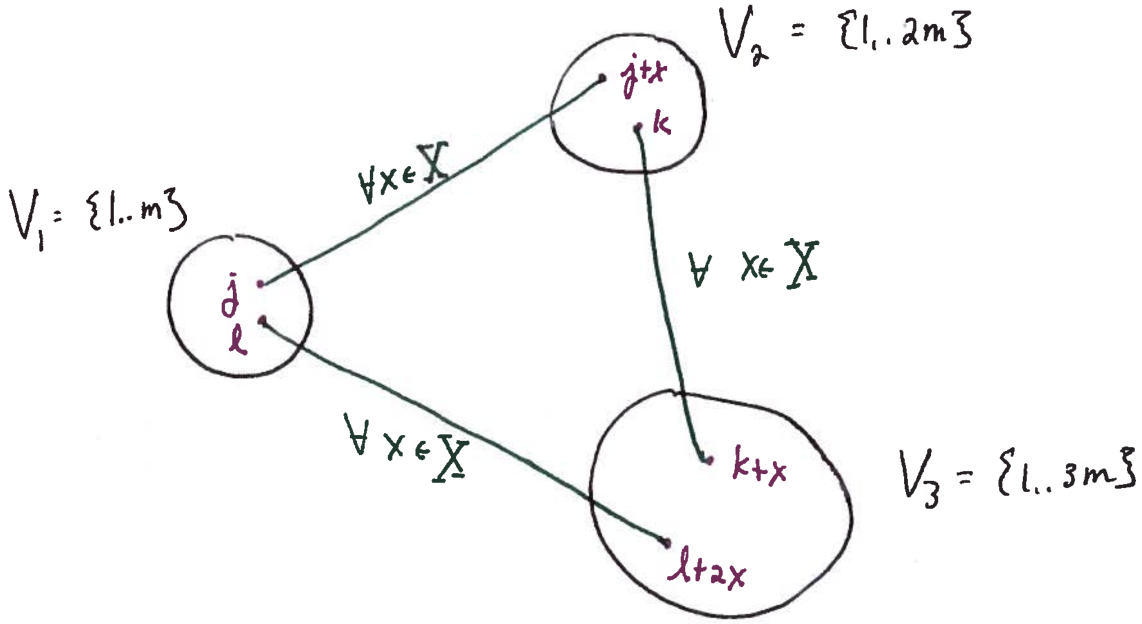
- $\exists\ B$ s.t. $|X_B| \geq \dfrac{\left(\frac{d}{2}\right)^k}{kd^2}$

- Use settings of d, k, get $|X_B| \geq \dfrac{m}{e^{10\sqrt{\log m}}}$

Not enough! need another idea, but won't do it here ∎

<u>Proof of Thm</u> (prop testing bound)

given sum-free $X \subseteq \{1..m\}$
construct a graph:



- will abuse notation:

  node should be (i, j)

  i in $\{1,2,3\}$, $j \in \{1..c \cdot m\}$

  will drop i if easy to see from context

- #nodes $= 6m$ so $m = \Theta(n)$
- #edges $= \Theta(m \cdot |X|) = \Theta(n^2 / e^{10\sqrt{\lg n}})$ ← not exactly dense

l.b. on Δ-free . . .

• # cycles :

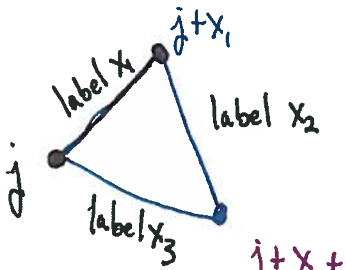
intended Δ's : $j, \; j+x, \; j+2x$

$\#$ intended Δ's is $m|X| = \Theta\left(n^2 / e^{10\sqrt{\log n}}\right)$

nonintended Δ's :

• no edges internal to V_1, V_2 or V_3

$\therefore$ any Δ has
$u \in V_1$
$v \in V_2$
$w \in V_3$



$$\left.\begin{array}{l} j+x_1+x_2 \\ = j+2x_3 \end{array}\right\} \Rightarrow x_1+x_2 = 2x_3$$

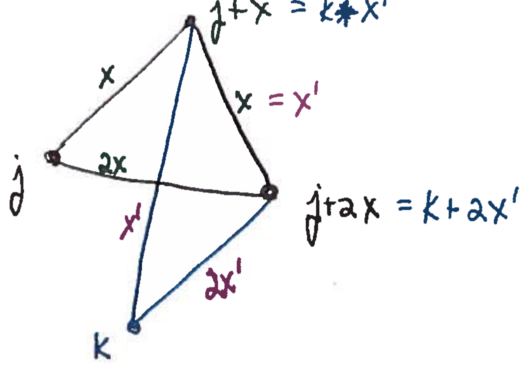
$$\Rightarrow \underbrace{x_1 = x_2 = x_3}_{\text{but these are intended!}} \quad \text{since } X \text{ is sum-free}$$

$\therefore$ no nonintended Δ's

- # disjoint cycles:

all intended Δ's are disjoint (share no edges at all)

suppose not:



since $x = x'$, $k = j$ $\rightarrow\leftarrow$

- distance to Δ-free:

must remove ≥ 1 edge from each Δ.

$\Downarrow$

"absolute" distance from Δ-free $= \Theta(\# \Delta\text{'s})$

$$= \Theta\left(\frac{n^2}{e^{10\sqrt{\log n}}}\right)$$

$$= \Theta(m/|X|)$$

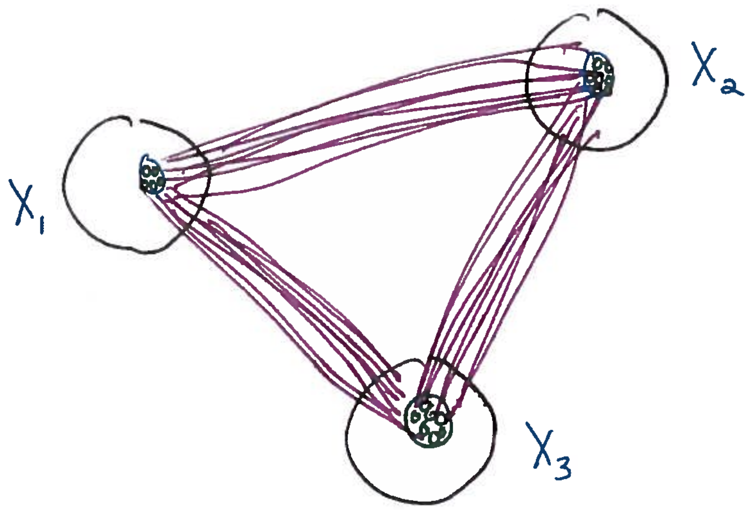
Problem need $\Omega(\varepsilon n^2)$ distance

Idea for fix

s-blow-up $\quad G \to G^{(s)}$

vertex in G $\to$ size s independent set in $G^{(s)}$

edge in G $\to$ complete bipartite graph in $G^{(s)}$



Note: Δ in G $\Rightarrow$ s^3 Δ's in $G^{(s)}$ $\Rightarrow$ likely to find one!

nodes in $G^{(s)}$ $\sim m \cdot s$ (actually $6ms$)

edges " " $\sim m|X| \cdot s^2$

triangles " " $\sim m|X| s^3$

Lemma dist of $G^{(s)}$ from Δ-free

$\geq$ # edge disjoint Δ's

$\geq m|X| s^2$

Proof show each triangle in G $\Rightarrow$ s^2 disjoint Δ's in $G^{(s)}$

l.b. on
Δ-free

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Given ε, pick m to be largest int

satisfying

$$\varepsilon \leq \frac{1}{e^{10\sqrt{\log m}}}$$

this m satisfies

$$m \geq \left(\frac{c}{\varepsilon}\right)^{c \log c/\varepsilon}$$

Pick $s = \left\lfloor \frac{n}{6m} \right\rfloor \approx n \left(\frac{\varepsilon}{c}\right)^{c \log c/\varepsilon}$

$\Rightarrow$ #edges $\sim$ distance $\sim \varepsilon n^2$ (since $\approx \frac{m|X|s^2}{m^2 s^2}$ $\leftarrow$ size of adj matrix

$$= \frac{|X|}{m} \geq \frac{1}{e^{10\sqrt{\log m}}} \geq \varepsilon)$$

$$|X| = \frac{m}{e^{10\sqrt{\log m}}}$$

triangles $\sim \left(\frac{\varepsilon}{c'}\right)^{c' \log c'/\varepsilon} n^3$

$$m|X| \cdot s^3 = \frac{m^2}{e^{10\sqrt{\log m}}} s^3$$

$$= \frac{1}{\varepsilon}\left(\left(\frac{c}{\varepsilon}\right)^{c \log c/\varepsilon}\right)^2 \cdot \left(\frac{\varepsilon}{c'}\right)^{c' \log c'/\varepsilon} n^3$$

Finally if take sample of size q

$$E[\#\Delta\text{'s in sample}] < \binom{q}{3}\left(\frac{\varepsilon}{c}\right)^{c' \log c'/\varepsilon}$$

$$\ll 1 \quad \text{unless} \quad q > \left(\frac{c'}{\varepsilon}\right)^{c' \log c'/\varepsilon}$$

by Markov's $\neq$ $\Rightarrow$ $\Pr[\text{see } \Delta] \ll 1$

But since 1-sided error,

must find Δ in order to fail $\blacksquare$