

Ph.D.

ONE SIZE DOES NOT FIT ALL

Streaming

OLAP

OLTP

Archiving

Log-processing

Web-search

Scan-oriented

Streaming

OLAP

OLTP

Archiving

Log-processing

Web-search

Scan-oriented

Streaming

OLAP

OLTP

Archiving

facebook

Log-processing

Web-search

Scan-oriented

amazon.com[®]

ebay[®]

Streaming

OLAP

OLTP

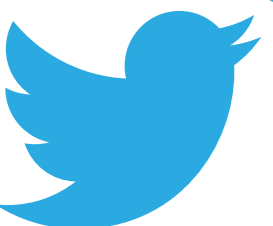
Archiving

facebook

Log-processing

Web-search

Scan-oriented



Streaming

amazon.com[®]

ebay[®]

OLAP

OLTP

Archiving

facebook[®]

Log-processing

Web-search

Scan-oriented

YAHOO![®]

OLTP



OLAP



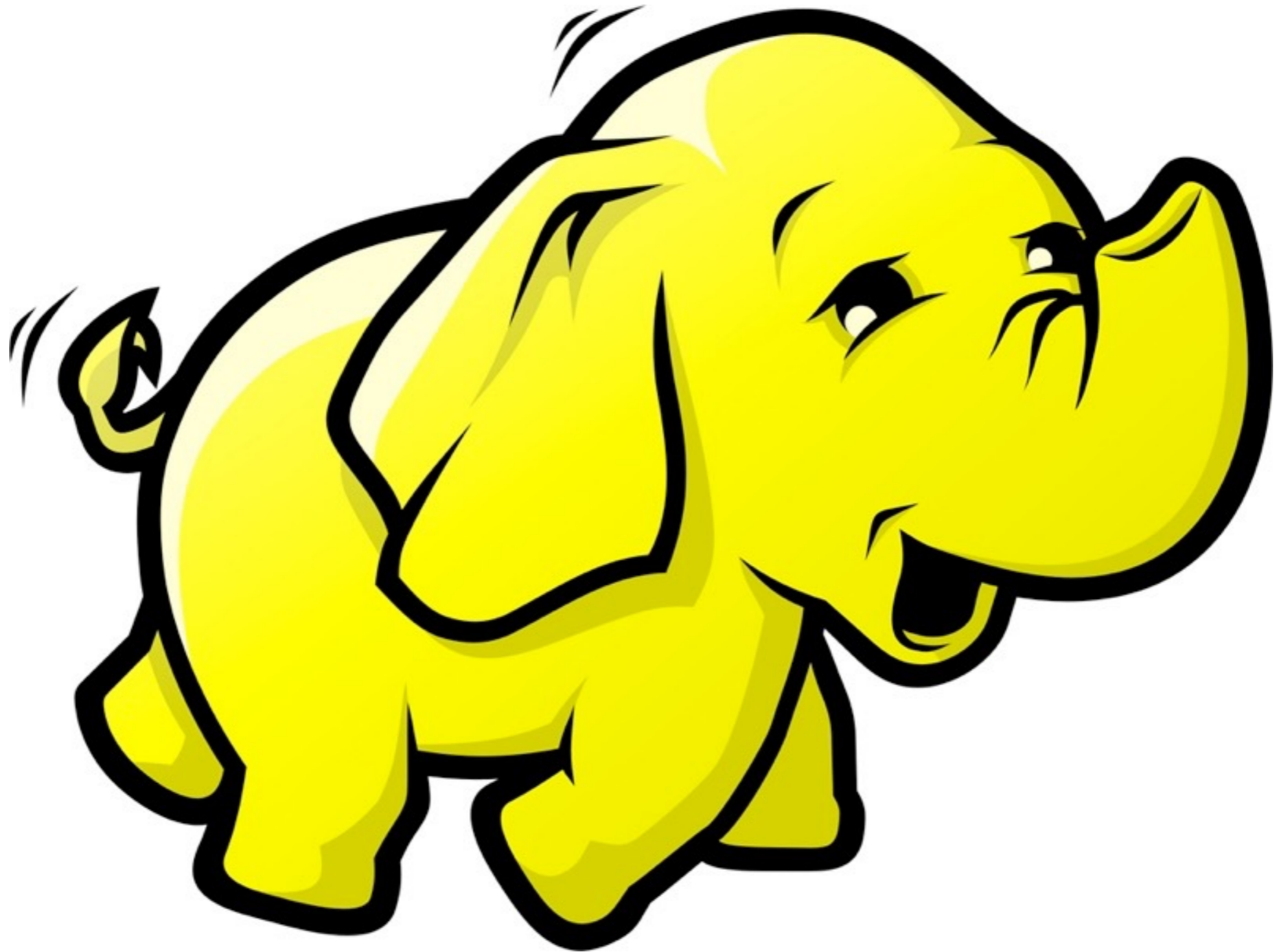
Archive



Streaming



Log-processing



Streaming

OLAP

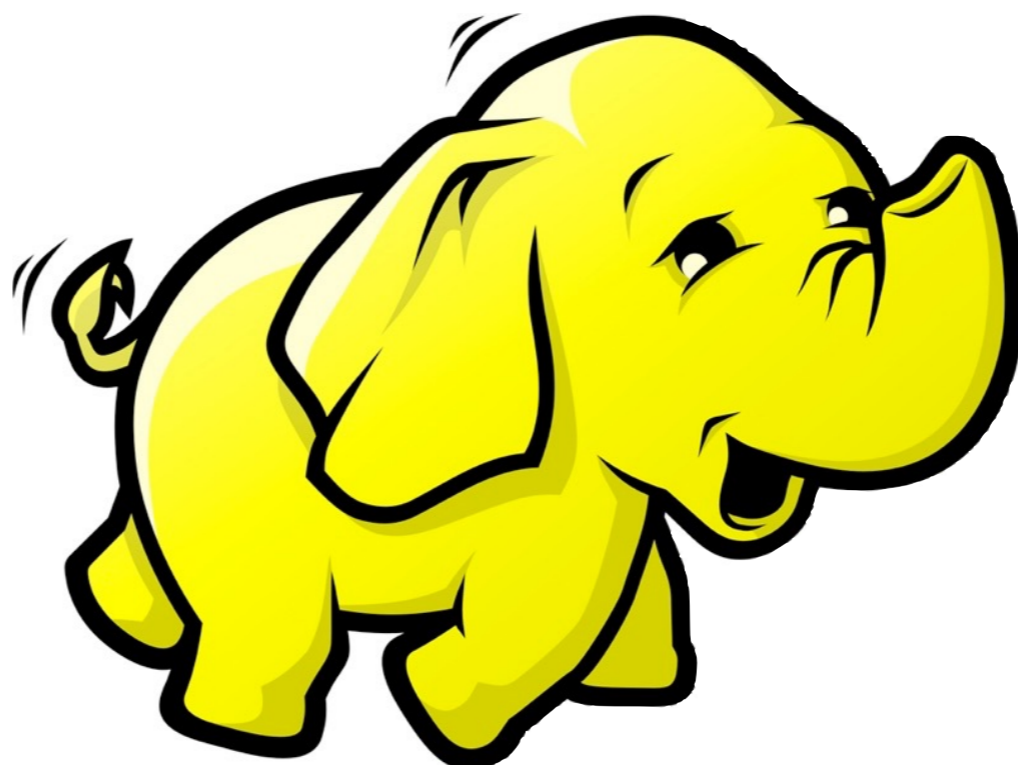
OLTP

Archiving

Log-processing

Web-search

Scan-oriented





Indexes



Column



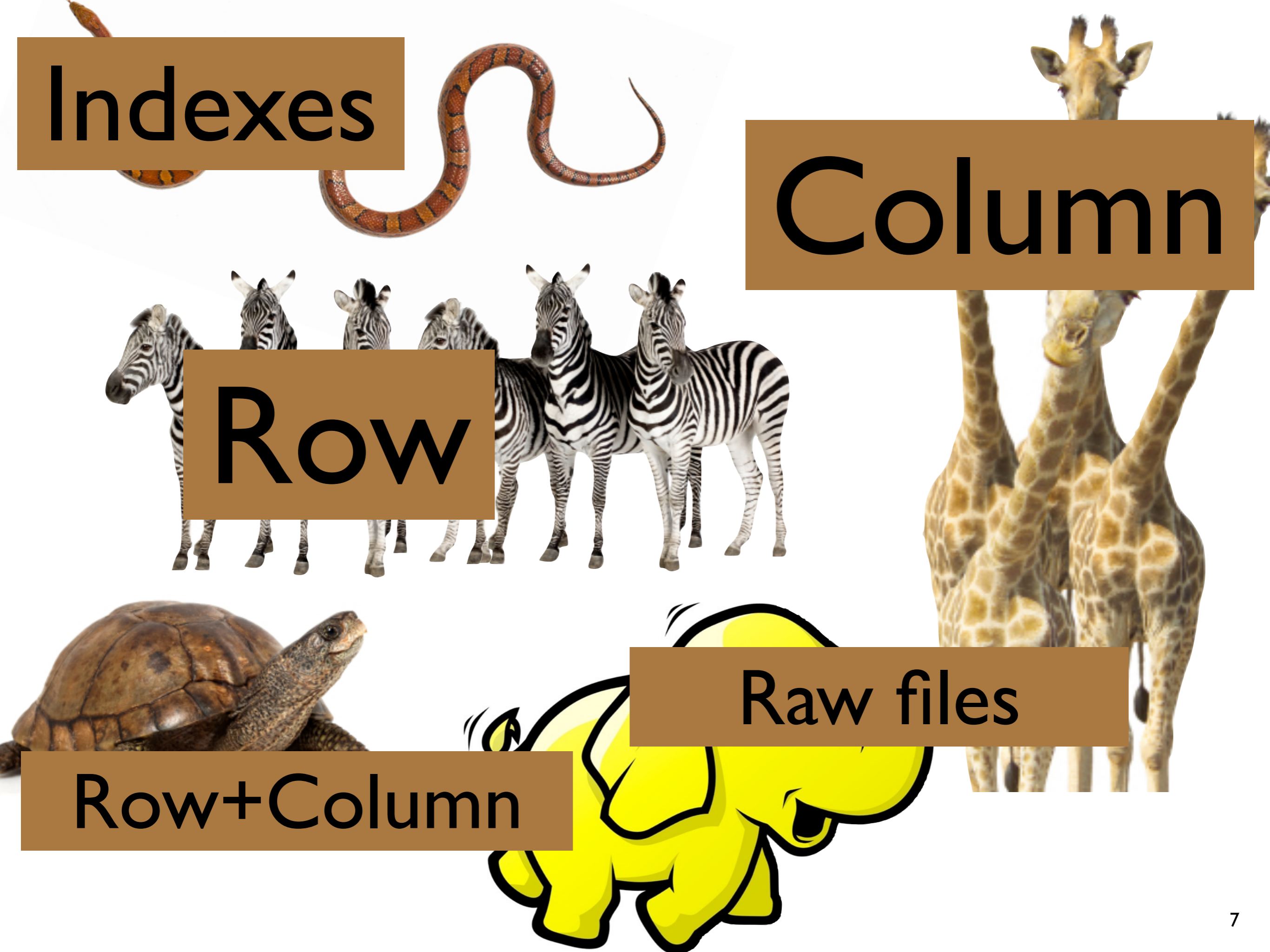
Row



Raw files



Row+Column



Storage Views

1	abc	56	887.9
2	fdg	89	445.35
3	poe	67	234.67
4	lkj	12	385.92
5	yui	17	612.13
6	omg	90	148.9

Storage Views

Log		56	887.9
		89	445.35
		67	234.67
		4	lkj
5	yui	17	612.13
6	omg	90	148.9

Storage Views

Log		56	887.9
		89	445.35
		67	234.67
4	Row		
5	yui	17	612.13
6	omg	90	148.9

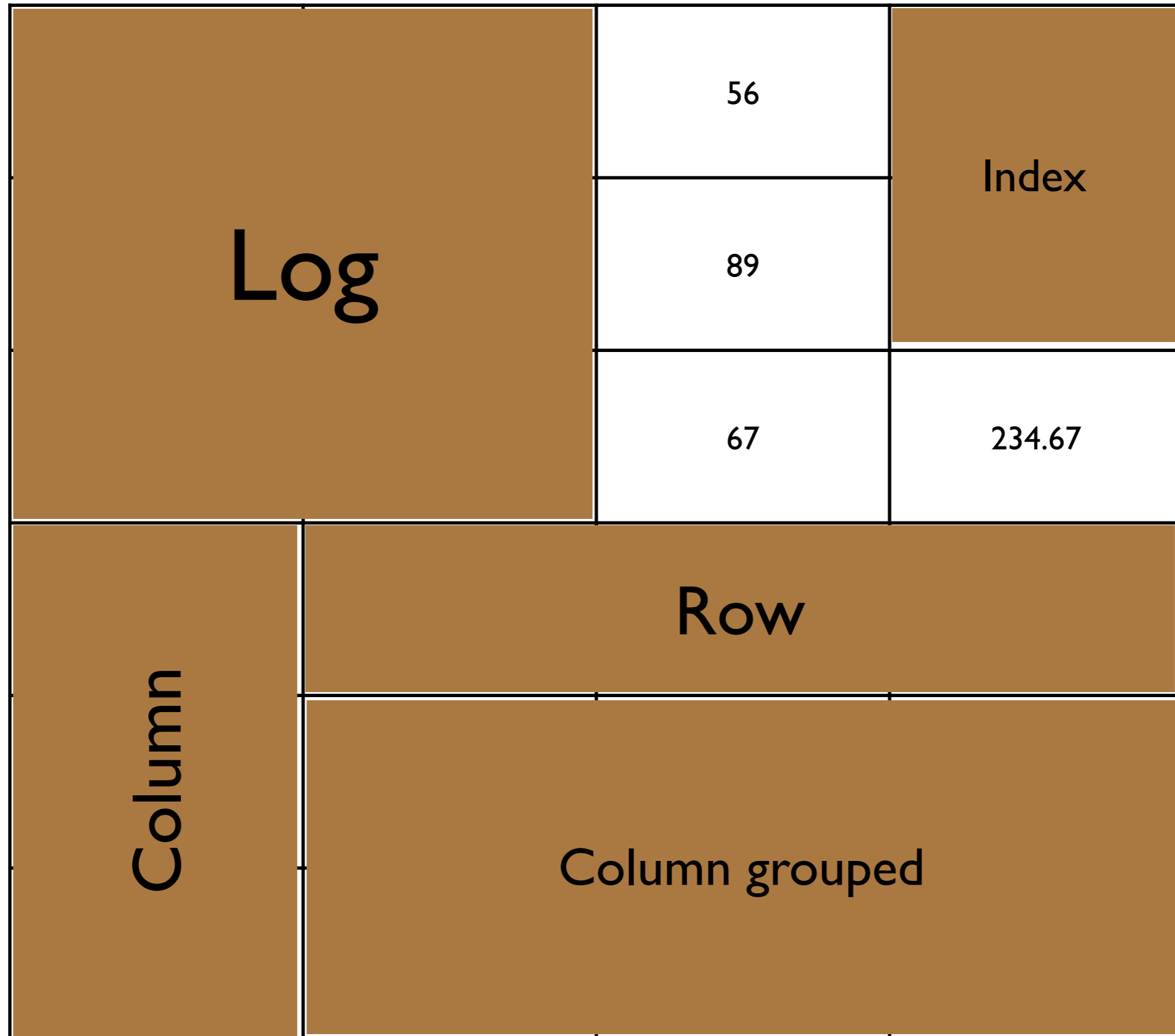
Storage Views

Log		56	887.9
		89	445.35
		67	234.67
Column	Row		
	yui	17	612.13
	omg	90	148.9

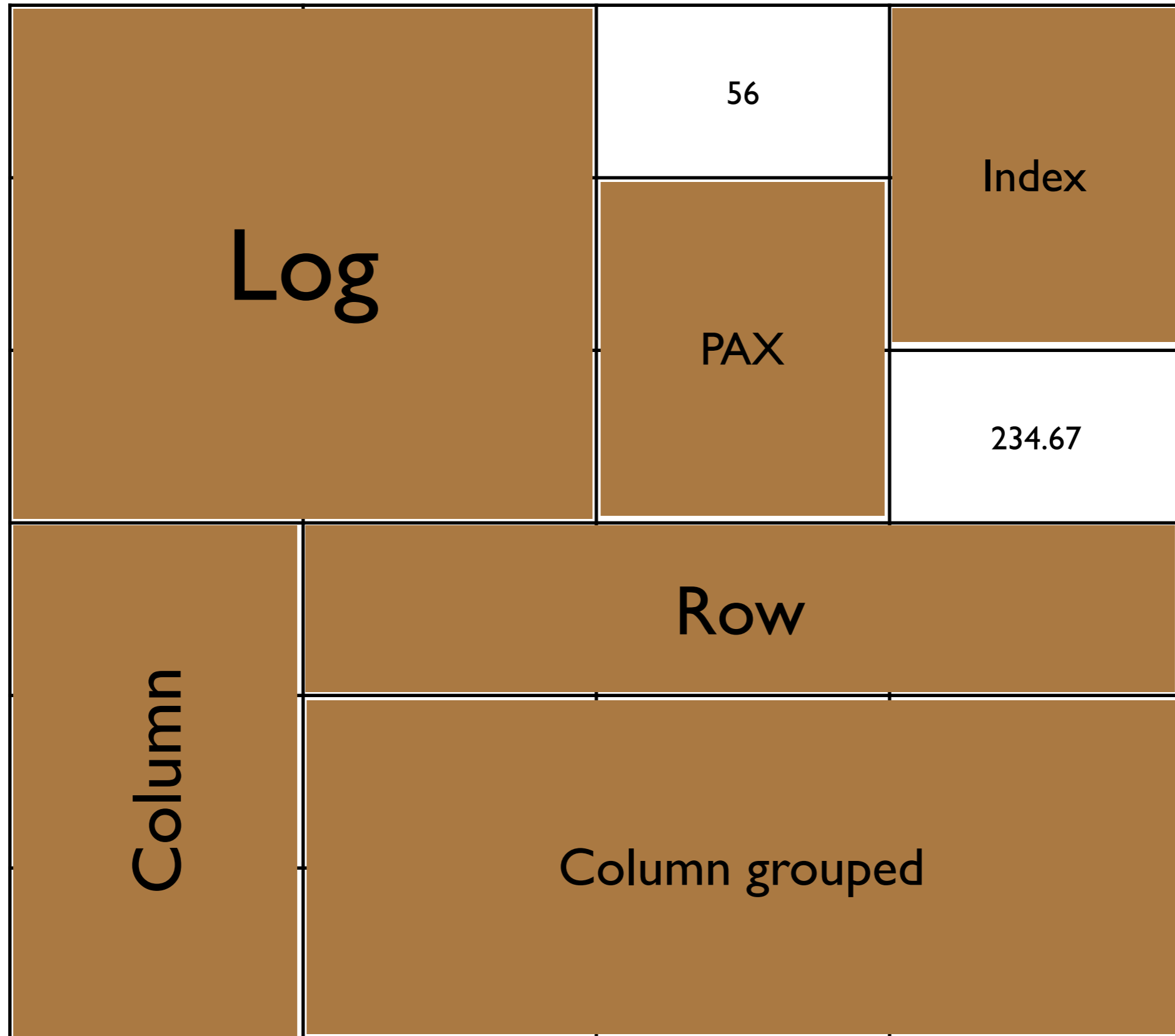
Storage Views

Log	56	887.9
	89	445.35
	67	234.67
Column	Row	
	Column grouped	

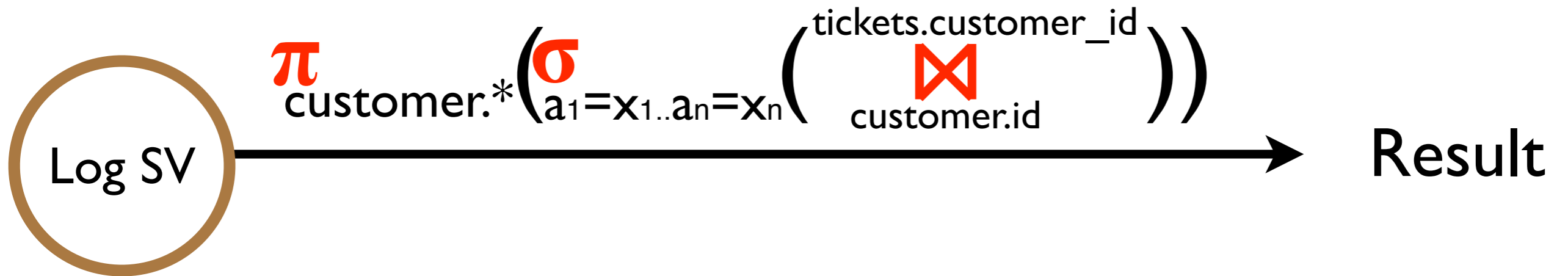
Storage Views



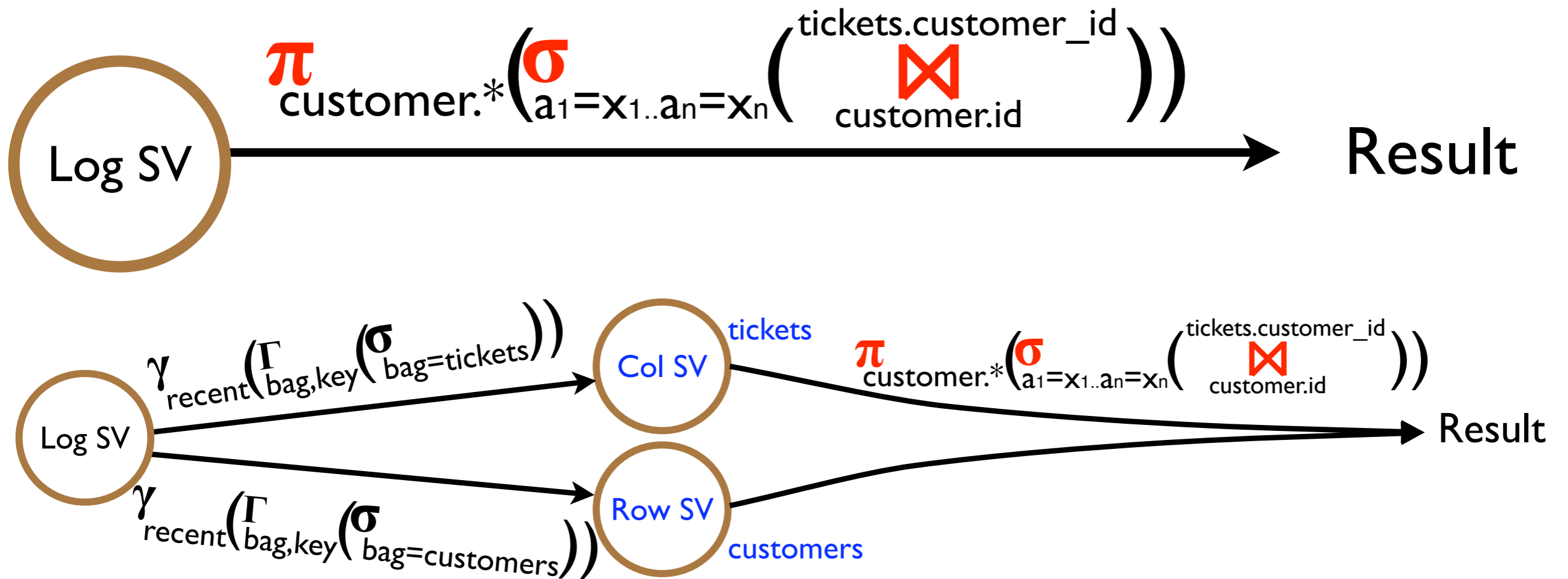
Storage Views



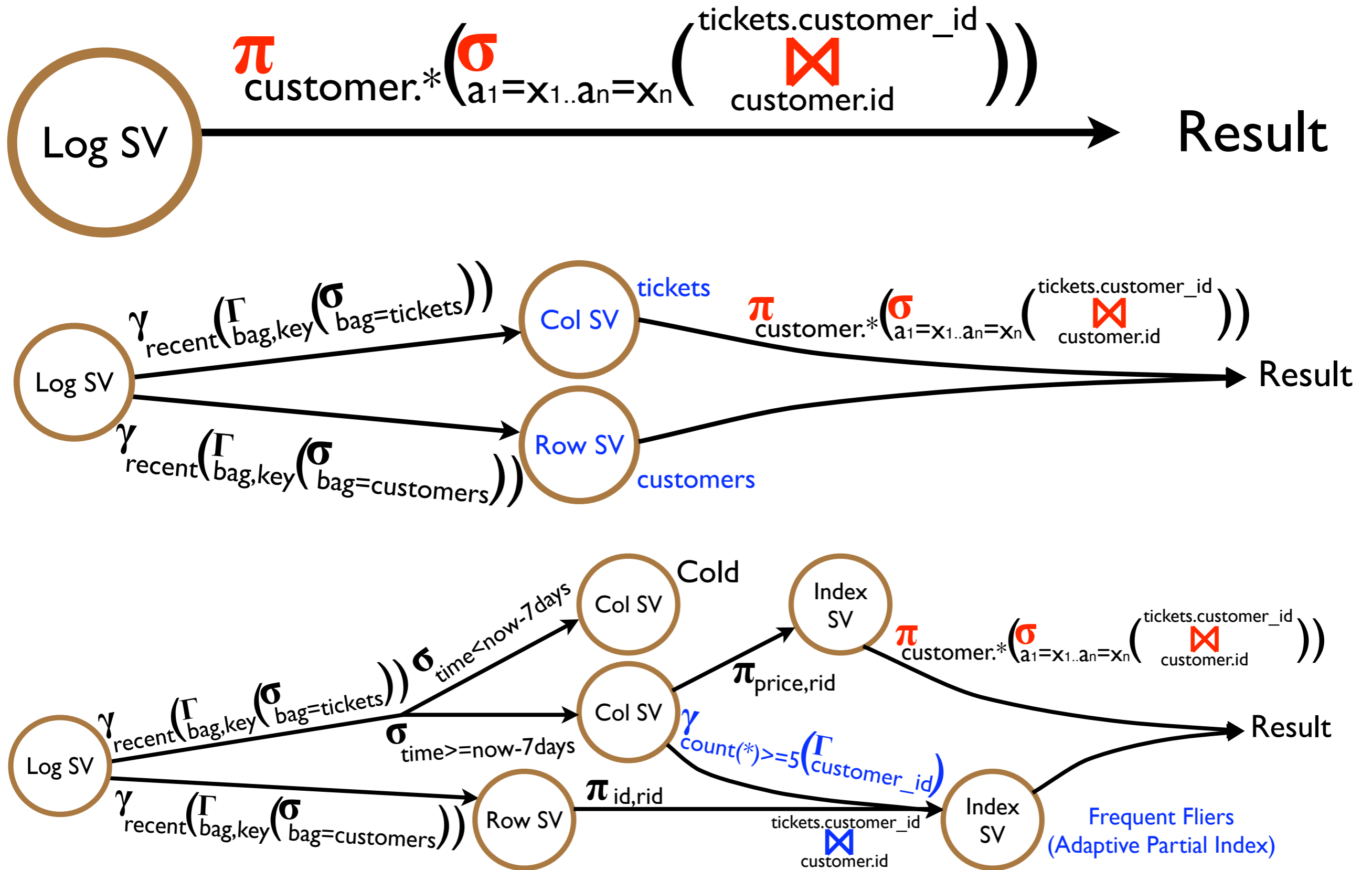
Example: Flight Tickets



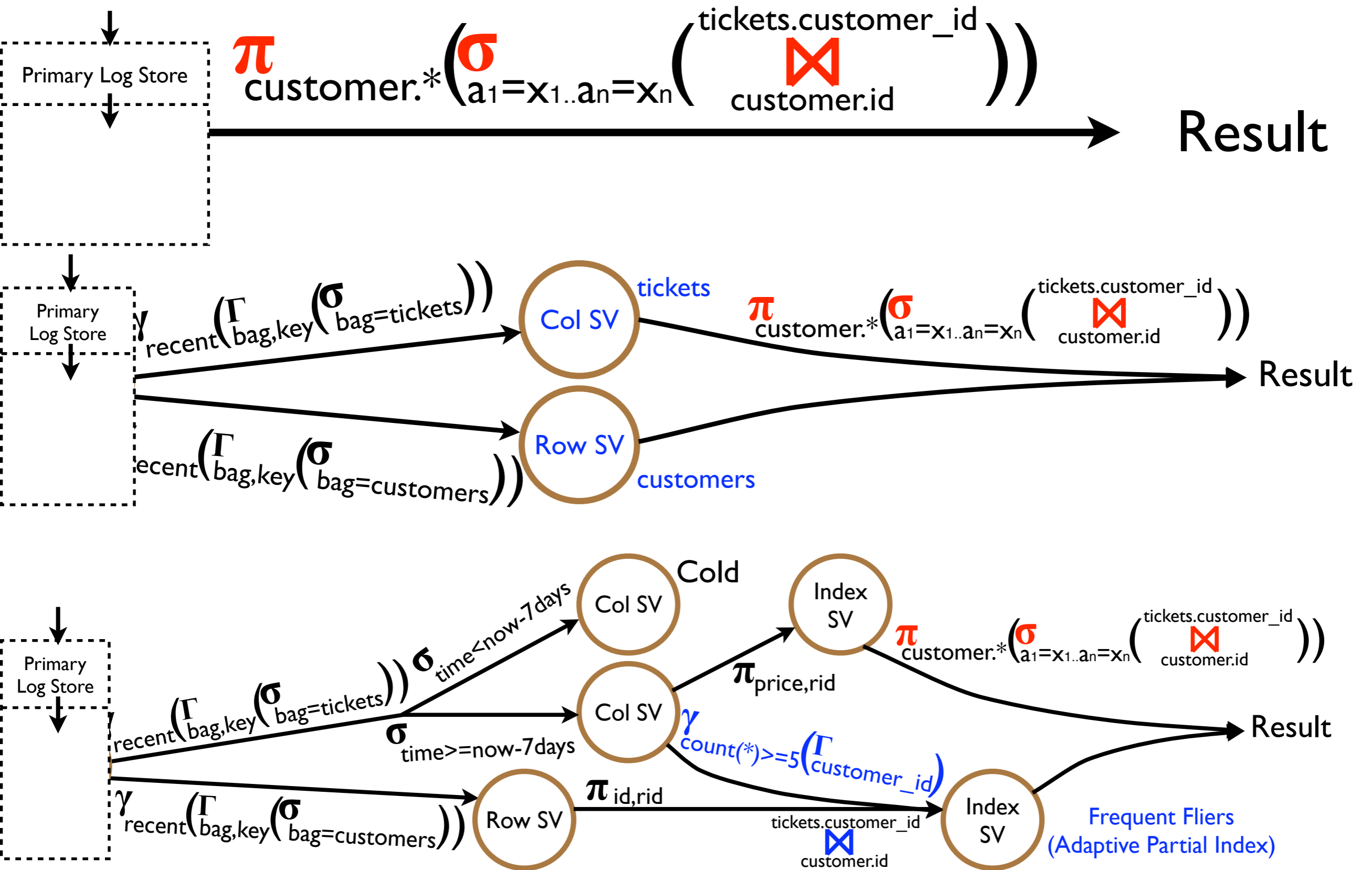
Example: Flight Tickets

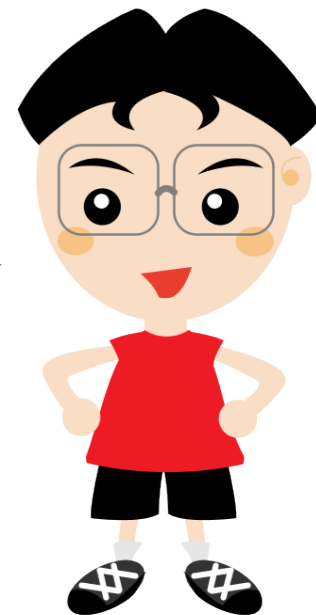


Example: Flight Tickets

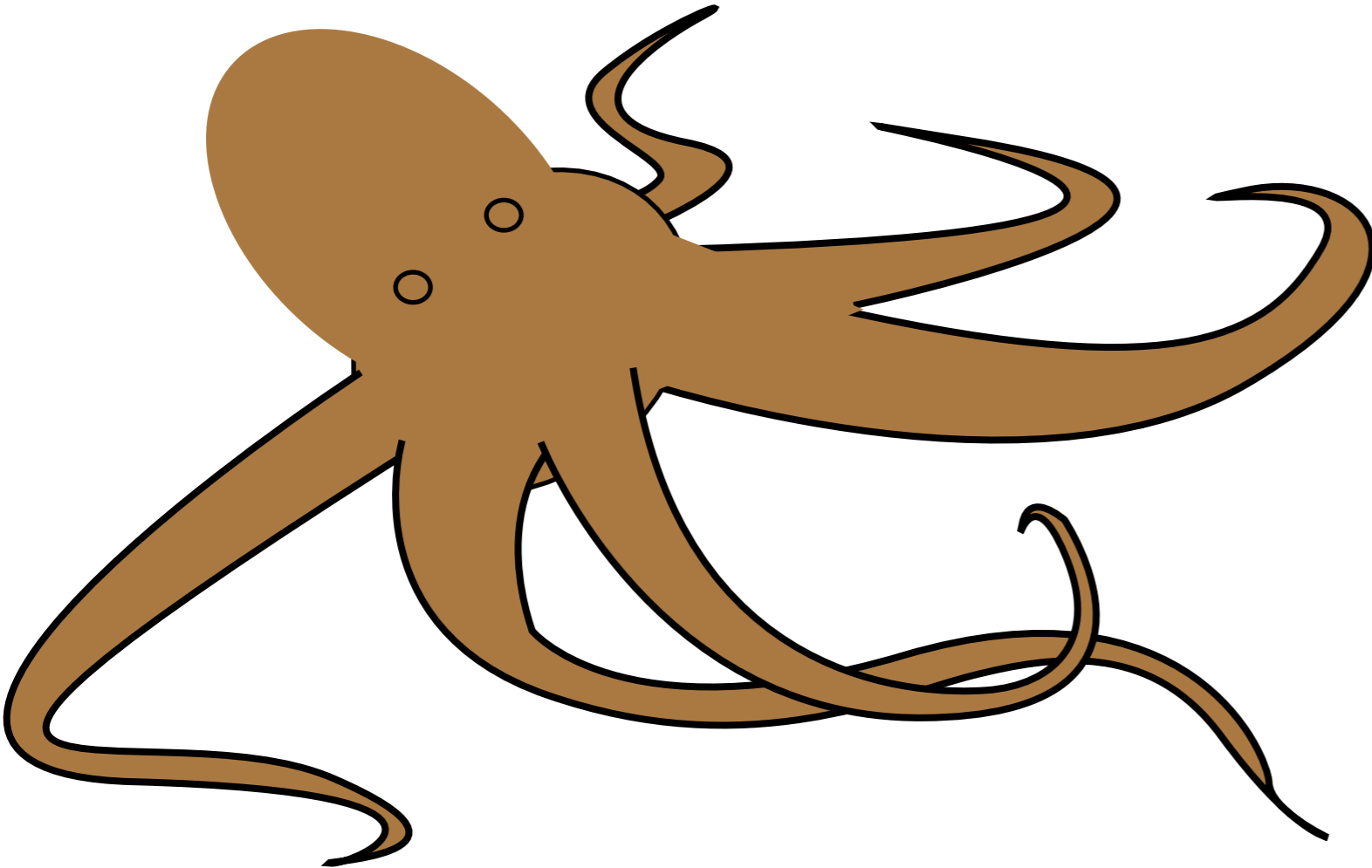


Example: Flight Tickets





WTF!
Where's The Food!



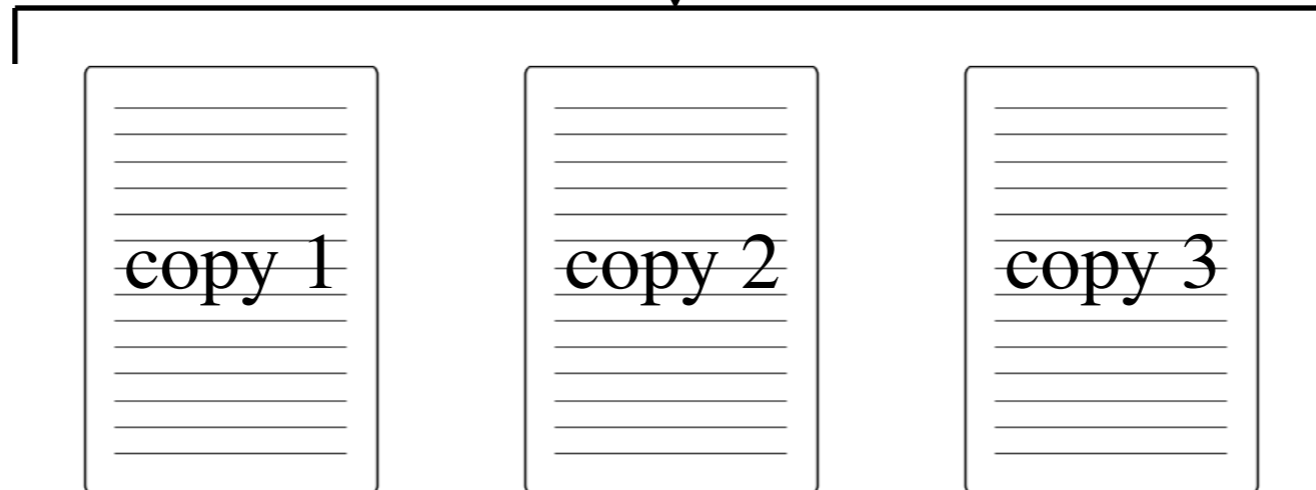
Rodent Store



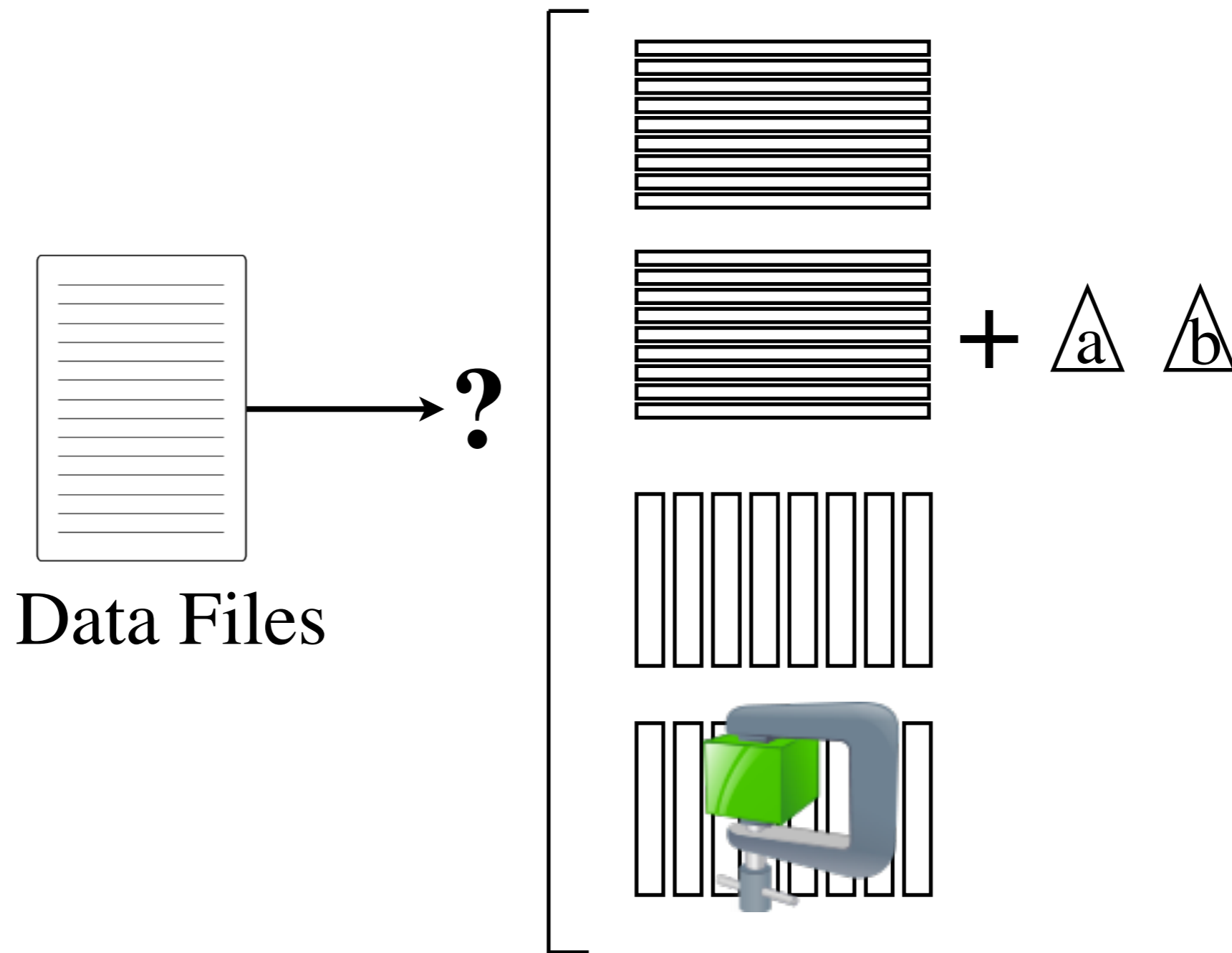
What to store?



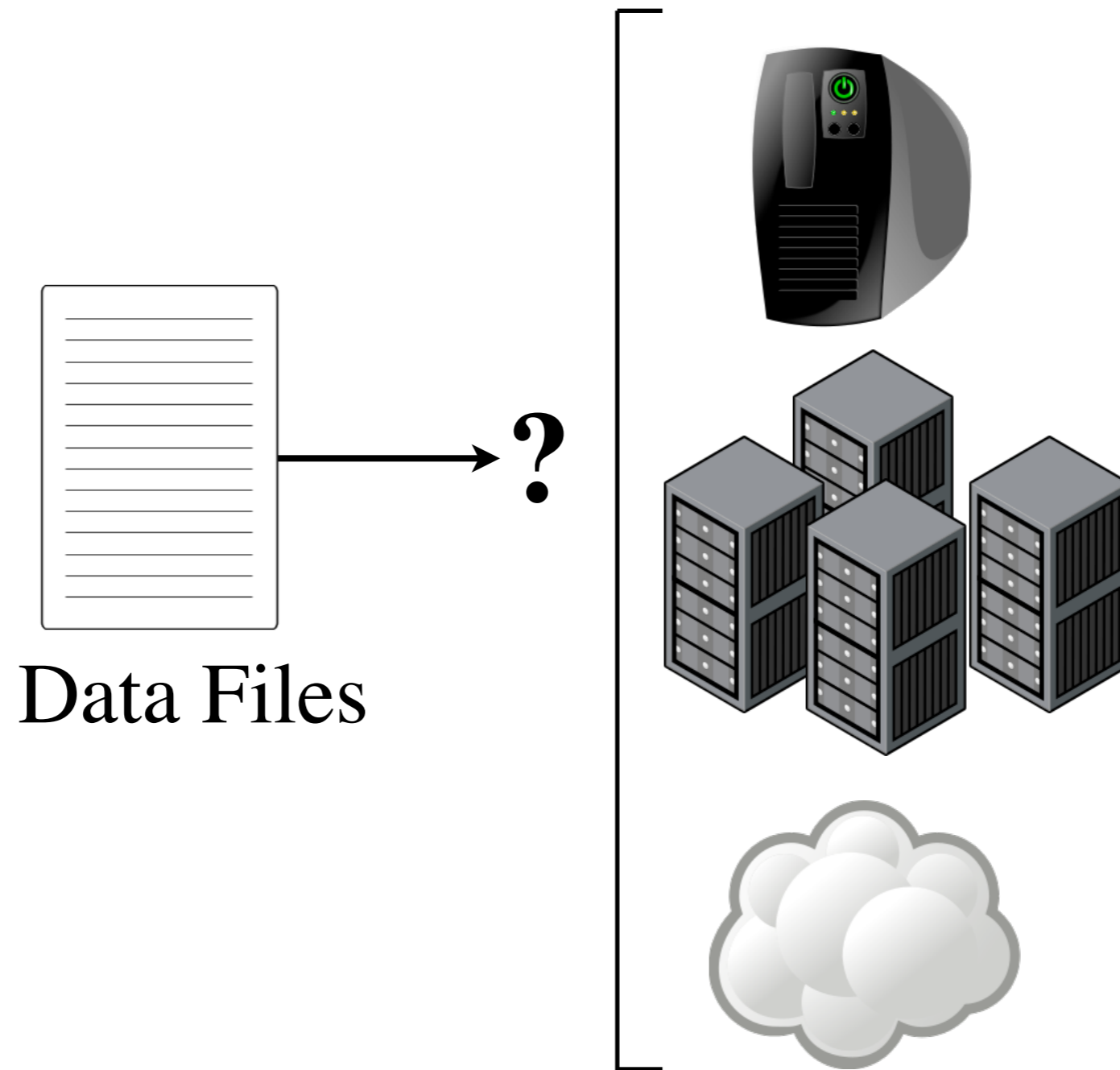
Data Files

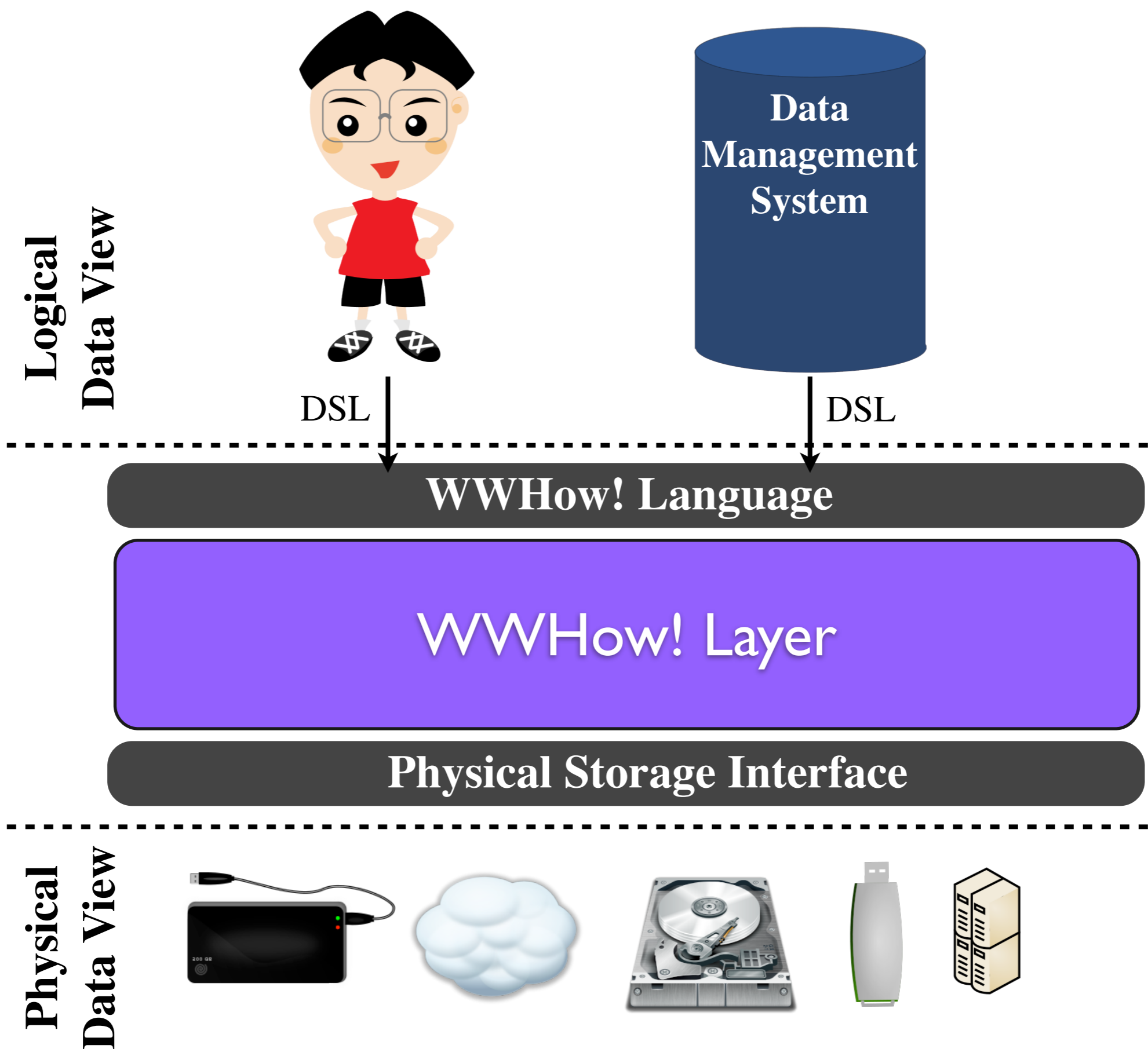


How to store?



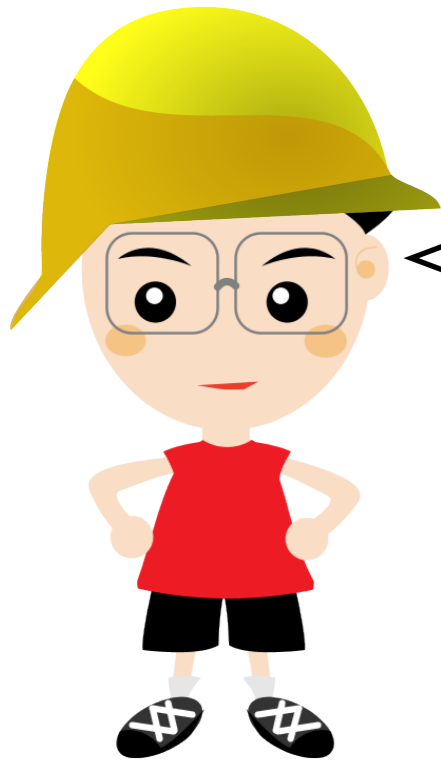
Where to store?





Example Use-cases

- WWHow! File System
- WWHow! RAID
- WWHow! Relational DBMS
- WWHow! Cloud



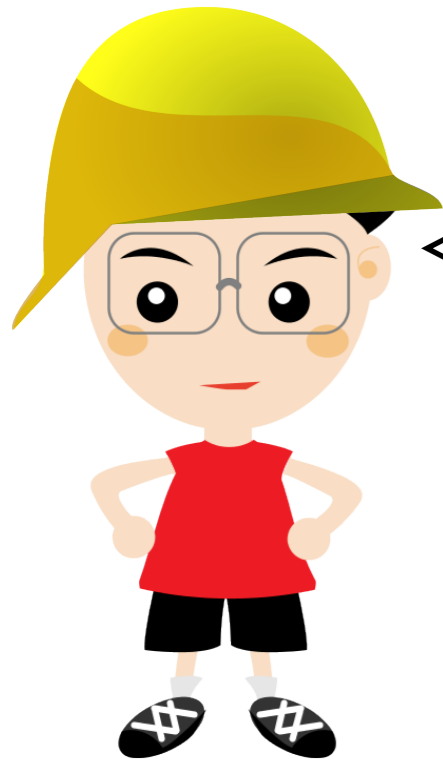
Store my conferences talks (PDFs 2x and PPTs 1x)
using RSA compression on University server

STORE '/Users/Bob/Conferences/Talks/*.*''

WHAT *.*(pdf | ppt), *.pdf

WHERE vise4

HOW encryption(rsa) FOR *;



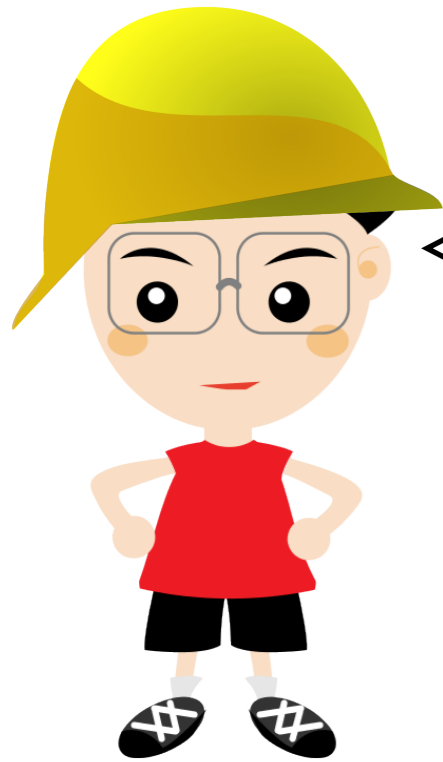
I want my conference talks to be highly available

STORE '/Users/Bob/Conferences/Talks/*.*'

WHAT *.(pdf | ppt), *.pdf

HOW encryption(rsa) FOR *

PREFERENCE Availability='high';



I want my conference talks to be highly available

STORE '/Users/Bob/Conferences/Talks/*.*'

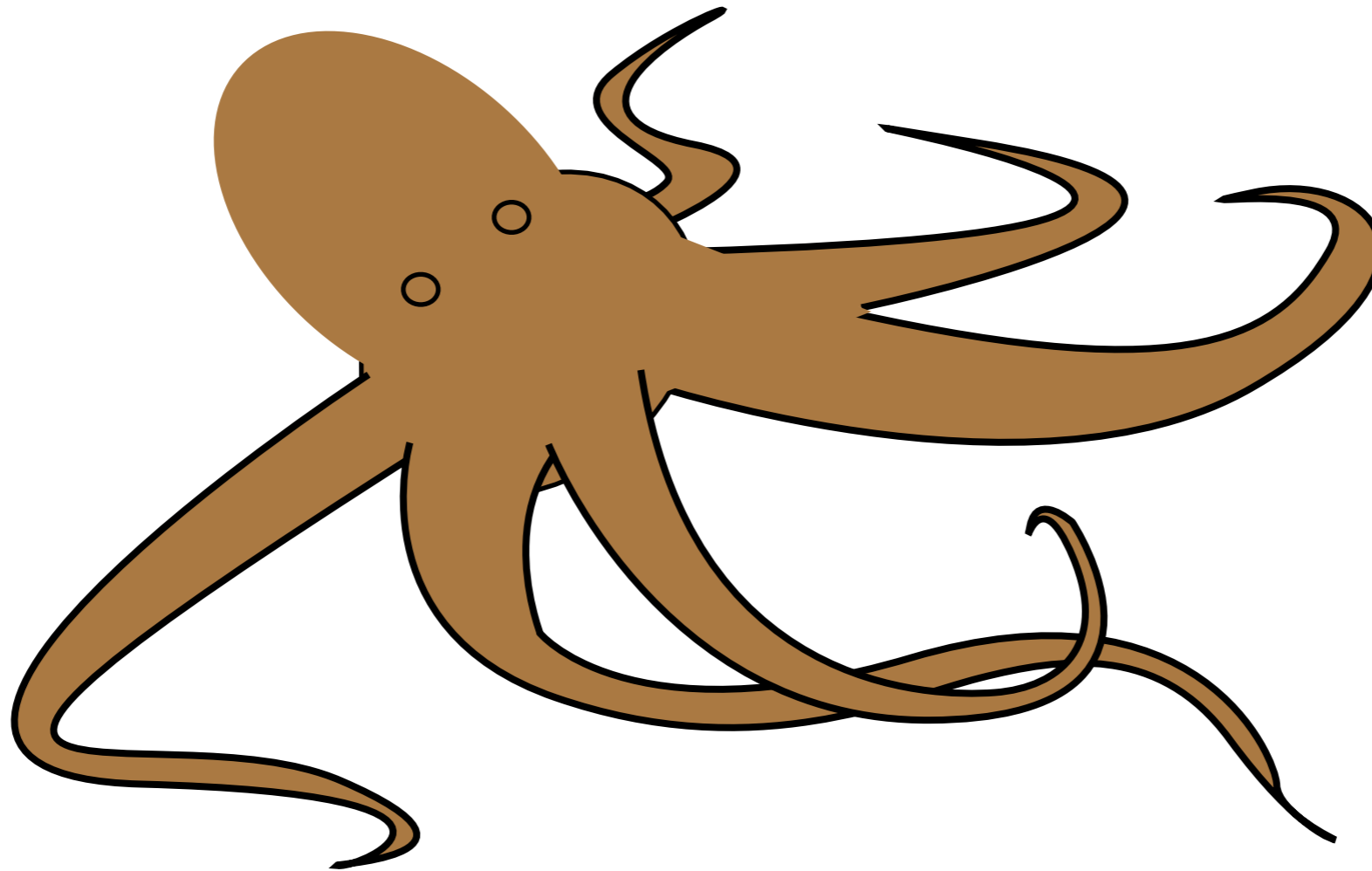
WHAT *.*(pdf | ppt), *.pdf

HOW encryption(rsa) FOR *

PREFERENCE Availability='high';

job for the
WWhow! data storage optimizer

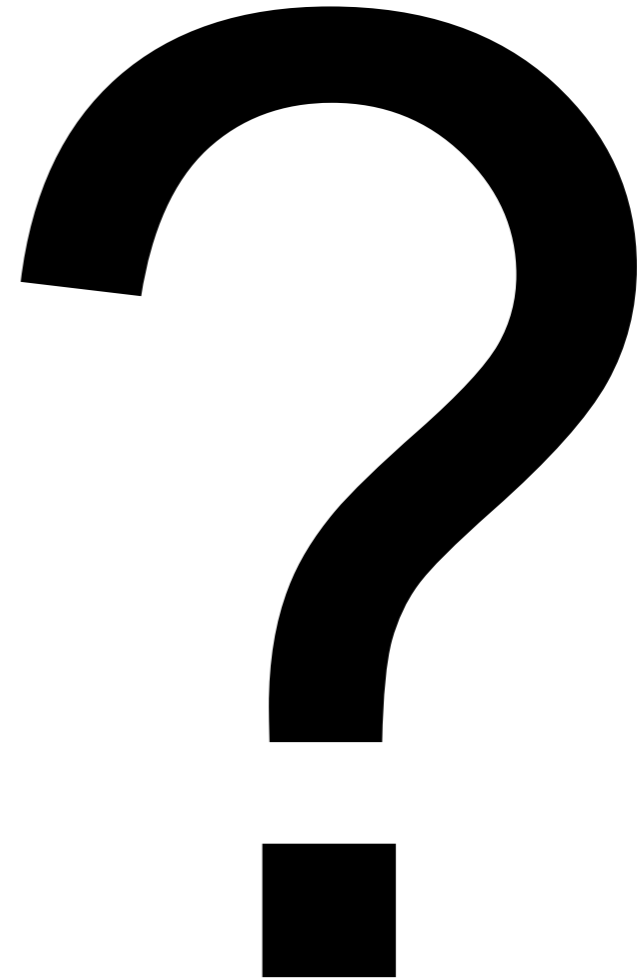
OctopusDB



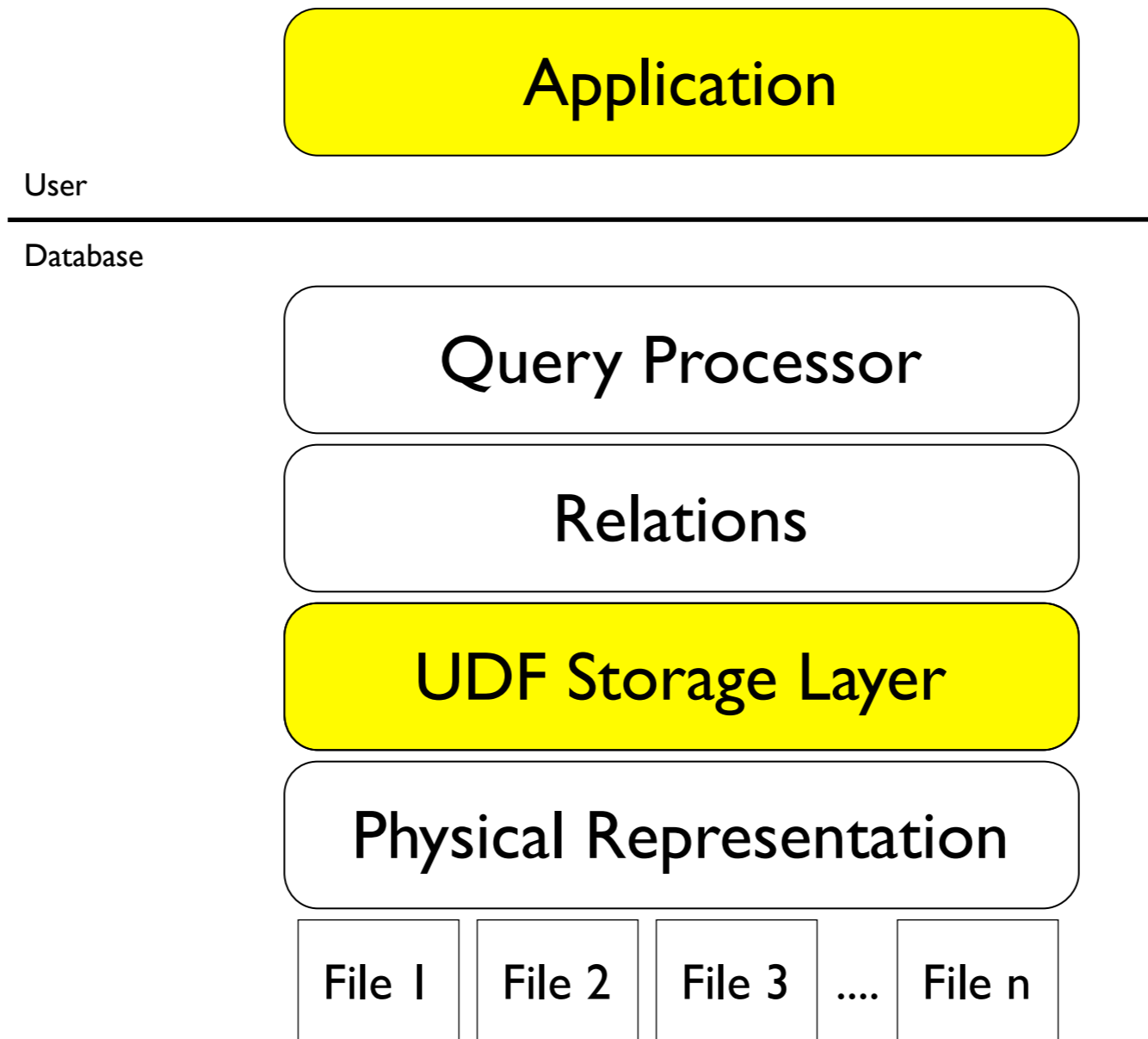
- Cool Vision
- Tough to realize



C-Store



Trojan Columns



Trojan Columns

Relation

Customer		
name	phone	market_segment
smith	2134	automobile
john	3425	household
kim	6756	furniture
joe	9878	building
mark	4312	building
steve	2435	automobile
jim	5766	household
ian	8789	household

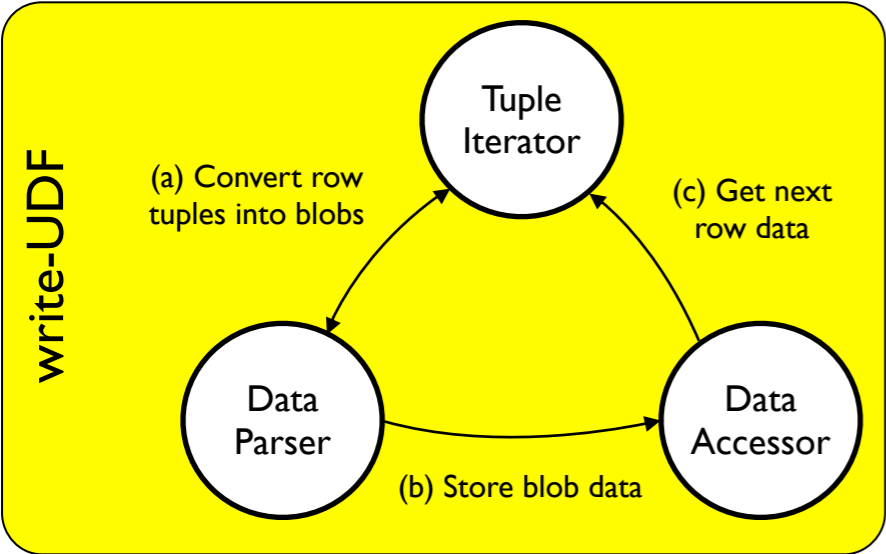
Physical Table

Customer_trojan		
segment_ID	attribute_ID	blob_data
1	name	smith, john, kim, joe
1	phone	2134, 3425, 6756, 9878
1	market_segment	automobile, household, furniture, building
2	name	mark, steve, jim, ian
2	phone	4312, 2435, 5766, 8789
2	market_segment	building, automobile, household, household

Trojan Columns

Relation

Customer		
name	phone	market_segment
smith	2134	automobile
john	3425	household
kim	6756	furniture
joe	9878	building
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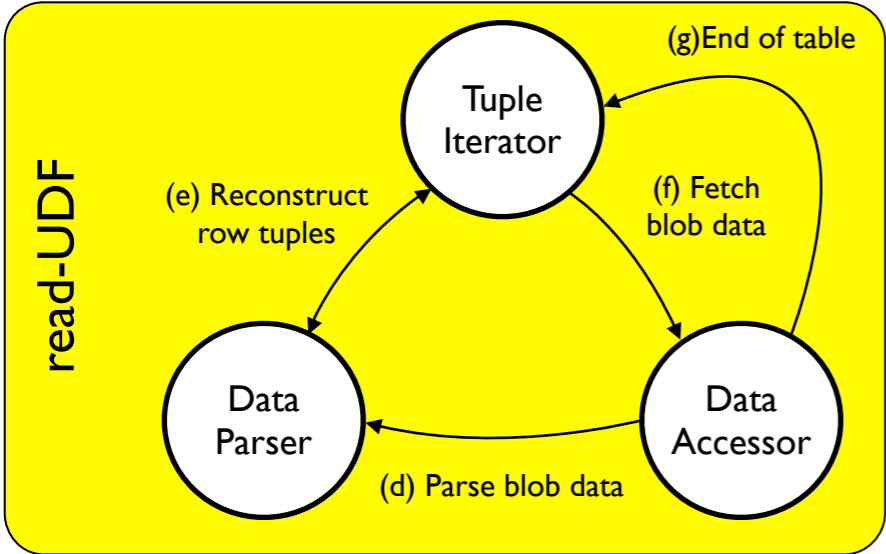


Physical Table

Customer_trojan		
segment_ID	attribute_ID	blob_data
1	name	smith, john, kim, joe
1	phone	2134, 3425, 6756, 9878
1	market_segment	automobile, household, furniture, building
2	name	mark, steve, jim, ian
2	phone	4312, 2435, 5766, 8789
2	market_segment	building, automobile, household, household

Trojan Columns

Relation



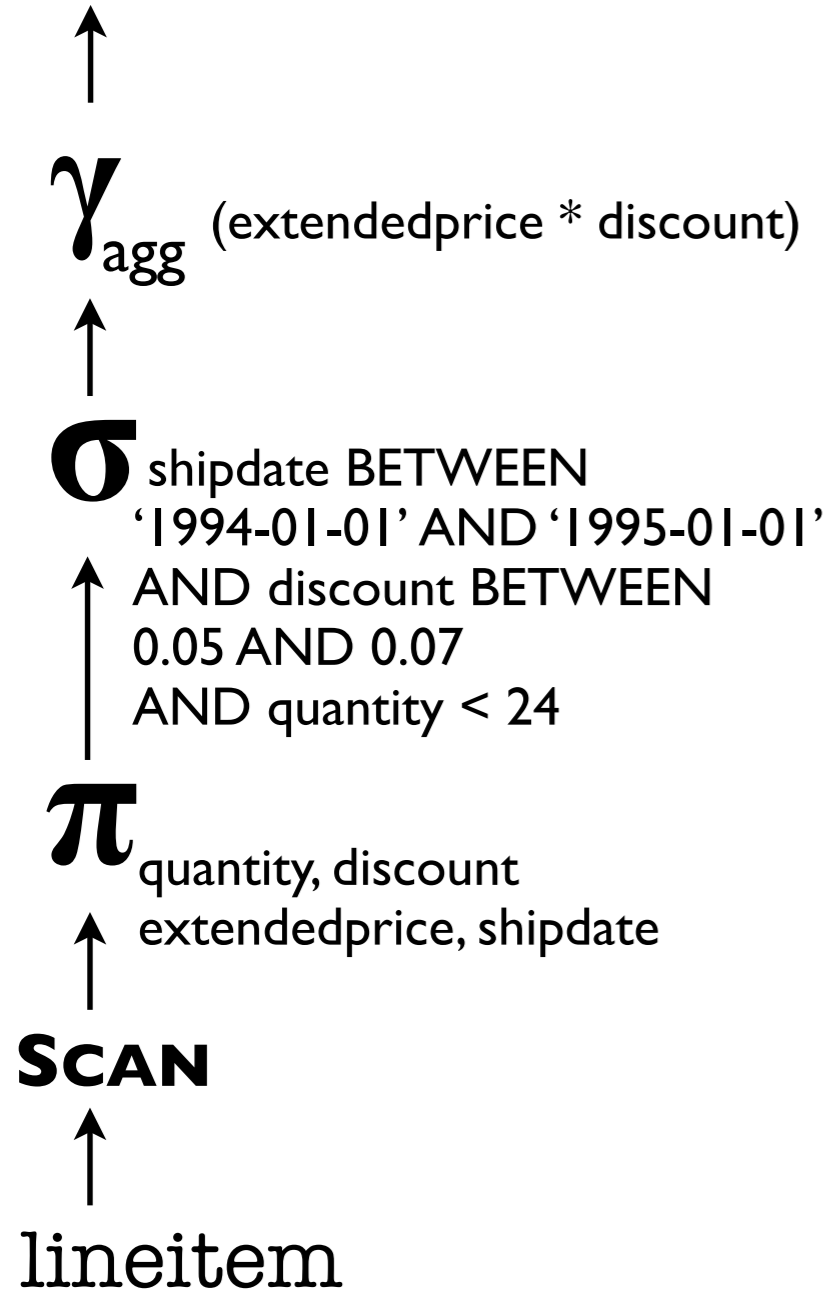
Customer		
name	phone	market_segment
smith	2134	automobile
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kim	6756	furniture
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Customer_trojan		
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2	phone	4312, 2435, 5766, 8789
2	market_segment	building, automobile, household, household

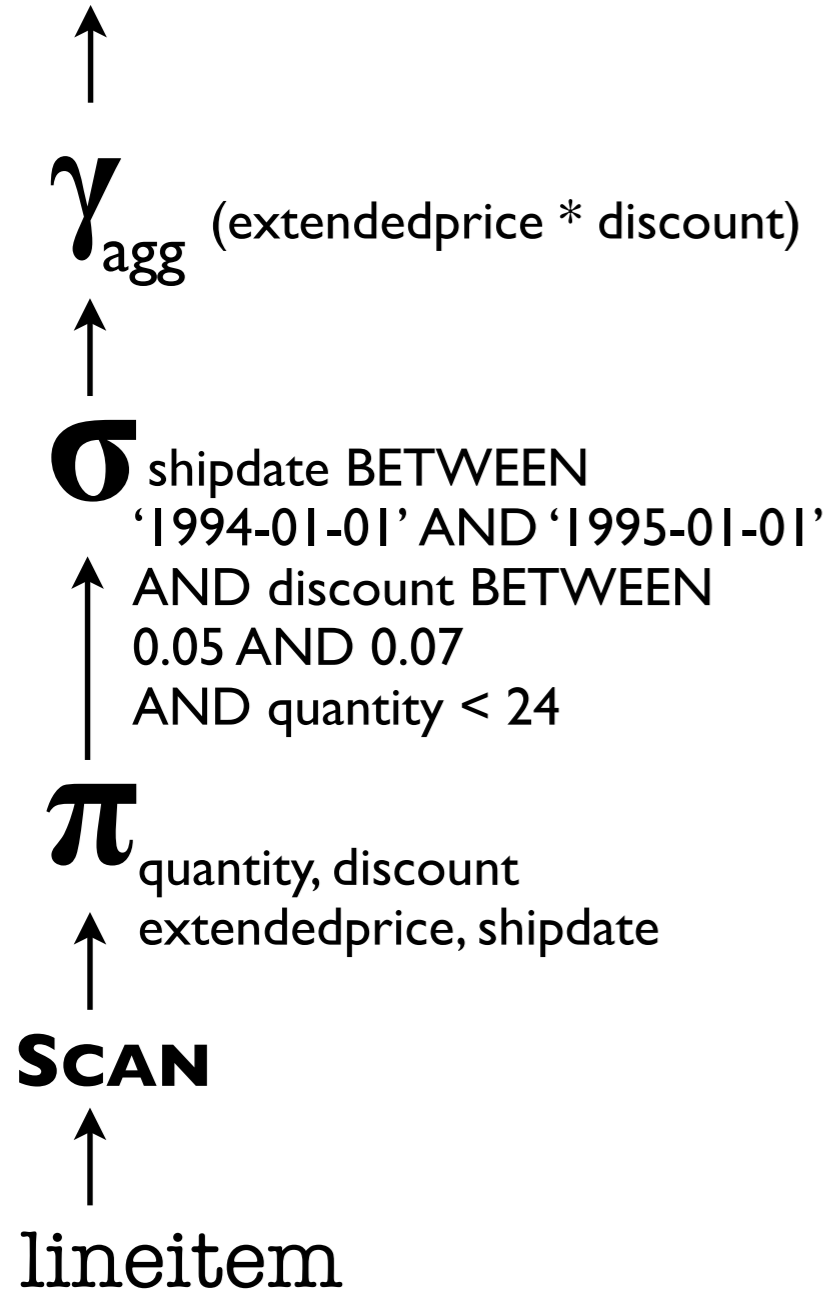
Example: TPC-H Query 6

Result

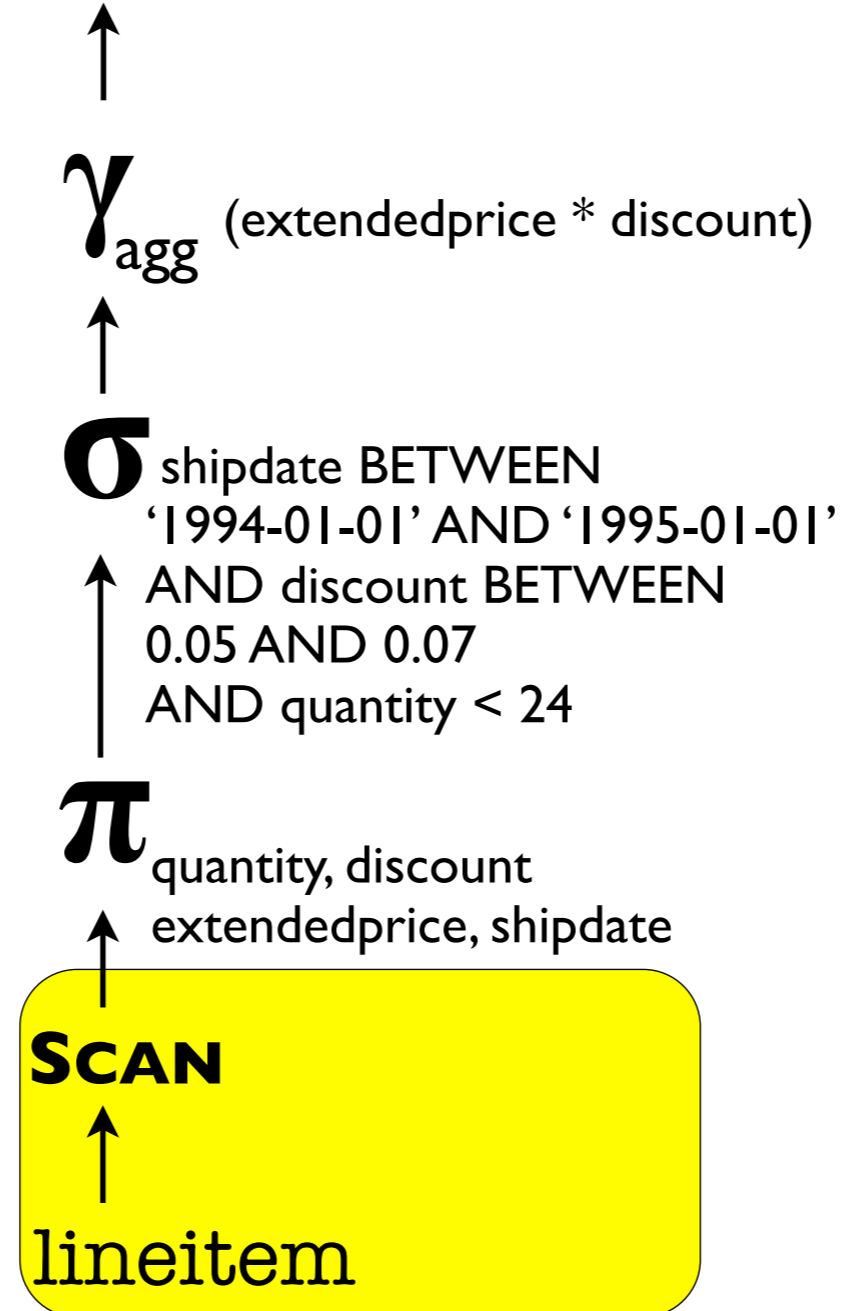


Example: TPC-H Query 6

Result



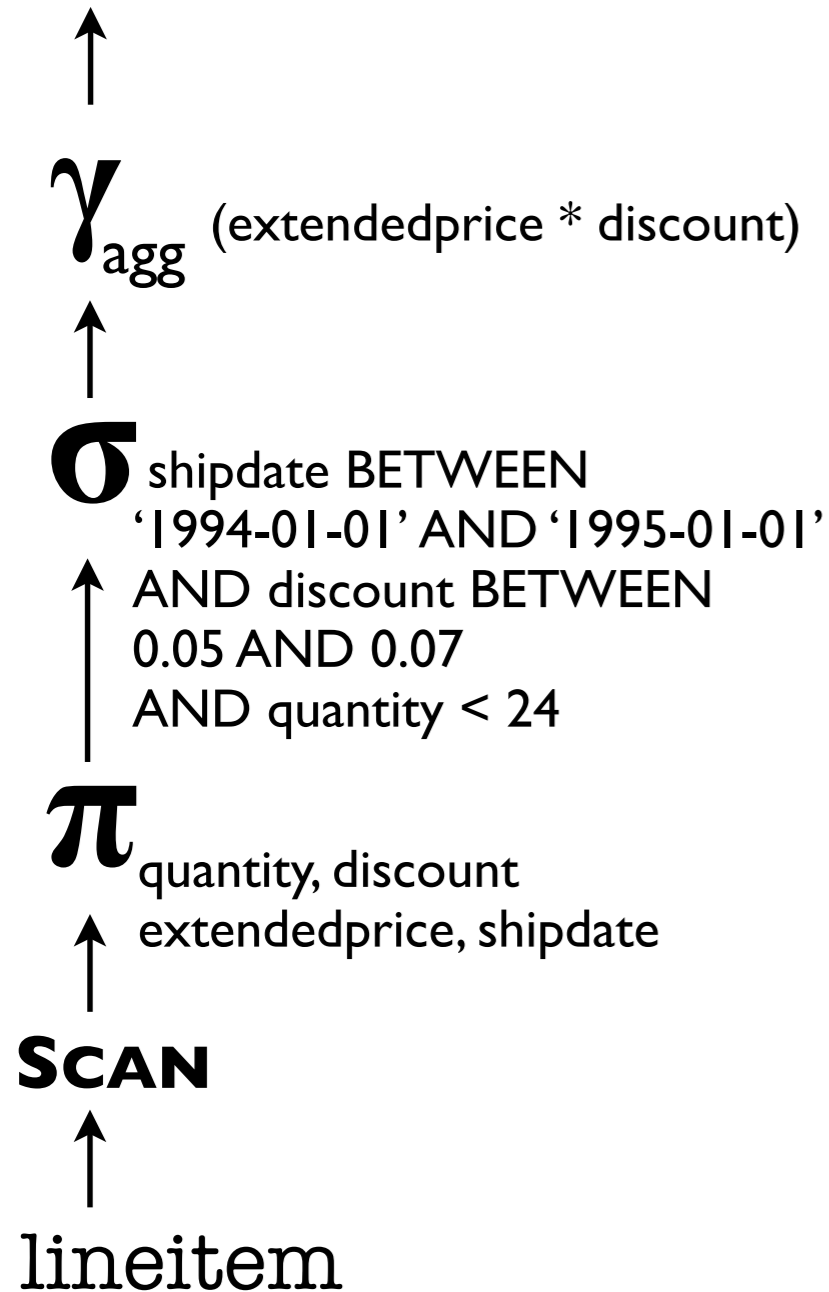
Result



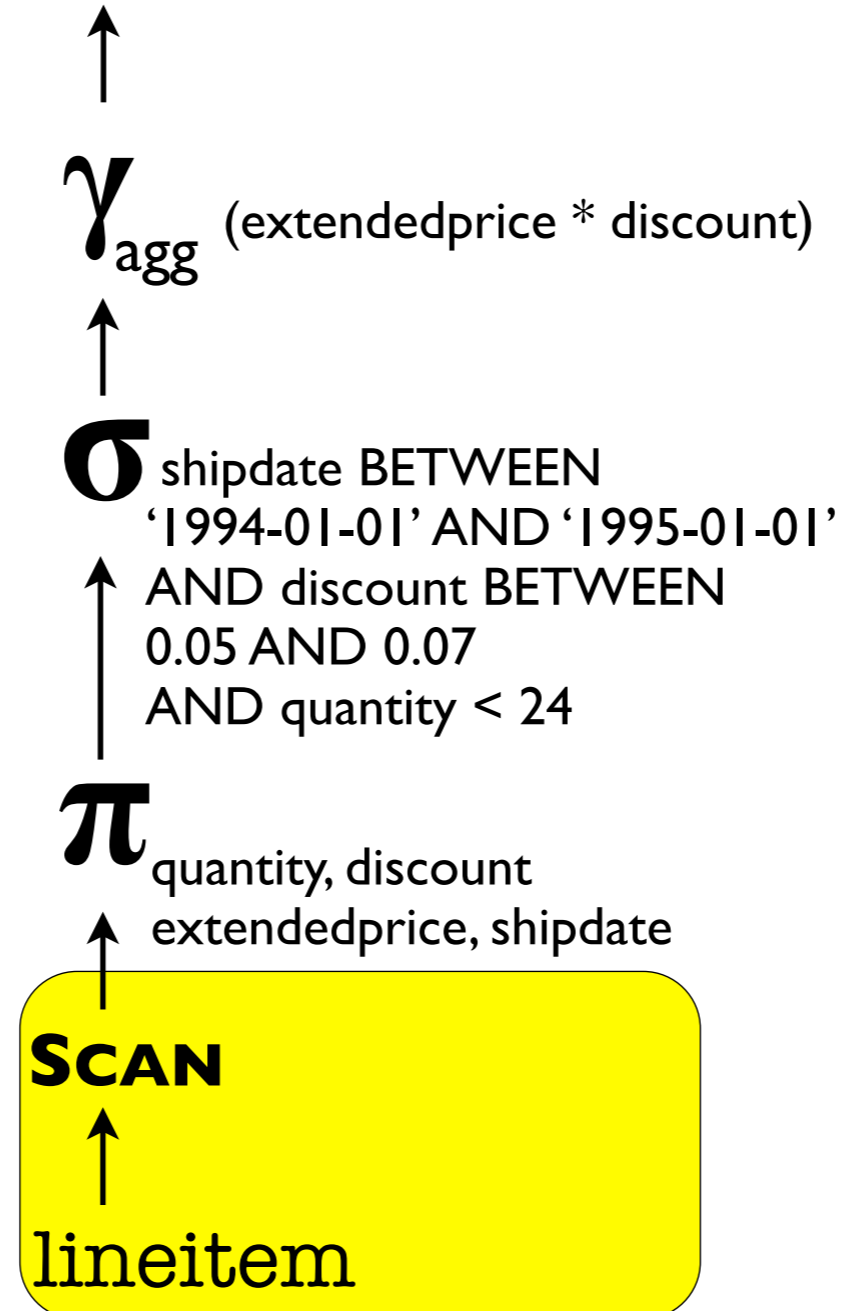
scanUDF

Example: TPC-H Query 6

Result

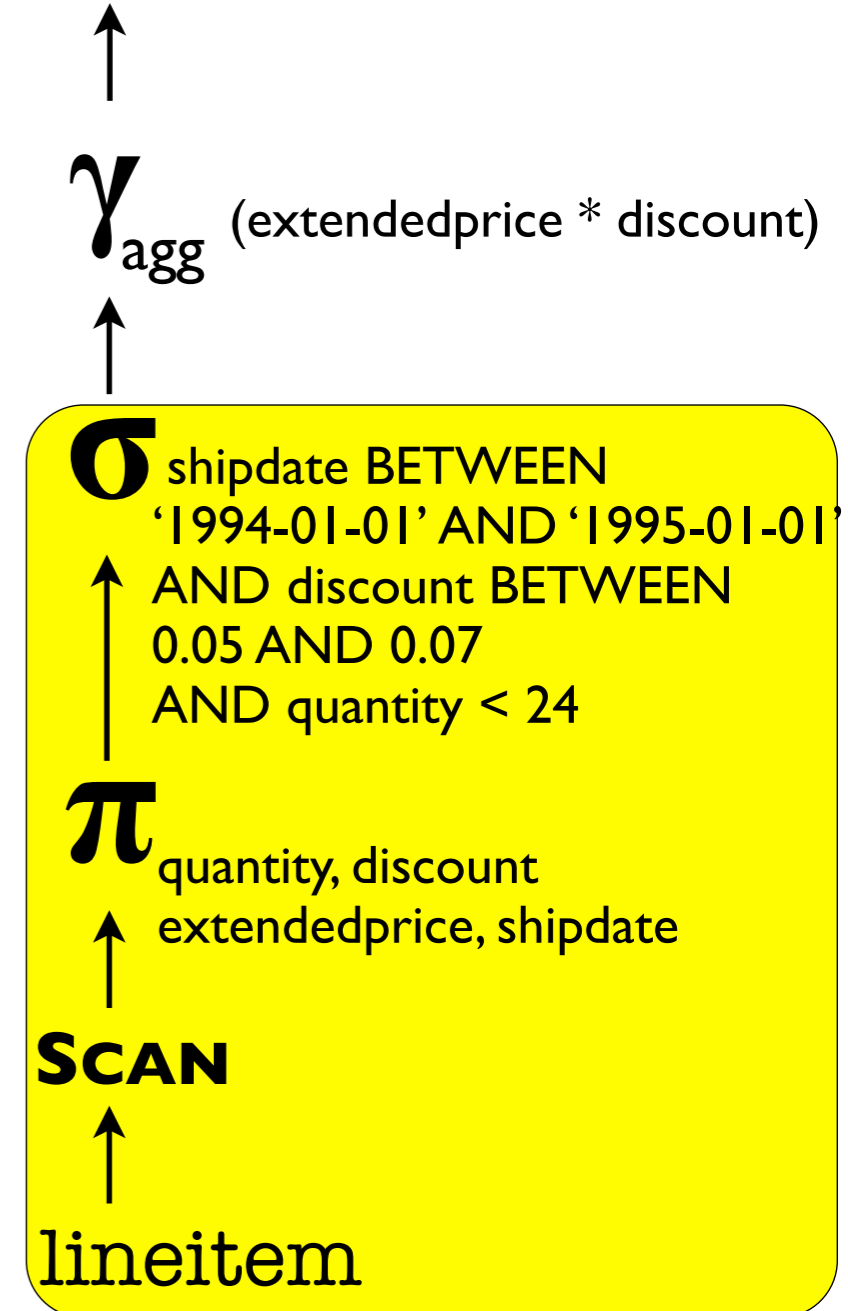


Result



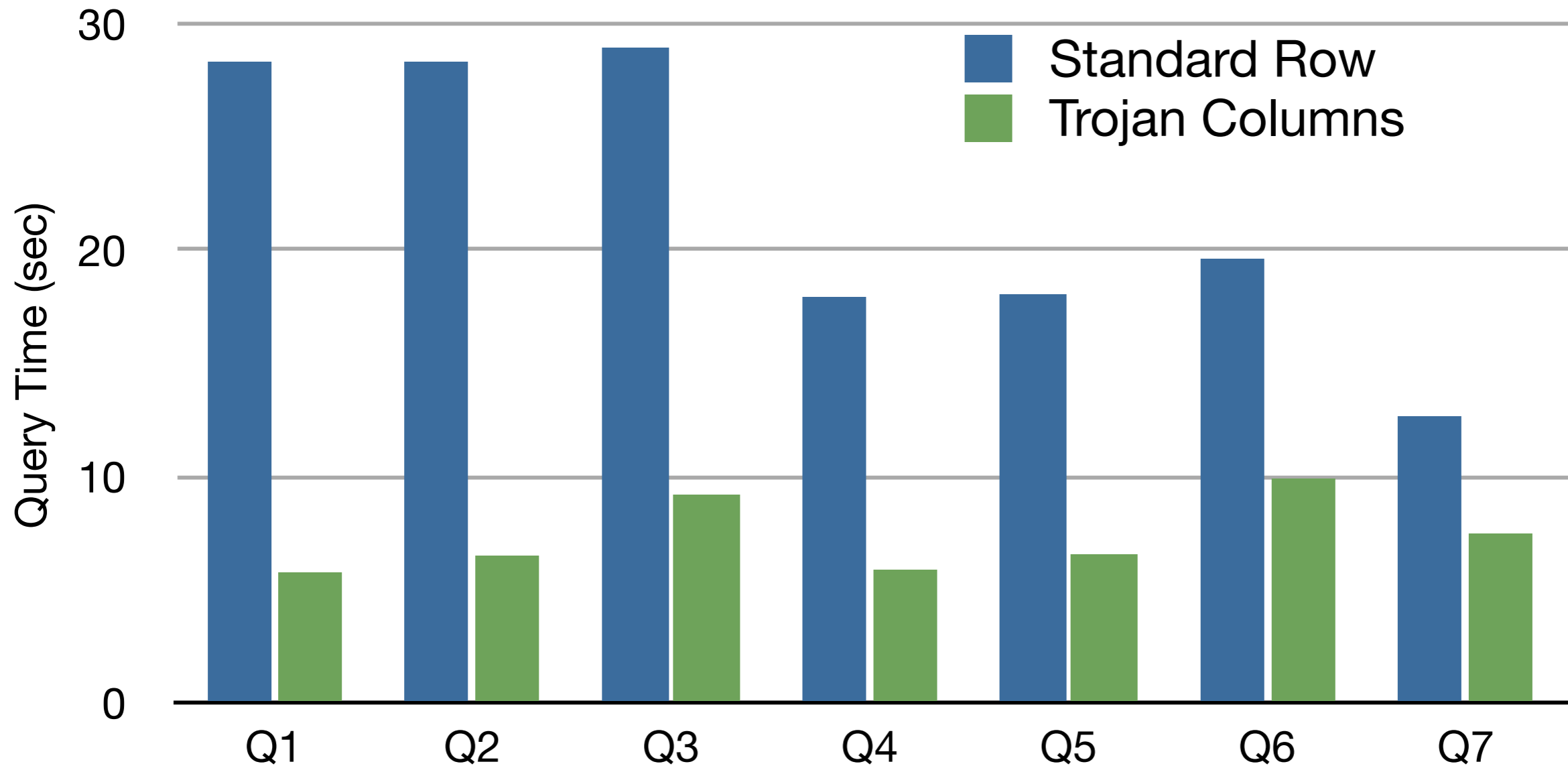
scanUDF

Result



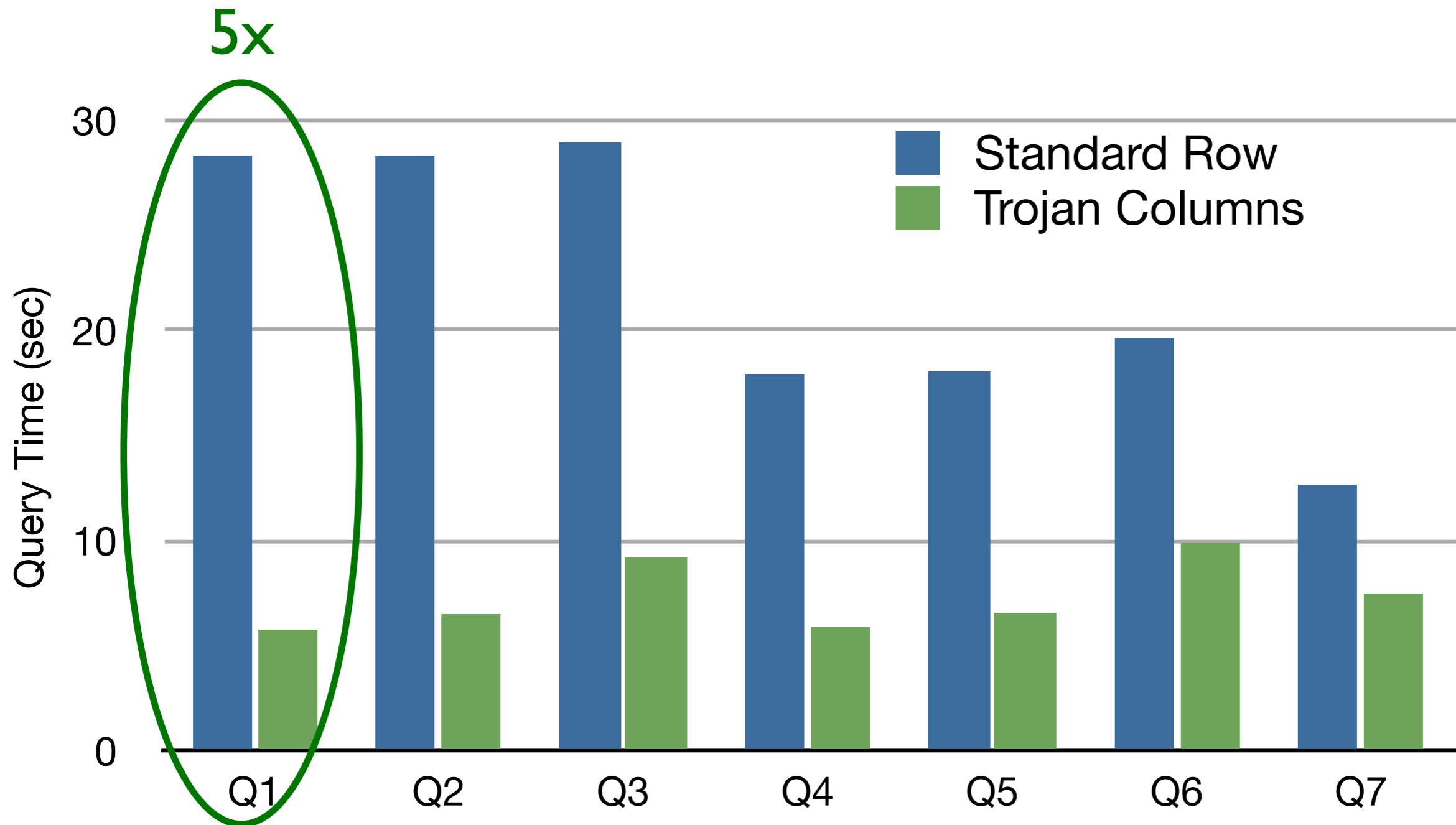
selectUDF

Benchmark Results *



* Mike Stonebraker et. al. C-Store: A Column Oriented DBMS. VLDB 2005

Benchmark Results *



* Mike Stonebraker et. al. C-Store: A Column Oriented DBMS. VLDB 2005

| 1980s

1990s

2000s

HYRISE

2010s

7 Vertical Partitioning Algorithms

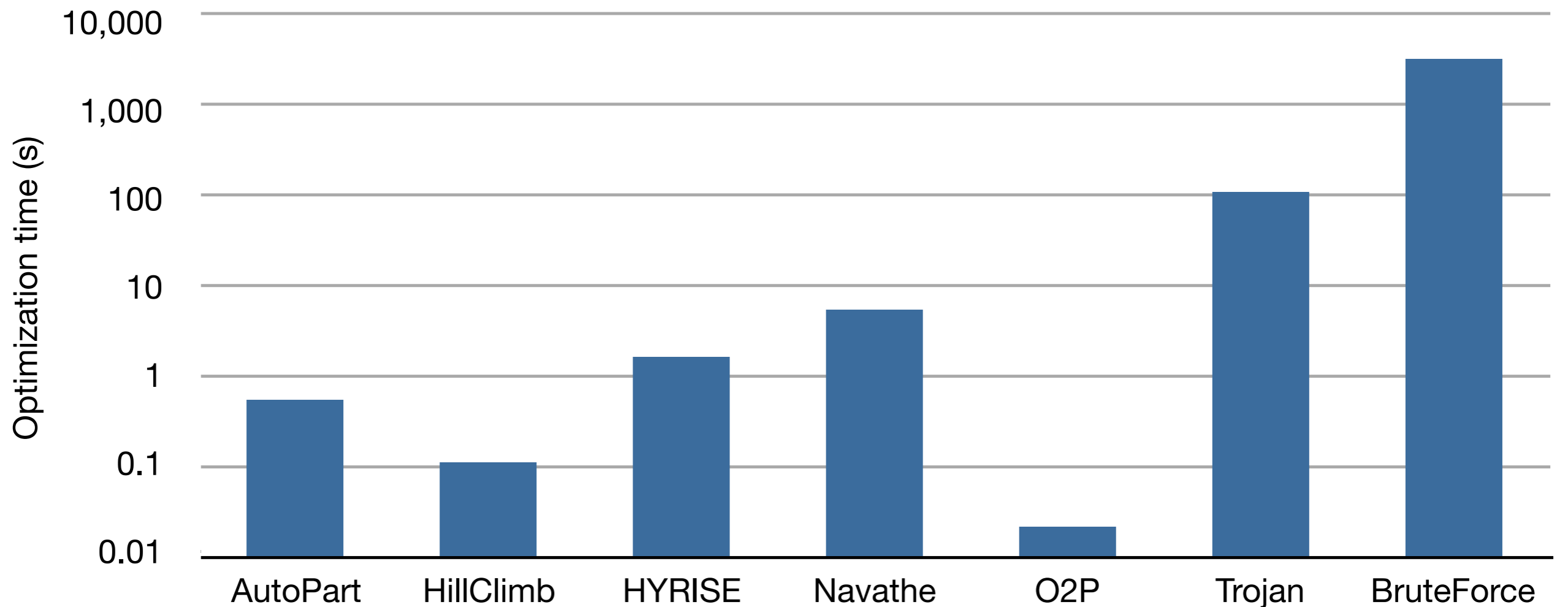
- Brute Force
- Navathe's Algorithm
- HillClimb
- AutoPart
- HYRISE
- O₂P
- Trojan



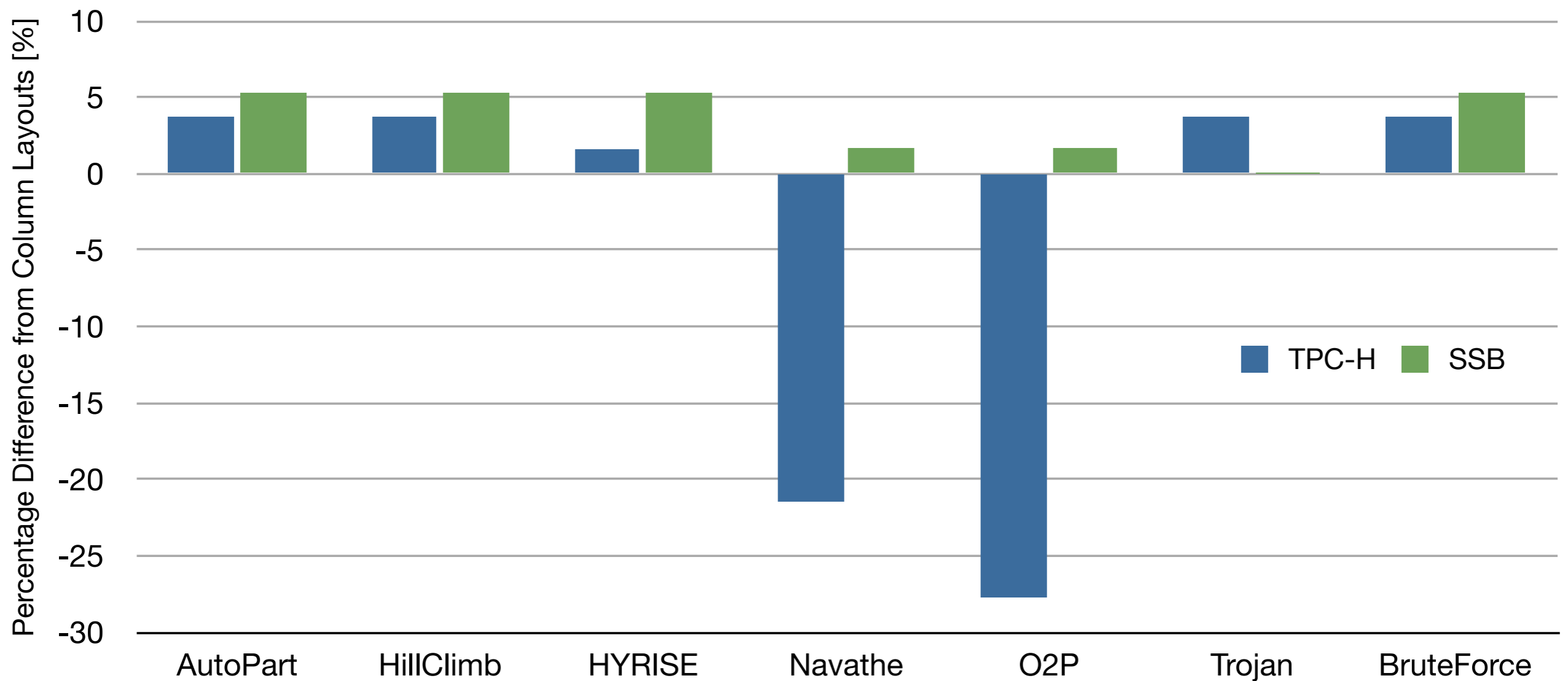
Four Comparison Metrics

- How Fast?
- How Good?
- How fragile?
- Where does it makes sense?

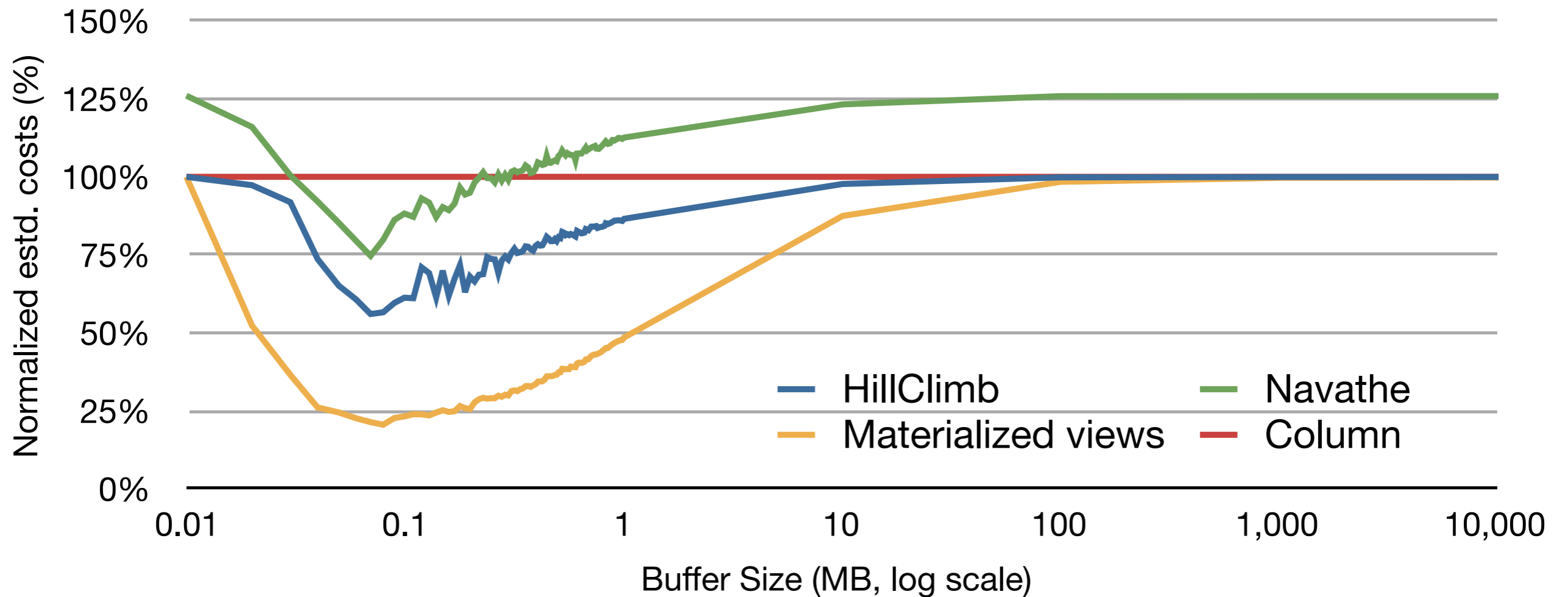
Optimization Runtime



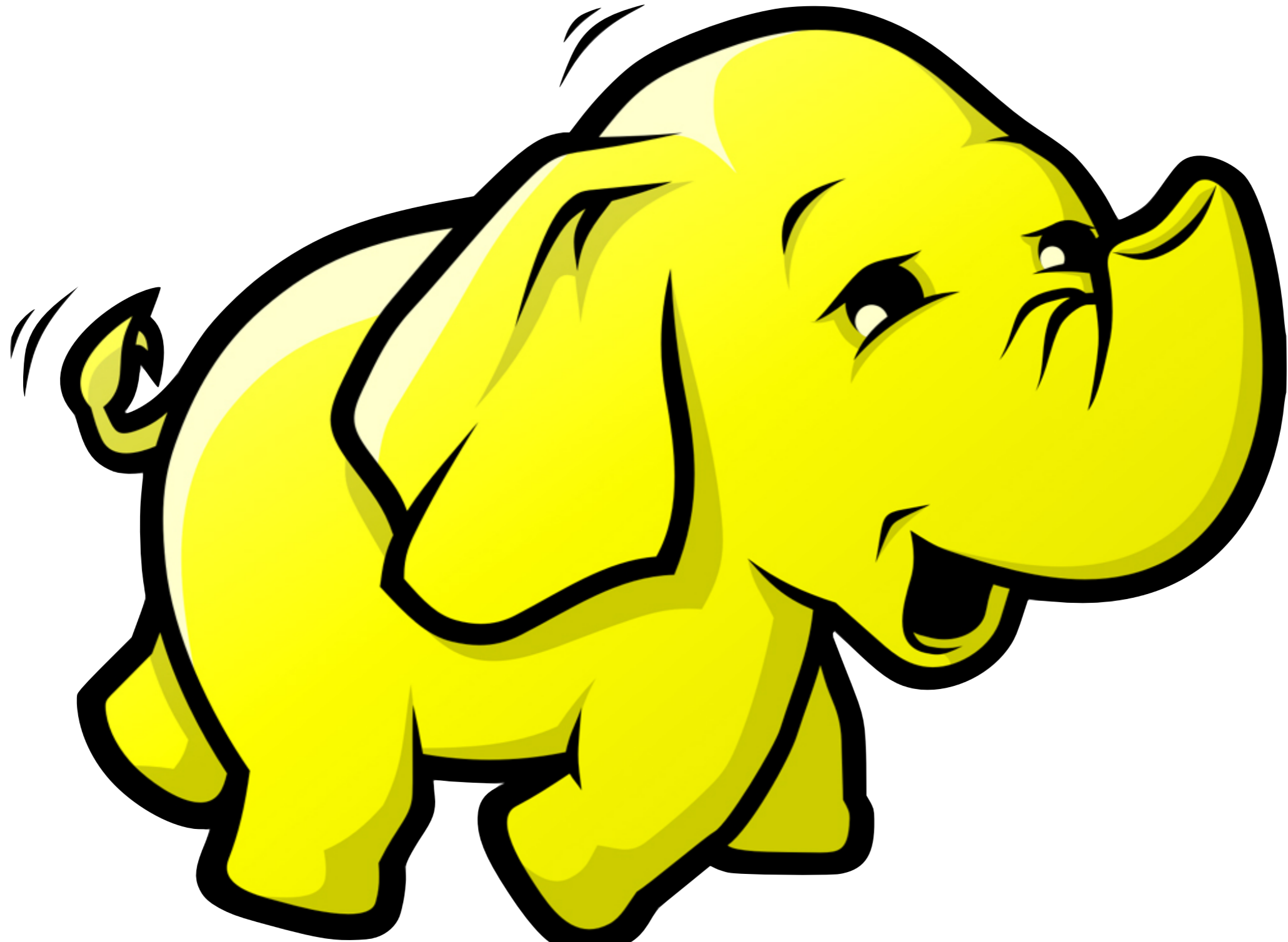
Distance from Column Layouts



Effect of Buffer Size



Comparison's Paper: Hadoop Vs PDBMS

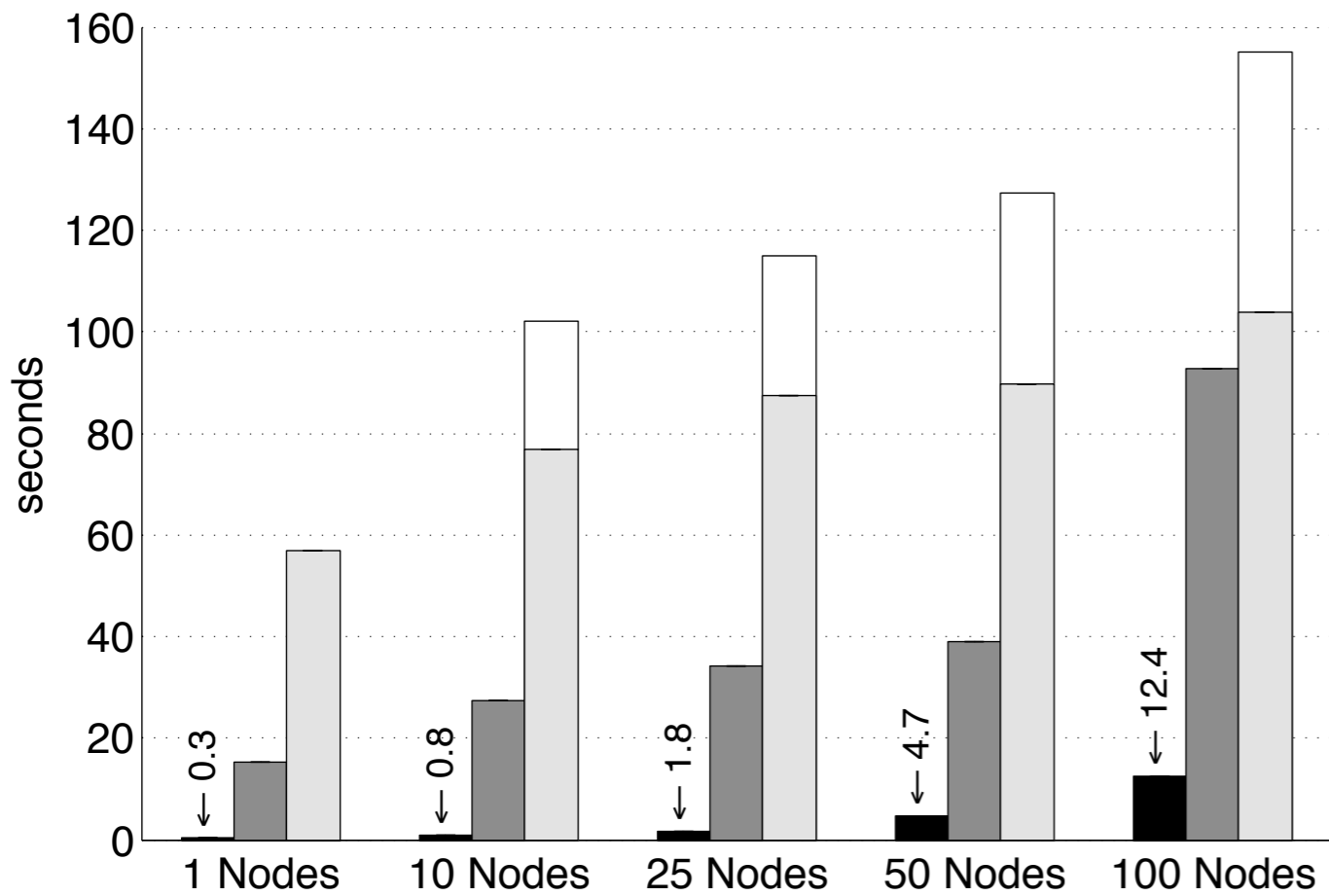


Comparison's Paper: Hadoop Vs PDBMS

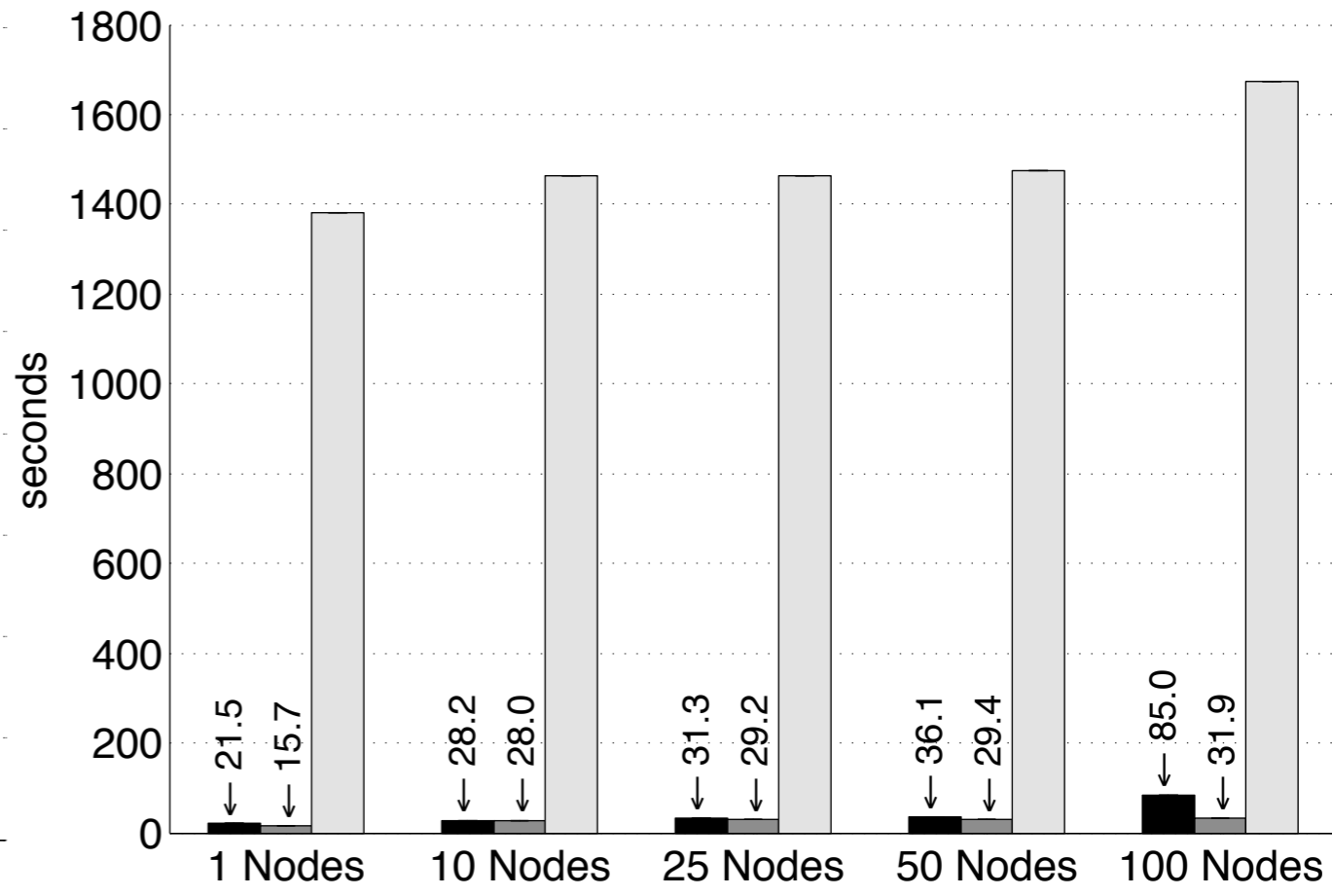


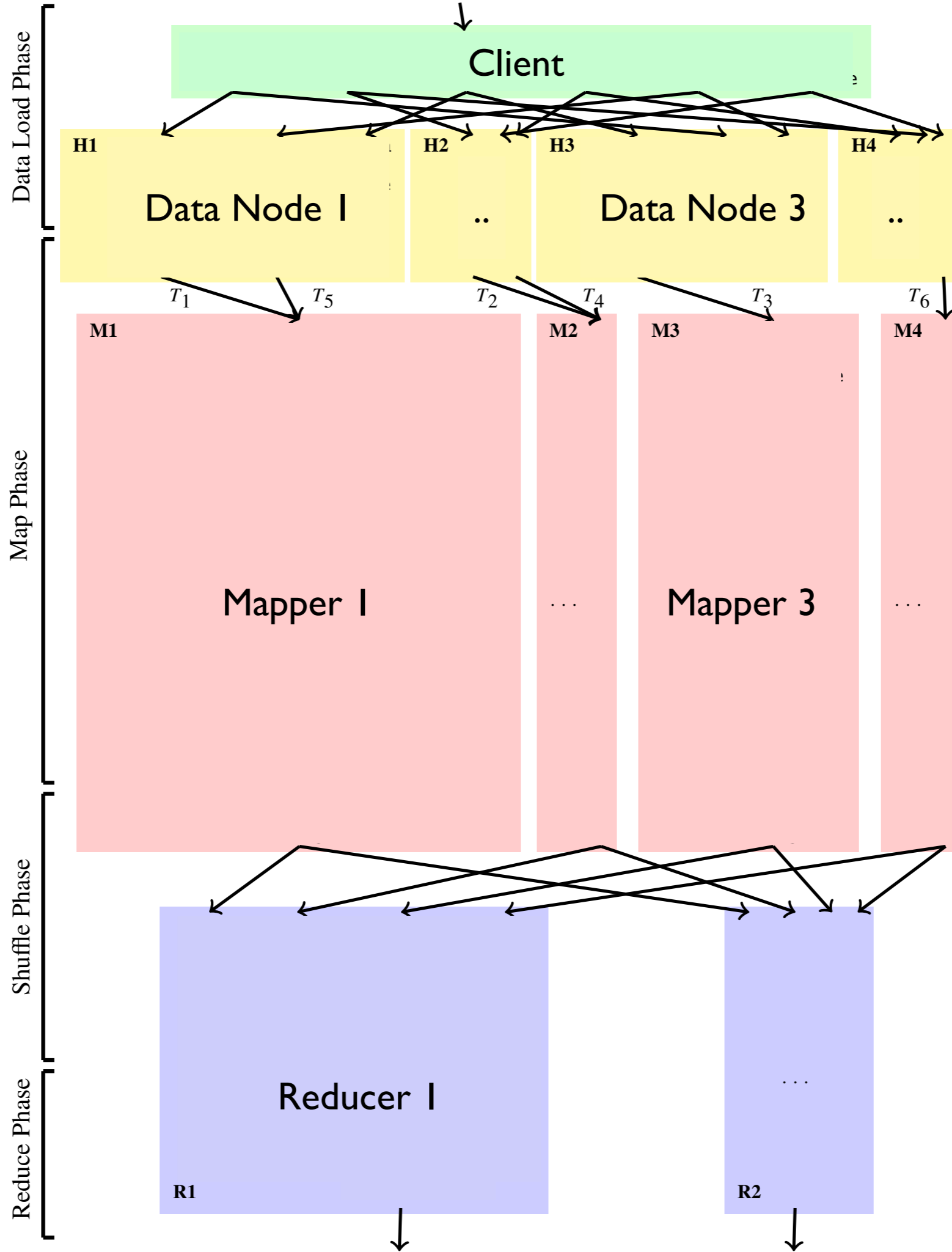
Analytical Query Performance

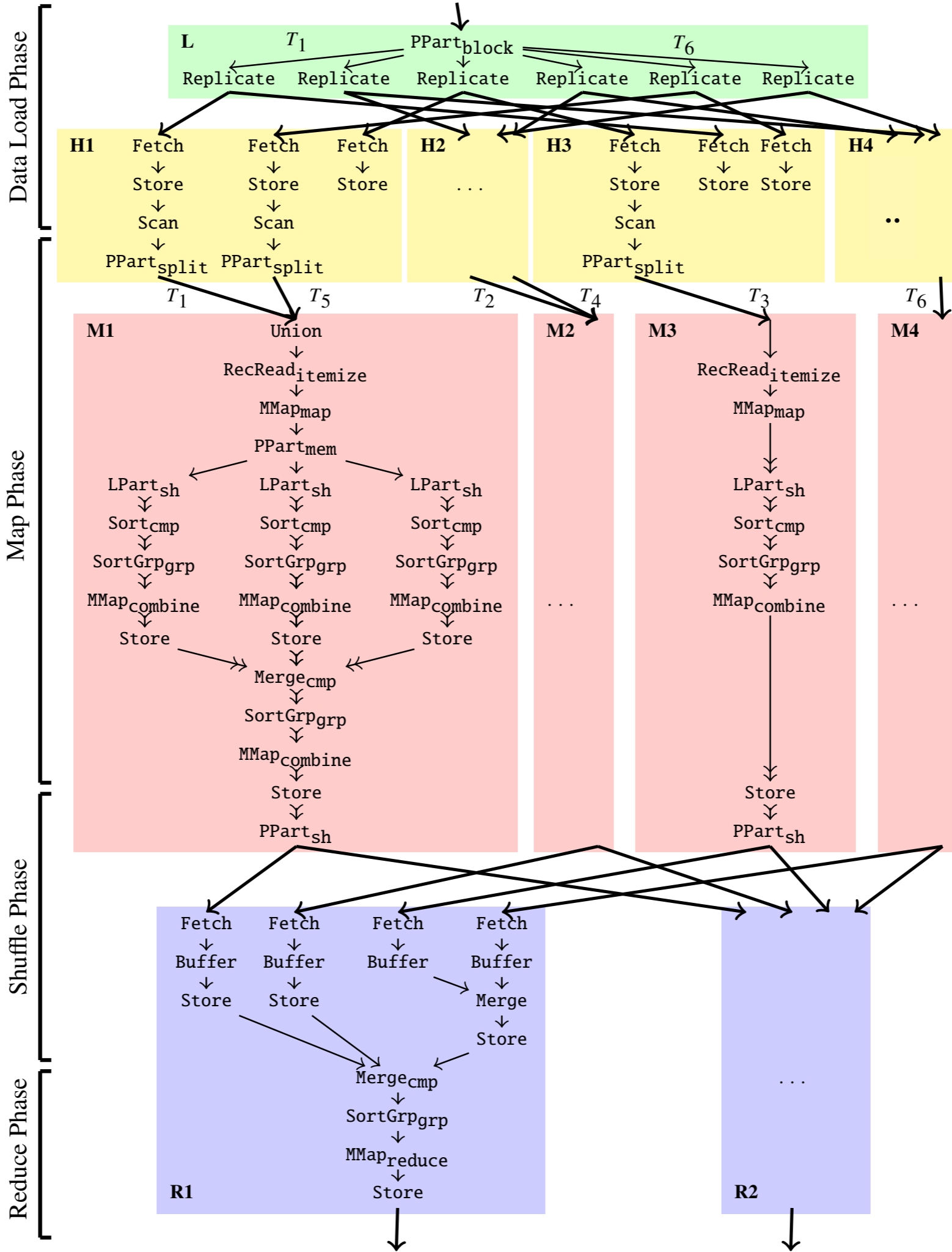
Selection Task

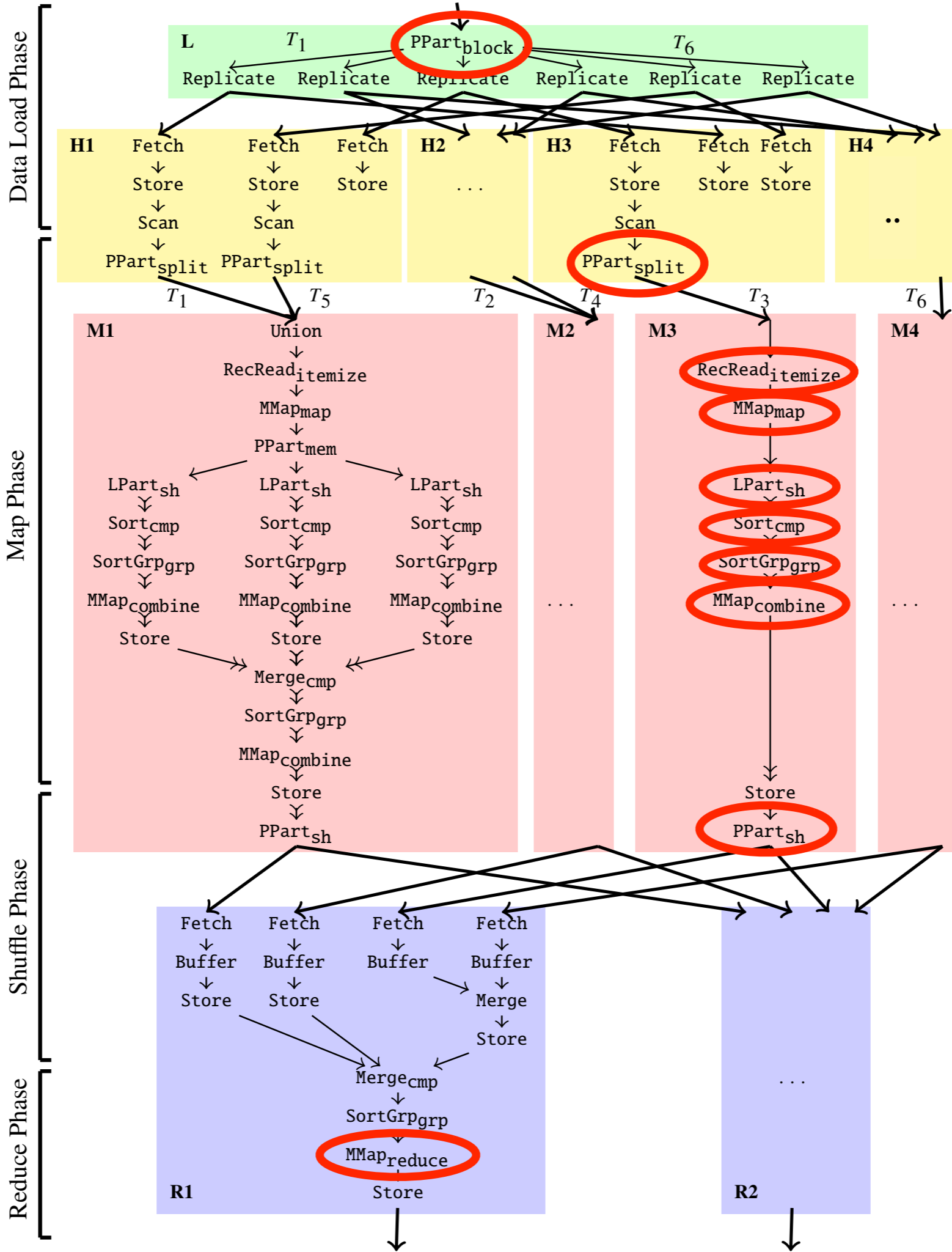


Join Task

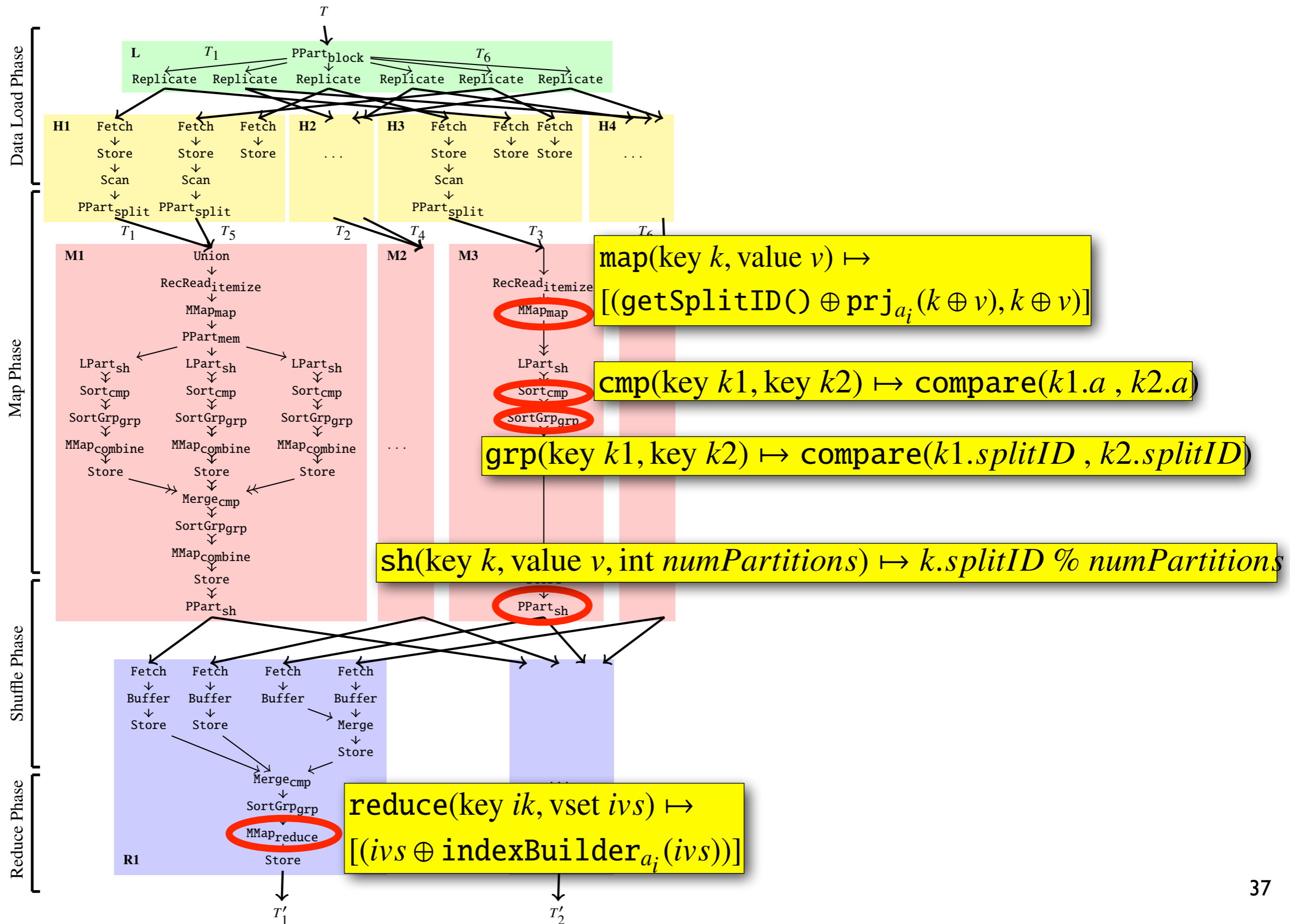




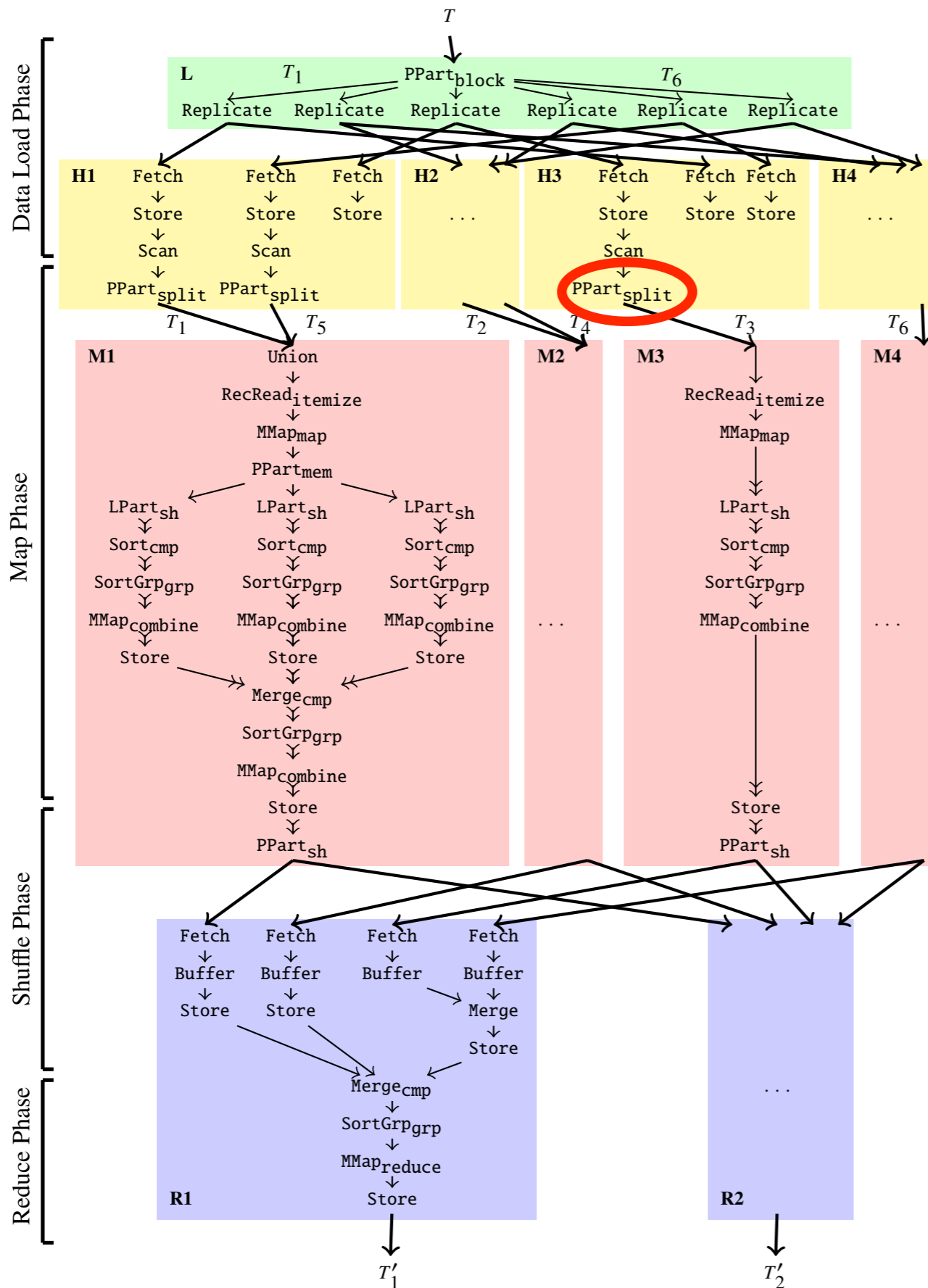




Trojan Index Creation



Trojan Index Access

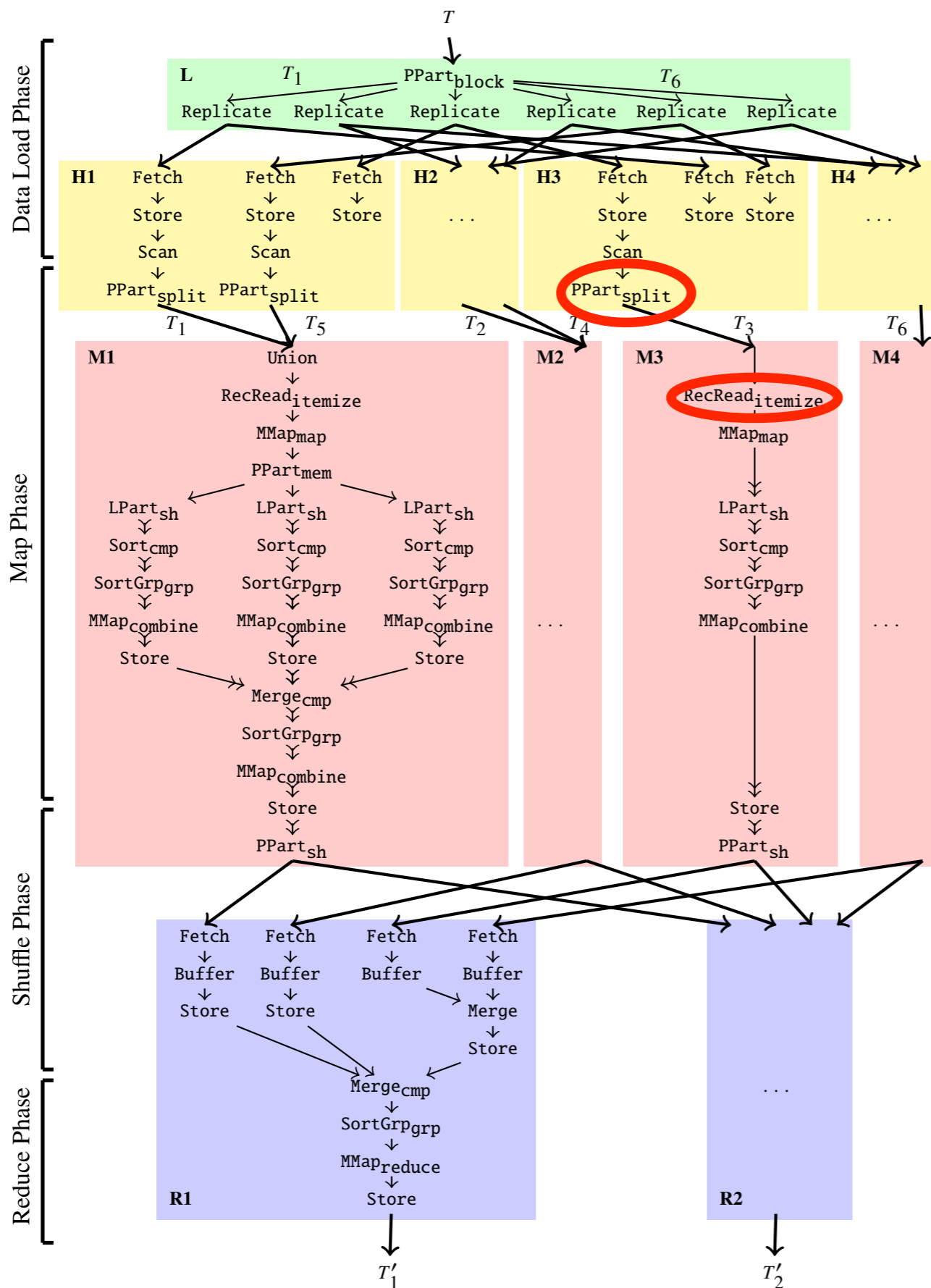


Algorithm 1: Trojan Index/Trojan Join split UDF

```

Input : JobConf job, Int numSplits
Output: logical data splits
1 FileSplit [] splits;
2 File [] files = GetFiles(job);
3 foreach file in files do
4   Path path = file.getPath();
5   InputStream in = GetInputStream(path);
6   Long offset = file.getLength();
7   while offset > 0 do
8     in.seek(offset-FOOTER_SIZE);
9     Footer footer = ReadFooter(in);
10    Long splitSize = footer.getSplitSize();
11    offset -= (splitSize + FOOTER_SIZE);
12    BlockLocations blocks = GetBlockLocations(path,offset);
13    FileSplit newSplit = CreateSplit(path,offset,splitSize,blocks);
14    splits.add(newSplit);
15  end
16 end
17 return splits;
  
```

Trojan Index Access

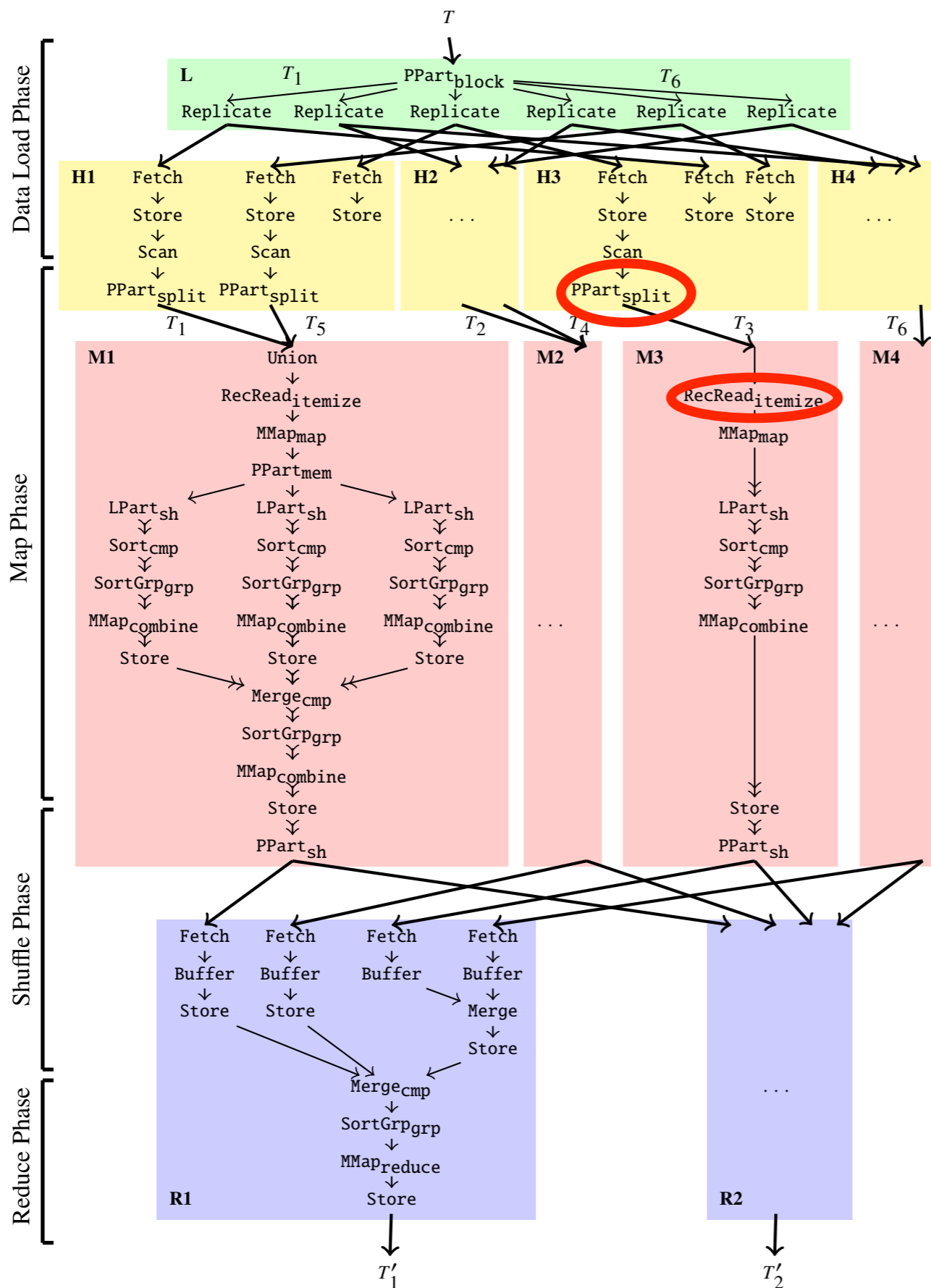


Algorithm 2: Trojan Index `itemize.initialize` UDF

```

Input: FileSplit split, JobConf job
1 Global FileSplit split = split;
2 Key lowKey = job.getLowKey();
3 Global Key highKey = job.getHighKey();
4 Int splitStart = split.getStart();
5 Global Int splitEnd = split.getEnd();
6 Header h = ReadHeader(split);
7 Overlap type = h.getOverlapType(lowKey, highKey);
8 Global Int offset;
9 if type == LEFT_CONTAINED or type == FULL_CONTAINED or type ==
POINT_CONTAINED then
10 | Index i = ReadIndex(split);
11 | offset = splitStart + i.lookup(lowKey);
12 else if type == RIGHT_CONTAINED or type == SPAN then
13 | offset = splitStart;
14 else
15 | // NOT_CONTAINED, skip the split;
16 | offset = splitEnd;
17 end
18 Seek(offset);
    
```

Trojan Index Access



Algorithm 3: Trojan Index `itemize.next` UDF

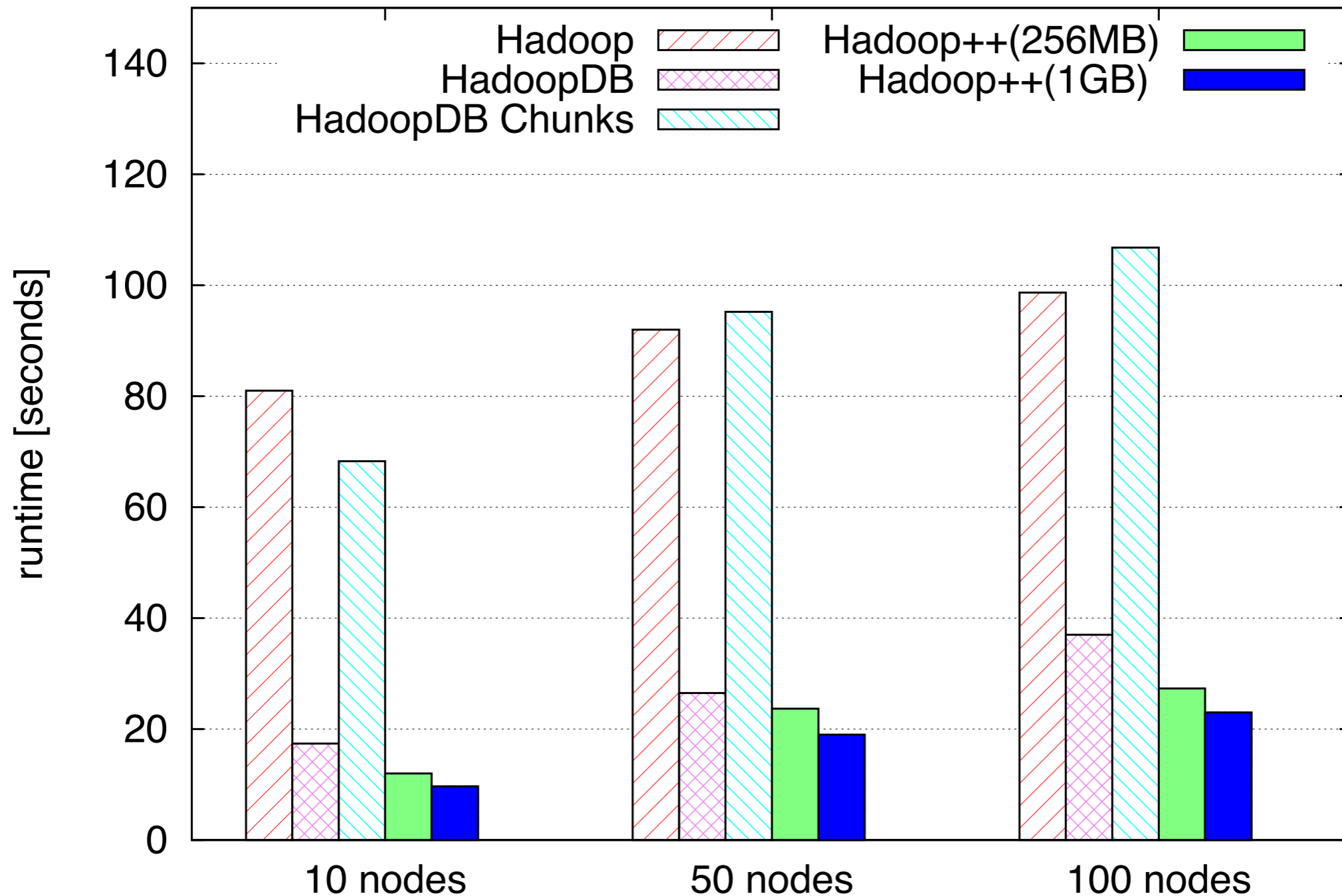
Input : KeyType key, ValueType value
Output: has more records

```

1 if offset < splitEnd then
2   Record nextRecord = ReadNextRecord(split);
3   offset += nextRecord.size();
4   if nextRecord.key < highKey then
5     SetKeyValue(key, value, nextRecord);
6     return true;
7   end
8 end
9 return false;

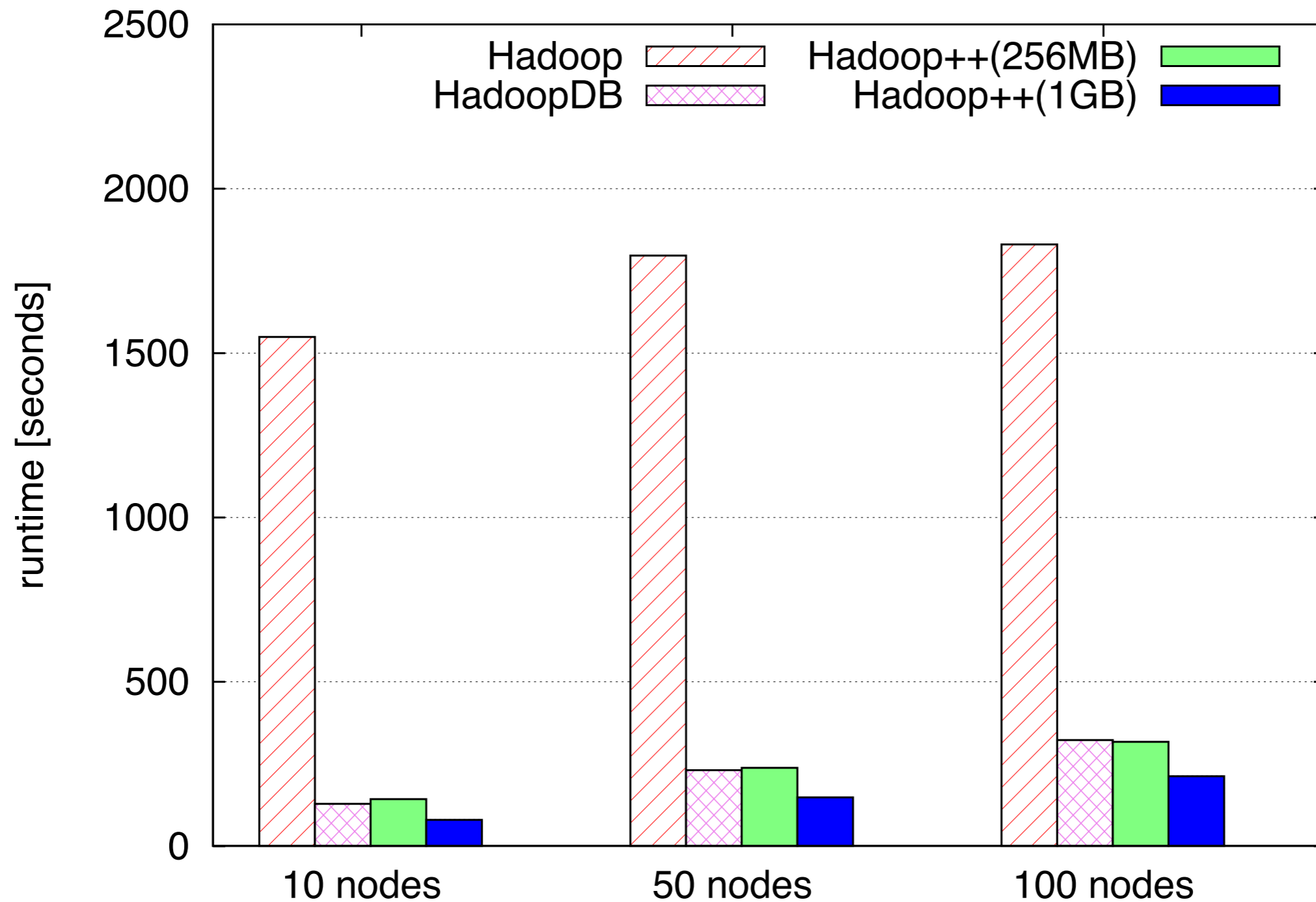
```

Selection Analytical Task *



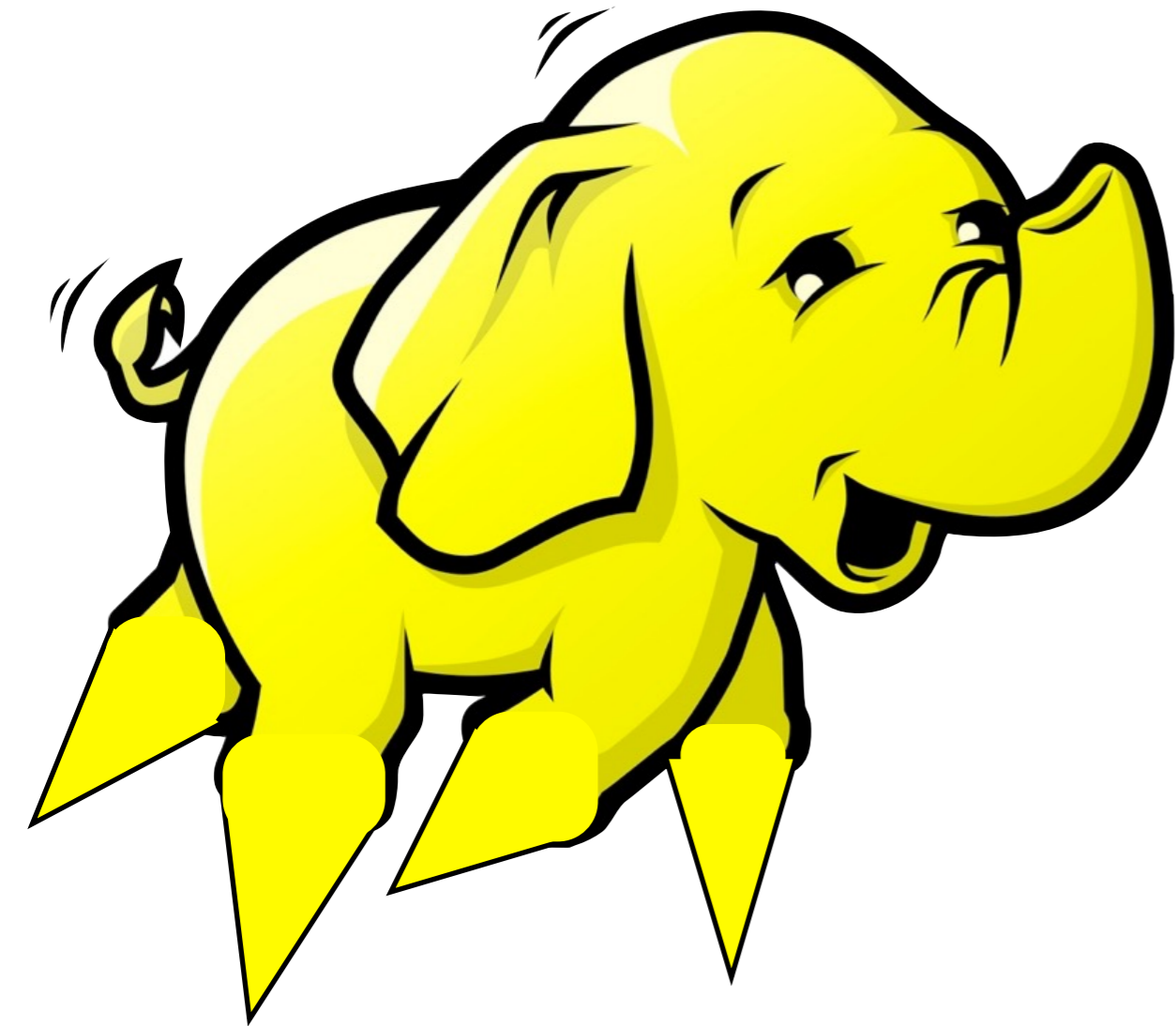
* Pavlo et. al. A Comparison of Approaches to large-Scale Data Analysis. SIGMOD 2009

Join Analytical Task *

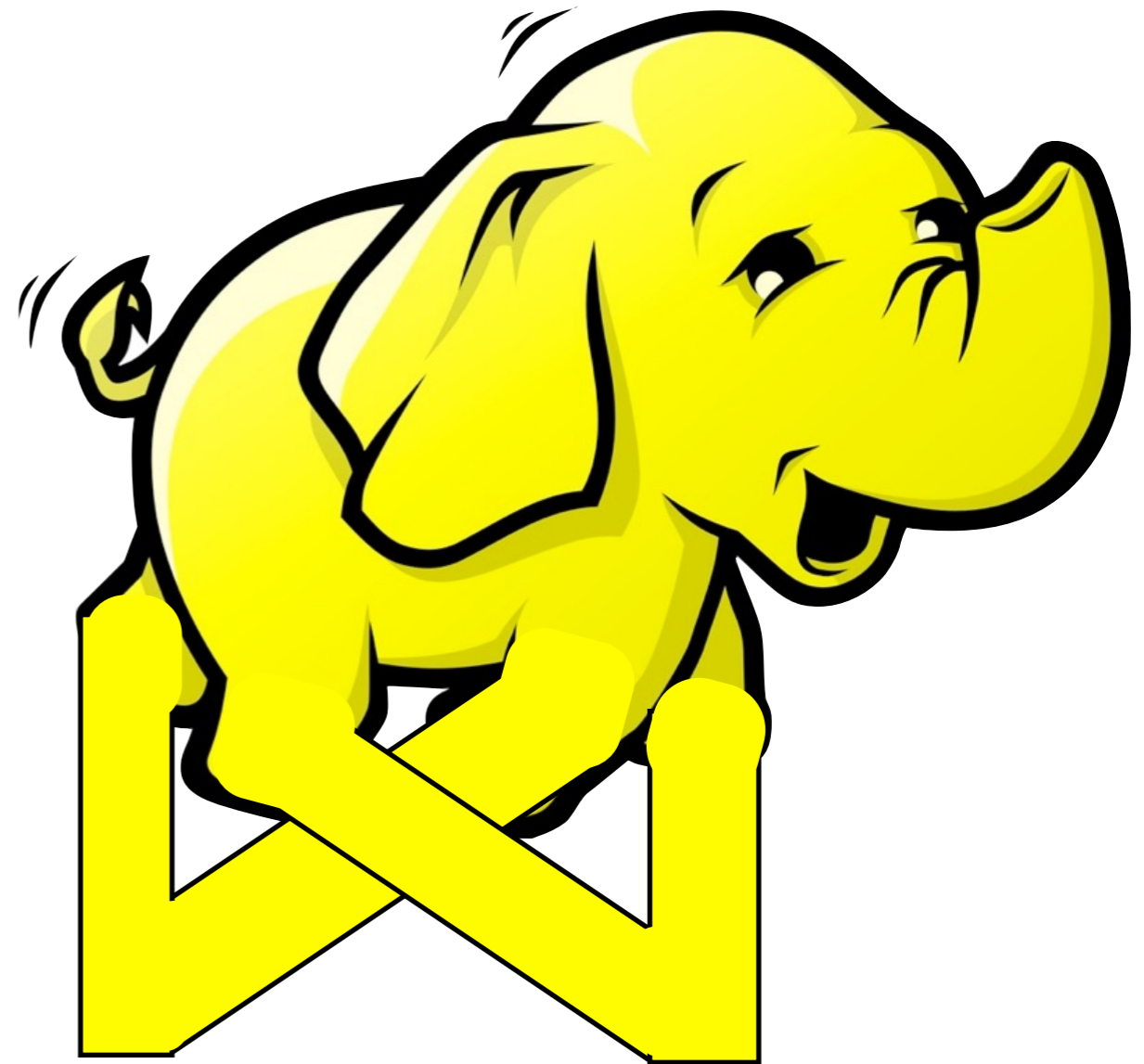


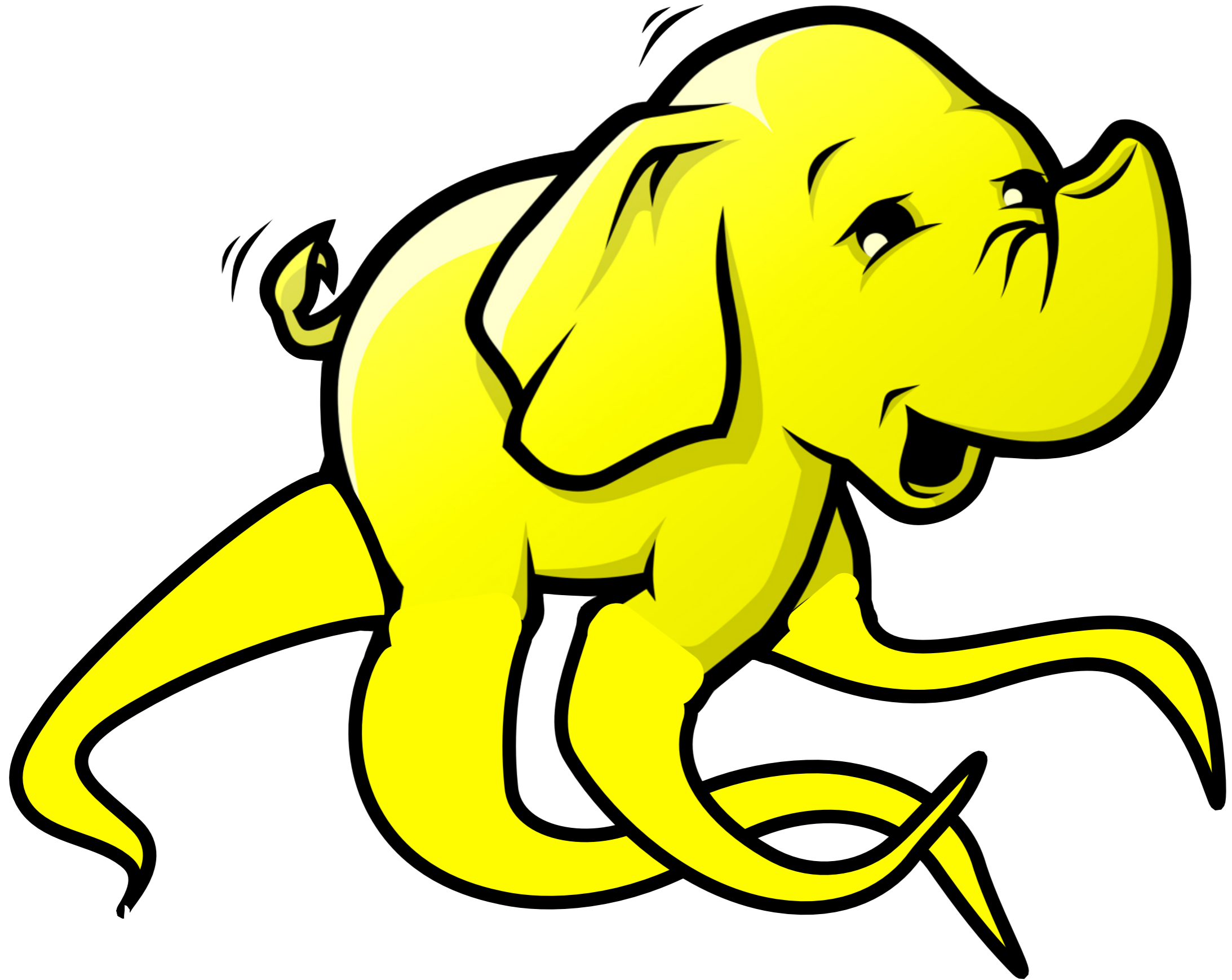
* Pavlo et. al. A Comparison of Approaches to large-Scale Data Analysis. SIGMOD 2009

Trojan Index



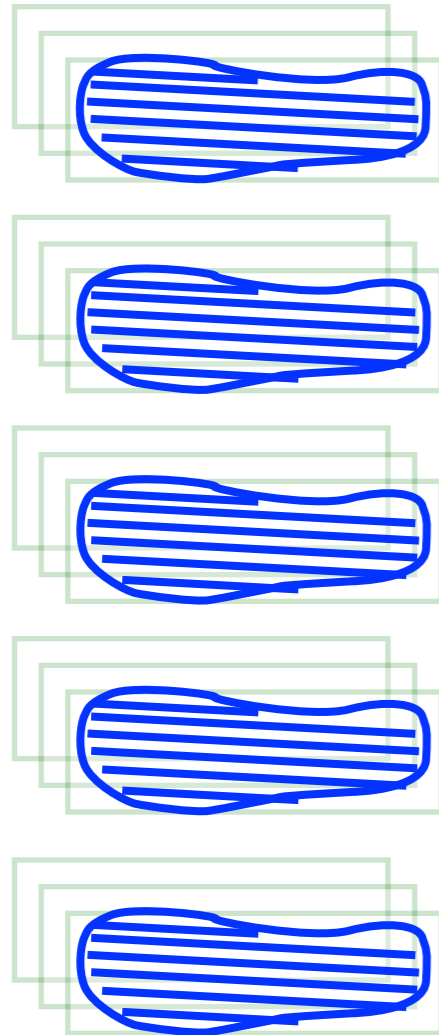
Trojan Join



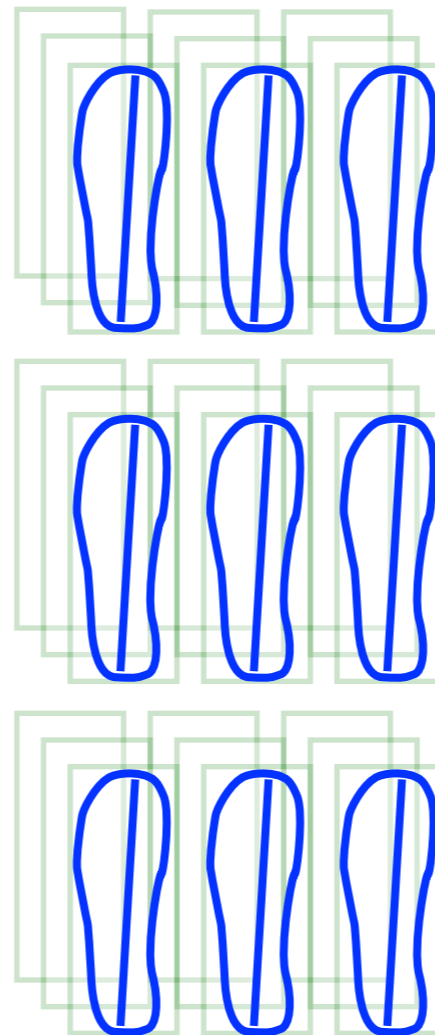


Traditional Layouts

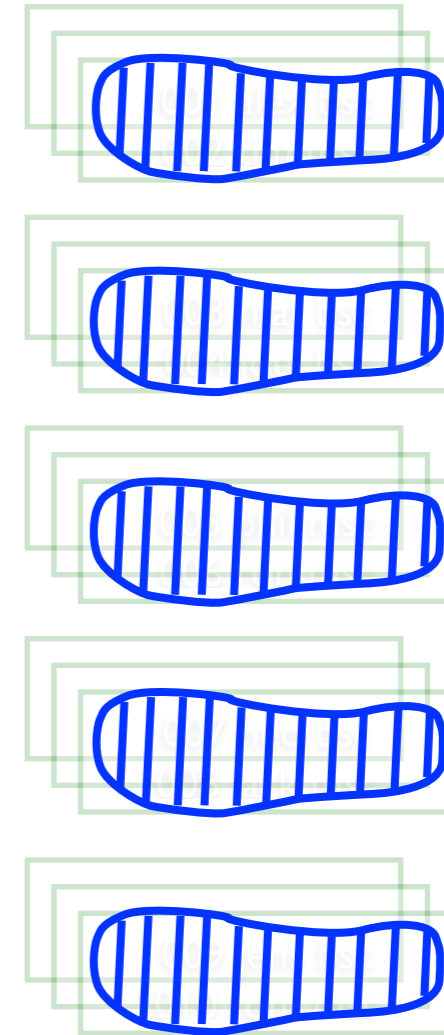
Row
(default)



Column*















PAX**



* A. Floratou et al. Column-Oriented Storage Techniques for MapReduce. PVLDB, April, 2011

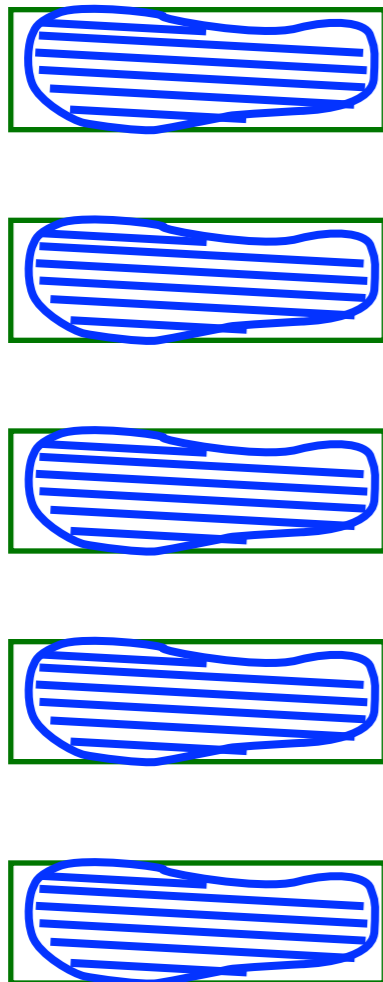
** Y. He et al. RCFile: A fast and space-efficient data placement structure in MapReduce-based warehouse systems. ICDE, 2011

Traditional Layouts

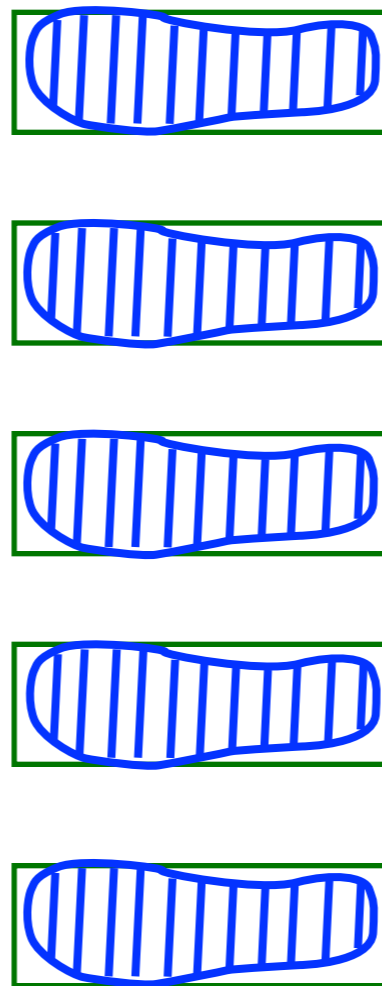
	Row	Column	PAX
Non-required Reads			
Network Costs			
Data Block Placement			
Tuple Reconstruction			

Trojan Data Layouts

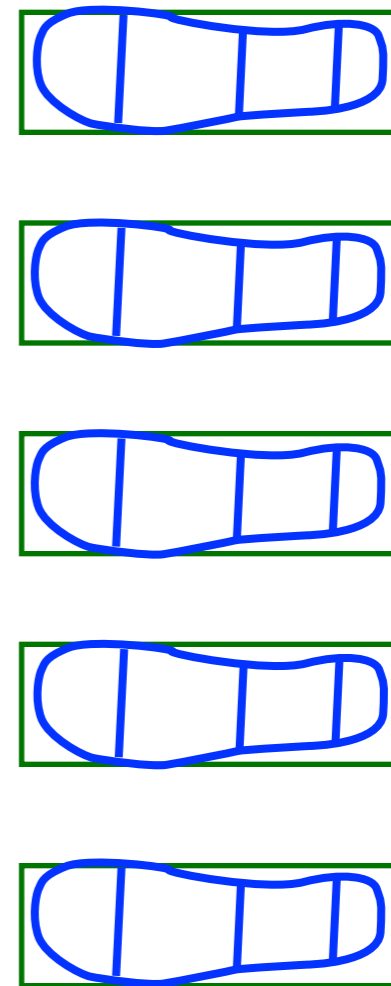
Replica 1



















Replica 2



Replica 3



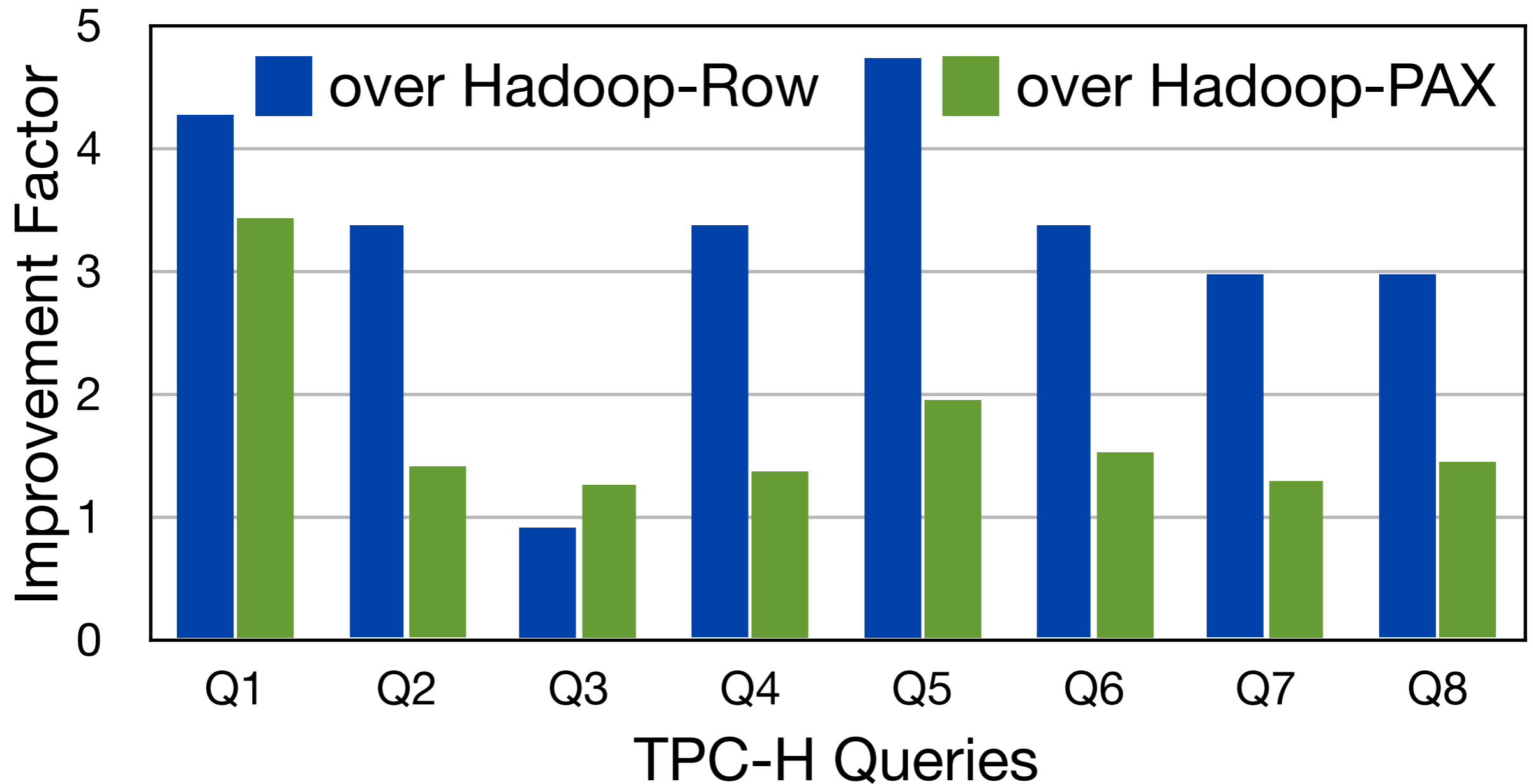
Trojan Data Layouts

	Row	Column	PAX	Trojan
Non-required Reads				
Network Costs				
Data Block Placement				
Tuple Reconstruction				

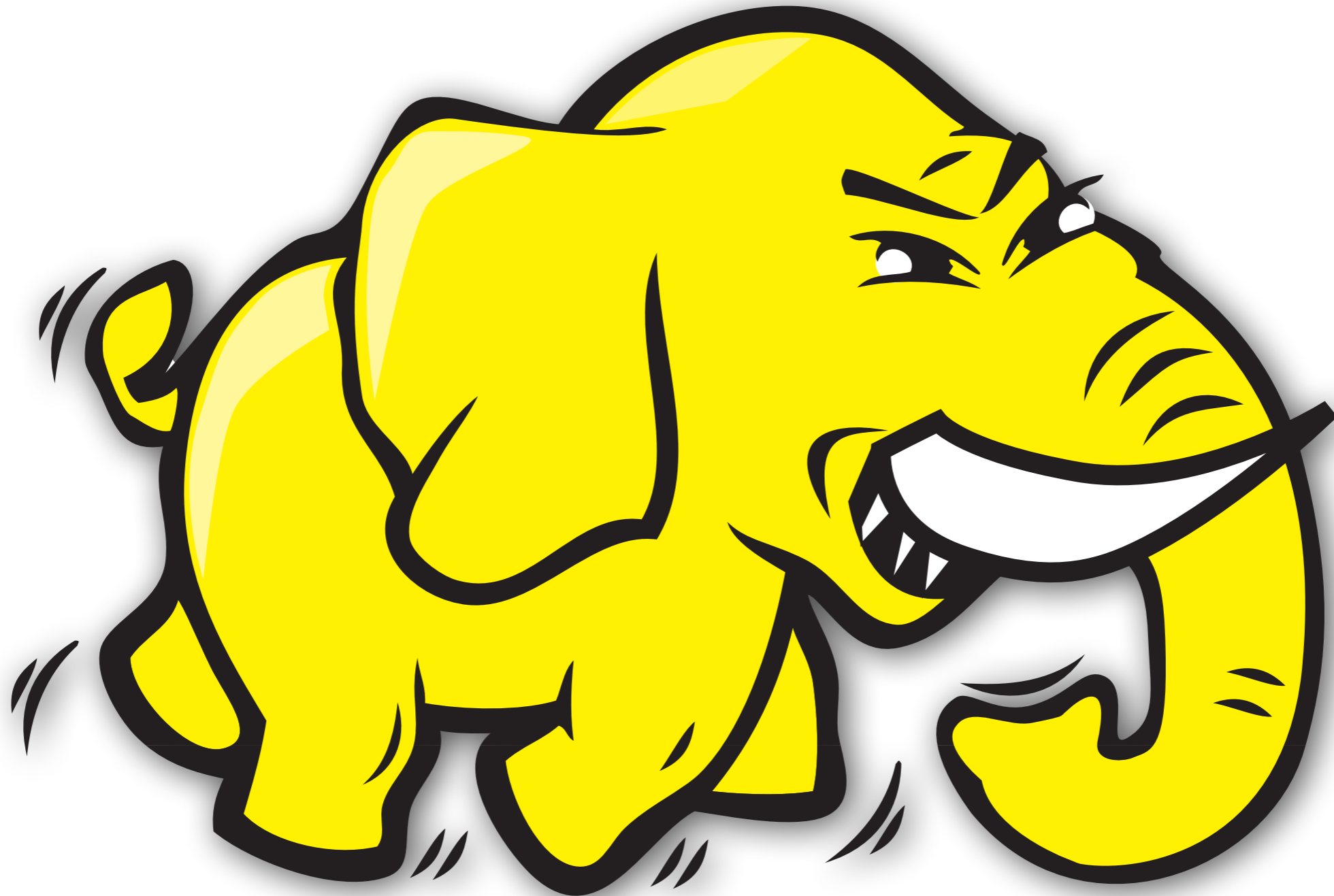
Layout Quality

	#Non-required Attributes Read	#Joins in Tuple Reconstruction
HADOOP-ROW	525	0
HADOOP-PAX	0	139
Trojan Layout	14	20

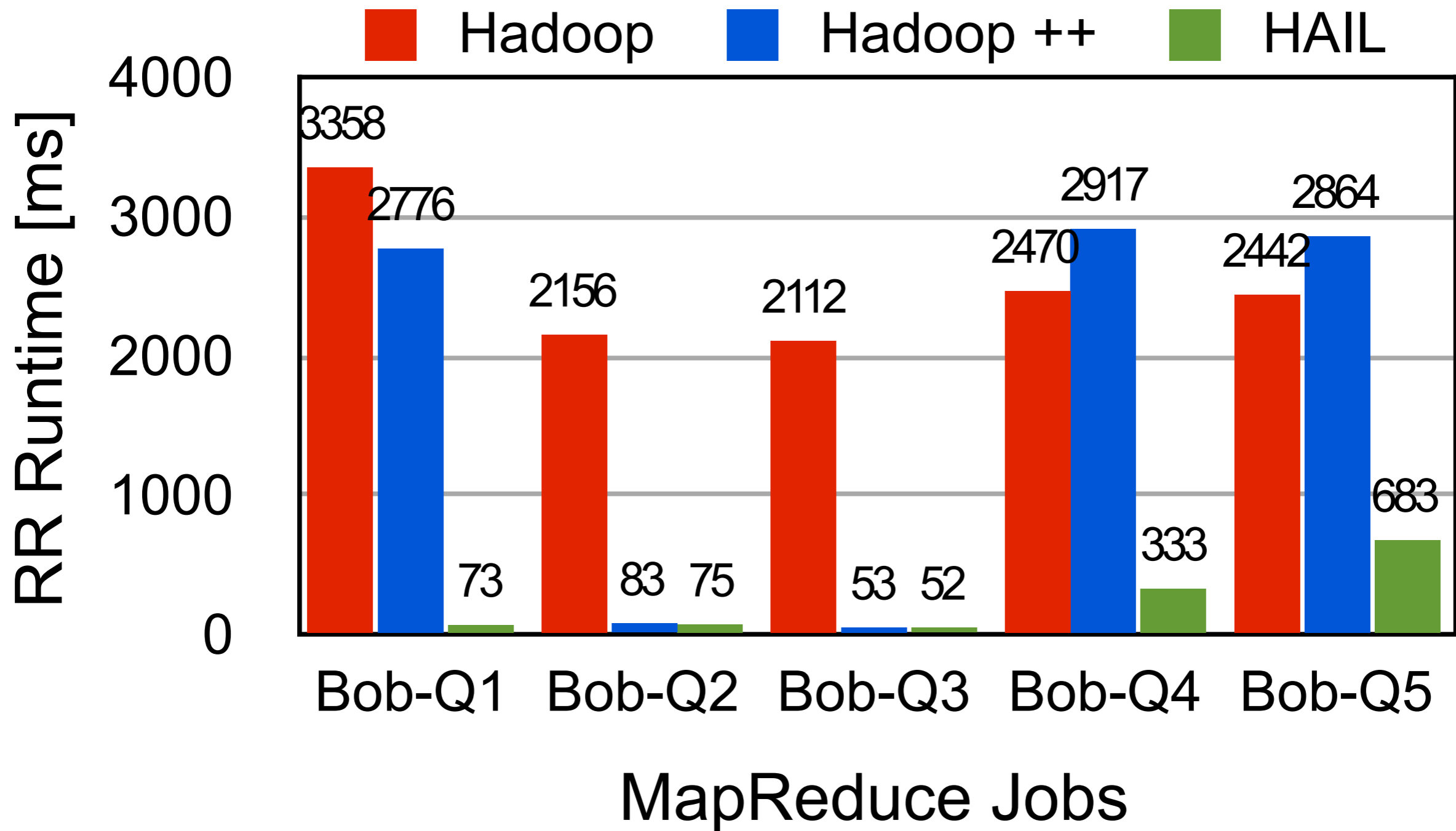
Projection Analytical Task



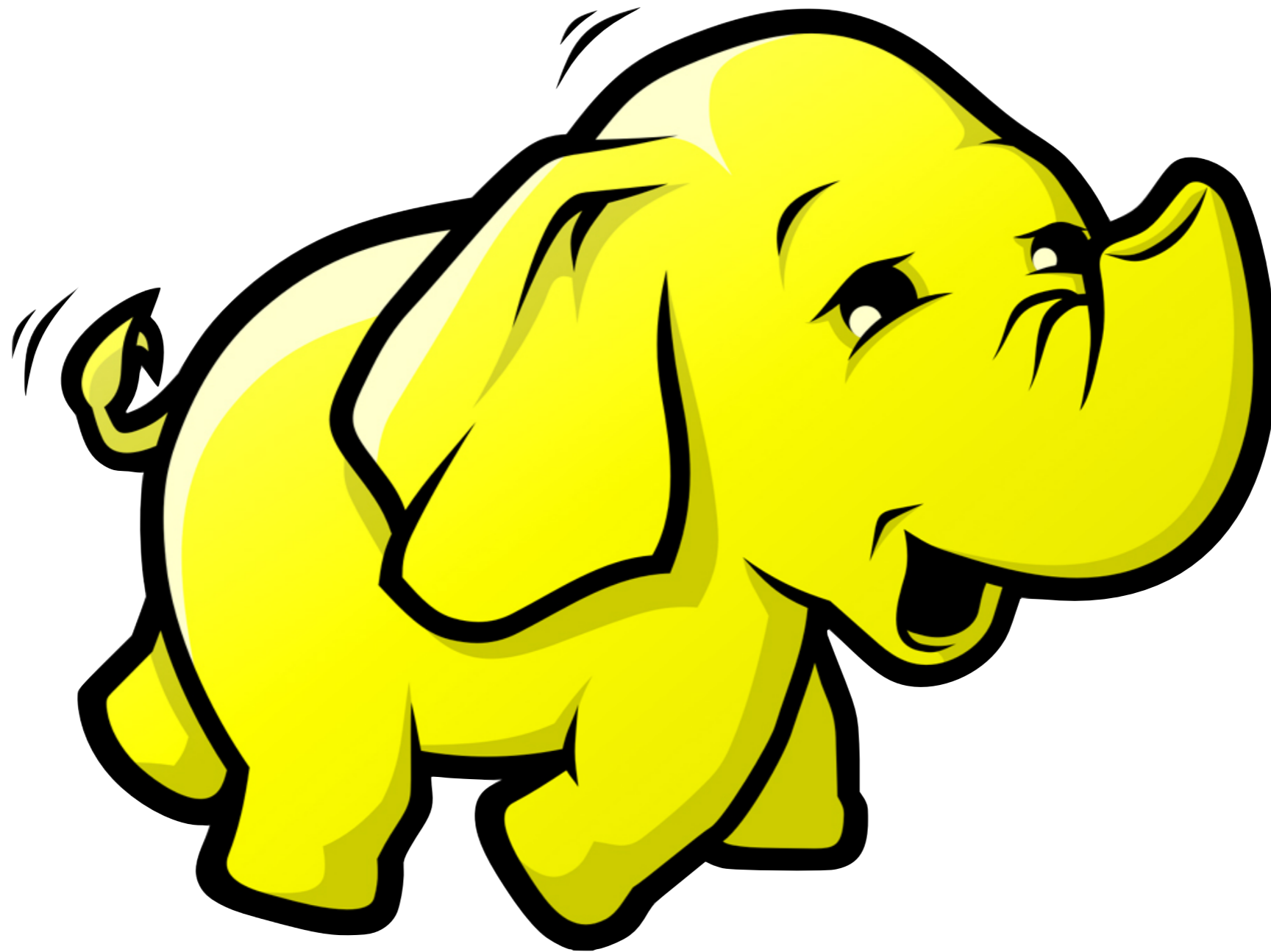
Hadoop Aggressive Indexing Library



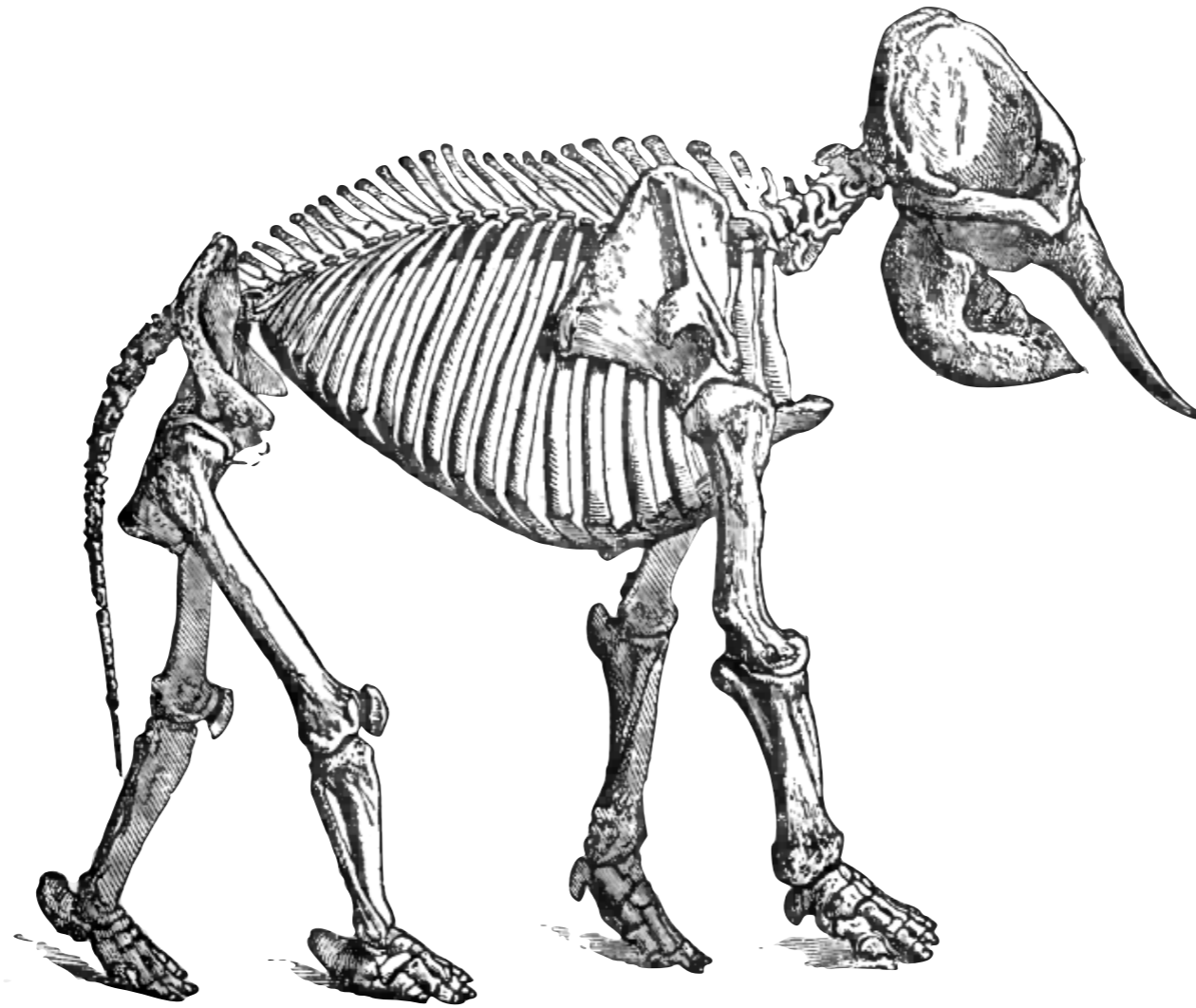
Individual Jobs: Weblog, RecordReader



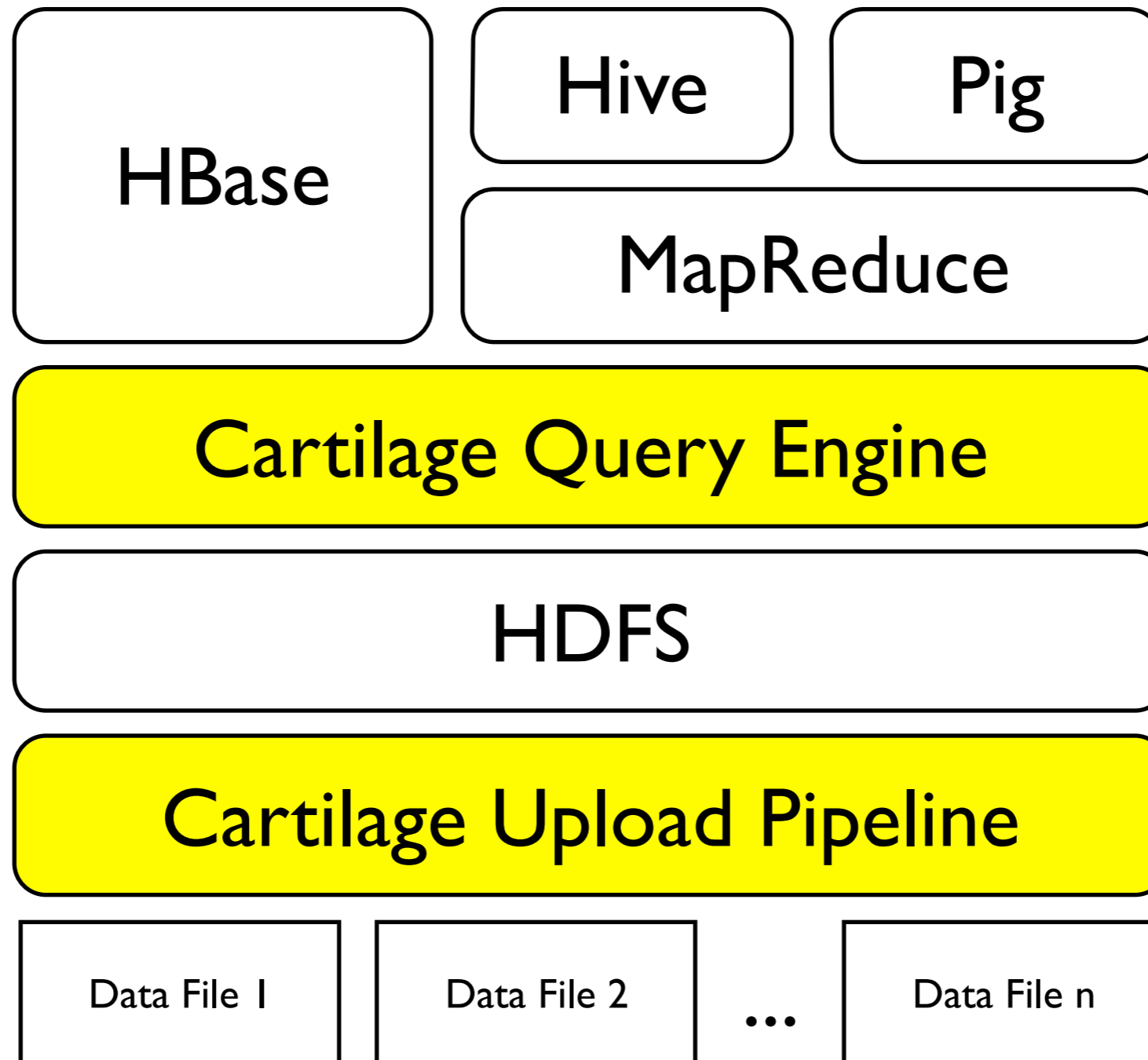
Cartilage



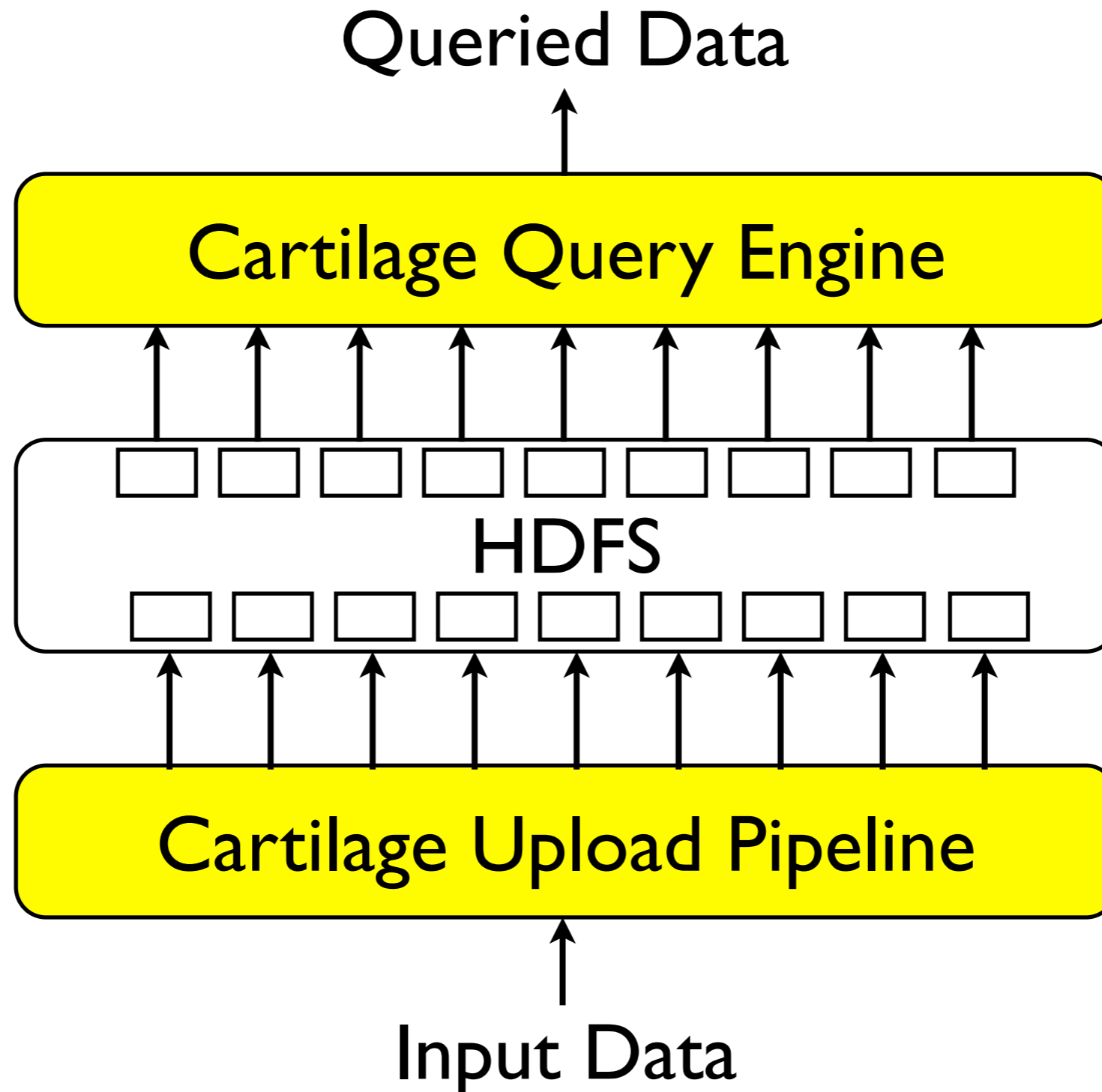
Cartilage



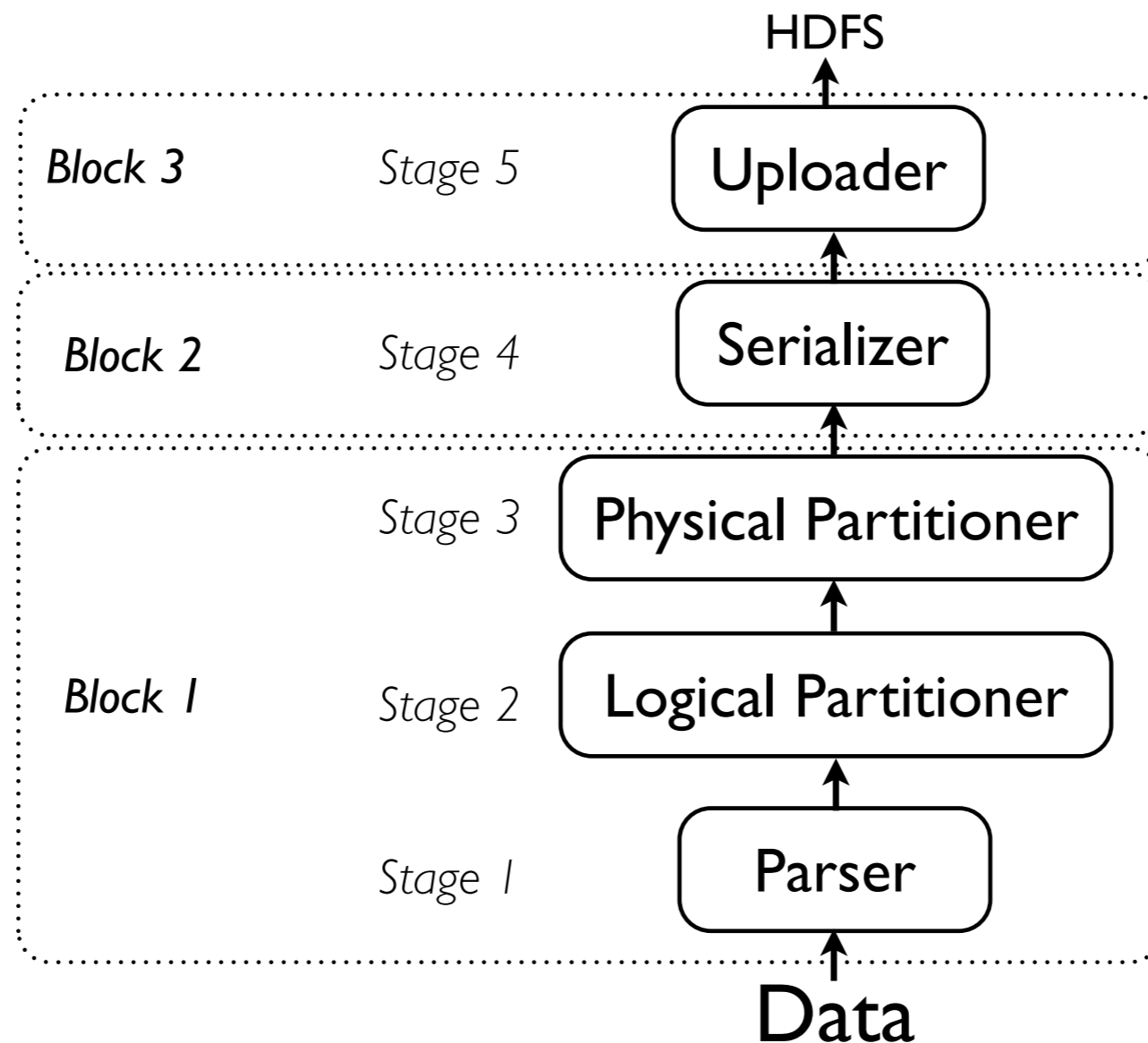
Hadoop Stack



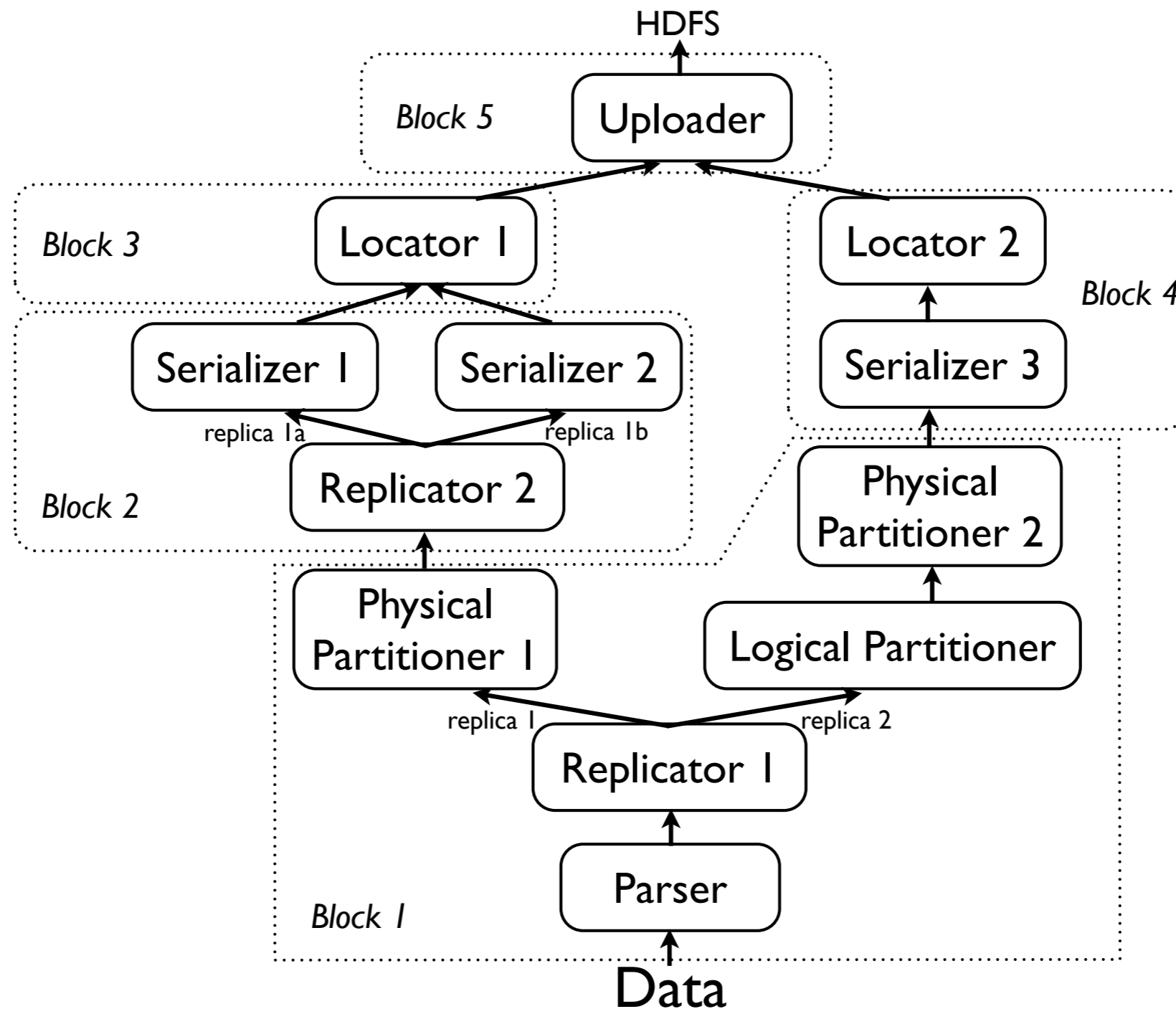
Hadoop Stack



Upload Plans

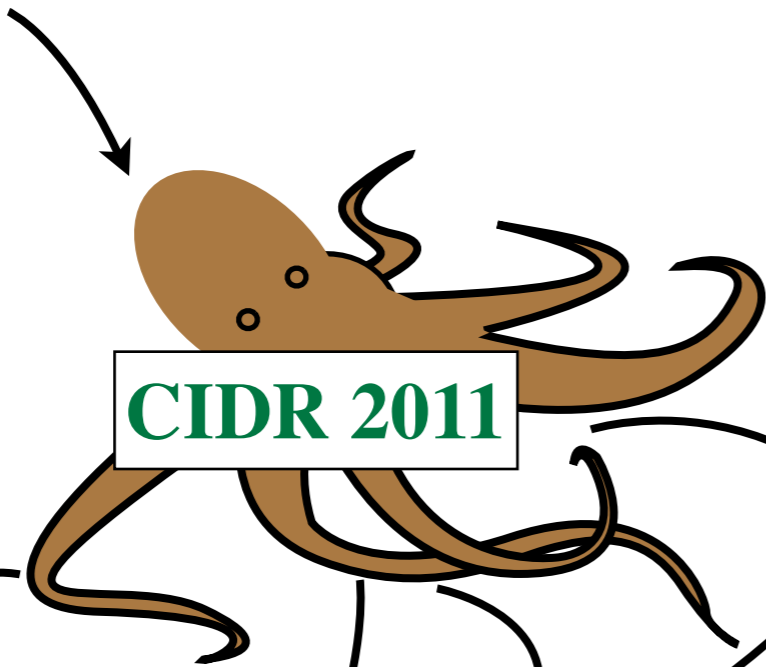


Upload Plans

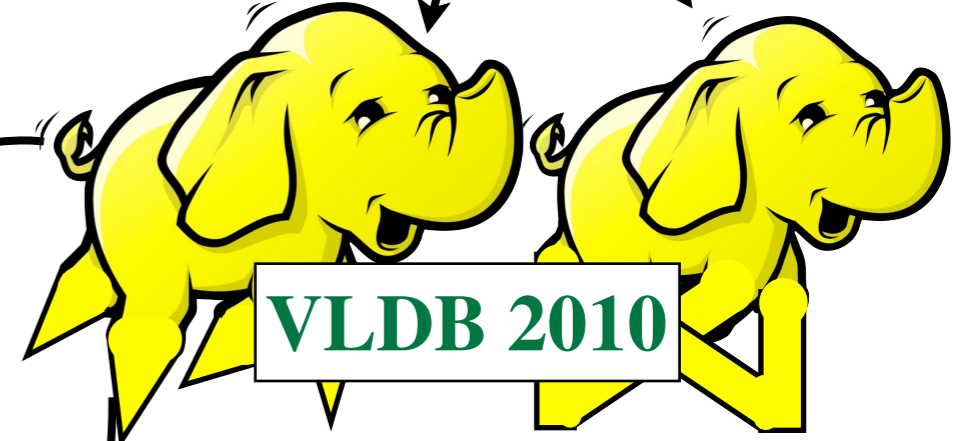


Summary

ONE SIZE DOES NOT FIT ALL



CIDR 2011

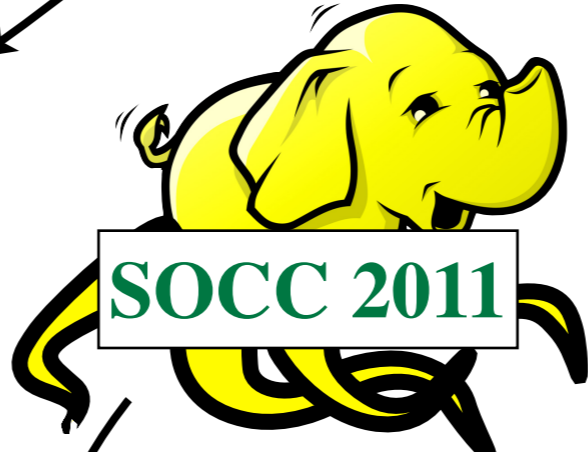


VLDB 2010

WVHow! Layer

CIDR 2013

HYRISE



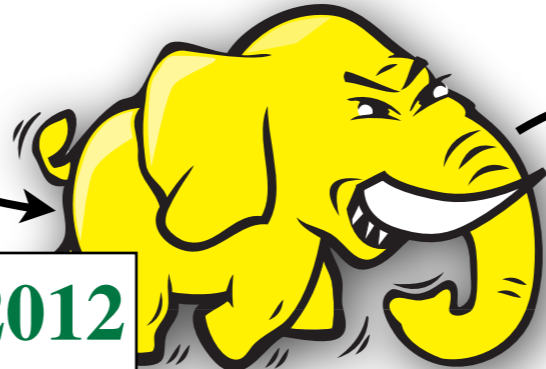
SOCC 2011



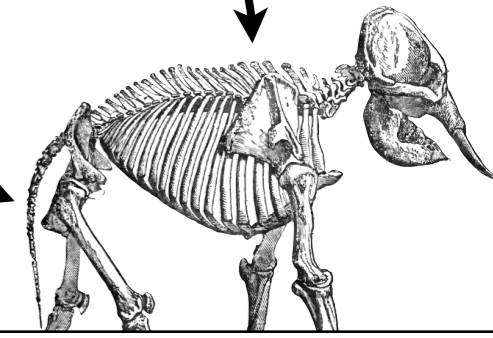
CIDR 2013



VLDB 2013



VLDB 2012



SIGMOD (demo)

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