

Today

FACTORIZATION OF POLYNOMIALS

- Quadratics
- Linear Factors
- General Case

x

Simple Case: $q = \text{odd}$;

Input: $x^2 + ax + b$ [= $(x - \alpha) \cdot (x - \beta)$]
"random"

Goal: $x - \alpha, x - \beta$

- Key to all factorization alg (over $\mathbb{F}_q[x]$)

$$x^q - x = x \cdot (x^{\frac{q-1}{2}} - 1) \cdot (x^{\frac{q-1}{2}} + 1)$$

- Why is this useful?

- Contains all linear factors
- Sparse

- Idea / Hope:

$$(x - \alpha) \mid x^{\frac{q-1}{2}} - 1$$

but not $x - \beta$

- Assuming above

$$\gcd(x^{\frac{q-1}{2}} - 1, x^2 + ax + b) = x - d$$

- Implementation

- Compute $x^{\frac{q-1}{2}} \pmod{x^2 + ax + b} = g(x)$
by repeated squaring

- Compute $\gcd(q-1, (x^2+ax+b))$

- Works if $\alpha = \mathbb{QR}$
 $\beta \neq \mathbb{QR}$

• General Case? ... shortly

Aside: $q = 2^l$

Use the identity

$$(x^q - x) = \text{Tr}(x) (\text{Tr}(x) - 1)$$

where

$$\text{Tr}(x) = x + x^2 + x^4 + \dots + x^{q/2}$$

Details omitted (easy, no surprises)

Rest of lecture - odd case

GENERAL QUADRATIC : $x^2 + ax + b = f$

Case 1: No linear factors?

Test: $\gcd(f, x^2 - x) = ?$

✓ • $\gcd = 1 \Rightarrow f$ is irreducible

? • $\gcd = f \Rightarrow f$ has two distinct linear factors.

✓ • $\gcd = (x - \alpha) \Rightarrow f = (x - \alpha)^2$

Only need to deal with middle case.

Remaining Case: $f = (x - \alpha)(x - \beta)$

$$\alpha \neq \beta.$$

Idea: Map roots to two random pts.

$$\alpha \longrightarrow \alpha \cdot c + d$$

$$\beta \longrightarrow \beta \cdot c + d$$

How? Let

$$h(x) = f\left(\frac{x-d}{c}\right).$$

$$\begin{aligned} \text{Then } h(x) &= \left(\frac{x-d}{c} - \alpha\right)\left(\frac{x-d}{c} - \beta\right) \\ &= \frac{1}{c^2} (x - (\alpha c + d))(x - (\beta c + d)) \end{aligned}$$

Key Observation:

$$c \in \mathbb{F}_q^* ; d \in \mathbb{F}_q \text{ random}$$

$\Rightarrow \alpha c + d$ & $\beta c + d$ are
random & independent.

$$\Rightarrow \text{w.p.} \geq \frac{1}{4} \quad \alpha c + d = QR$$

$$\beta c + d \neq QR$$

\Rightarrow Reduces to simple case!

Algorithm dates back to [Berlekamp '72]

Algorithm Summarized

FACTOR ($f(x) = x^2 + ax + b$)

- if $\text{gcd}(f, x^2 - x) = 1 \Rightarrow$ irreducible
- " " = $(x - \alpha) \Rightarrow f = (x - \alpha)^2$
- Else

Repeat till success

- Pick $c \leftarrow \mathbb{F}^*$, $d \leftarrow \mathbb{F}$ randomly

- if $\text{gcd}\left(f\left(\frac{x-d}{c}\right), x^{\frac{q-1}{2}} - 1\right) = (x - \delta)$

output $\left(f_1 = x - \left(\frac{\delta - d}{c}\right); f/f_1\right) \leftarrow \text{STOP}$

- continue

Towards higher-degree polynomials

Case 1: f splits into linear factors.

- Above algorithm still makes progress.

Case 2: f has irreducible factors of distinct degrees.

- Use the fact that

f_i irred. of deg d

$\Leftrightarrow d$ is the least integer s.t.

$$f_i \mid x^{q^d} - x$$

- Fact above + gcd's give non-trivial factorization ... makes progress.

Case 3: All n ^{irred.} factors of f have same degree d , but some maybe repeated.

Calculus in Algebra:

Definition

- Derivative of $x^i = i \cdot x^{i-1}$
- Extend linearly $(\alpha x^i)' = i \alpha x^{i-1}$
- $(f+g)' = f' + g'$

Prop.

- Product rule holds: $(f \cdot g)' = f'g + g'f$
- $g^2 \mid f \Rightarrow g \mid f'$
 $\Rightarrow g \mid \gcd(f, f')$
- $\left. \begin{array}{l} g^i \nmid f \\ \& g^{i-1} \mid f \end{array} \right\} \Rightarrow g^{i-1} \nmid \gcd(f, f')$.
(for irred. g)

Use Calculus / gcd (f, f') to eliminate repeated factors.

Case 4: (the interesting case)

f has distinct factors all of same degree d .

Inspiration:

- should be able to factor f over extension field of size q^d .
- f will have linear factors there.
- Use idea of minimal polynomials to get back solution over \mathbb{F}_q

Hurdles:

- Don't have \mathbb{F}_{q^d} explicitly.
- Computing minimal poly? How?

Direct approach: $f = f_1 \cdot f_2$
 $\deg(f_i) = d.$

- Pick random poly g of $\deg \leq 2d-1$

- Chinese Remainder Theorem

$$\Rightarrow (g \bmod f_1), (g \bmod f_2)$$

uniquely specify $g = (g \bmod f)$

$\Rightarrow g_1 = g \bmod f_1 \ \& \ g_2 = g \bmod f_2$
are random & independent.

- W.P. $\geq \frac{1}{4}$

g_1 is a QR in $\mathbb{F}_q[x]/f_1$

g_2 is ^{not} a QR in $\mathbb{F}_q[x]/f_2$

$\Rightarrow \gcd(g^{\frac{q-1}{2}} - 1, f) = \text{non-trivial}$
 \boxtimes

Summary

- Can factor poly of deg n in \mathbb{F}_q using $\text{poly}(n, \log q)$ operations over \mathbb{F}_q
- # operations depends on $q!$
- Next lectures:
 - ① Deterministic algorithms running in time $\text{poly}(n, l, p)$ if $q = p^l$
 - ② Factorization of bivariate polynomials
 - ③ Factorization of rational " " .