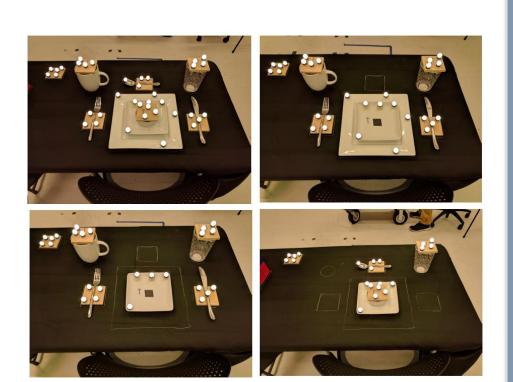


Motivation

Specifications for setting a dinner table:

- Necessity of object placements
- Correct object positions
- No collisions
- Placement orders

Not Markovian in object poses



Expressible in Linear temporal logic (LTL), a flexible specification language used in:

- Synthesis of verifiable controllers^[1]
- Reinforcement learning^[2]
- Goal description in symbolic planning^[3]

Aim: Infer task specifications from demonstrations

Approach: Bayesian specification inference for task specifications as LTL formula

Bayesian Formulation

$$P(\varphi|\boldsymbol{D}) = \frac{P(\varphi)P(\boldsymbol{D}|\varphi)}{\sum_{\varphi \in \boldsymbol{\varphi}} P(\varphi)P(\boldsymbol{D}|\varphi)}$$

- $P(\varphi)$ must have positive support over all relevant formulas.
- $P(\mathbf{D}|\varphi)$ is the likelihood distribution that honors the size principle:
 - Large likelihood for complex formula.
 - Small likelihood for simple formula
 - Number of conjunctions a measure of formula complexity
- Probabilistic programming languages for sampling based inference^[4]

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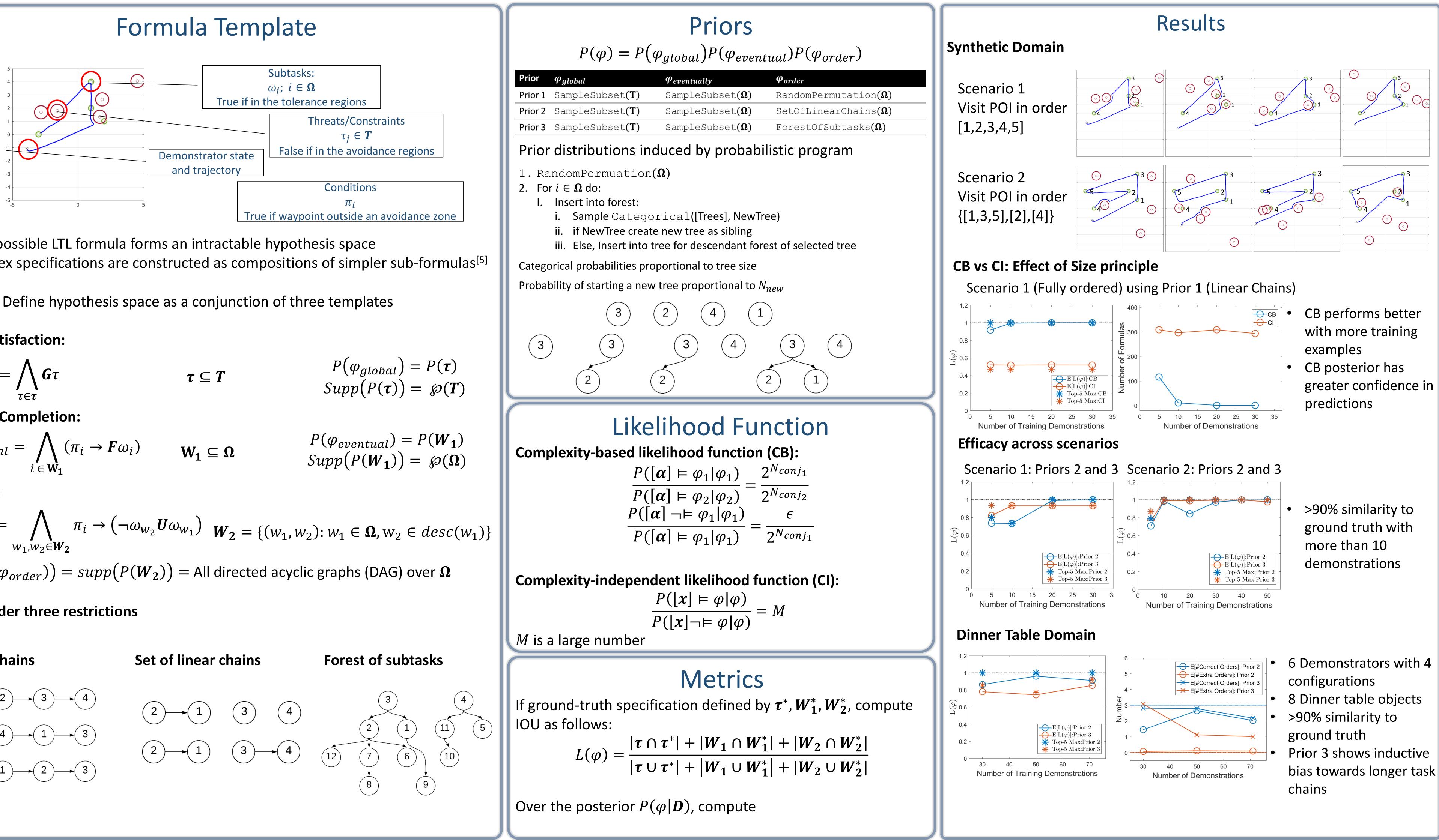
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Bayesian Inference of Temporal Task Specifications from Demonstrations

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- Every possible LTL formula forms an intractable hypothesis space
- Complex specifications are constructed as compositions of simpler sub-formulas^[5]

Key Idea: Define hypothesis space as a conjunction of three templates

Global satisfaction:

$$\varphi_{global} = \bigwedge_{\tau \in \tau} \mathbf{G}\tau \qquad \mathbf{\tau} \subseteq \mathbf{T} \qquad P(\varphi_{global}) = P(\tau) \\ Supp(P(\tau)) = \mathcal{P}(\tau)$$

Eventual Completion:

 $\varphi_{eventual} = / (\pi_i \to F \omega_i)$

Ordering:

$$\varphi_{order} = \bigwedge_{w_1, w_2 \in W_2} \pi_i \to \left(\neg \omega_{w_2} U \omega_{w_1}\right) \quad W_2 = \{(w_1, w_2) : w_1 \in \Omega, w_2 \in des\}$$

 $Supp(P(\varphi_{order})) = supp(P(W_2)) = All directed acyclic graphs (DAG) over <math>\Omega$

We consider three restrictions

Linear Chains $1 \longrightarrow 2 \longrightarrow 3 \longrightarrow 4$ $2 \rightarrow 4 \rightarrow 1 \rightarrow 3$ (4)→→(1)→→(2)→→(3)



- Prior 3 shows inductive bias towards longer task