

# Block Coordinate Descent and Exact Minimization

Jelena Diakonikolas

**Boston University** 

joint work with Lorenzo Orecchia (BU)

Workshop on Local Algorithms June 2018

#### Full-Gradient First-Order Convex Optimization

Unconstrained convex minimization:

$$\begin{array}{ll} \min & f(\mathbf{x}) \\ \text{s.t.} & \mathbf{x} \in \mathbb{R}^N \end{array}$$



First-order blackbox (oracle) model:



#### History

$$\begin{array}{ll} \min & f(\mathbf{x}) \\ \text{s.t.} & \mathbf{x} \in \mathbb{R}^N \end{array}$$

- Methods with optimal iteration complexity in various settings are well-known:
  - Gradient descent folklore
  - Nemirovski's mirror descent [Nemirovski, Yudin'83]
  - Nesterov's accelerated method (AGD) [Nesterov'83]
  - Frank-Wolfe methods [Frank, Wolfe'56]
  - ... and many more books: [Bubeck'14], [Sra, Nowozin, Wright'11]
- Typical complexity of an iteration is near-linear in the input size, few iterations
- Particularly attractive for large-scale problems; broad applications in machine learning and TCS

#### Block Coordinate Descent: Setting

• Fix a partition of the vector of variables into *n* blocks:



#### Block Coordinate Descent: Setting

• Fix a partition of the vector of variables into *n* blocks:

Block 1	Block 2	Block 3	Block 4
1 1 1	2 2 2	3 3 3	4 4 4

# Block Coordinate Descent: Setting

• Fix a partition of the vector of variables into *n* blocks:



• Assume access to two types of oracles:



#### Assumptions about the Problem

- Assumptions:
  - Function is differentiable and *L*-smooth:

$$\|\nabla f(\mathbf{x}) - \nabla f(\mathbf{y})\|_* \le L \|\mathbf{x} - \mathbf{y}\|, \ \forall \mathbf{x}, \mathbf{y}$$

• Each block *i* is *L<sub>i</sub>*-smooth:

$$\|\nabla_i f(\mathbf{x}) - \nabla_i f(\mathbf{y})\|_* \le L_i \|\mathbf{x}_i - \mathbf{y}_i\|, \ \forall \mathbf{x}, \mathbf{y}, \ \text{where} \ \mathbf{y}_k = \mathbf{x}_k \ \text{for} \ k \ne i$$

• Block *n* is "least" smooth, possibly with  $L_n = \infty$ :

$$L_n = L_{\max} = \max_{1 \le i \le n} L_i$$

Note that if  $L_n = \infty$ , then it must be  $L = \infty$ !

Basic (Nonaccelerated) Methods

- Almost a folklore method [Ortega & Rheinboldt, 1970]
- Fix a (possibly random) permutation of the blocks
- Take either exact minimization or a gradient step over a block selected in the cyclic order



- Almost a folklore method [Ortega & Rheinboldt, 1970]
- Fix a (possibly random) permutation of the blocks
- Take either exact minimization or a gradient step over a block selected in the cyclic order



- Almost a folklore method [Ortega & Rheinboldt, 1970]
- Fix a (possibly random) permutation of the blocks
- Take either exact minimization or a gradient step over a block selected in the cyclic order



- Almost a folklore method [Ortega & Rheinboldt, 1970]
- Fix a (possibly random) permutation of the blocks
- Take either exact minimization or a gradient step over a block selected in the cyclic order



- Almost a folklore method [Ortega & Rheinboldt, 1970]
- Fix a (possibly random) permutation of the blocks
- Take either exact minimization or a gradient step over a block selected in the cyclic order



Example order: 1, 3, 2, 4

Dependence of the optimality gap on smoothness parameters:

$$L_n + \frac{\min(nL^2, (\sum_{i=1}^n L_i)^2)}{L_{\min}}$$
 [Sun, Hong'15],  $Ln^3$  [Hong, Wang, Razaviyayn, Luo'17]

#### **Alternating Minimization**

- A special case of cyclic BCD when there are only two blocks
- Exact minimization on the less smooth block; exact minimization or gradient descent step on the other block



Block 2					
2	2	2	2	2	2

#### **Alternating Minimization**

- A special case of cyclic BCD when there are only two blocks
- Exact minimization on the less smooth block; exact minimization or gradient descent step on the other block



Block 2					
2	2	2	2	2	2

#### Alternating Minimization

- A special case of cyclic BCD when there are only two blocks
- Exact minimization on the less smooth block; exact minimization or gradient descent step on the other block





Dependence of the optimality gap on smoothness parameters:

 $\min(L_1, L_2)$  [Beck'15]

- Introduced by [Nesterov, 2012]
- Fix a probability distribution  $\{p_i\}_{i=1}^n$  over the blocks

i		
	1	

Block 1	Block 2	Block 3	Block 4
1 1 1	2222	3 3 3	4 4 4

- Introduced by [Nesterov, 2012]
- Fix a probability distribution  $\{p_i\}_{i=1}^n$  over the blocks

i		
	1	

Block 1	Block 2	Block 3	Block 4
1 1 1	2222	3 3 3	4 4 4

- Introduced by [Nesterov, 2012]
- Fix a probability distribution  $\{p_i\}_{i=1}^n$  over the blocks

i		
	3	

Block 1	Block 2	Block 3	Block 4
1 1 1	222	3 3 3	4 4 4

- Introduced by [Nesterov, 2012]
- Fix a probability distribution  $\{p_i\}_{i=1}^n$  over the blocks



Block 1	Block 2	Block 3	Block 4
1 1 1	222	3 3 3	4 4 4

- Introduced by [Nesterov, 2012]
- Fix a probability distribution  $\{p_i\}_{i=1}^n$  over the blocks



• Select block i with probability  $p_i$  and do a gradient descent step or exact minimization over it



Dependence of the optimality gap on smoothness parameters:

 $\sum_{i=1}^{n} L_i$  for  $p_i \sim L_i$  [Nesterov'12]

#### Dependence of the optimality gap on smoothness parameters:

• Cyclic block coordinate descent (*n* blocks):

$$L_n + \frac{\min(nL^2, (\sum_{i=1}^n L_i)^2)}{L_{\min}}$$
 [Sun, Hong'15],  $Ln^3$  [Hong, Wang, Razaviyayn, Luo'17]

• Randomized block coordinate descent (*n* blocks):

 $\sum_{i=1}^{n} L_i$  for  $p_i \sim L_i$  [Nesterov'12]

• Alternating minimization (2 blocks):

 $\min(L_1, L_2)$  [Beck'15]

So far, only alternating minimization (two blocks) can avoid paying for the less-smooth block

**Q**. Can we avoid paying for the least-smooth block when there are more than two blocks?

- ... or how to avoid paying for the least smooth (non-smooth) block
- Fix a probability distribution  $\{p_i\}_{i=1}^{n-1}$  over blocks 1, 2, ..., n-1



i

• Do a gradient descent (or exact min) step over block *i*, then exact minimization over block *n* 



- ... or how to avoid paying for the least smooth (non-smooth) block
- Fix a probability distribution  $\{p_i\}_{i=1}^{n-1}$  over blocks 1, 2, ..., n-1

i

• Do a gradient descent (or exact min) step over block *i*, then exact minimization over block *n* 



- ... or how to avoid paying for the least smooth (non-smooth) block
- Fix a probability distribution  $\{p_i\}_{i=1}^{n-1}$  over blocks 1, 2, ..., n-1



i

• Do a gradient descent (or exact min) step over block *i*, then exact minimization over block *n* 



- ... or how to avoid paying for the least smooth (non-smooth) block
- Fix a probability distribution  $\{p_i\}_{i=1}^{n-1}$  over blocks 1, 2, ..., n-1

i

• Do a gradient descent (or exact min) step over block *i*, then exact minimization over block *n* 



- ... or how to avoid paying for the least smooth (non-smooth) block
- Fix a probability distribution  $\{p_i\}_{i=1}^{n-1}$  over blocks 1, 2, ..., n-1



• Do a gradient descent (or exact min) step over block *i*, then exact minimization over block *n* 



- ... or how to avoid paying for the least smooth (non-smooth) block
- Fix a probability distribution  $\{p_i\}_{i=1}^{n-1}$  over blocks 1, 2, ..., n-1



• Do a gradient descent (or exact min) step over block *i*, then exact minimization over block *n* 



Generalizes randomized BCD and alternating minimization

Dependence of the optimality gap on smoothness parameters:  $\sum_{i=1}^{n-1} L_i$  for  $p_i \sim L_i$  (no dependence on  $L_n$ !)

**Possible to accelerate**: gives the same convergence time as the fastest known accelerated BCD – NUACDM [Allen-Zhu, Qu, Richtárik, Yuan'16], except without any dependence on  $L_n$ 

#### Convergence Analysis: Main Ideas

• Extension of Approximate Duality Gap Technique [D, Orecchia, 2017]



 $f(\mathbf{x}^k) - f(\mathbf{x}^*) \le G_k$ . Let  $A_k$  be an increasing (rate) function of iteration count k. Then, if  $A_k G_k \le A_{k-1} G_{k-1}$ ,  $\forall k$ 

$$f(\mathbf{x}^k) - f(\mathbf{x}^*) \le G_k \le \frac{A_0 G_0}{A_k}$$

#### Convergence Analysis: Main Ideas

• Lower bound uses the following:

•

by conversity (and differentiability) W

• by convexity (and differentiability), 
$$\forall \mathbf{x}_k$$
:  

$$f(\mathbf{x}^*) \ge f(\mathbf{x}^k) + \langle \nabla f(\mathbf{x}^k), \mathbf{x}^* - \mathbf{x}^k \rangle$$

$$= f(\mathbf{x}^k) + \sum_{i=1}^n \langle \nabla_i f(\mathbf{x}^k), \mathbf{x}^*_i - \mathbf{x}^k_i \rangle$$

$$= f(\mathbf{x}^k) + \sum_{i=1}^{n-1} \langle \nabla_i f(\mathbf{x}^k), \mathbf{x}^*_i - \mathbf{x}^k_i \rangle$$
because  $\nabla_i f(\mathbf{x}^k) = \mathbf{0}$ 
Upper bound uses the (block) gradient descent progress:  

$$f(\operatorname{Grad}_i(\mathbf{x}^k)) \le f(\mathbf{x}^k) - \frac{1}{2L_i} \|\nabla f(\mathbf{x}_k)\|_*^2$$
can sample only over the first  $n - 1$  (i.e., "smoother") blocks

Numerical Experiments

#### Experiments: Linear Regression on BlogFeedback Dataset



#### Experiments: Linear Regression on BlogFeedback Dataset



#### Summary

- A novel block coordinate descent method (and its accelerated version) that can handle a completely non-smooth block (structured non-smoothness)
- The method outperforms existing methods if one block has much worse (but finite) smoothness parameter than the remaining ones

#### • Ongoing work:

- Extension to the smooth and strongly convex setting
- Extension to the composite non-smooth setting
- Improved convergence bounds for randomized BCD with exact minimization
- Open question:
  - We need to know which block is the least smooth to not pay for it. Is it possible to relax this?

ielena@ielena-diakonikolas.com www.ielena-diakonikolas.com

