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MASSACHUSETTS INSTITUTE OF TECHNOLOGY Department of Electrical Engineering and Computer Science **6.821** Programming Languages

*Lisp Reference Manual

Version 5.0 October 1988

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Preface

Objectives of This Manual

The *Lisp Reference Manual describes the essential constructs of the *Lisp language and explains the concepts used in programming the Connection Machine in *Lisp.

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Intended Audience

The reader is assumed to have a working knowledge of Common Lisp, as described in Common Lisp: The Language, and a general understanding of the Connection Machine system. The Connection Machine Front-End Subsystems provides the necessary background information on the Connection Machine system.

Revision Information

This manual is a revision of the *Lisp Reference Manual, Version 4.0, published October 1987. This revision corrects information presented in that version and updates descriptions to account for the implementation of *Lisp Version 5.0. It does not fully describe *Lisp Version 5.0; the Supplement to the *Lisp Reference Manual provides information on language features new with Version 5.0.

Overview of Manual

Chapter 1. Introduction

Chapter 2. Overview of *Lisp

This chapter provides an overview of the Connection Machine computer and of *Lisp, including:

- A description of the language's basic concepts, such as parallel variables and the selection of particular processors
- The parts of a typical *Lisp program
- Code examples that illustrate using parallel variables and processor selection, defining parallel functions, and performing interprocessor communication

Preface

Names that stand for pieces of code (metavariables) appear in italics. In function or macro descriptions, the names of the arguments appear in italics.

Argument names can restrict the type of an argument; argument names that end in the suffix *pvar* must be parallel variables. For example, the name *integer-pvar* restricts the argument to a parallel variable whose fields in the currently selected set of processors must all contain integers.

Braces followed by a star (as in {symbol}*) are used as in Common Lisp: The Language to indicate the symbol may appear zero or more times.

Related Documents

- *Lisp Release Notes, Version 5.0 The current release notes provide a succinct overview of the changes made to to *Lisp since the release of Version 4.3. These are essential reading.
- Supplement to the *Lisp Reference Manual, Version 5.0 This manual updates the *Lisp Reference Manual, adding descriptions of all features new with the release of *Lisp Version 5.0.
- *Lisp Compiler Guide, Version 5.0 This manual describes the current implementation of the *Lisp compiler.
- Connection Machine Front-End Subsystems The manuals in this volume should be read before the *Lisp Reference Manual. It explains the configuration of the Connection Machine system, and how to access the Connection Machine from a front-end computer.
- Connection Machine Parallel Instruction Set The *Lisp Reference Manual explains how to call Paris from *Lisp. Users who wish to do so should refer also to the Paris manual.
- Common Lisp: The Language by Guy L. Steele Jr. Burlington, Mass.: Digital Press, 1984.

This book defines the de facto industry standard Common Lisp.

Customer Support

Thinking Machines Customer Support encourages customers to report errors in Connection Machine operation and to suggest improvements in our products.

When reporting an error, please provide as much information as possible to help us identify and correct the problem. A code example that failed to execute, a session transcript, the record of a backtrace, or other such information can greatly reduce the time it takes Thinking Machines to respond to the report.

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To contact Thinking Machines Customer Support:

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For Symbolics users only:

The Symbolics Lisp machine, when connected to the Internet network, provides a special mail facility for automatic reporting of Connection Machine system errors. When such an error occurs, simply press CTRL-M to create a report. In the mail window that appears, the To: field should be addressed as follows:

To: bug-connection-machine@think.com

Please supplement the automatic report with any further pertinent information.

Chapter 1

Introduction

*Lisp (pronounced star lisp) is a data parallel language designed to program the Connection Machine.

A Connection Machine (CM) data parallel computer consists of a large number of simple processors. Each has some associated local memory and is integrated into a highly connected communications network. A CM configuration can have up to 65,536 processors, each with 4K bits (in model CM-1) or 64K bits of memory (in model CM-2). Typical applications use data types that have components spanning many Connection Machine processors. *Lisp provides the means for creating and manipulating these parallel types.

*Lisp is an extension of Common Lisp. *Lisp adds a new data structure and extensions of many Common Lisp functions that execute in many Connection Machine processors in parallel.

*Lisp has several important features:

- Many *Lisp language features map directly onto Connection Machine instructions; therefore, users quickly develop an intuition for predicting program performance.
- *Lisp includes an interface that permits direct calls to Paris from within a *Lisp program.
- A *Lisp compiler is provided (described in the *Lisp Compiler Guide).
- A *Lisp simulator is provided for preliminary program testing and debugging. Executing entirely on a serial front end, it simulates the Connection Machine operations.

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Chapter 2 Overview of *Lisp

This chapter introduces the main concepts and terminology of *Lisp. It then provides a brief overview, with code examples, of *Lisp operations. All operations appearing in this chapter are described more fully in subsequent chapters.

The primary concepts of the *Lisp language follow:

- *Lisp programs execute on a front-end computer, typically a Symbolics Lisp machine or a Digital Equipment Corporation VAX. As a side effect of running the *Lisp program, the front-end computer generates Paris instructions for the Connection Machine processors to execute. Every so often, the front-end computer transfers data or results of computations to or from the Connection Machine.
- *Lisp programs refer to and manage memory in the Connection Machine processors through Lisp objects called *pvars* (pronounced *pee-vars*, for *parallel variables*). These objects contain information about memory locations in the Connection Machine processors and the possible types of values stored at those locations. A pvar looks like a large vector of Common Lisp values, with each value stored in a different Connection Machine processor. These values may be integers, floating-point numbers, booleans, or any other Lisp object (for example: 0, 102, -5, 10.333, t, nil, hi-there). As with Common Lisp, the *Lisp programmer need not be concerned with type coercion, since it is done automatically.

 *Lisp programs control the set of Connection Machine processors that are executing instructions. This set may range from all processors in the machine to none of the processors.

Given these concepts, a *Lisp program typically consists of these parts:

- Permanent Connection Machine storage declarations.
- Code for creating static data structures in Connection Machine memory. This
 often involves substantial transfers of data from the front-end computer to the
 Connection Machine.

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- **Pvar Component** A single instance among the collection of values represented by a pvar. A pvar component may be any Lisp value; the set of components represented by a pvar need not be all of the same data type.
- Field The memory occupied by all the components of a pvar. A field is a string of contiguous bits in the same memory location in *each* processor. The components of a pvar all occupy the same amount of memory (even if they are of different types), and they are all stored at the same memory address in the respective processors.
- Pvar Contents The set of values (components) represented by a pvar. These values are stored in the field in the Connection Machine that is described by the pvar.
- Currently Selected Set Most *Lisp operations are only carried out in a subset of the Connection Machine processors. This subset is called the *currently selected set* and is specified by using *Lisp special forms, such as *all, *when, *cond, and *if.
 - The names of functions and macros that return pvars as their values end with !!. This suffix, pronounced *bang-bang*, is meant to look like two parallel lines. We recommend that user-defined functions follow this convention (although nothing enforces it), because it helps ensure that pvars are produced only in contexts where they can be used. It is an error to produce pvars in contexts where they cannot be used (see chapter 8).

There are a few *Lisp macros whose names do not end in !!, such as *when, *all and *let, that, nonetheless, may optionally return a pvar.

All *Lisp functions that perform parallel computation and do not end in !! begin with * (pronounced *star*); hence, the name *Lisp.

Parallel Equivalent of This phrase is used to describe a *Lisp function with reference to a Common Lisp function. For example, "mod!! is the parallel equivalent of the Common Lisp mod." This means that mod!! performs the same calculation as mod, only mod!! performs the operation in parallel using each component of an argument pvar.

2.2 *Lisp Concepts

This section contains sample code that illustrates some common *Lisp expressions. All the functions used are described fully in later sections of this manual. The above sets the contents of pvar **a** to the sum of the contents of pvar **b** and pvar **c**. Notice that because **c** contains floating-point values, the integers contained in **b** are properly coerced to a floating-point value, and the result is a floating-point value as well.

((((d à []+) # ####")

(Citd a 11+) a 202")

i, 2, or 3 after dividing by 4, the following might be use

Expressions can be nested:

(*set a (-!! b (*!! a (!! 2))))

This sets **a** to the difference of **b** and 2 times **a**. This simple expression causes thousands of operations to go on simultaneously.

2.2.2 Predefined Pvars

Two pvars are predefined in *Lisp. The pvar t!! contains the Lisp symbol t in each processor; the symbol t!! is equivalent to (!! t). Similarly, the pvar nill! contains the Lisp symbol nll in each processor.

2.2.3 Selection

When the Connection Machine is initialized, every processor will be in a state to execute all *Lisp instructions in parallel. However, it may be necessary to execute instructions in some subset of the Connection Machine processors.

One way of temporarily selecting a subset is to wrap the *when macro around a body of forms. For example, to select the set of all processors whose cube addresses (contained in the pvar returned by the function self-address!!) end in 1, one might use the following:

```
;; Create a pvar that is True in all odd processors
(*defvar odd-address-p
                     (=!! (!! 1) (mod!! (self-address!!) (!! 2))))
;; Now select all processors with odd cube addresses
(*when odd-address-p ...)
```

In another case, it may be desirable to perform an operation in the processors in which the cube address is even and the pvar a contains zero. Two natural ways to do this are to (1) use the logic functions to select the correct set:

2.2.4 *Defun

To define functions that can take pvars as arguments or return them as values, use *defun instead of defun. For example, to define a function that takes two pvar arguments and returns their sum, difference, product, or quotient (depending on whether the processor's address has remainder 0, 1, 2, or 3 when divided by 4 in all processors in the currently selected set), use the following:

MEWS for short). It is possible to sum the value contained in a in

This may now be used like any other !! function, as in:

(*set a (four-function!! (+!! a (!! 4)) (-!! a b)))

To pass a *Lisp function as an argument, use *funcall. For example, the following:

```
(defun *compose (*f *g x)
        (*funcall *f (*funcall *g x)))
(*set a (*compose 'sqrt!! '1+!! (!! 8)))
```

acts like:

```
(*set a (sqrt!! (1+!! (!! 8))))
```

2.2.5 Communication

This section demonstrates how to cause the processors to communicate with one another.

rate advantage of this flexibility. In addition, it is desirable to write software that mus

One connectivity pattern that can be specified upon initialization is a two-dimensional grid in which each processor has a neighbor on the north, east, west, and south (or

PRUMPIN POWERSHIP NETT.

The Pvar Data Structure

The basic abstraction in *Lisp is the pvar. A pvar is a Lisp object that references a field of memory in the Connection Machine system. It contains everything necessary to describe the field. In *Lisp, the contents of pvars may be any valid Lisp object. As in Common Lisp, coercion between data types and allocation of memory is handled automatically.

non degriteries proc-initial-value nome type

Chapter 3 because deletes the prompted for each prior over declared and and the second states of the second states

sta-*defvars &reet pronuncs

3.1 Creating New Pvars

To create a permanent, named pvar, use *defvar (analogous to the Lisp defvar). To create a permanent, unnamed pvar, use allocate!!.

*defvar symbol &optional initial-value-pvar documentation-string [Macro]

This creates a new pvar that is permanently allocated. Symbol contains the allocated pvar. The optional argument *initial-value-pvar* may be any pvar or pvar expression. *defvar creates a new pvar, initializes it to the contents of *initial-value-pvar*, and sets the symbol to that new pvar using setq. If no *initial-value-pvar* argument is given, the symbol contains a pvar whose values are uninitialized. Note that *cold-boot resets the values of all pvars allocated by *defvar. This form returns symbol.

Some example uses of *defvar are:

```
(*defvar a)
(*defvar b (!! 5))
(*defvar c (+!! b (!! 6)))
(*defvar d t!!)
(*defvar e (self-address!!))
(*defvar f c)
```

expressions stems from its maintenance of this stack. While this automatic allocation takes care of many situations, there are times when it is desirable to explicitly allocate a temporary variable. The *Lisp macros for performing this operation are *let and *let*.

"let ({(symbol & optional pvar-expression)}") & rest body [Macro]

The first expression following the ***let** should be a list of lists, each specifying one temporary pvar. The elements of each sublist should consist of a Lisp symbol whose value will be the temporary pvar, followed by an optional pvar or pvar expression that will be copied into the new one.

These pvars survive only for the extent of the form. It is an error to try to refer to these pvars outside the body of the ***let**. In other words, the *symbols* have lexical scope (as in Common Lisp), whereas the pvars themselves have dynamic extent that terminates when the ***let** form is exited.

*let returns the value of the last form of the body, regardless of whether that value is a pvar. It is legitimate to return a temporarily bound pvar. *let is *not* able to return multiple values.

"let" ({(symbol & optional pvar-expression)}") & rest body [Macro]

This macro behaves in the same manner as "let except that, as in Common Lisp, the defining expressions are evaluated in sequence, so that previous bindings affect the evaluation of future initialization forms.

Some example expressions are:

is for processor 17 to 16

3.3 Setting the Values of Pvars

The *set macro allows the contents of one pvar to be set to the contents of another. The destination field is set in those processors that are currently selected. A *set expres-

This function returns, as a Lisp value, the component of the field specified by *pvar* in the processor whose grid address is given by *addresses*. There must be as many *address* values as there are dimensions in the processors' current configuration (as specified previously with *cold-boot). *setf may be used with pref-grid to write a value into a single processor of a pvar.

(pref-grid bar 4 7)

The above returns the component of pvar bar from processor (4,7) (on a two-dimensionally configured machine).

```
(*setf (pref-grid bar 4 7 8) (* 19 89))
```

The above sets the component of pvar bar for processor (4,7,8) (on a three-dimensionally configured machine) to 1691.

3.5 Declaring Pvar Types

*Lisp does not require that the programmer declare the type of a pvar's contents, nor does it require that all of a given pvar's values be of the same type. A pvar can have an integer in one processor and a floating-point number in the next. This flexibility comes at the cost of decreased efficiency.

Type declarations are a method to reduce runtime overhead for *Lisp code running either interpreted or compiled. Using pvars of defined types results in faster interpreted code and allows the *Lisp compiler to translate *Lisp code into Lisp/Paris.

For more information on pvar types, see chapter 8 of the Supplement to the *Lisp Reference Manual. For more information on how the *Lisp Compiler uses pvar types see the *Lisp Compiler Guide.

3.5.1 Syntax of Declarations

The pvar types supported by *Lisp are signed and unsigned integers, floating-point numbers of varying precision, complex numbers containing floating point data of varying precision, characters, string-chars, and booleans. Pvar type declarations presently are processed only by *proclaim, *let, *let*, *defun, *locally and the. These declarations have the same syntax as declarations in Common Lisp. The type of a pvar is specified in the following manner:

(pvar pvar-type-specifier)

) ...

```
(*proclaim '(type (pvar boolean) finished-p))
(*defvar finished-p nil!!)
(*let* ((temp1 (load-byte!! foo (!! 0) (!! 9)))
        (temp2 (sqrt!! temp1)))
(declare (type (pvar (unsigned-byte 9)) temp1)
        (type (pvar single-float) temp2))
;; temp1 may contain values between 0 and 511.
;; temp2 may contain single-floats.
...)
```

(setq xx (allocate!! (!! 1.23) 'xx '(pvar double-float)))

*Lisp allows certain elements of declarations to be computed at run time as opposed to compile time. All the *length* arguments to the type declarations may be either constants or run-time expressions such as global configuration variables. Following is a typical example using a global configuration variable:

```
(*let ((temp (self-address!!)))
  (declare
   (type (pvar (unsigned-byte *current-send-address-length*))
        temp))
```

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processors in which pror is non-nit. Even if there are no selected processors, ALE forms to the body are evaluated. The value of the final expression in the body is tothread whether it is a List value or a over.

Chapter 4

Processor Selection

When the Connection Machine is initialized, every processor is in a state to execute all *Lisp instructions in parallel. However, it may be necessary to execute instructions in some subset of processors. In fact, most operations are executed in a subset of Connection Machine processors, which is known as the *currently selected set*.

sors of the currently selected set in which the poor is all. Even if there are

rs, both thereform and che-form are evaluated. Unlike Lisp's R

Some of the macros in *Lisp that change the currently selected set are *all, *when, *cond, and *if. These macros select processors based on the result of a pvar expression. Any processor in which the pvar expression evaluates to nil is eliminated from the selected set. Although these macros may modify the currently selected set, they all obey the discipline of restoring the currently selected set to its previous state upon completion. Also, they may be nested as deeply as desired.

It is sometimes useful for user functions to have a *all surrounding their bodies to ensure that they are starting out with the complete machine selected. Using the functions described in this section, the selected set is whittled down to select only the processors that should perform a given operation. The body of these forms is *always* executed, even if there are no selected processors.

Note: In the current implementation, of the forms below that return values, none are configured to allow the return of multiple values. It is an error to attempt to return multiple values from any of these forms.

*all &body body

[Macro]

This form selects all processors. Its body is executed with the currently selected set equal to the entire machine. The value of the final expression in the body is returned whether it is a Lisp value or a pvar.

*when pvar &body body

[Macro]

This form subselects from the currently selected set. Thus every processor that is unselected when *when is called remains unselected in the body of the *when. It selects

the use of these macros are

```
(*when (=!! a b) (*set e (+!! c d)))
(*cond ((=!! a (!! 1)) (*set e (+!! b c)))
       ((not!! (=!! c d)) (*set f (*!! b c)))
       (t!! (*set f (!! 9))))
(*if (=!! c d) (*set e f) (*set g h))
(*when a (*set b (-!! b))
(*defun f (x y)
   "Returns y divided by x for y greater than 0.
   Returns nil if any x is 0. The return value is
   undefined a processors where y<0"
   (block foo
      (with-css-saved
         (*when (>!! y (!! 0))
            (if (*or (=!! (!! 0) x))
               (return-from foo nil)
               (/!! y x))
              ))))
```

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(Puniction

The post returned by this predicate contains (for each processor where the value of the argument number-poor to lets than zero, and all in all others. N 5 ranged of -0.0)) is always faire.

Computations on Pvars

This chapter introduces a variety of functions that work within each processor and that return a pvar containing all the processors' results. Recall that in *Lisp, it is conventional for functions that return a pvar to have the suffix !! on their names.

5.1 Predicate Operations

The *Lisp predicate functions are used in the same way as Common Lisp predicates, for instance, in conditional expressions. They return a pvar that contains a t in all processors of the currently selected set in which the predicate holds, and a nil in those in which it does not.

oddp!! integer-pvar

The pvar returned by this predicate contains t for each processor where the value of the argument *integer-pvar* is odd, and nil in all others. It is an error if any component of *integer-pvar* is not an integer.

This is the parallel equivalent of the Common Lisp function zeros.

evenp!! integer-pvar

The pvar returned by this predicate contains t for each processor where the value of the argument *integer-pvar* is even, and nil in all others. It is an error if any component of *integer-pvar* is not an integer.

plusp!! number-pvar

The pvar returned by this predicate contains t for each processor where the value of the argument *number-pvar* is greater than zero, and nil in all others.

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[Function]

[Function]

[Function]

/=!! numeric-pvar &rest numeric-pvars

This returns a pvar that contains t in each processor where the argument pvars contain unequal values and nil elsewhere. If only one argument is given, the returned pvar is t!!.

<!! numeric-pvar &rest numeric-pvars

This returns a pyar that contains t in each processor where the argument pyars contain values which are in strictly increasing order and nil elsewhere. If only one argument pvar is given, the returned pvar is t!!.

This returns a for all processors to which over is nit, and nit othered >!! numeric-pvar &rest numeric-pvars

This returns a pvar that contains t in each processor where the argument pvars contain values which are in strictly decreasing order and nil elsewhere. If only one argument pvar is given, the returned pvar is tll. the more evaluates to all in a processor, that processor is removed from the curves

remainder of the and/i. and/i returns the val

<=!! numeric-pvar &rest numeric-pvars

This returns a pvar that contains t in each processor where the argument pvars contain values which are in non-decreasing order and nil elsewhere. If only one argument pvar is given, the returned pvar is t!!.

>=!! numeric-pvar &rest numeric-pvars

This returns a pvar that contains t in each processor where the argument pvars contain values which are in non-increasing order and nil elsewhere. If only one argument pyar is given, the returned pvar is tll.

5.1.1 Predefined Pvars

nil!!

[Constant]

This is a pvar whose contents in each processor is the Lisp symbol t.

r this returns t if there are an odd number of argumentitiens

[Constant]

This is a pvar whose contents in each processor is the Lisp symbol nil.

It is an error to use t!! or nil!! as the destination for "set, "pset or any other form which modifies its argument.

[Function]

[Function]

[Function]

[Function]

[Function]

ogloal Operations

lognot!! integer-pvar

[Function]

This returns a pvar whose bits are the logical complement of the bits in integer-pvar.

logior!! &rest integer-pvars

[Function]

This returns a pvar whose bits are the logical inclusive or of the bits in integer-pvars. If there are no pvars, then (!! 0) is returned.

logxor!! &rest integer-pyars

[Function]

This is the parallel equivalent of the Common Lisp function logxor. If there are no integer-pvars, then (!! 0) is returned.

logand!! &rest integer-pvars

[Function]

This returns a pvar whose bits are the logical and of the bits in integer-pvars. If no integer-pvars are given, then (!! -1) is returned.

and shares and and the state

logeqv!! &rest integer-pvars [Function]

This is the parallel equivalent of the Common Lisp function logeqv. If no integer-pvars are given, then (!! -1) is returned. if that argument has any field whose value is

month is is an error if any argument but the first has any field whose value is 0.

Numerical Operations 5.4

This section describes the elementary numerical functions. As with Common Lisp, the results of these functions are always numerically correct. For example, the result of an addition is never truncated, no matter how much memory is required to represent the result. If not enough memory is available, an error is signaled. However, the numerical accuracy of certain arithmetic operations on floating point data is subject to restrictions noted in the Version 5.0 *Lisp Release Notes. A few arithmetic operations are also restricted when operating on integer data to a maximum number of bits for each argument. The Release Notes also describe these limitations.

These functions each return results of the same type as the most expensive of their arguments (e.g. if all arguments are integers, the result is generally an integer; but if any argument is a float, the result is a float).

mod!! numeric-pvar integer-pvar

This is the parallel equivalent of the Common Lisp function mod. It is an error if *inte-ger-pvar* contains zero in any processor.

ash!! integer-pvar count-pvar

[Function]

[Function]

This is the parallel equivalent of the Common Lisp function ash.

truncate!! numeric-pvar & optional divisor-numeric-pvar [Function]

This is the parallel equivalent of the Common Lisp function truncate, except that only one value (the first) is computed and returned.

round!! numeric-pvar & optional divisor-numeric-pvar [Function]

This is the parallel equivalent of the Common Lisp function round, except that only one value (the first) is computed and returned.

ceiling!! numeric-pvar & optional divisor-numeric-pvar [Function]

This is the parallel equivalent of the Common Lisp function ceiling, except that only one value (the first) is computed and returned.

floor!! numeric-pvar & optional divisor-numeric-pvar

This is the parallel equivalent of the Common Lisp function floor, except that only one value (the first) is computed and returned.

sqrt!! non-negative-or-complex-pvar

This returns the non-negative square root of its argument, if the argument is not complex. If the argument is complex, the principal square root is returned. Unlike Common Lisp., it is an error to provide a negative, non-complex value to sqrt!!.

isgrt!! non-negative-integer-pvar

This is the parallel equivalent of the Common Lisp function isqrt.

random!! limit-pvar

This returns a pvar whose contents is a random value between 0 inclusive and *limit-pvar* exclusive for each processor.

.

[Function]

....

[Function]

[Function]

[Function]

least significant bit. In any processor in which zero bits are extracted, the resulting field contains zero. This operation is especially fast when both position-pvar and sizepvar are constants, as in (!! lisp-value). from-pvar must be a pvar containing integers, while position-pvar and size-pvar must be pvars containing non-negative integers. Outof-range bits are treated as zero for positive integers (for example, (load-byte!! (!! 1) (!! 2) (!! 3)) returns a pyar that contains zero in each processor), and one for negative integers (for example, (load-byte!! (!! -1) (!! 2) (!! 3)) returns a pvar that contains 7 in each processor).

[Function] deposit-byte!! into-pvar position-pvar size-pvar byte-pvar

[Macro]

This returns a pvar whose contents are a copy of into-pvar with the low order size-pvar bits of byte-pvar inserted into the bits starting at location position-pvar.

When the into-pvar is positive (negative), zeros (ones) are appended as high order bits of byte-pvar as needed. The returned value may have more bits than into-pvar if the inserted field extends beyond the most significant bit of into-pvar. For example, (deposit-byte!! (!! 3) (!! 1) (!! 2) (!! 2)) returns (!! 5). This function is especially fast when both position-pvar and size-pvar are constants, as in (11 lisp-value). Into-pvar and byte-pvar must contain integers, while position-pvar and size-pvar must be pvars containing non-negative integers only.

if!! pvar then-pvar & optional else-pvar

This returns a pvar that contains the contents of the then-pvar in all processors in which pvar is non-nil, and the contents of else-pvar in all processors in which pvar is nil. The else-pvar argument defaults to nill!. For the execution of the then-pvar expression, the currently selected set is set to all processors that passed the predicate, whereas for the execution of the *else-pvar* the currently selected set is set to all the selected processors that failed the predicate. (See also "if, which is executed only for side effect.)

plet to the value of numeric-year in proc

d its the symbol cell, and its social analog is the predice

This is equivalent to:

```
(*let ((result)
         (temp-pred pvar))
     (*when temp-pred
        (*set result then-pvar))
     (*when (not!! temp-pred)
        (*set result else-pvar))
     result
1
```

An example that demonstrates the usefulness of if!! is the following function to take the absolute value:

*defun name arg-list &body body [Macro]

This is analogous to the Common Lisp defun and can be used in place of it in defining user functions that might take as an argument a pvar or that might return a pvar as a result. Using "defun is only required if a function is to take pvar arguments and possibly return a non-pvar result. "defun returns, as a symbol, the name of the function being defined. Like the Common Lisp defun, the body may contain declarations and a documentation string. In particular type declarations for pvar arguments may be provided within the body of a "defun.

If, in a given file, a function foo defined by ***defun** is called before it is defined textually in the file, or is called but is not defined in the current file, then the user must declare that foo is actually a function defined by ***defun** and is not a regular function defined by **defun**. One makes such a declaration with the ***Lisp macro *proclaim**. For example:

(*proclaim '(*defun foo))

Failure to make such declarations results in incorrectly compiled code.

*funcall function &rest arguments [Macro]

This is used just like Common Lisp's funcall, but with functions defined using *defun. One may not use funcall with a function defined using *defun.

provider the default value for the lowword are

*apply function arg &optional more-args

[Macro]

[Macro]

This is used just like Common Lisp's apply, but with functions defined using *defun.

5.7 Debugging Tools

pretty-print-pvar pvar

&key (:mode *ppp-default-mode*)
 (:format *ppp-default-format*)
 (:per-line *ppp-default-per-line*)
 (:start *ppp-default-start*)
 (:end *ppp-default-end*)

This prints out the contents of a pvar in all processors. If :per-line is nil, no newlines are ever printed between values; otherwise, :per-line values are printed out and then a newline is output. The keyword :mode can have the value :cube or :grid; in the latter

33

list-of-active-processors

[Function]

This simply returns a list of cube addresses of all the currently selected processors. The order of this list is not specified. Since this function is so useful, an alias, loap, is also defined. This could be written as:

```
(defun list-of-active-processors ()
  (let ((return-list nil))
     (do-for-selected-processors (processor)
        (push processor return-list))
     return-list))
```

pretty-print-pvar-in-currently-selected-set pvar &key format start end

[Function]

This function prints out the the cube address and value of *pvar* for all processors in the currently selected set. Since this function is so useful, an alias, **ppp-css**, is also defined. *format* defaults to "-s". This function returns no values.

6.4 Global Operations

The following functions reduce the contents of a pvar in all selected processors into a single Lisp value, which is then returned:

*logior integer-pvar

This returns a Lisp value that is the bitwise logical inclusive or of the contents of *inte*ger-pvar in all selected processors. This returns the Lisp value 0 if there are no selected processors.

*logand integer-pvar

This returns a Lisp value that is the bitwise logical and of the contents of *integer-pvar* in all selected processors. This returns the Lisp value -1 if there are no selected processors.

*min numeric-pvar

This returns a Lisp value that is the minimum of the contents of *numeric-pvar* in all selected processors. This returns the Lisp value nil if there are no selected processors.

*max numeric-pvar

This returns a Lisp value that is the maximum of the contents of *numeric-pvar* in all selected processors. This returns the Lisp value nil if there are no selected processors.

*or pvar

This returns a Lisp value of t if the contents of *pvar* is non-nil in any selected processor; otherwise, it returns nil. If there are no selected processors, this function returns nil. For example, to determine if there are any processors currently selected, use (* or t!!), which returns t only if there are selected processors.

*and pvar [*Defun]

This function translates a grid address consisting of (possibly) several a

This returns a Lisp value of t if the contents of *pvar* is non-nil in every selected processor; otherwise, it returns nil. If there are no selected processors, this function returns t.

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[*Defun]

00003013.

[*Defun]

[*Defun]

[*Defun]

[*Defun]

Chapter 6. Communication

grid-from-cube-address!! cube-address-pvar dimension-pvar [Function]

This function takes a *cube-address-pvar* and returns a pvar containing the grid address for the specified *dimension-pvar* for each selected processor.

cube-from-grid-address!! address-pvar &rest address-pvars [Function]

This function translates a grid address consisting of (possibly) several address-pvars into a cube address for each selected processor.

off-grid-border-p!! &rest grid-address-pvars [Function]

This function returns a boolean pvar that is true if the grid-address-pvars specify an address that is invalid given the current dimensions, and false otherwise. It is an error for any component of grid-address-pvar to be a non-integer.

off-grid-border-relative-p!! & rest relative-grid-address-pvars [Function]

This function is identical to off-grid-border-p!! except that the relative-grid-addresspvars specify relative addresses. "Ling Reference Manual

Machine hardware. All "dubur pyars are reallocated and their initial values are recomputed.

Chapter 7

Using the Connection Machine

The *Lisp language resides in a package named *LISP. To use the language, the user must either be in that package:

der normal circomstances, this need only be called at the start of a session

```
(in-package '*LISP)
```

or else make that package available to the package the user is in:

(use-package '*LISP)

On Symbolics Lisp machines, the user should put the following package attribute in the attribute list of any file that contains functions to be put in the *LISP package:

Package: (*LISP COMMON-LISP-GLOBAL)

For instructions on how to load the *Lisp language into your Lisp machine, please refer to the Connection Machine Front-End Subsystems.

7.1 Using the Connection Machine Hardware

This section describes the two *Lisp functions (*cold-boot and *warm-boot) that allow the user to use the hardware.

Note: The macro *cold-boot has been enhanced to work with n-dimensional NEWS, *Lisp interpreter safety, and geometry objects. These changes and the new features to which they are related are all documented in the Supplement to the *Lisp Reference Manual.

*cold-boot &key : initial-dimensions initial-dimensions [Macro]

This function initializes *Lisp and must be called immediately after loading in the *Lisp software. It resets the internal state of the *Lisp system and of the Connection

7.1.1 Initialization Lists for *cold-boot and *warm-boot

Users can define a set of forms to be executed automatically before and after each execution of ***cold-boot** and ***warm-boot**. These user-defined initialization lists are stored in one or more of these variables:

*before-*cold-boot-initializations*

*after-*cold-boot-initializations*

*before-*warm-boot-initializations*

*after-*warm-boot-initializations*

New forms are added using the function add-initialization, and removed using deleteinitialization.

add-initialization name-of-form form variable

j-jorni jorni variable

The argument *name-of-form* is a character string that names the form being added. The argument *form* may be any executable Lisp form. Adding two forms with the same name is permissible only if the forms are the same according to the function equal; otherwise an error is signaled. The *variable* should be one of the initialization-list variables above, or it may be a list of such variables, in which case the *form* is added to each initialization list named. The *form* and *variable* arguments must be quoted so that they are not evaluated during the call to add-initialization. For example:

delete-initialization name-of-form variable

[Function]

This function deletes the form named by *name-of-form* from the initialization list (or lists) specified by *variable*. The arguments are specified in the same manner as the first and third arguments for add-initialization. For example:

55

[Variable]

[Variable]

[Variable]

[Variable]

[Function]

boot and *warm-boot work in the same manner as for the interpreter. The simulator is more lenient than the interpreter with respect to dimensioning the machine: the simulator allows arbitrary extent in each dimension, while the interpreter currently supports only extents that are powers of two. Warnings are issued if one tries to configure the simulator in a way that the interpreter cannot handle.

Since the simulator and interpreter implementations of *Lisp are very different, it is unfortunately necessary to recompile code when switching from one system to the other. It is not possible to load the simulator once the interpreter version of *Lisp has been loaded, and vice versa.

The symbols "lisp-hardware and "lisp-simulator are appended to the Common Lisp "features" list when running on the interpreter and on the simulator, respectively. These symbols can be used to perform read-time conditionalization of code.

(DURSE long times (and)

7.2 Interfacing Paris Code to *Lisp

It is sometimes necessary to explicitly call Paris instructions from within *Lisp programs.

Paris instructions require pvars' memory addresses and lengths as their arguments. While a pvar is commonly—and appropriately—regarded as a "parallel variable," that is, a program variable that has a value in each CM processor, a pvar actually exists in the front end as a Lisp object that points to and describes a field in each CM processor's memory. These fields in the CM contain the values that comprise the parallel variable. The front-end Lisp object also contains pvars' addresses and lengths, information that may be needed as arguments to Lisp/Paris functions. The following functions return this information.

The stock is used for fast allocation of pyacs returned

```
(pvar-location pvar)
(pvar-length pvar)
```

These functions return the location (address) and length of a pvar.

bas words we betrees more its asbulant and abate a we bestaper relate that and (pvar-type pvar)

This function returns the type of a pvar. The type of a pvar may be one of :general, :field, :signed, :float, :boolean, :complex, :character, :string-char, :array, or :structure. These correspond to pvar types t (for general), unsigned-byte, signed-byte, defined-float, boolean, complex, character, string-char, array, and types defined by *defstruct, respectively. "Lity Referance Manual

Chapter 8

Avoiding Potential Difficulties

This chapter describes potential difficulties in using *Lisp and ways the user can avoid them.

8.3 Using Mapoar on Functions that Return pyars

8.1 Pvar Values in Non-Selected Processors

It is an error to depend on the value of a pvar in a processor that was not in the currently selected set at the time the pvar was created. For instance, the following is incorrect:

```
(*when (<!! (self-address!!) (!! 10))
 (*let ((foo (self-address!!)))
  (print (pref foo 20))
))</pre>
```

The *Lisp language definition does not define the value printed in the above example because the pvar foo was only given values in the active processors, 0 through 9.

8.2 The Extent of Pvars

Unlike Common Lisp, pvars defined using "let or "let" have dynamic extent; that is, it is an error to reference the value of a pvar once the body of its defining "let or "let" has been exited. For example:

```
(*defun will-not-work (pvar constant)
  (funcall
    (*let ((xyzzy (!! constant)))
    #* (lambda (x) (*sum (+!! (!! x) xyzzy)))
```

8.5 *cold-boot

*cold-boot calls cm:coldboot to initialize Connection Machine state and also initializes the *Lisp run-time systems. This function *must* be called when first entering the *Lisp environment (hardware or simulator) if the user intends to evaluate any code. In addition, we strongly recommend that *cold-boot be called between runs of application programs. Finally, *cold-boot should be called if *warm-boot has failed to clear an error state.

a Grid and Cube Addresses

*cold-boot is intended to be used as a top-level form. Under no circumstances should it be used inside of a *defun form, either lexically or in such a manner that it might be executed while a *defun function is being evaluated.

8.6 pref!!, pref-grid!! and pref-grid-relative!!

These *Lisp functions (which are actually macros) are defined to evaluate their source argument(s) in the context of the set of addresses defined by evaluating their address pvar(s). Therefore, writing a function foo that calls one of these functions with a source argument s passed into foo may not work because s will have already been evaluated by the Lisp evaluator before the body of foo is evaluated.

The solution is to define such functions as macros. For example:

```
(*defun foo (condition source-pvar address-pvar)
  (if condition
        (pref!! source-pvar address-pvar)
        (!! 0)
  ))
```

The above may not work as intended and should be rendered as

```
(defmacro foo (condition source-pvar address-pvar)
`(if ,condition
        (pref!! ,source-pvar ,address-pvar)
        (!! 0)
))
```

If source-pvar is always a symbol and not a pvar expression, then this modification is not necessary.

8.7 Multiple Values

The current implementation does not support the return of multiple values from any *Lisp form. It is error to attempt such an operation.

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Appendixes

Appendix A *Lisp Symbols

Lisp Reference Manual

This appendix lists all *Lisp symbols and pvar types described in this manual.

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