

On the Existence of Hazard-Free Multi-Level Logic

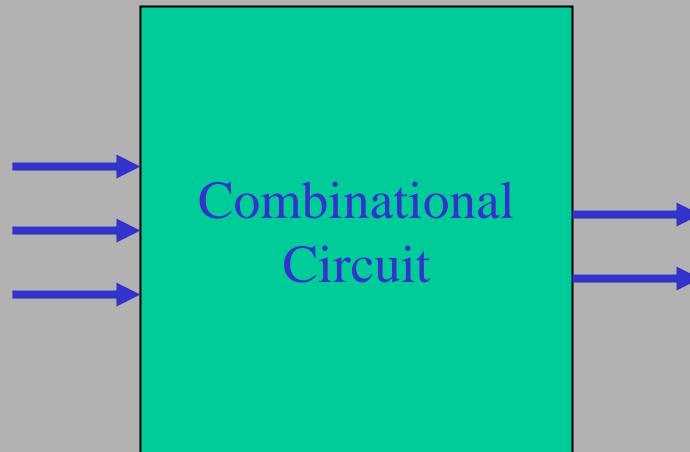
Steven M. Nowick

Charles W. O'Donnell

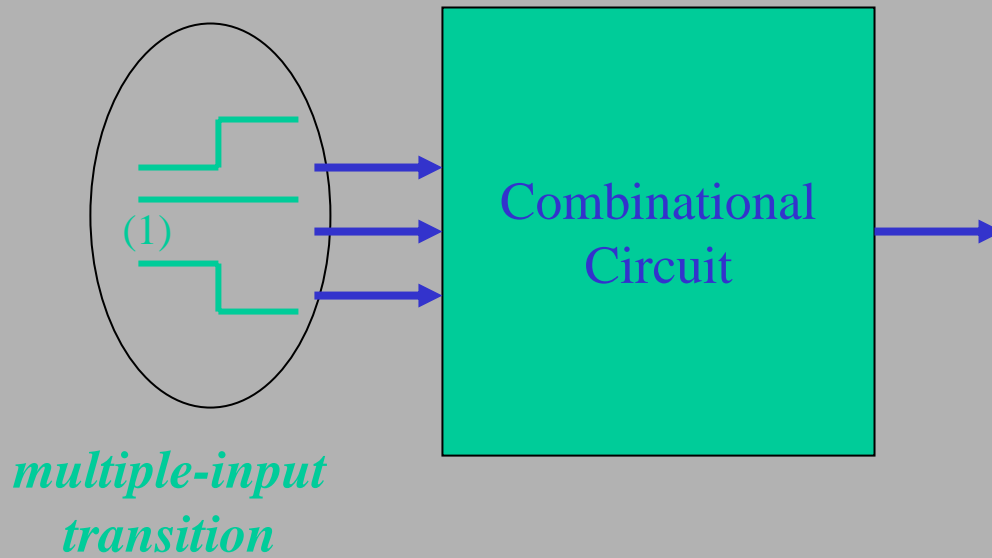
Department of Computer Science

Columbia University

Preliminaries

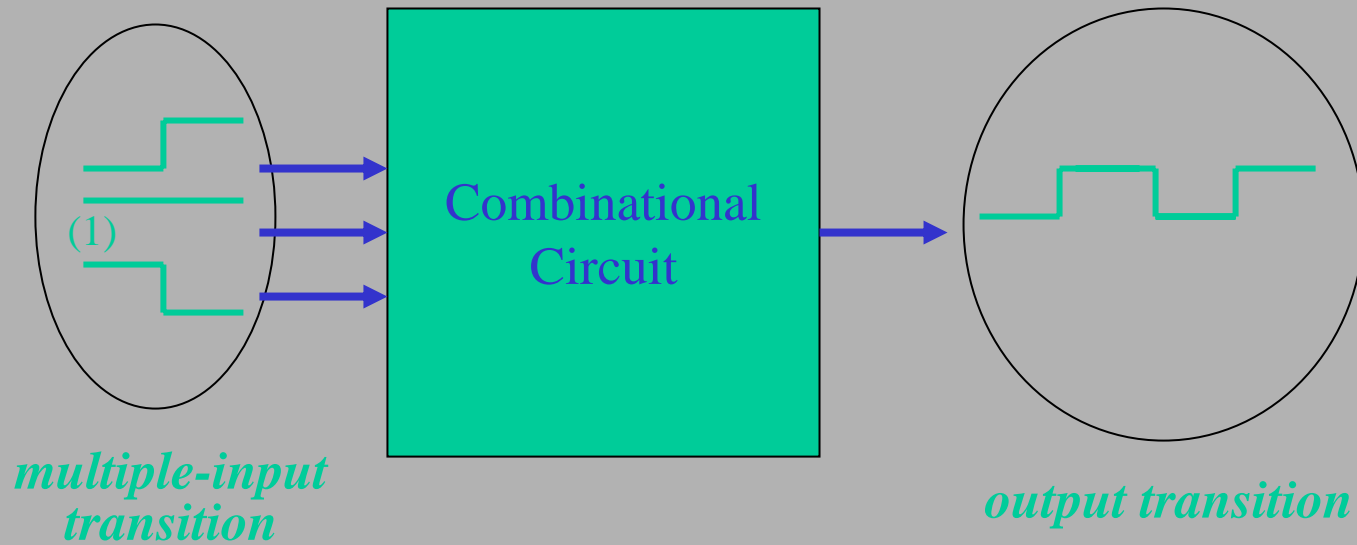


Preliminaries



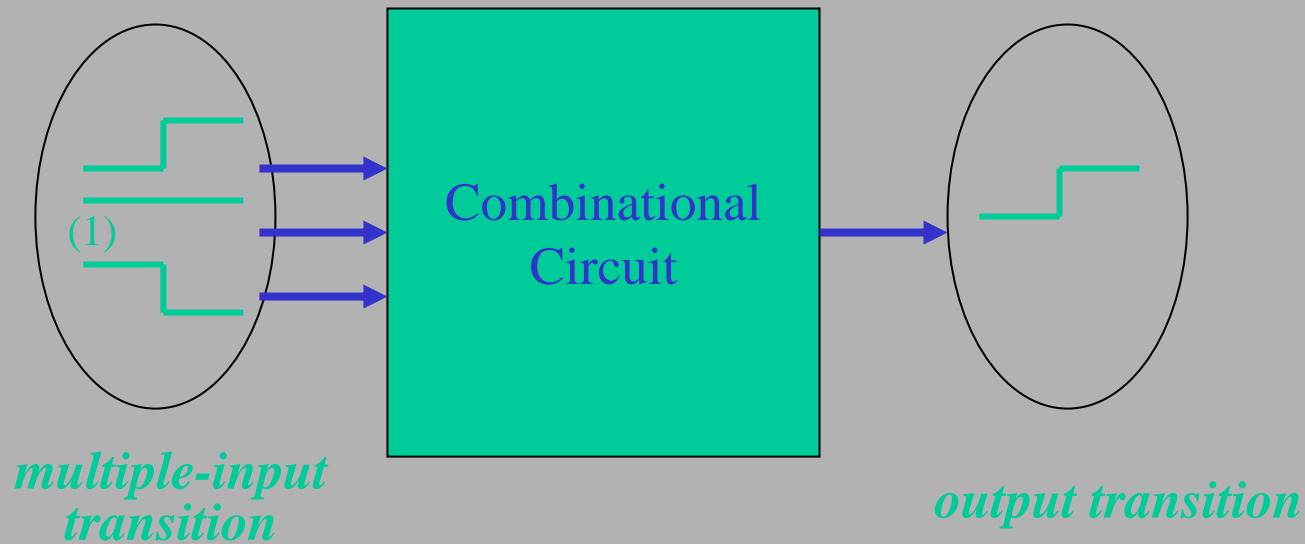
Preliminaries

hazard = potential for a glitch



Preliminaries

hazard-free = *no potential for a glitch*

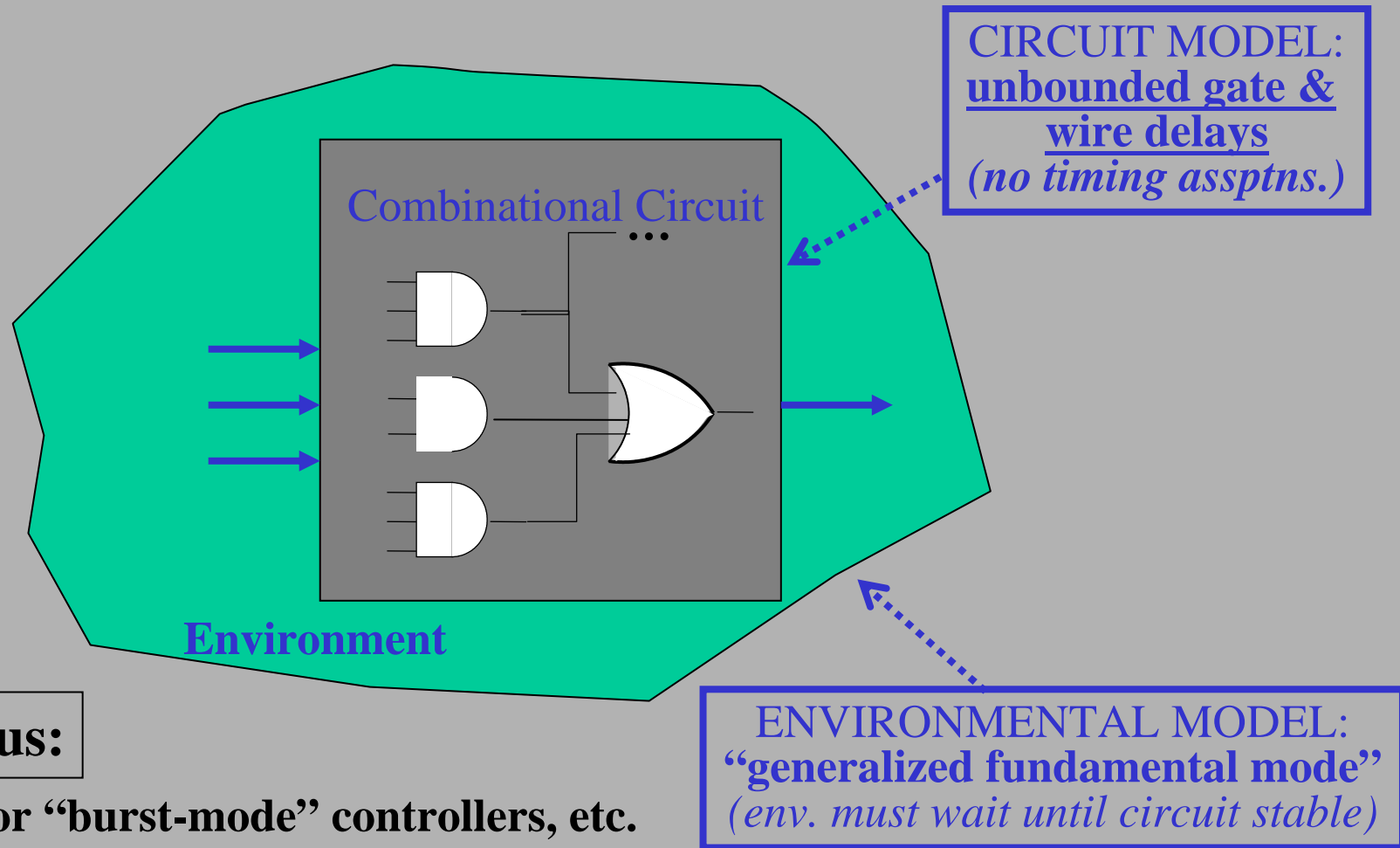


Goal: design hazard-free combinational circuits

Outline

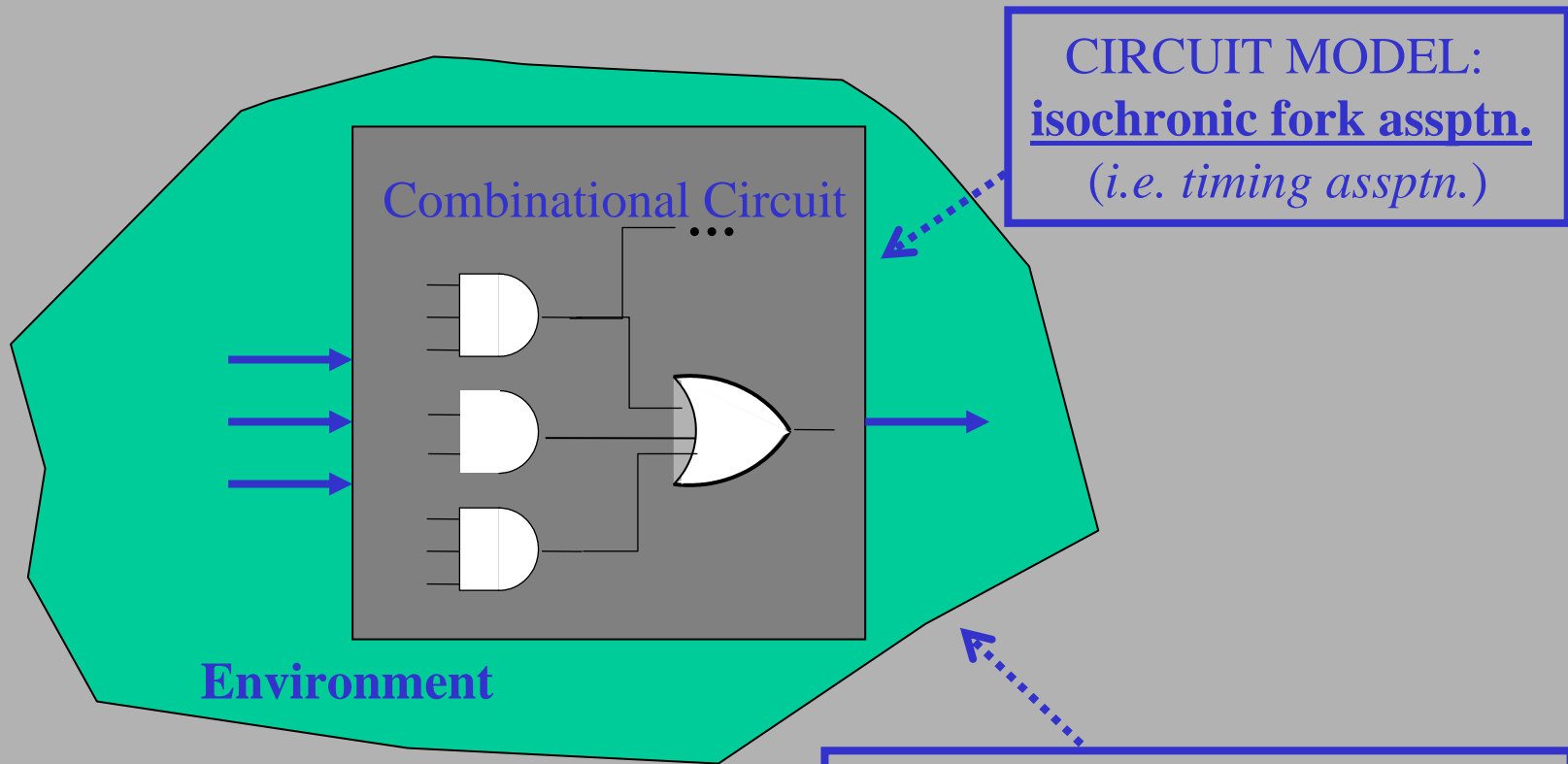
- • **Background:** Hazard-Free 2-Level Logic
- **Overview of Multi-Level Method**
 - Hazard-Free Decomposition
 - Algorithm + Examples
 - Standard 3-Level Circuit Structures
 - Canonicity
- **Related Work:** Bredeson ['74]
- **Experimental Results**
- **Conclusions**

Circuit & Environmental Models: #1



Goal: synthesize *hazard-free circuit* under this model

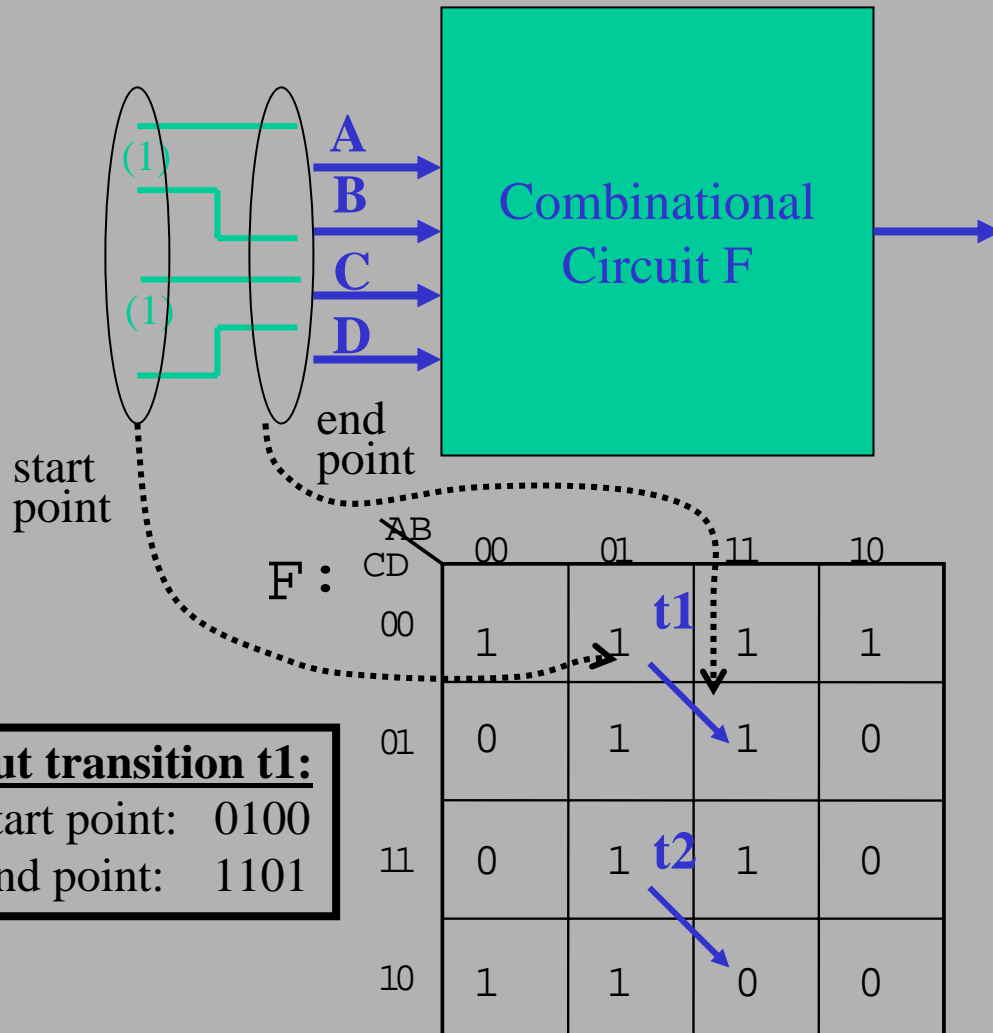
Circuit & Environmental Models: #2



Alternative Model:

... used for "QDI" controllers, etc.

Combinational Hazards in 2-Level Circuits: Notation



input transition t1:

start point: 0100

end point: 1101

Function F: represented using Karnaugh Map (= truth table)

Function Hazards

- An input transition which is not monotonic in function F itself [Unger69]
- **Intuition:** occurs when user specifies glitchy behavior
 - circuit implementations cannot be guaranteed glitch-free!

F :

AB \ CD	00	01	11	10
00	1	1	1	1
01	0	1	1	0
11	0	1	1	0
10	1	1	0	0

t4: function hazard-free

t3: function hazard

... only consider function hazard-free input transitions [Unger 69]

Logic Hazards: Conditions for Hazard-Free Transitions

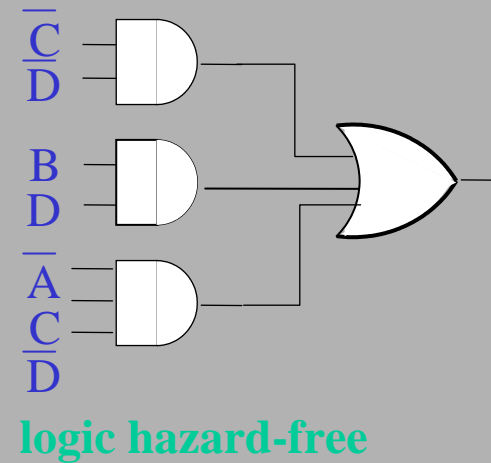
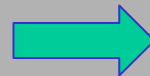
- Property of a **given circuit implementation**
- An input transition has a **logic hazard** if circuit impltn. may glitch

Case 1: $0 \rightarrow 0$ Transition

F :

AB \ CD	00	01	11	10
00	1	1	1	1
01	0	1	1	0
11	0	1	1	0
10	1	1	0	0

A blue arrow labeled t_5 points from the cell (01, 10) to the cell (11, 10).

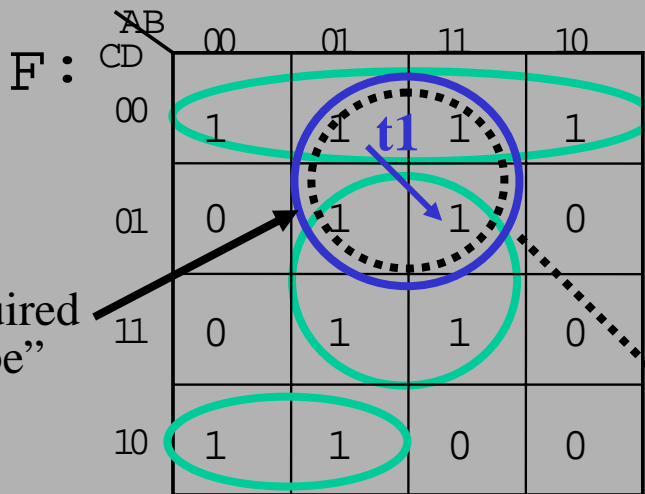
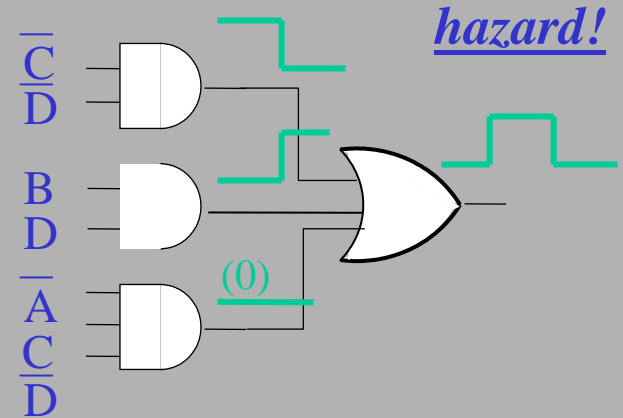
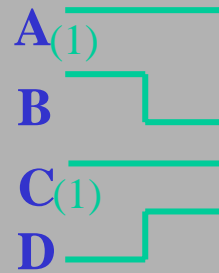
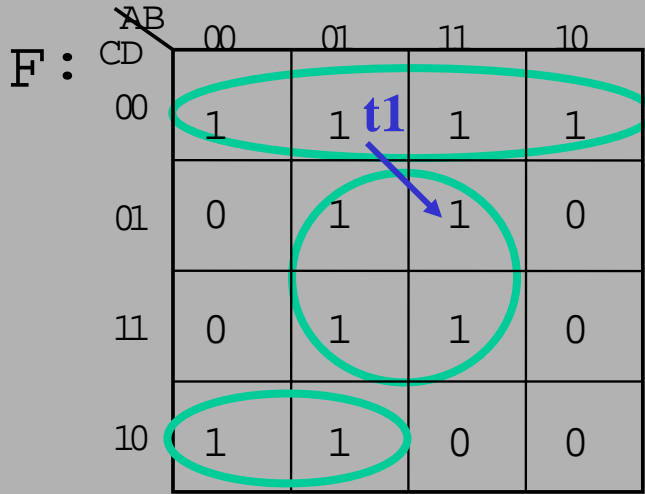


...every AND-OR circuit is logic hazard-free for $0 \rightarrow 0$ transitions [Unger 69]

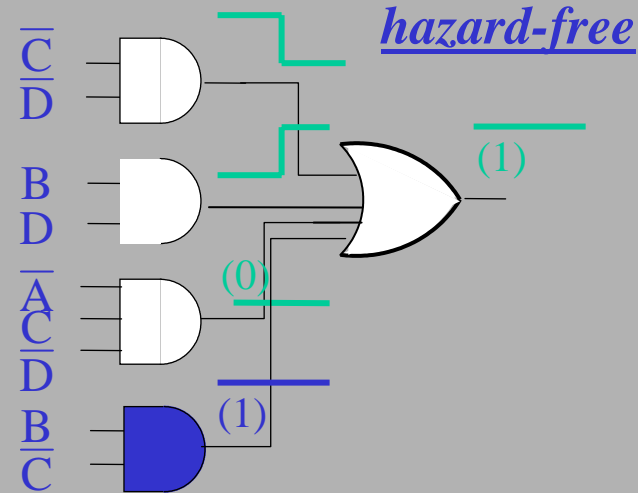
Logic Hazards (cont.)

Case 2: 1→1 Transition

Example: input transition t1



new gate

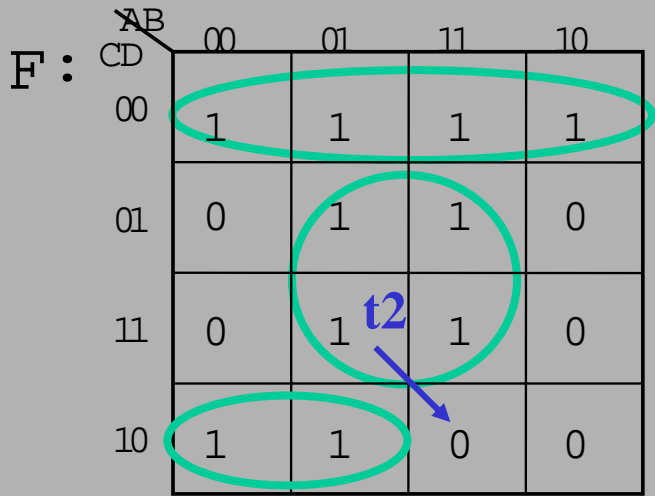


Transition t1's "required cube" must be covered

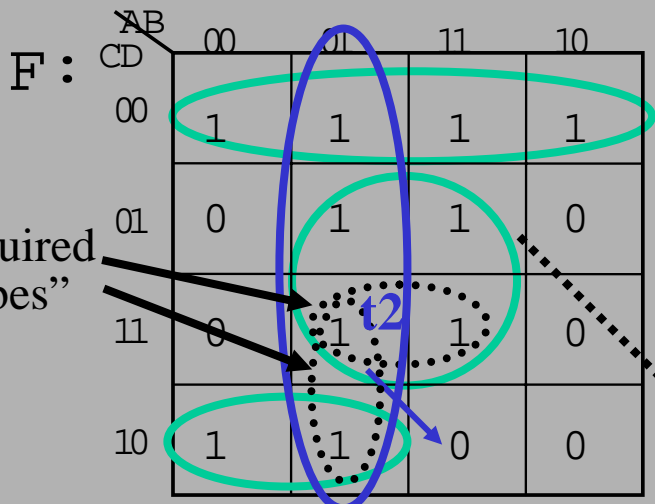
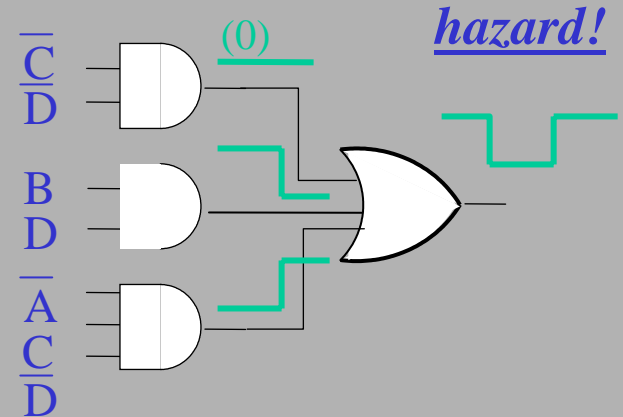
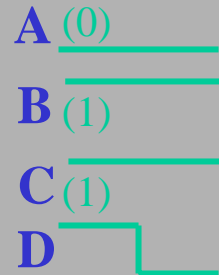
Logic Hazards (cont.)

Case 3: $1 \rightarrow 0$ ($0 \rightarrow 1$) Transition

Example: input transition t2

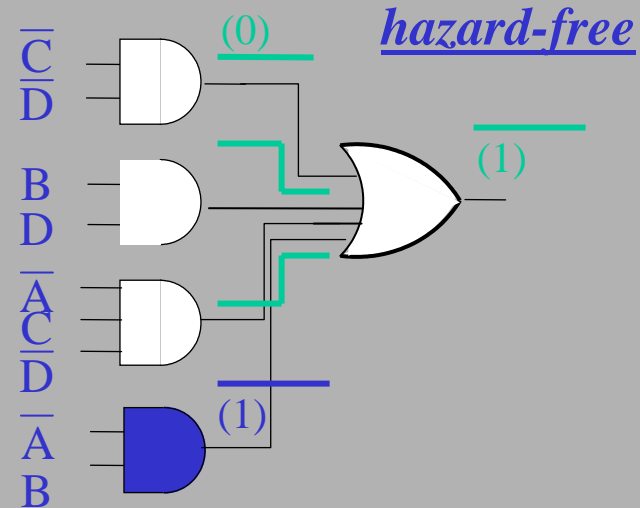


partial transition



“required cubes”

new gate

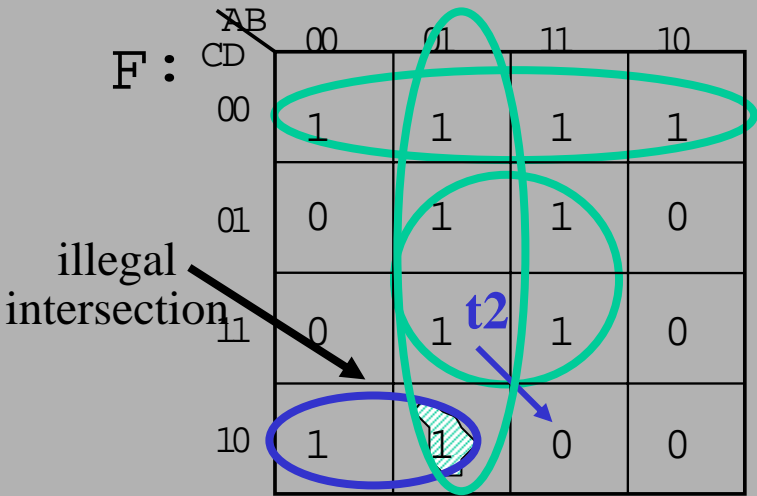


Transition t2: each max $1 \rightarrow 1$ subtransition = “required cube” which must be covered

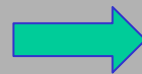
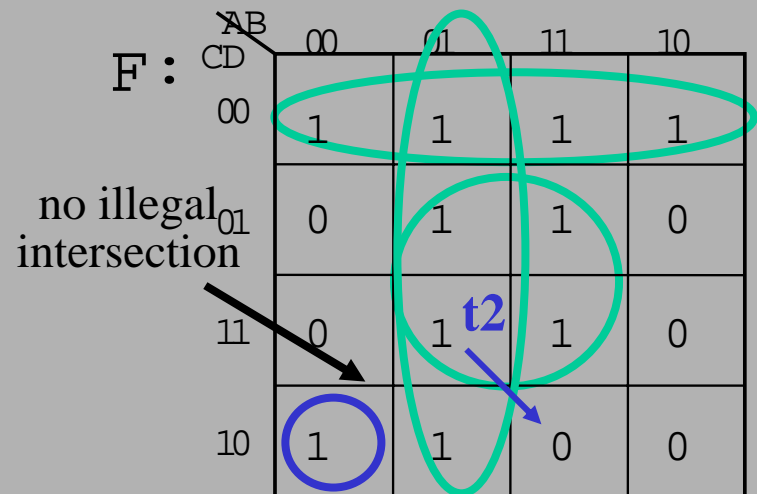
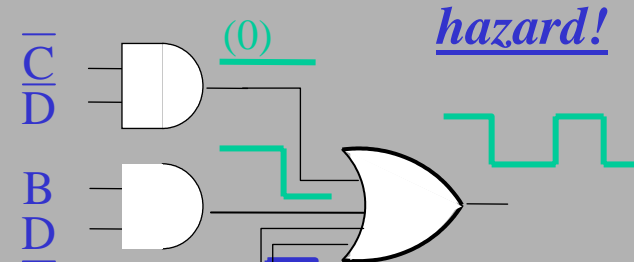
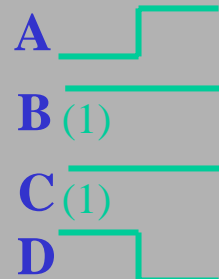
Logic Hazards (cont.)

Case 3: 1→0 (0→1) Transition (cont.)

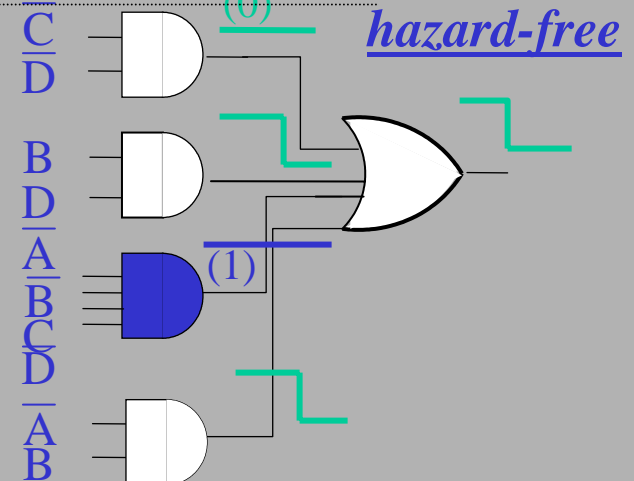
Example: input transition t2



complete transition



modify gate



Transition t2: avoid "illegal intersections" of 1→0 (0→1) transitions

Hazard-Free Two-Level Logic Minimization

Summary: “Hazard-Free Covering Rules”

In Karnaugh map for function F ...:

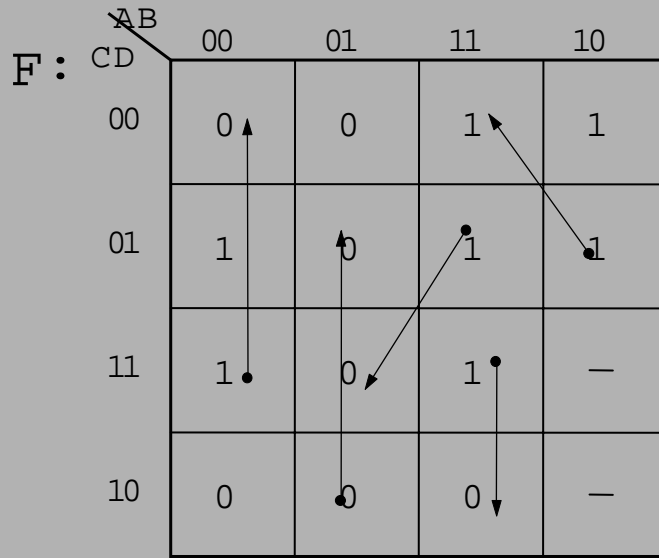
- each **required cube**: must be covered by some product
- no **illegal intersections**: with any $1 \rightarrow 0$ ($0 \rightarrow 1$) transition

Hazard-Free Minimization Algorithm [Nowick/Dill ICCAD92,TCAD95]

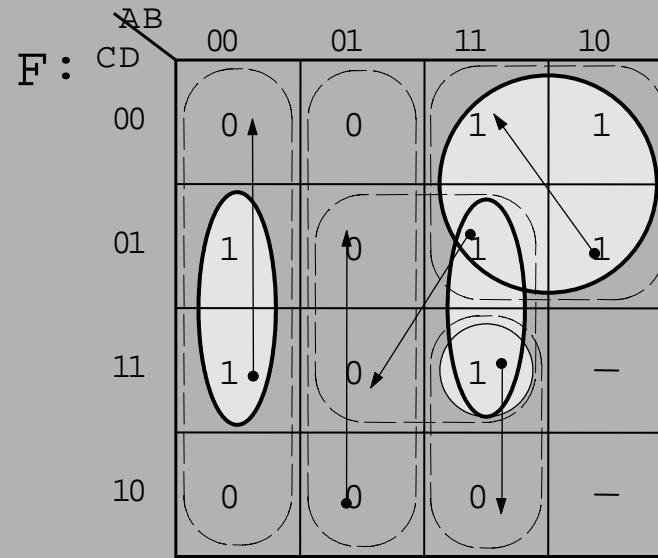
- produce a minimum-cost hazard-free AND-OR circuit

Existence of a Two-Level Solution: *not always!*

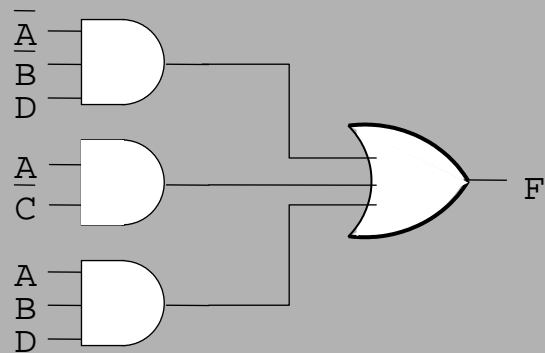
Example #1: Hazard-Free Two-Level Solution



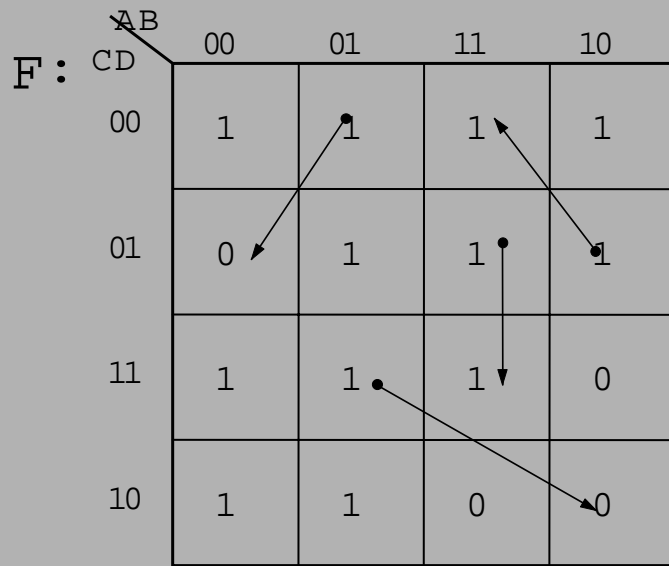
(a)



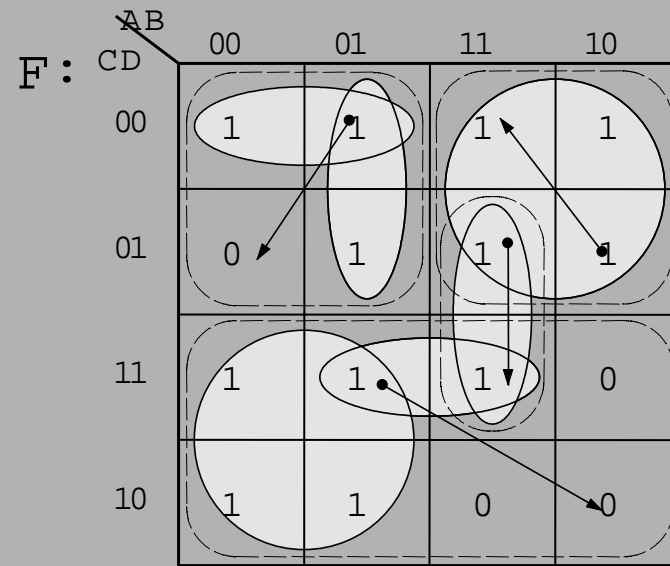
(b)



Example #2: No Hazard-Free 2-Level Solution Exists



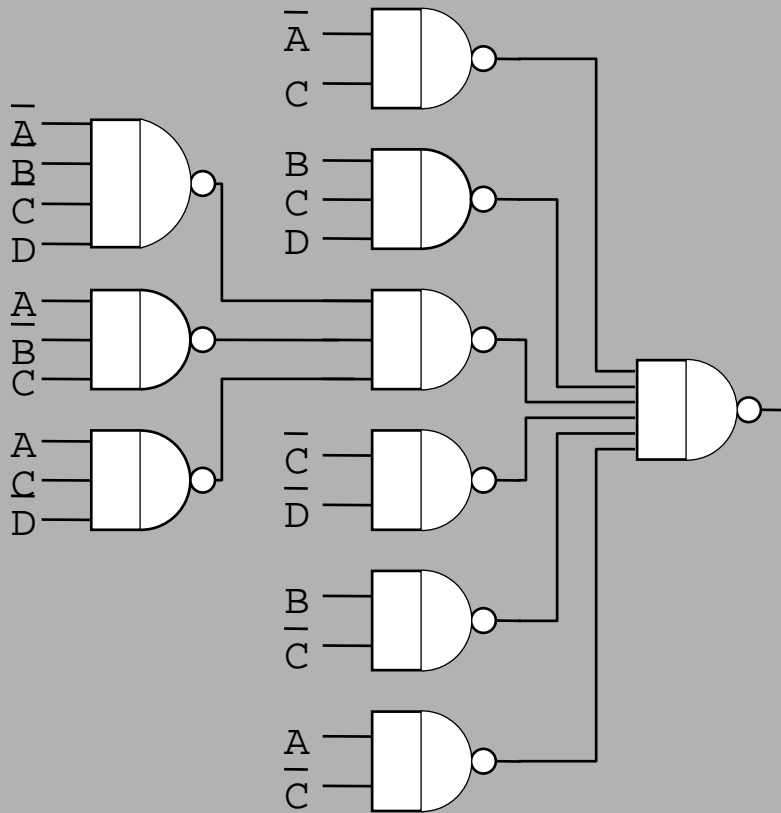
(a)



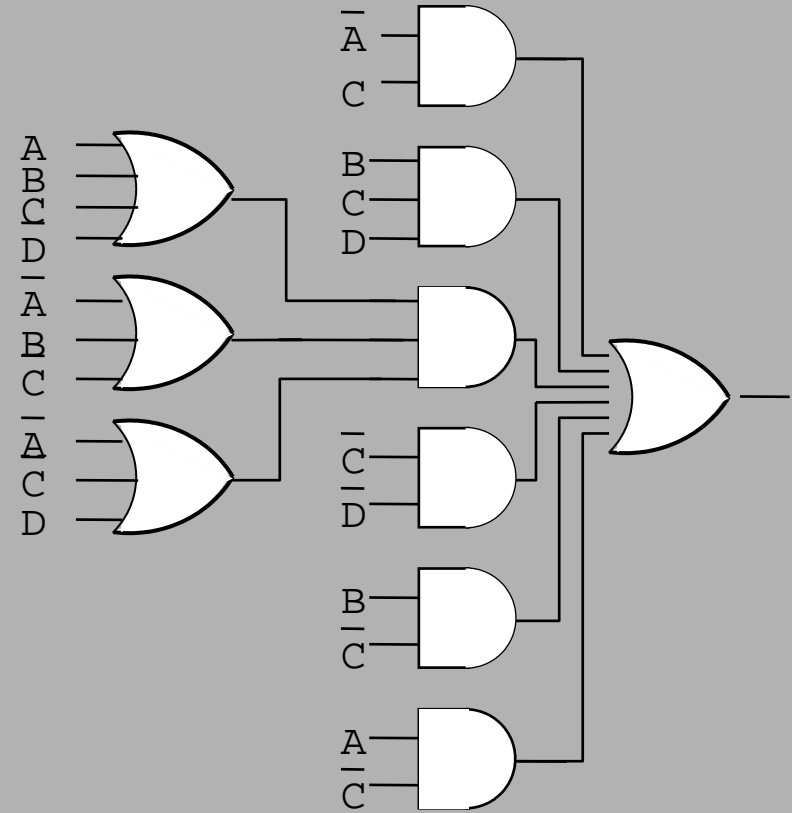
(b)

[Nowick/Dill, ICCAD-92]

Example #2 (cont.)



(a)



(b)

However, several hazard-free multi-level implementations do exist!

Contribution

Problem. Given a Boolean function F and specified set of input transitions T :

Does any multi-level hazard-free implementation exist?

Assumptions. Consider only networks of simple combinational gates:
AND, OR, INV, NAND, NOR

General Solution: *2 new algorithms, always find a solution if one exists.*

1. Basic Recursive Method
2. Iterative Method

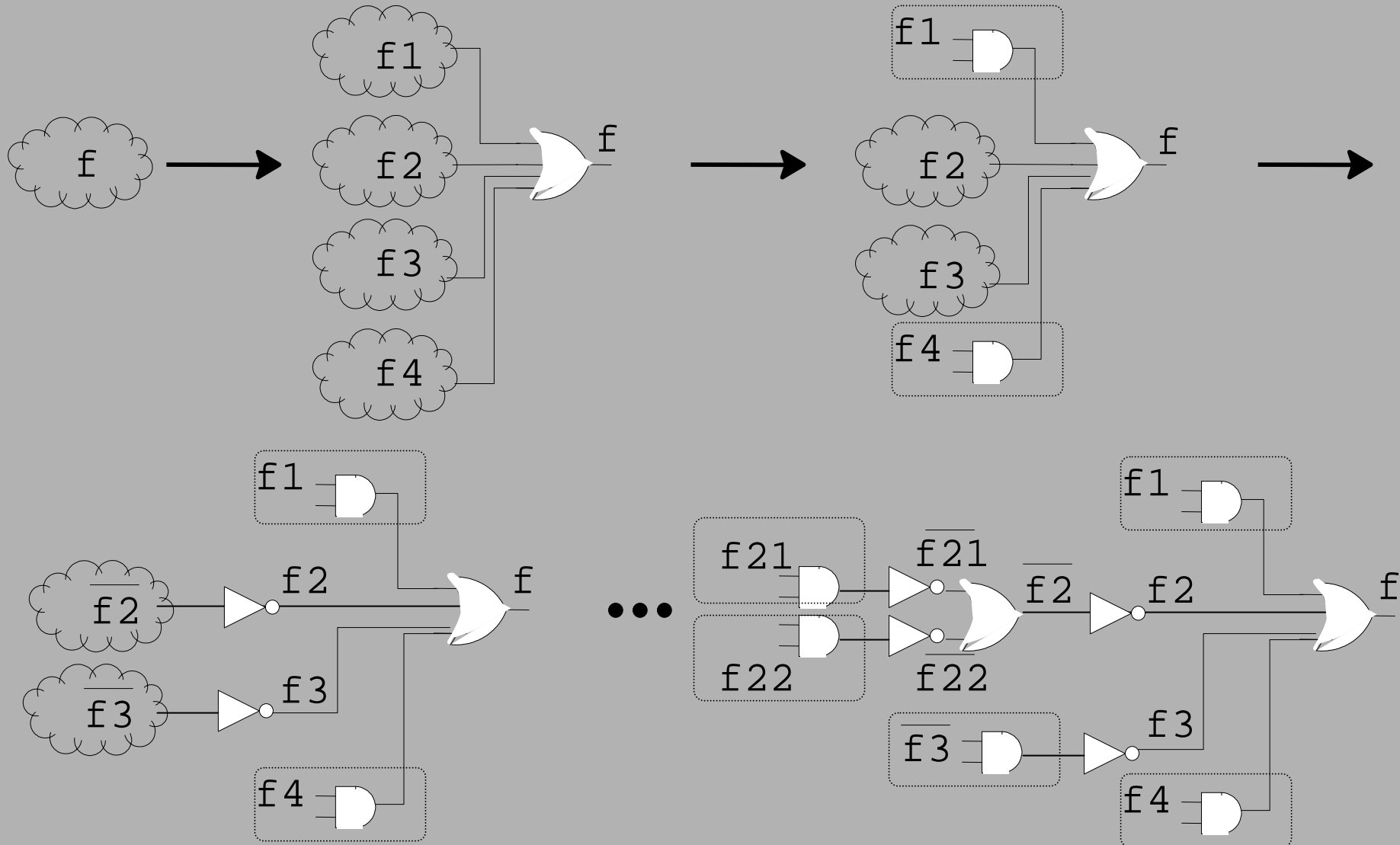
Key Contribution: show that, if some hazard-free multi-level implementation exists, then there is *always a hazard-free 3-level implementation:*

- **3-Level NAND**
- **OR-AND-OR**

Outline

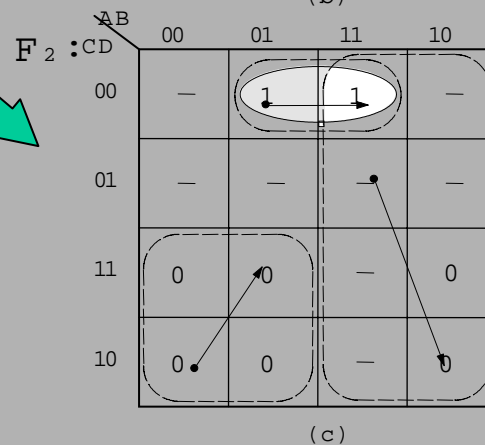
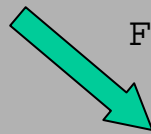
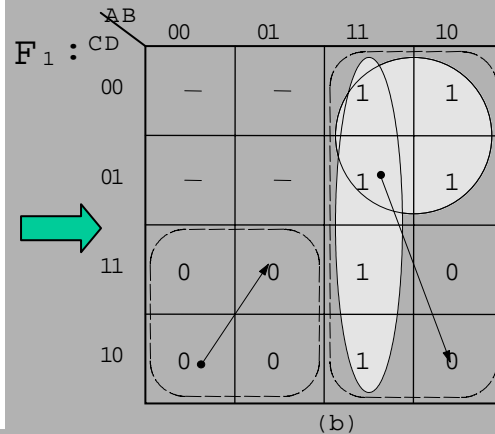
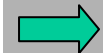
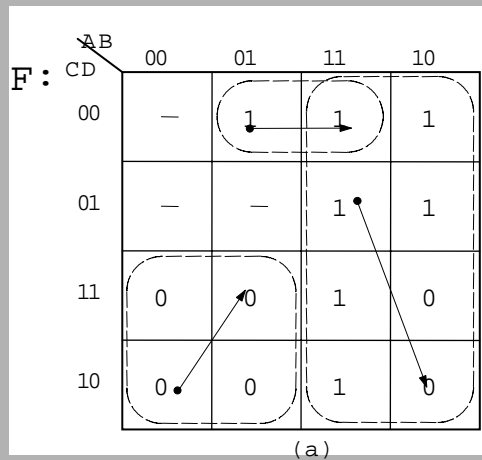
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- **Experimental Results**
- **Conclusions**

Basic Overview of Approach: Use Recursive “Hazard-Free Decomposition”



First Hazard-Free Decomposition: Disjunction

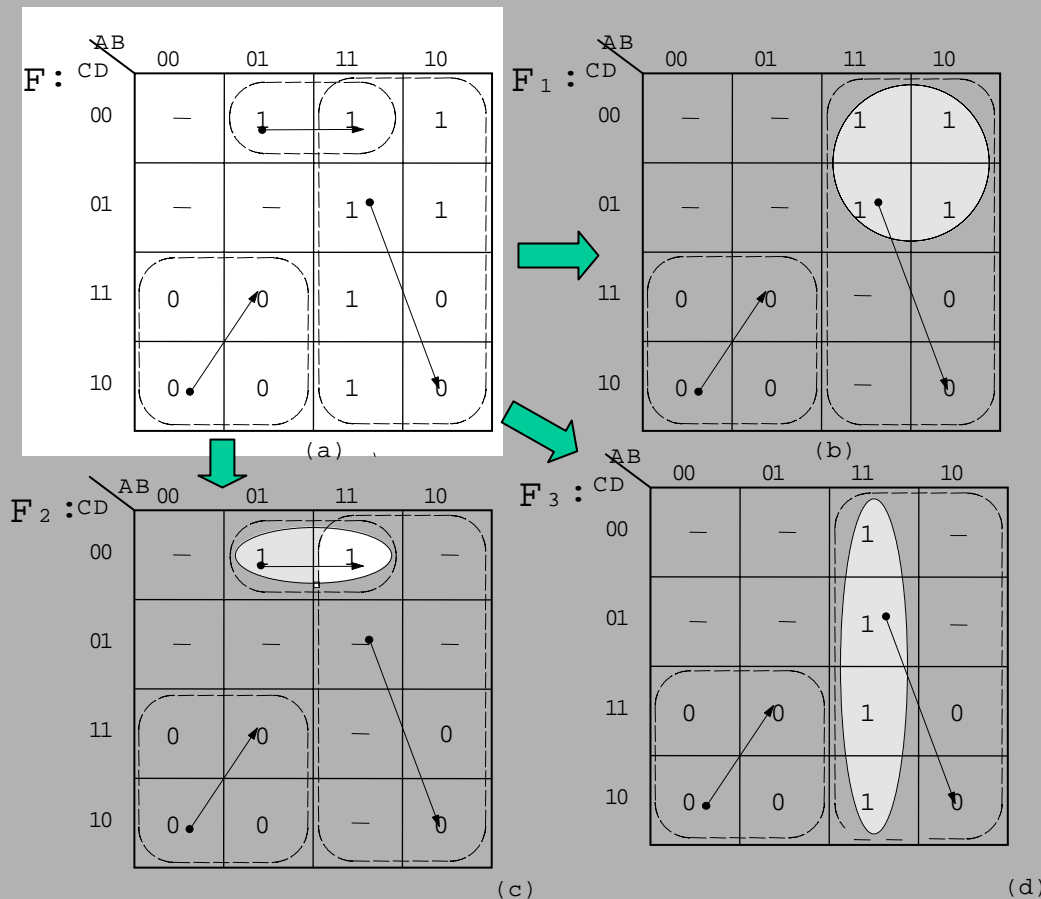
General disjunctive decomposition:



“Required cubes” are divided among the new subfunctions

First Hazard-Free Decomposition: **Disjunction** (cont.)

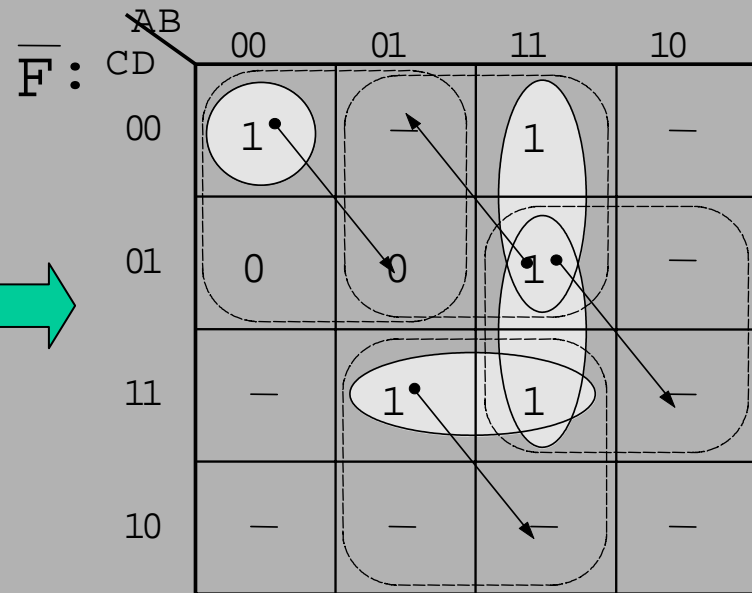
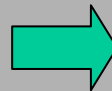
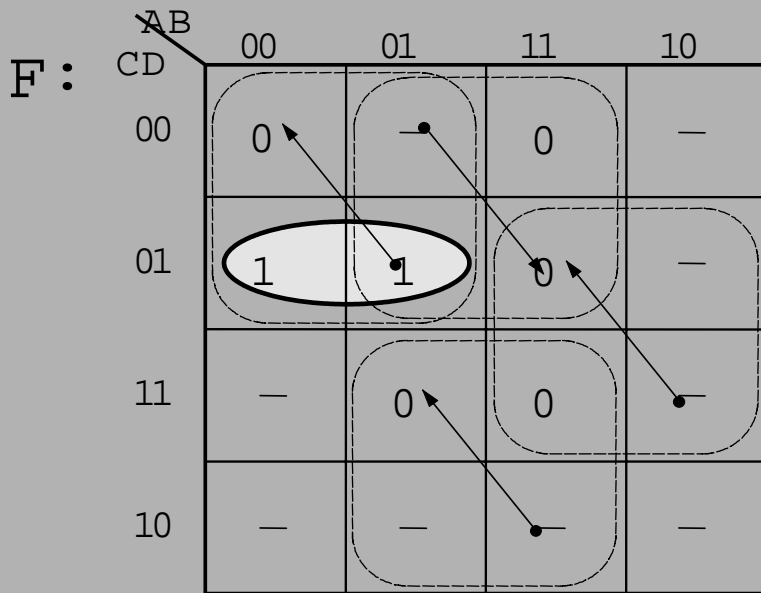
Maximal disjunctive decomposition: a special case



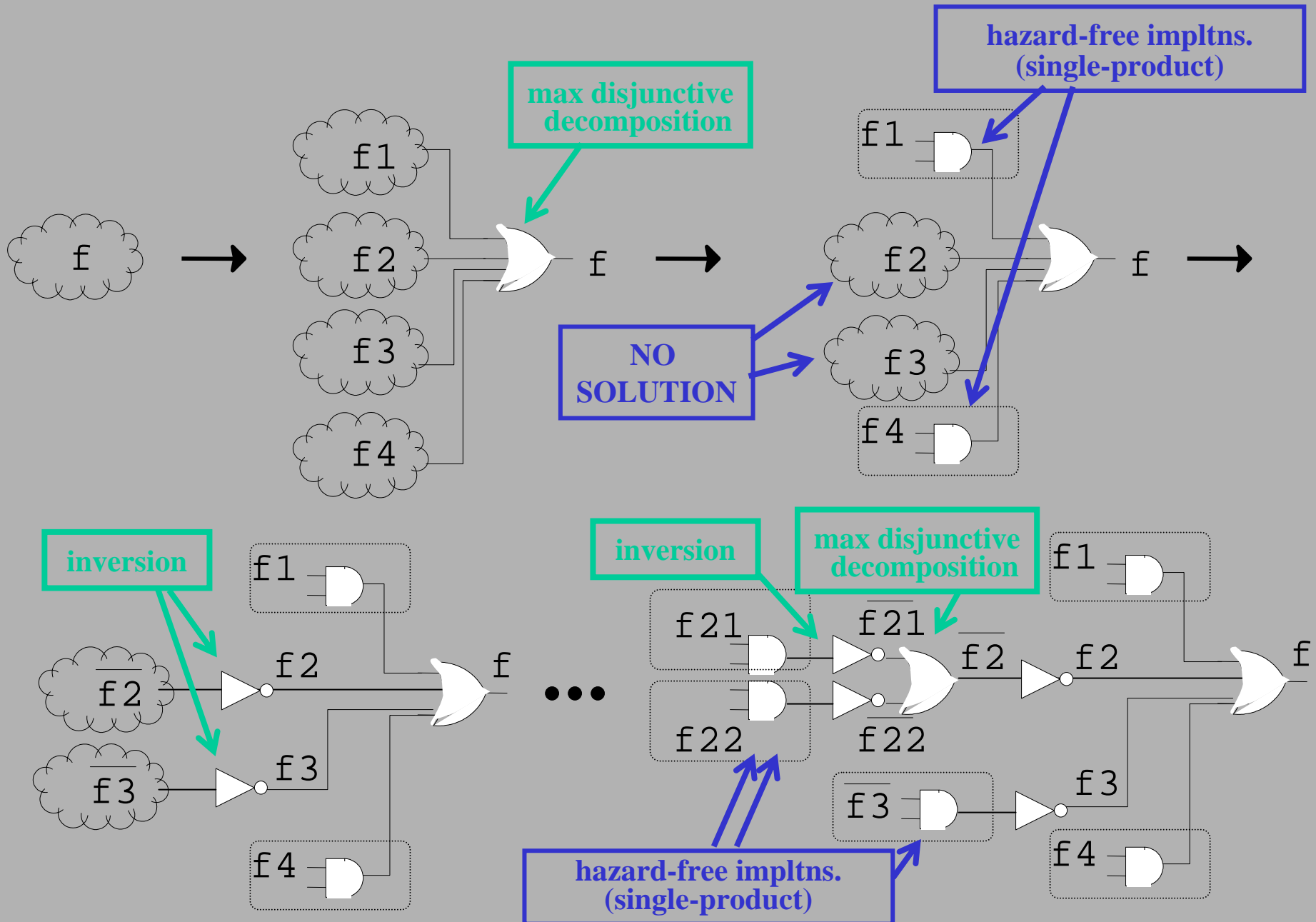
Each new subfunction gets exactly one (distinct) “required cube”

Second Hazard-Free Decomposition: **Inversion**

Swap OFF-set and required cubes



Basic Summary of Approach: *alternate* disjunction + inversion



First Complete Example:

Finding a Hazard-Free Multi-Level Implementation

F :

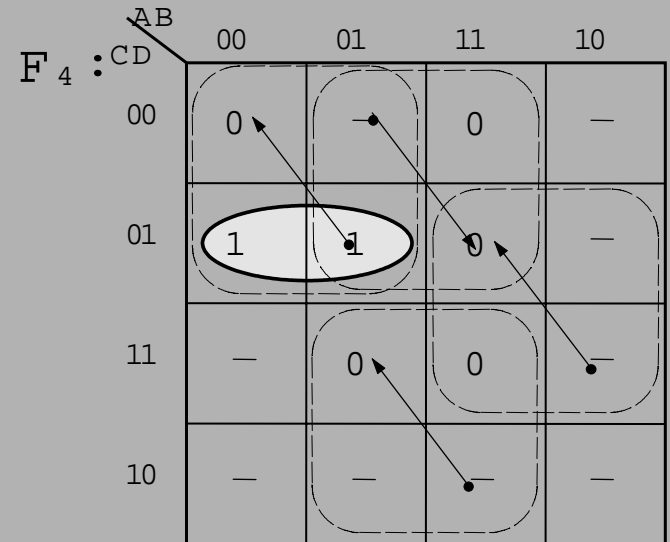
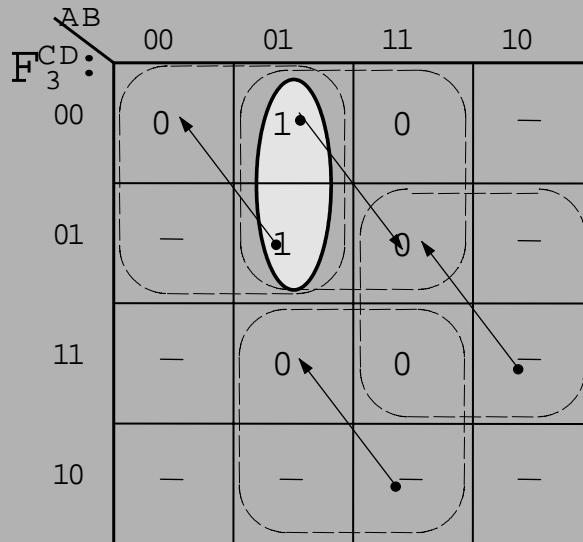
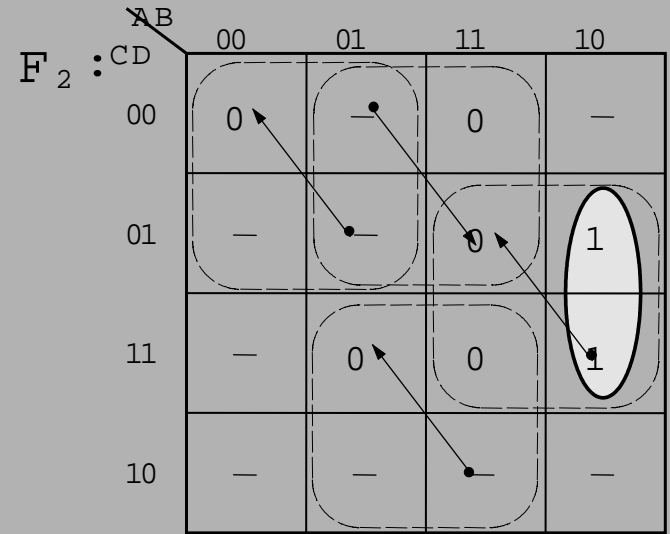
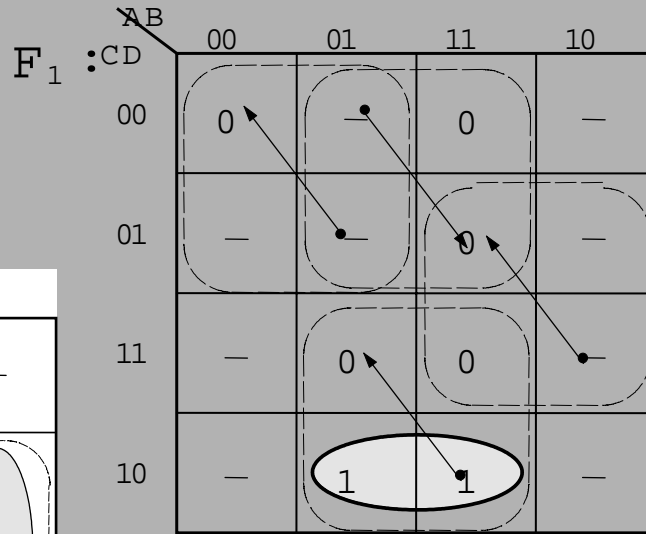
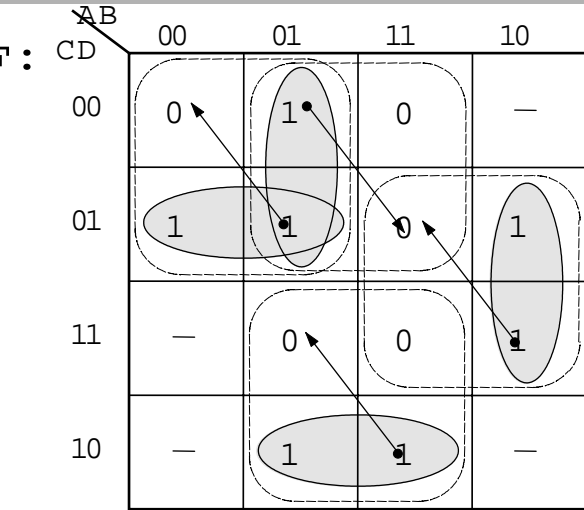
AB \ CD	00	01	11	10
00	0	1	0	—
01	1	1	0	1
11	—	0	0	1
10	—	1	1	—

F :

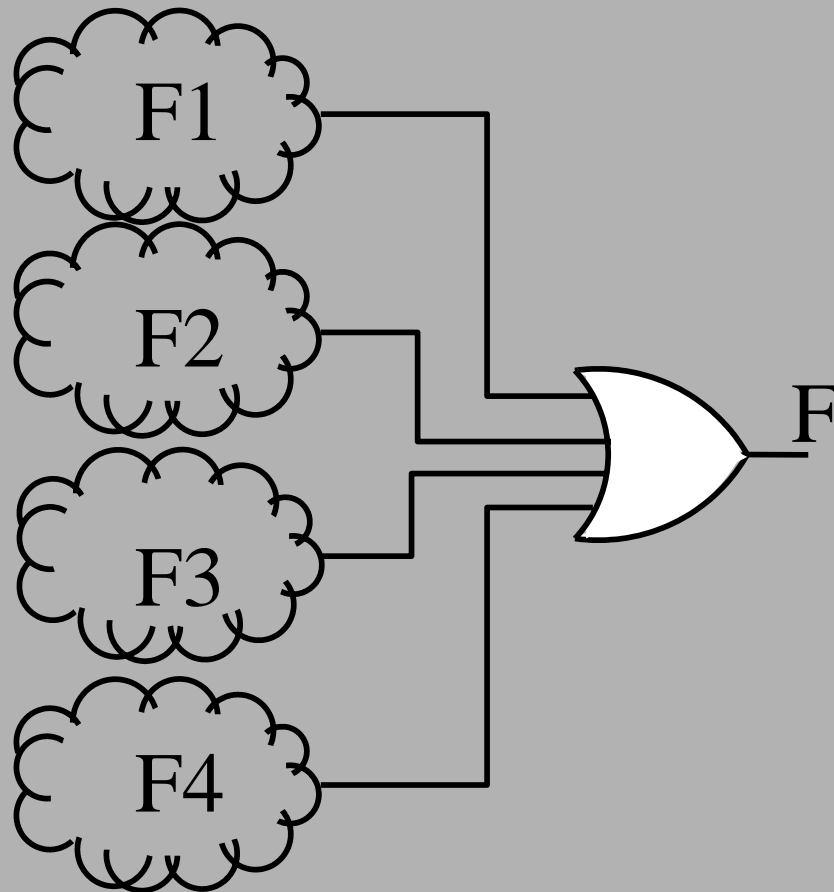
AB \ CD	00	01	11	10
00	0	1	0	—
01	1	1	0	1
11	—	0	0	1
10	—	1	1	—

Given Boolean Function + Specified Transitions

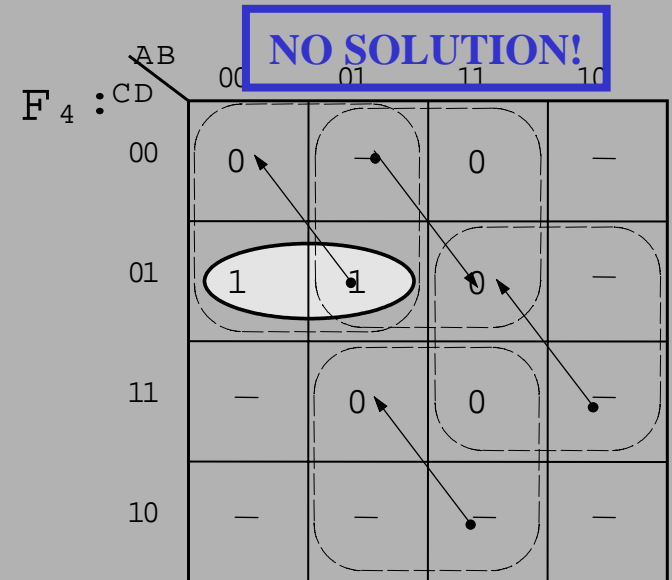
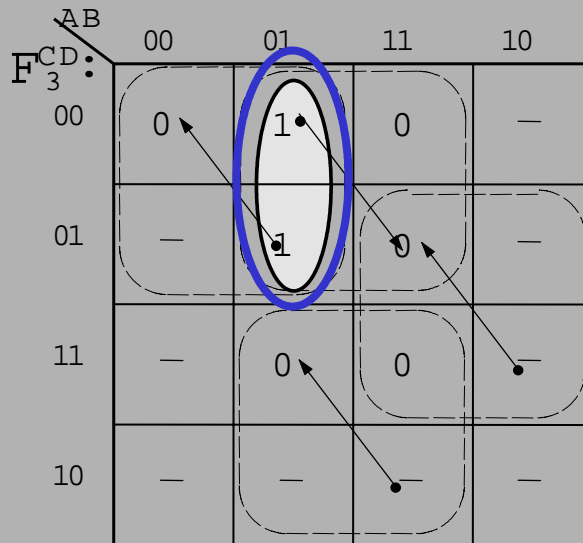
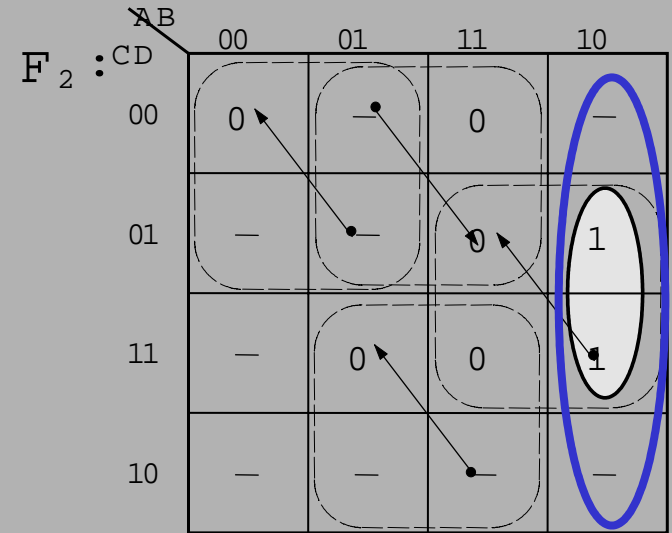
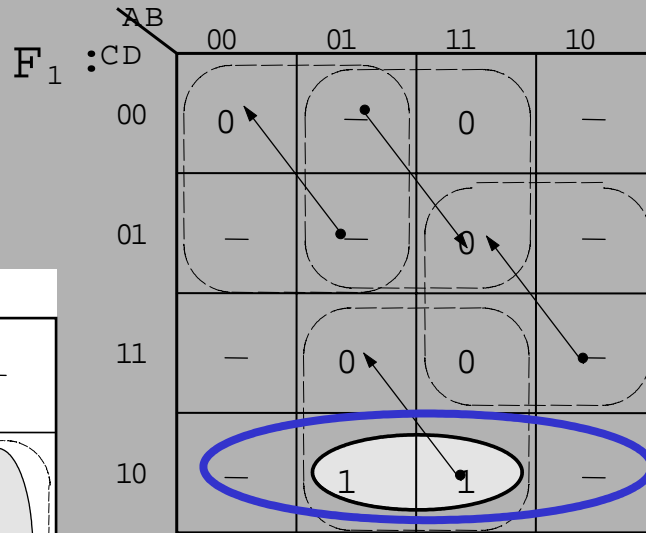
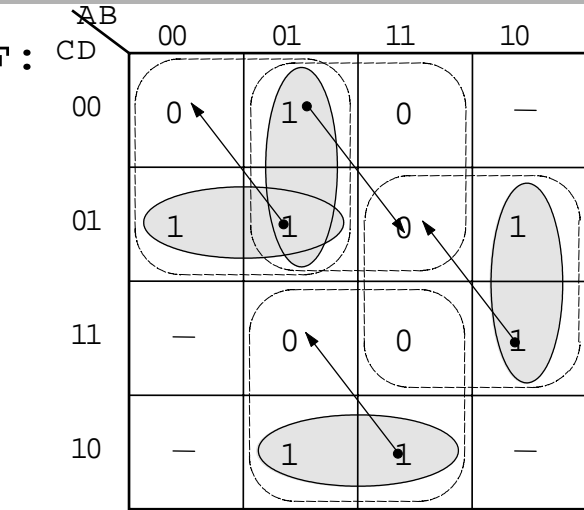
Level-0: Maximal Disjunctive Decomposition



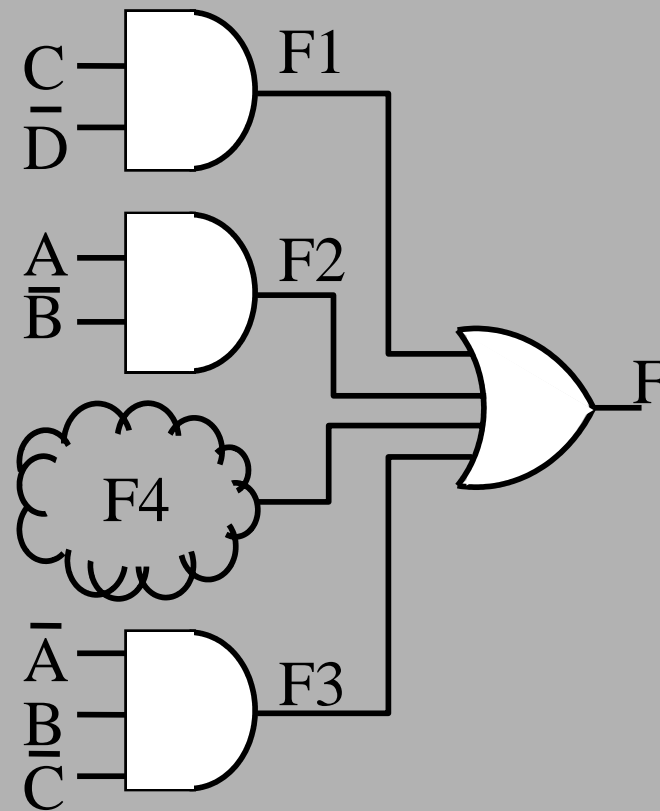
Level-0: Maximal Disjunctive Decomposition (cont.)



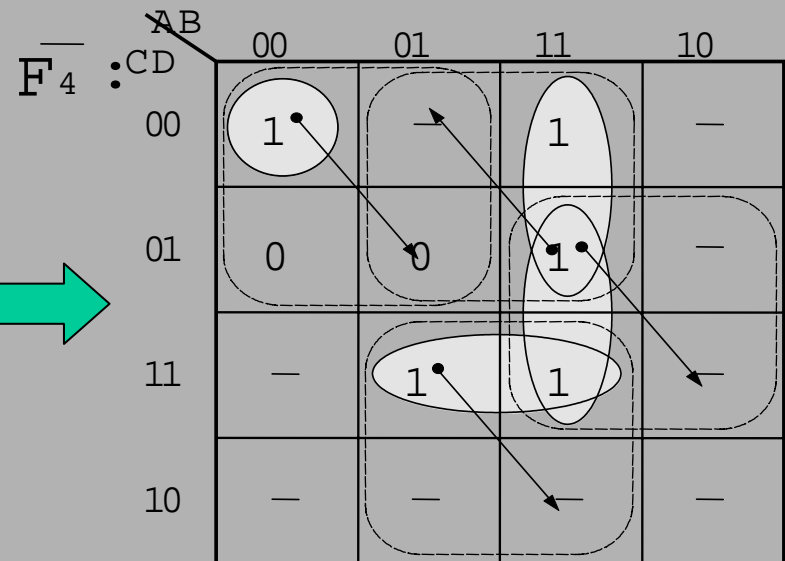
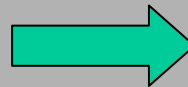
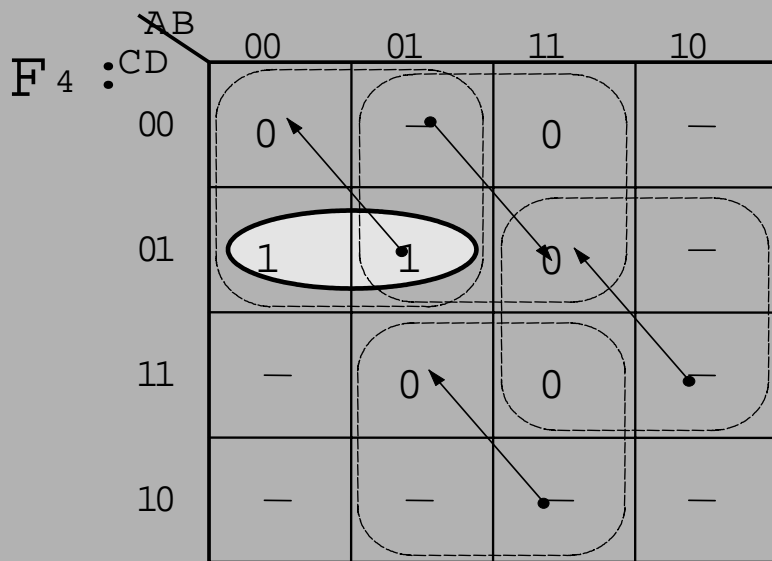
Level-0: Maximal Disjunctive Decomposition (cont.)



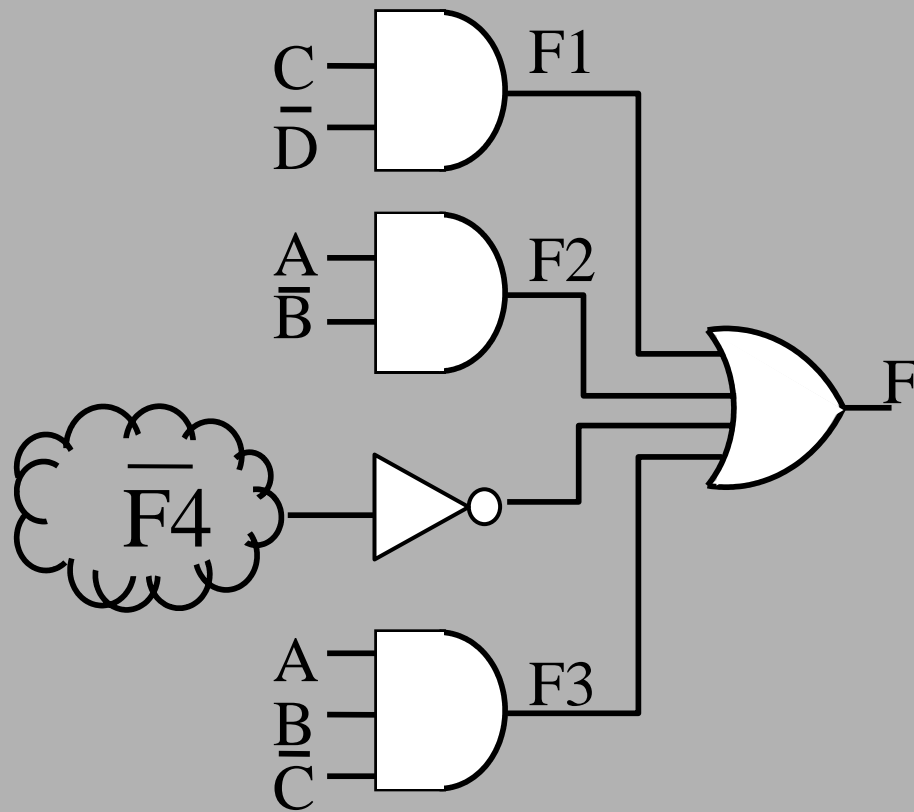
Level-0: Maximal Disjunctive Decomposition (cont.)



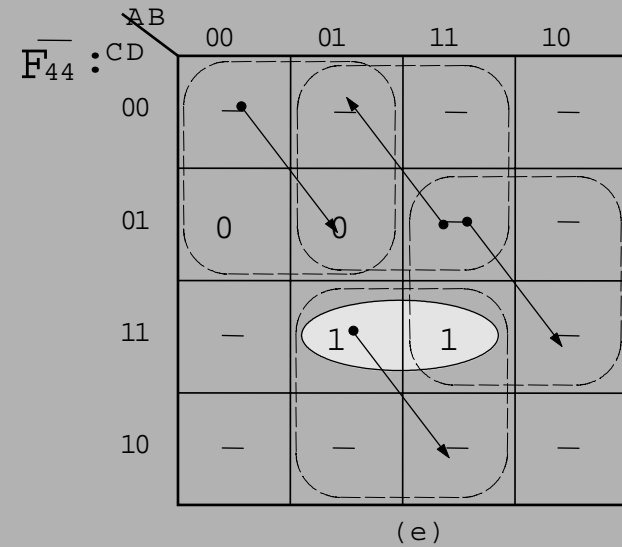
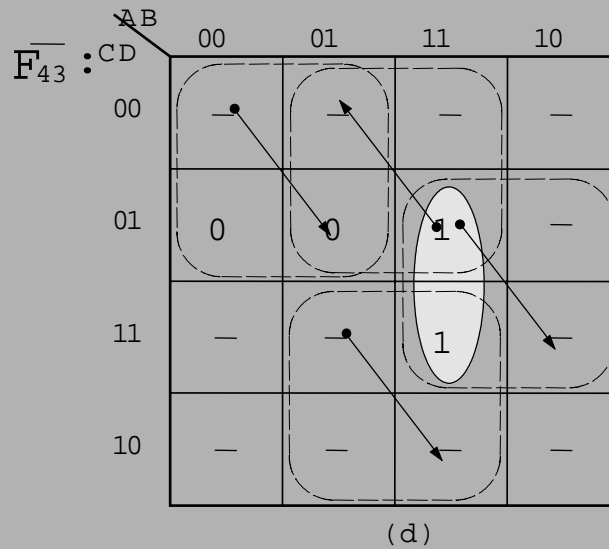
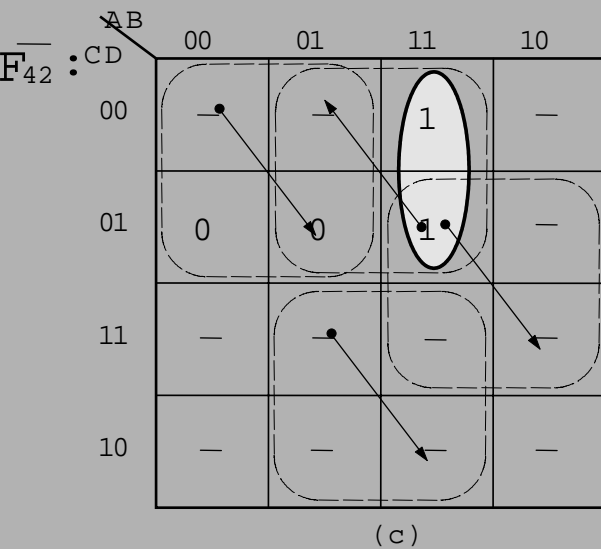
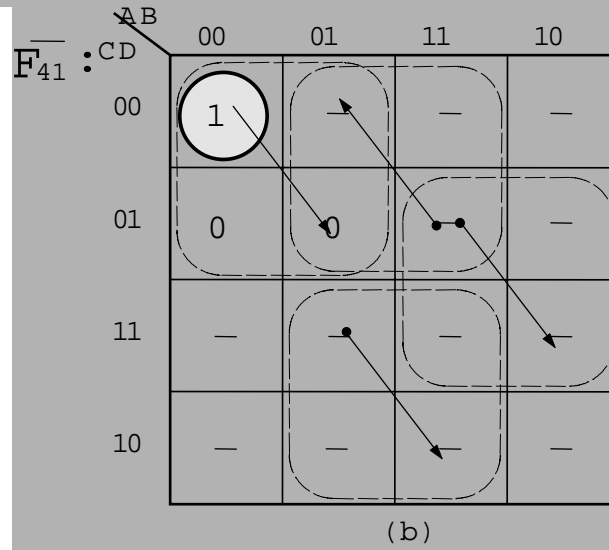
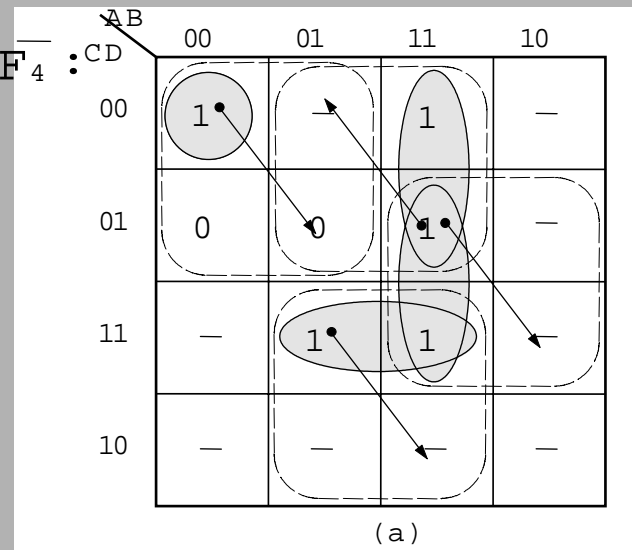
Level-1: Inversion



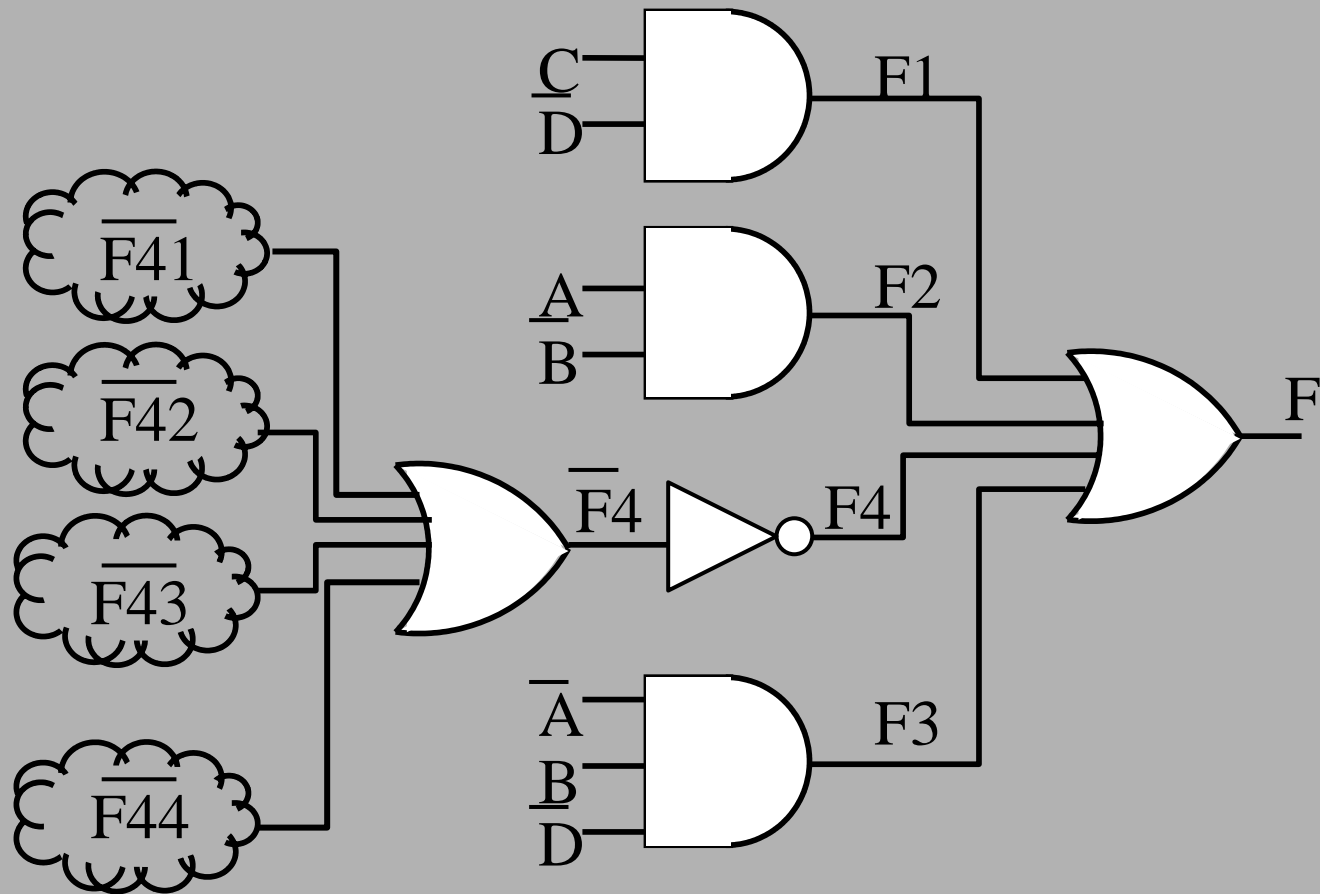
Level-1: Inversion (cont.)



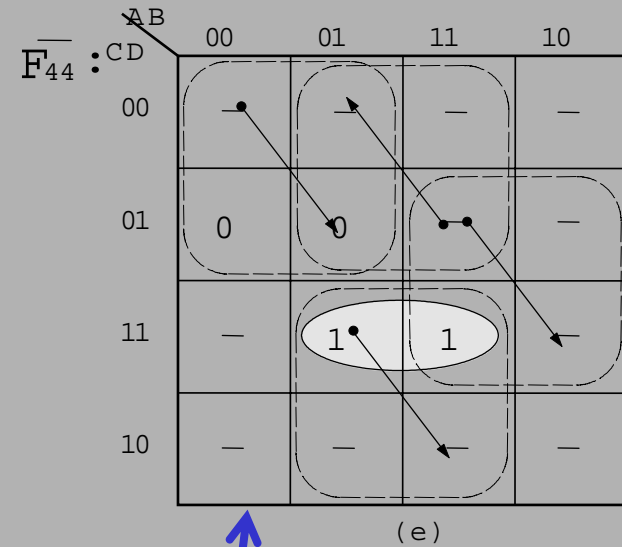
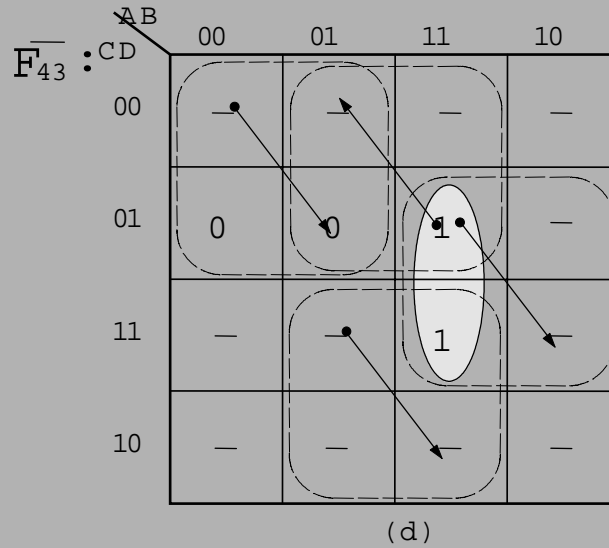
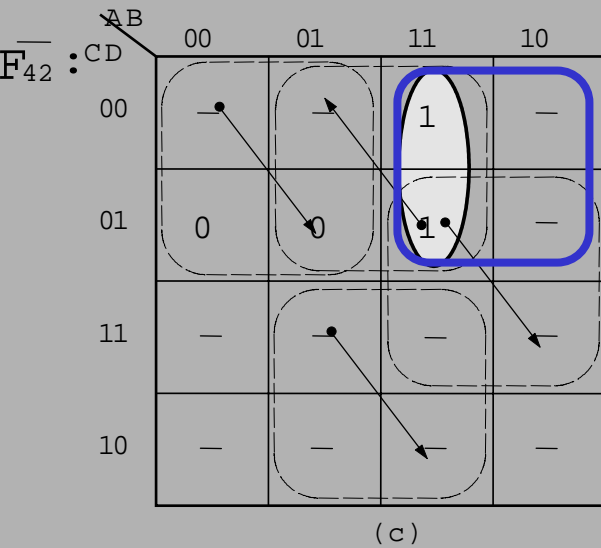
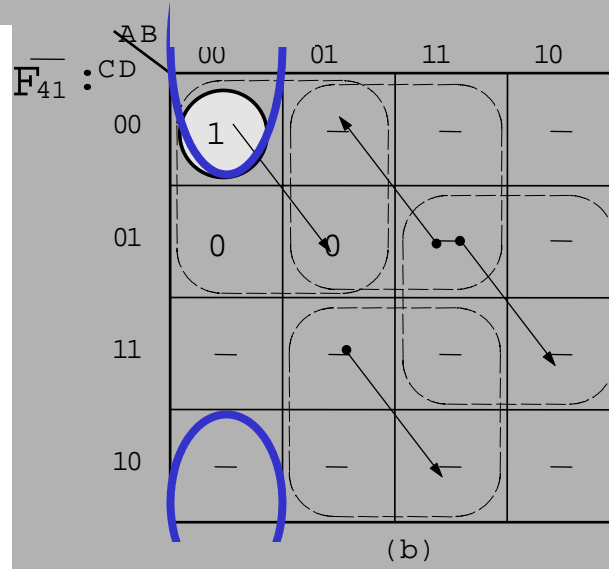
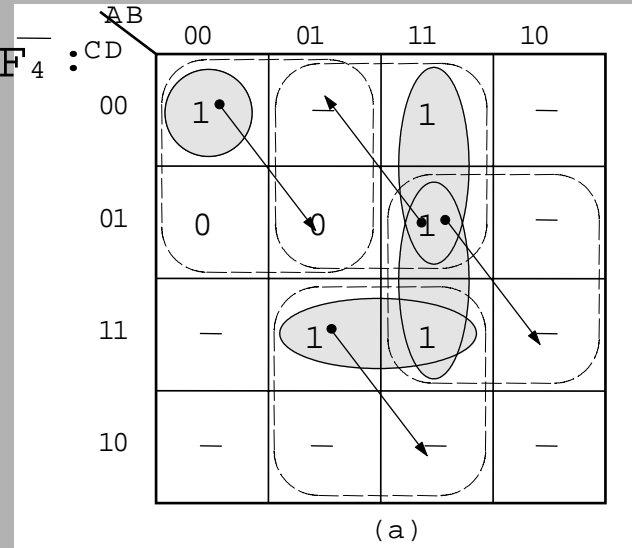
Level-1: Maximal Disjunctive Decomposition



Level-1: Maximal Disjunctive Decomposition (cont.)

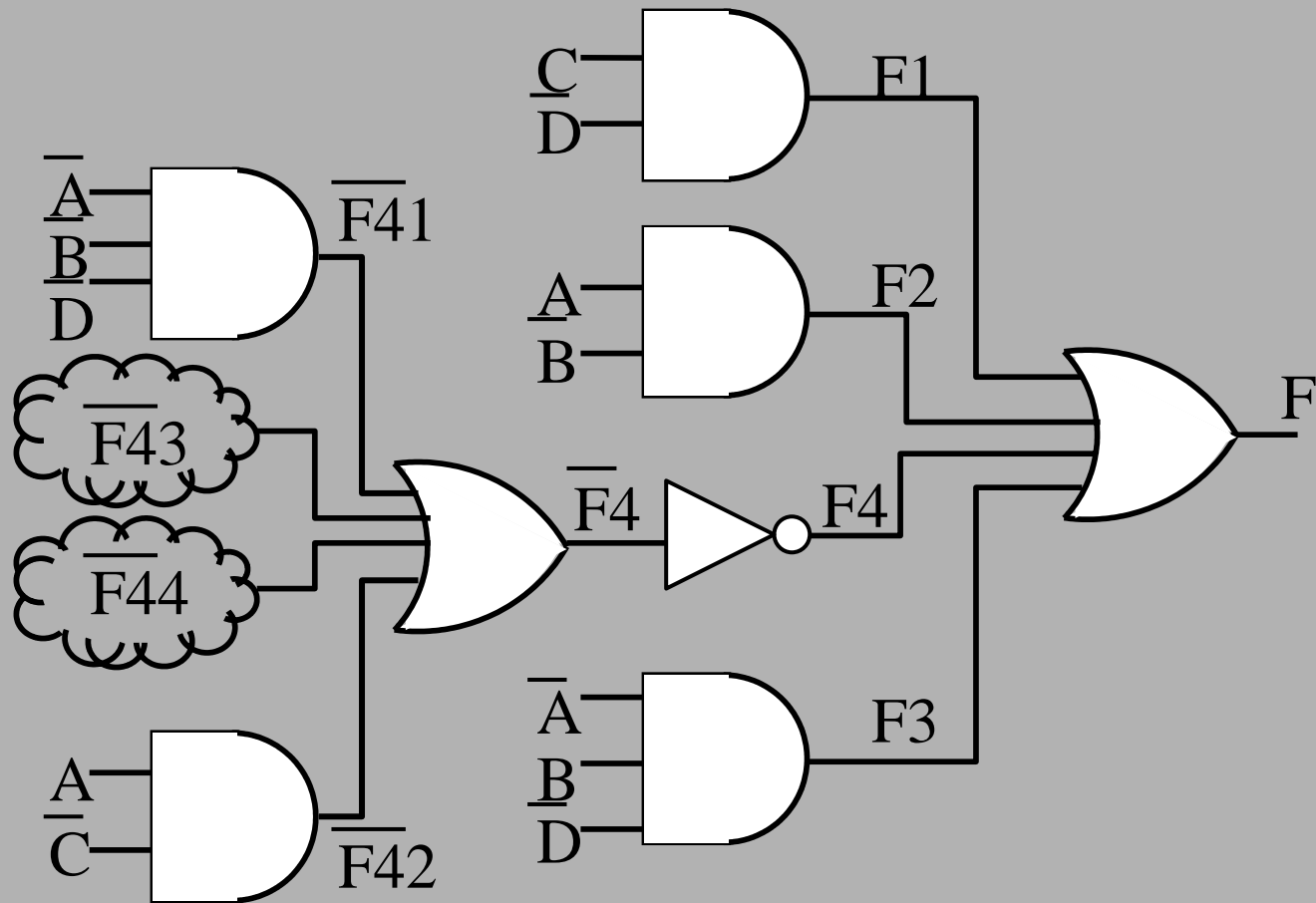


Level-1: Maximal Disjunctive Decomposition (cont.)

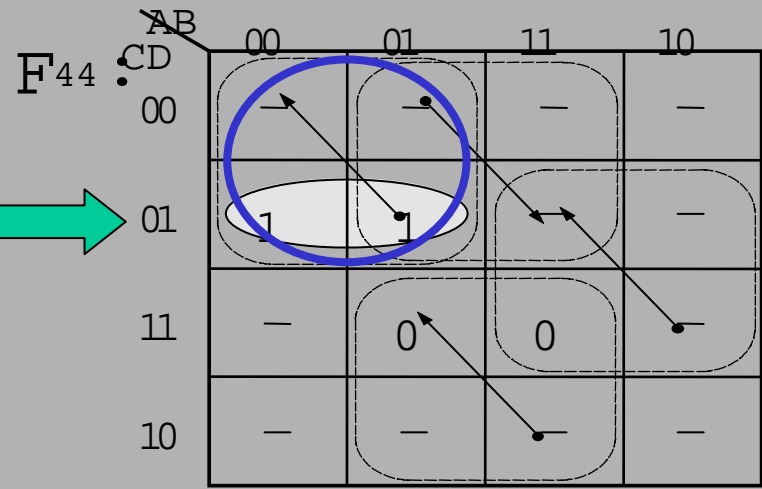
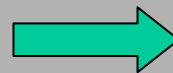
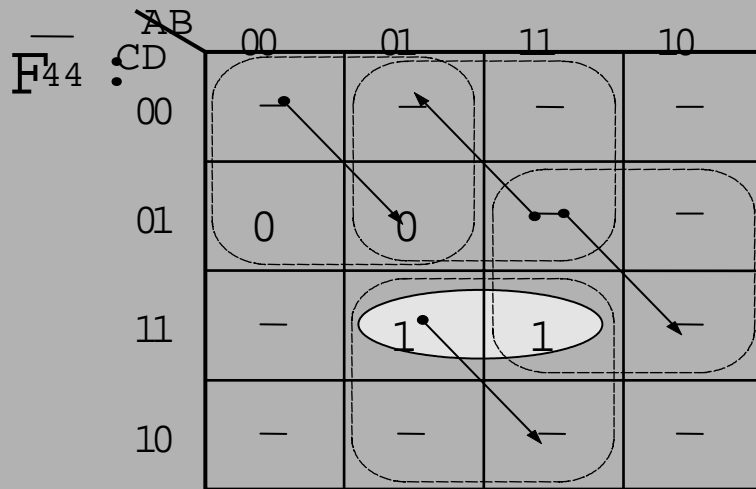
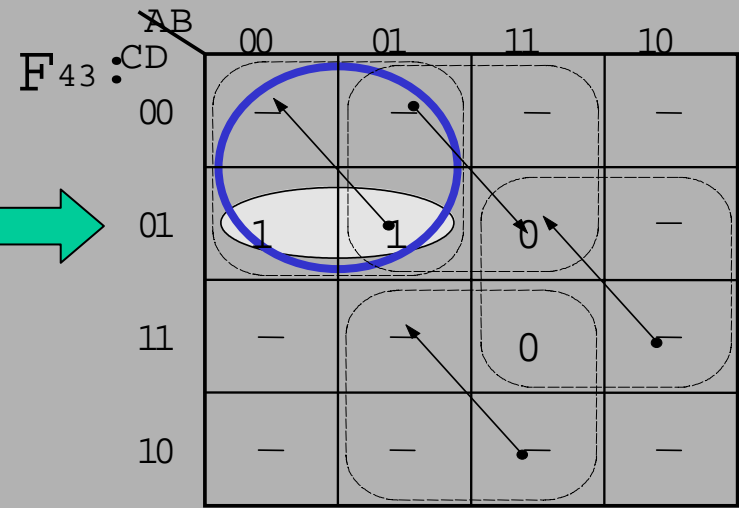
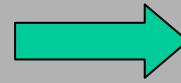
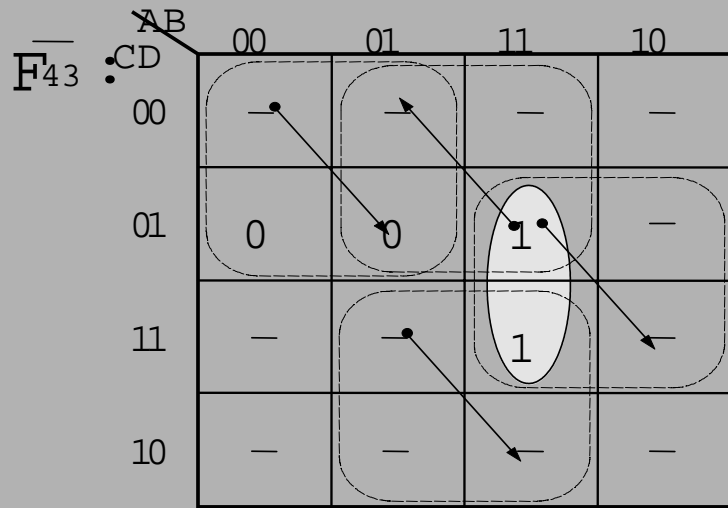


NO SOLUTION

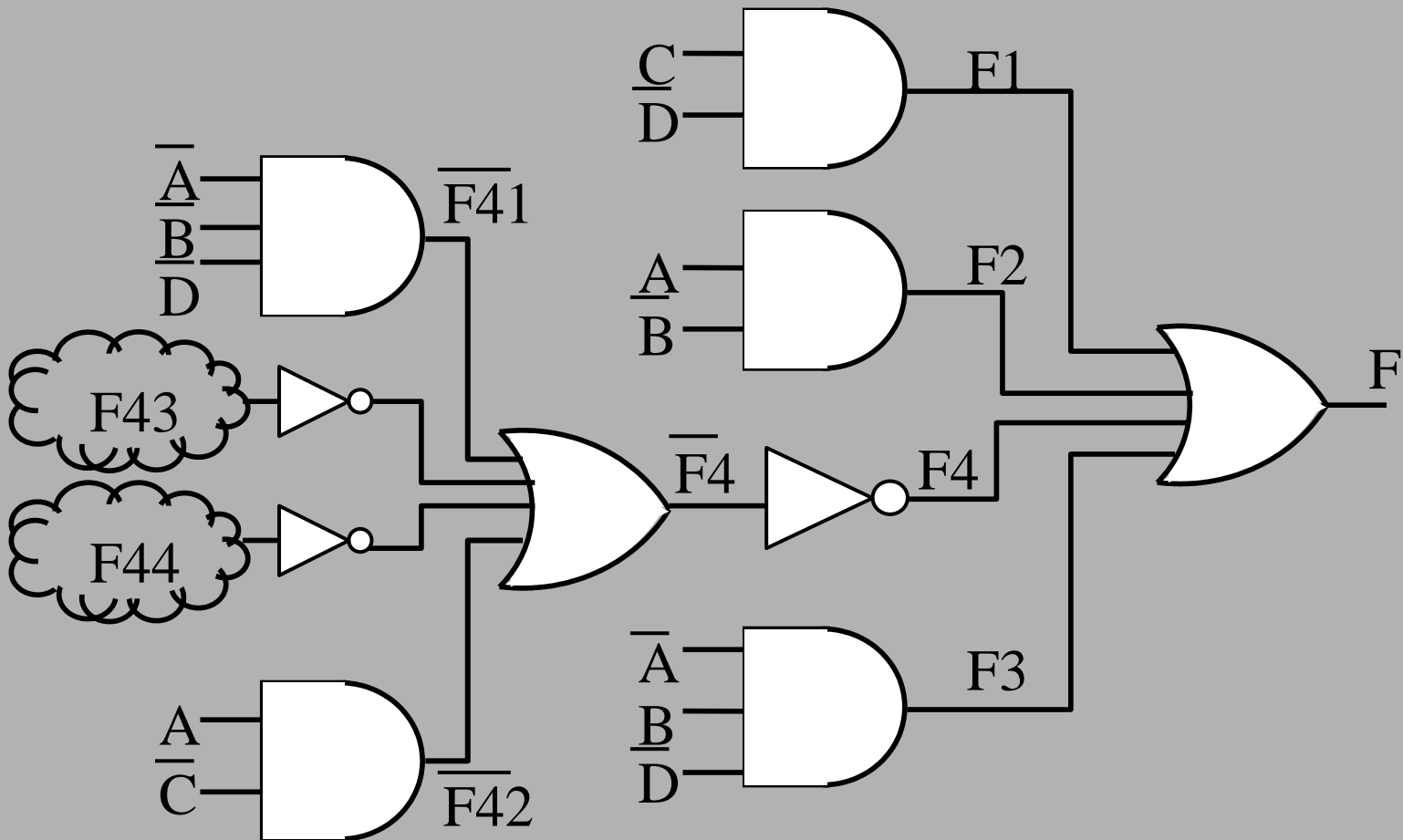
Level-1: Maximal Disjunctive Decomposition (cont.)



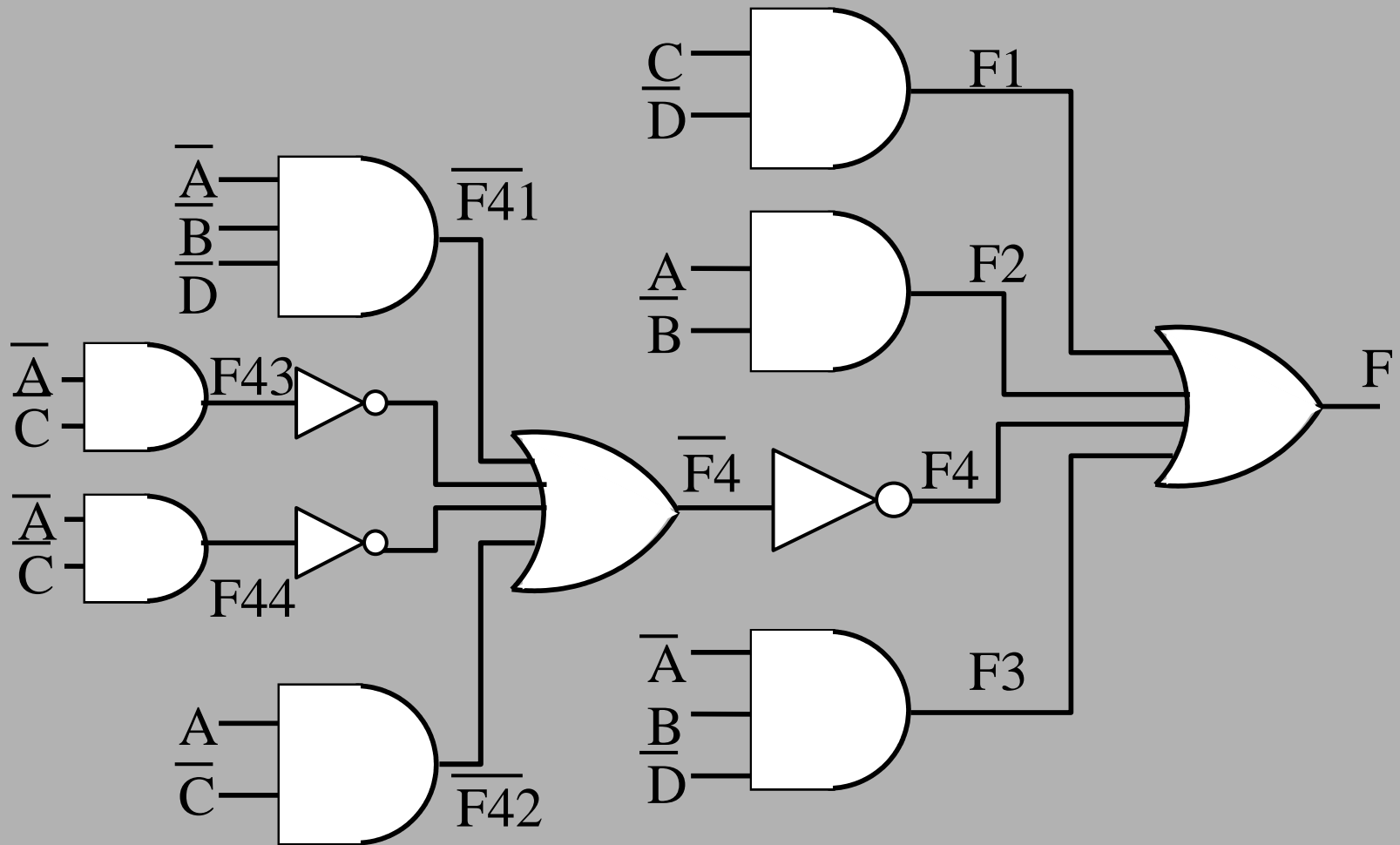
Level-2: Inversion



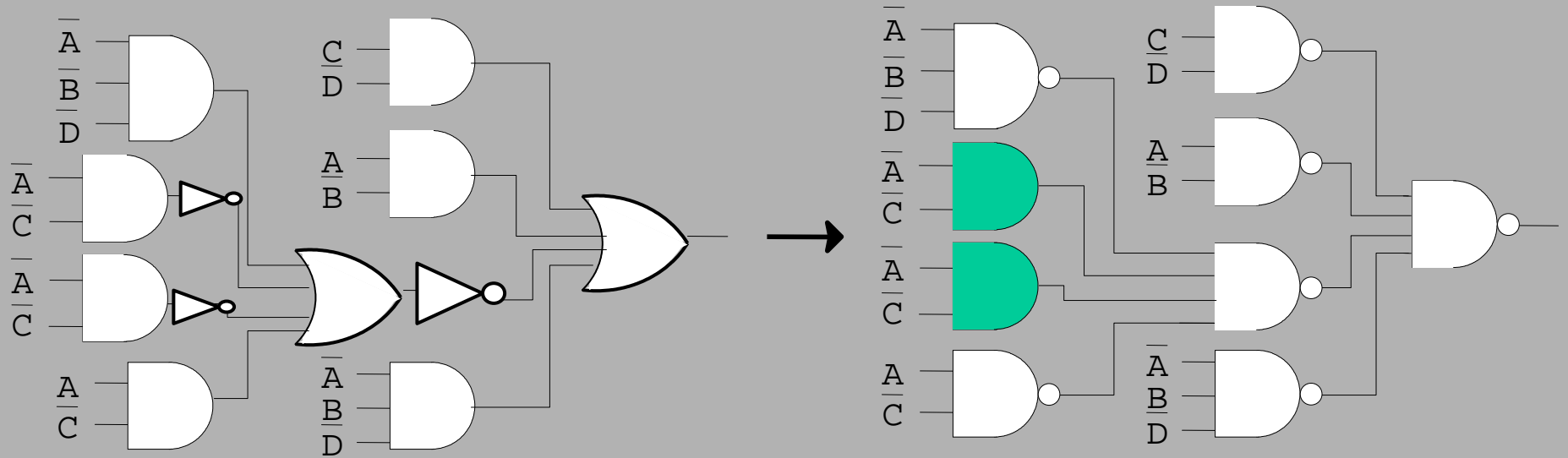
Level-2: Inversion (cont.)



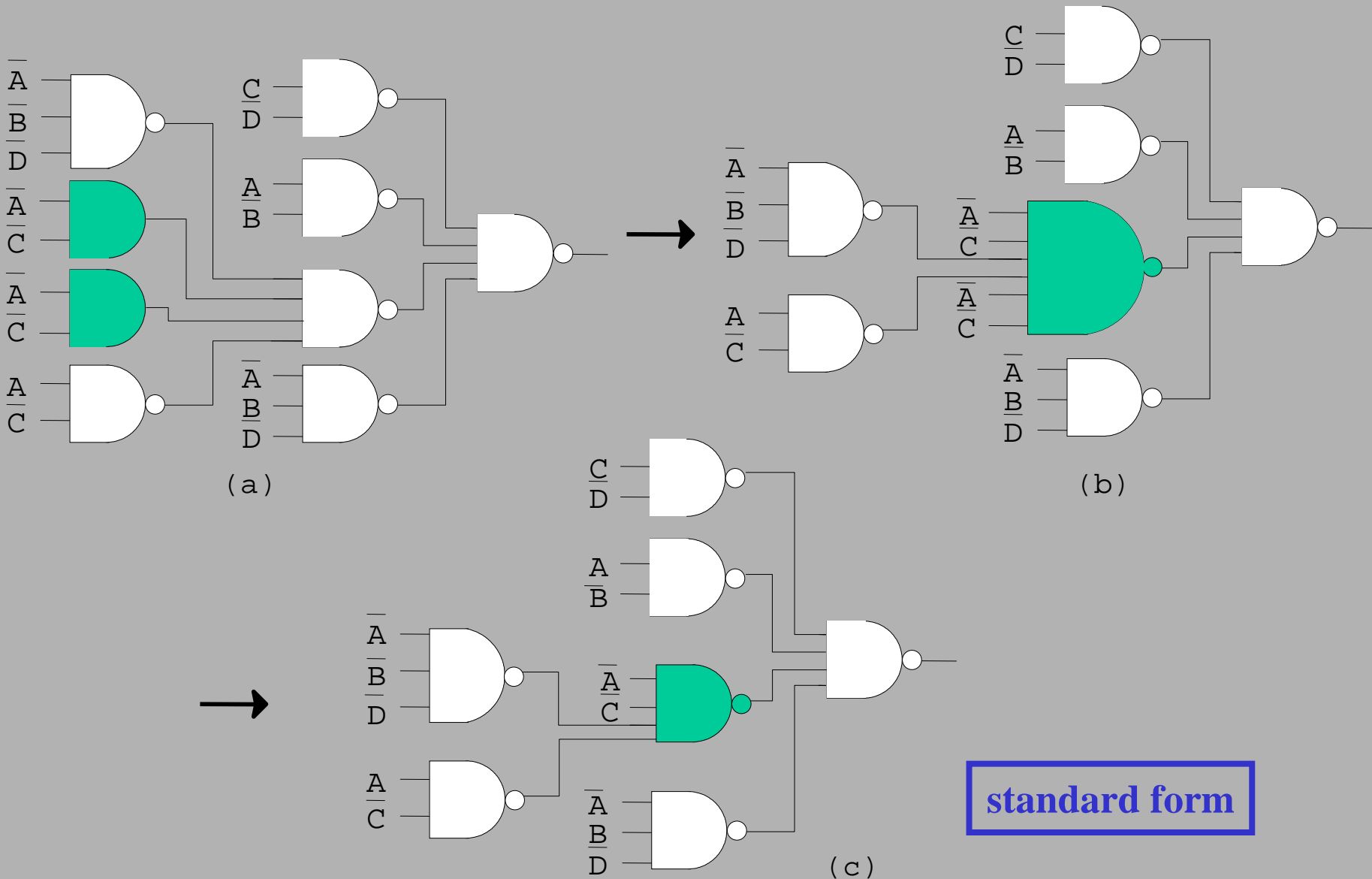
Final Hazard-Free Multi-Level Implementation



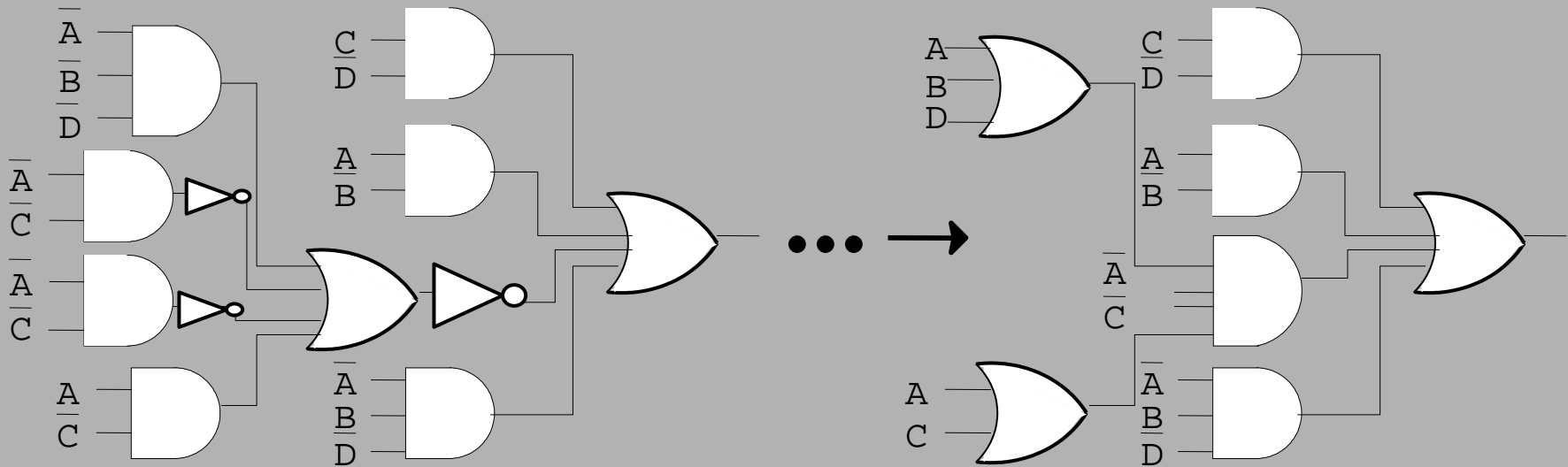
Deriving a 3-Level NAND Implementation: First Step



Deriving a 3-Level NAND Implementation: Final Step

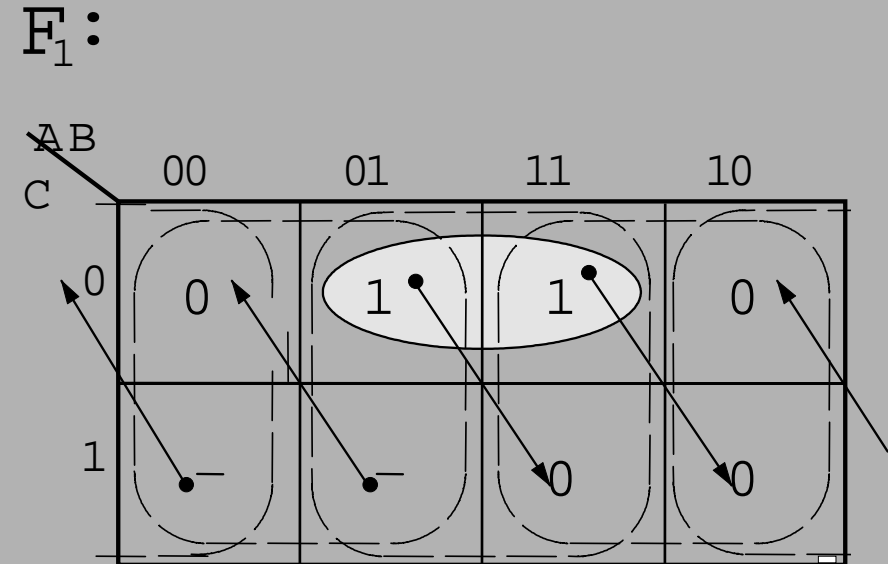
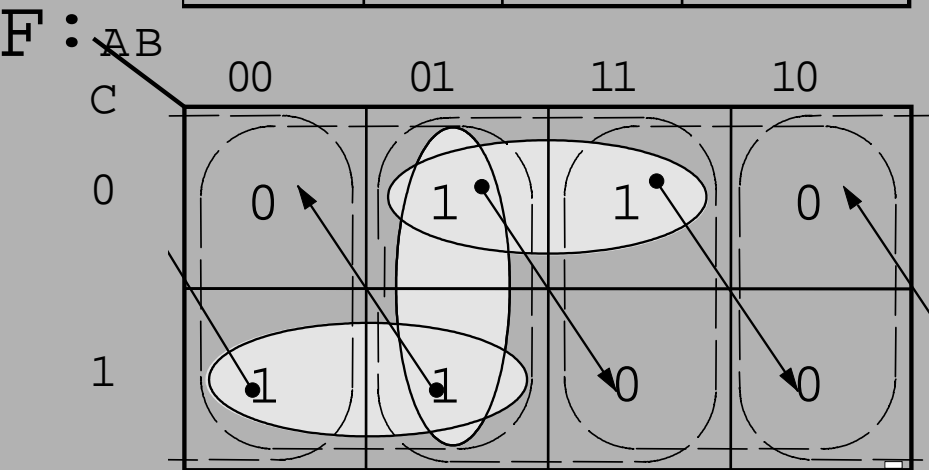
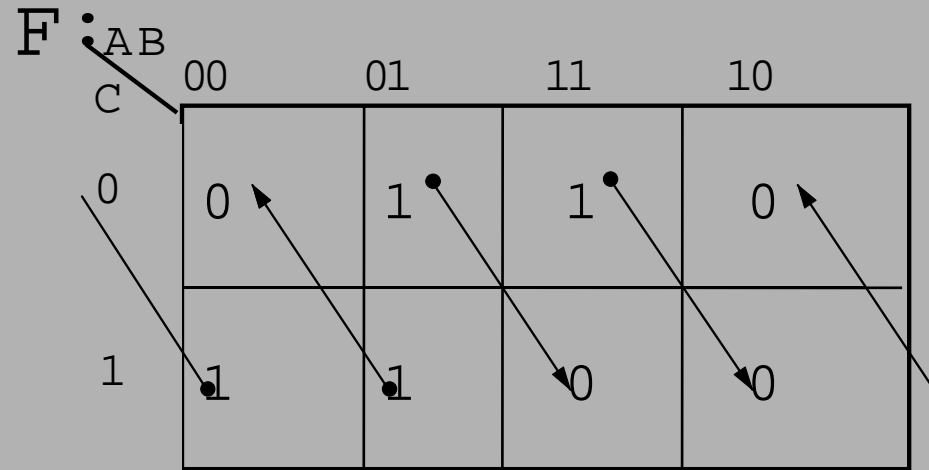


Deriving a 3-Level OR-AND-OR Implementation



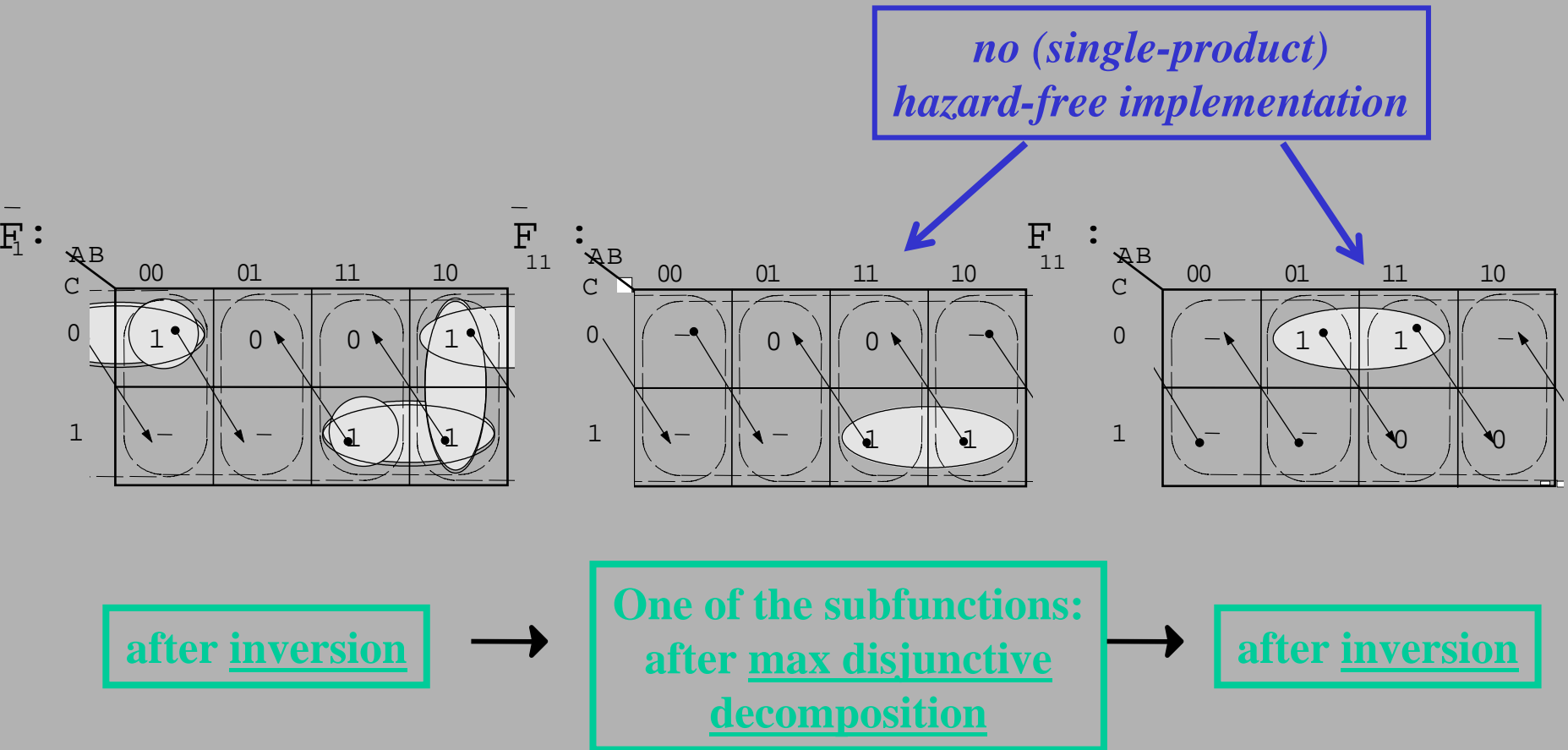
standard form

Second Example: No Hazard-Free Multi-Level Solution



One of the subfunctions:
after max disjunctive decomposition

Second Example: No Multi-Level Solution (cont.)



Conclusion: no hazard-free multi-level solution exists!

Other Contributions

Correctness of Method:

- **Soundness:** only produces hazard-free implementations
- **Completeness:** if any hazard-free multi-level solution exists, method will always produce a solution

Canonicity:

- variant of algorithm: identifies unique “representative” solution

Iterative Algorithm: (...see paper)

- **Observation:** after 3 recursive levels, sub-problems repeat
→ can terminate recursion
- Propose 3-level iterative algorithm

Summary

Given: a Boolean function F + specified set of transitions T .

If (F,T) has any hazard-free multi-level solution, then it

always will have a *hazard-free 3-level implementation*:

- 3-Level NAND
- OR-AND-OR

The proposed iterative method will always find such a solution... if one exists.

Outline

- **Background:** Hazard-Free 2-Level Logic
- **Overview of Multi-Level Method**
 - Hazard-Free Decomposition
 - Algorithm + Examples
 - Standard 3-Level Circuit Structures
 - Canonicity
- • **Related Work:** Bredeson ['74]
- **Experimental Results**
- **Conclusions**

Related Work: Bredeson [‘74]

First known approach to tackle the problem:

- **uses recursive algorithm:** alternates disjunction + inversion

Limitations:

- **does not identify that solutions are always 3-level**
- **does not address canonicity**
- **highly-restricted formulation:**
 - tight restrictions on allowed input transitions
 - limited notion of hazard-free decomposition
 - cannot handle incompletely-specified functions

Experimental Results

Prototype CAD Tool:

- Reads in Berkeley PLA format + input transition list
- Produces hazard-free multi-level implementation
 - or indicates if no solution exists

Function Name	# Inputs	Solution Exists?	Total # Subfunctns	Total # Inversions	Max Inverters per Level	Recursion Depth	CPU Time (ms)
bredeson	3	No	4	3	1	2	0.36
simple_ex	4	Yes	3	0	0	0	0.11
nowick_dill	4	Yes	8	1	1	1	0.59
z1	4	Yes	19	3	3	1	2.28
z2	4	No	12	10	4	2	1.42
z3	4	Yes	8	2	2	1	0.68
z4	4	Yes	4	0	0	0	0.18
z5	4	No	11	14	2	2	2.23
z6	4	Yes	7	3	2	2	0.99

Conclusions

- **General Constructive Method:**

- given function **F** and set of input transitions **T**, determine if any hazard-free multi-level solution exists
- based on “hazard-free decomposition”: alternate **disjunction** + **inversion**

- **Key Observation:**

- if any solution exists, then so does a 3-level solution
- target two forms: (a) **3-level NAND**, (b) **OR-AND-OR**

- **Potential Practical Benefits:**

- can use as starting point of hazard-free logic synthesis
- can simplify conservative restrictions in sequential synthesis
- provides deeper understanding of nature of hazard-free logic....