

Lecture 10

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Today, we will continue our approach to factoring bivariate polynomials. We will first focus on the tool of Hensel's Lifting; and then describe how we perform the factoring.

1 Hensel's Lifting

Suppose $f(x, y) = g(x)h(x) \pmod{y}$. We wanted a factorization for higher powers of y : $f(x, y) = \tilde{g}(x, y)\tilde{h}(x, y) \pmod{y^{2k}}$. Hensel's Lifting says we can obtain this if g and h are "relatively prime," and that the \tilde{g} and \tilde{h} we obtain are essentially unique. Formally:

Lemma 1 (Hensel's Lifting) *If R is a ring, $I \subseteq R$ is an ideal, and there exist f, g, h, a, b in R such that*

$$(H1) \quad f = gh \pmod{I}$$

$$(H2) \quad ag + bh = 1 \pmod{I}.$$

Then, for every positive integer s that is a power of 2, there exist \tilde{g}, \tilde{h} and \tilde{a}, \tilde{b} in R , such that

$$(C1) \quad f = \tilde{g}\tilde{h} \pmod{I^s}$$

$$(C2) \quad \tilde{g}\tilde{a} + \tilde{b}\tilde{h} = 1 \pmod{I^s}$$

$$(C3) \quad g = \tilde{g} \pmod{I} \text{ and } h = \tilde{h} \pmod{I}$$

Furthermore, the solution satisfying the above three conditions is "unique" in the following sense:

Uniqueness: *We say that an ideal J is special if for all all integer k , and a, b such that $ab \in J^k$, there is an integer l such that $a \in J^l$ and $b \in J^{k-l}$.*

Assume that I is special. Then, for every two solutions g_1, h_1 and g_2, h_2 satisfying conditions C1 to C3, there exists $u \in I^t$, such that:

$$g_2 = g_1(1 + u) \pmod{I^s}, \quad h_2 = h_1(1 - u) \pmod{I^s}$$

As we will see later, towards factoring a bivariate polynomial $f \in \mathbb{F}_q[x, y]$, we will apply the Hensel's lifting to $R = \mathbb{F}_q[x, y]$, $I = (y)$ and a factorization of $f \pmod{I}$. When $I = (y)$, $f \pmod{y}$ is simply a univariate polynomial on x over \mathbb{F}_q , which we know how to factor from previous lectures. Next we proceed to prove Lemma 1.

Proof We prove this lemma by induction.

Base Case $s = 2$: We first show that there exists \tilde{g}, \tilde{h} and \tilde{a}, \tilde{b} satisfying C1 to C3, and then establish the uniqueness of the solution. By condition H1 and H2, we have

$$f = gh + q, \text{ for some } q \in I$$

$$ag + bh = 1 + r \text{ for some } r \in I$$

Write $\tilde{g} = g + g_1$, $\tilde{h} = h + h_1$, for some $g_1, h_1 \in I$ to be set. Then

$$\tilde{g}\tilde{h} = gh + g_1h + h_1g + h_1g_1$$

Since $h_1g_1 \in I^2$, in order to satisfy condition C1, we want $g_1h + h_1g + h_1 = q \pmod{I^2}$. To satisfy this, set $g_1 = bq$, $h_1 = aq$, and get that $g_1h + h_1g = q(bh + ag) = q(1+r)$, which equals to q modulo I^2 as required. By construction $\tilde{g} = g \pmod{I}$ and $\tilde{h} = h \pmod{I}$, satisfying condition C3. To show that \tilde{g} and \tilde{h} are also relatively prime, observe that $a\tilde{g} + b\tilde{h} = ag + bh + r' = 1 + r + r'$, for some $r' \in I$. Let $r'' = r + r' \in I$. Now we can take $\tilde{a} = a(1 - r'')$ and $\tilde{b} = b(1 - r'')$, and get that:

$$\tilde{a}\tilde{g} + \tilde{b}\tilde{h} = (1 - r'')(a\tilde{g} + b\tilde{h}) = (1 - r'')(1 + r'') = 1 - r''^2 = 1 \pmod{I^2}$$

Now it remains to show that \tilde{g}, \tilde{h} is the unique solution satisfying C1 to C3. That is, if g^*, h^* is a different solution satisfying C1 to C3, then there is $u \in I$ such that $g^* = \tilde{g}(1 + u)$ and $h^* = \tilde{h}(1 - u)$. By condition C3, we have $g^* = \tilde{g} + g_2$ and $h^* = \tilde{h} + h_2$ for some $g_2, h_2 \in I$ (because, modulo I , we know that $g^* = g = \tilde{g}$ and $h^* = h = \tilde{h}$). Therefore, we have:

$$g^*h^* = \tilde{g}\tilde{h} + g_2\tilde{h} + h_2\tilde{g} + g_2h_2$$

By condition C1, we know that $g^*h^* = f = \tilde{g}\tilde{h} \pmod{I^2}$. Thus, the above equation modulo I^2 gives,

$$g_2\tilde{h} + h_2\tilde{g} = 0 \pmod{I^2}$$

By condition C2, we have $\tilde{a}\tilde{g} + \tilde{b}\tilde{h} = 1 \pmod{I^2}$. Therefore,

$$\begin{aligned} \tilde{b}(g_2\tilde{h} + h_2\tilde{g}) &= 0 \pmod{I^2} \\ g_2\tilde{b}\tilde{h} + \tilde{b}h_2\tilde{g} &= 0 \pmod{I^2} \\ g_2(1 - \tilde{a}\tilde{g}) + \tilde{b}h_2\tilde{g} &= 0 \pmod{I^2} \\ g_2 &= (\tilde{a}g_2 - \tilde{b}h_2)\tilde{g} \pmod{I^2} \end{aligned}$$

Let $u = \tilde{a}g_2 - \tilde{b}h_2$. Since g_2 and h_2 are all elements in I , so is u . Furthermore we have $g^* = \tilde{g} + g_2 = \tilde{g}(1 + u)$. Similarly, by symmetry, we obtain that

$$h_2 = (\tilde{b}h_2 - \tilde{a}g_2)\tilde{h} \pmod{I^2}$$

Therefore, $h^* = \tilde{h}(1 - u)$. This concludes the proof for the base case.

Induction Step: Assume that for the case of $s = t$, there exist $g_0, h_0 \in R[x]$ and $a_0, b_0 \in I^t$ satisfying conditions C1 to C3, and the solution to the three conditions is “unique”. We show that for the case of $s = 2t$, we can construct $g_1, h_1 \in R[x]$ and $a_1, b_1 \in I^{2t}$ satisfying conditions C1 to C3, and the solution is also unique.

The existence of g_1, h_1, a_1, b_1 satisfying conditions C1 to C3 follows exactly the same proof as in the base case. Therefore, we focus on the proof of uniqueness. Let g_1, h_2 be a different solution from g_1, h_1 . Then both $g_1 \pmod{I^t}$, $h_1 \pmod{I^t}$ and $g_2 \pmod{I^t}$, $h_2 \pmod{I^t}$ are solutions satisfying C1 to C3 for the case of $s = t$. Then by the induction hypothesis, we have that there is a $u_0 \in I^{t/2}$ such that,

$$\begin{aligned} (g_2 \pmod{I^t}) &= (g_1 \pmod{I^t})(1 + u_0) \pmod{I^t} \\ (h_2 \pmod{I^t}) &= (h_1 \pmod{I^t})(1 - u_0) \pmod{I^t} \end{aligned}$$

This implies that

$$\begin{aligned} g_2 &= g_1(1 + u_0) \pmod{I^t} \\ h_2 &= h_1(1 - u_0) \pmod{I^t} \end{aligned}$$

Notice that this is different from the condition in the base case where any two solutions must equal modulo I . Nevertheless, following the same argument, we can derive that there is an element $u \in I^t$ such that,

$$g_2 = g_1(1 + u_0)(1 + u) \pmod{I^{2t}}$$

$$h_2 = h_1(1 - u_0)(1 - u) \pmod{I^{2t}}$$

Below we show that u_0 is in fact an element in I^t , then

$$g_2 = g_1(1 + u_0 + u + u_0u) = g_1(1 + u_0 + u) \pmod{I^{2t}}$$

$$h_2 = h_1(1 - u_0 - u + u_0u) = h_1(1 - u_0 - u) \pmod{I^{2t}}$$

Thus $g_2 = g_1(1 + u')$ and $h_2 = h_1(1 - u')$ for $u' = u_0 + u \in I^t$ as desired.

To show that $u_0 \in I^t$, consider:

$$\begin{aligned} g_2h_2 &= g_1h_1(1 - u_0^2)(1 - u^2) \pmod{I^{2t}} \\ g_2h_2 &= g_1h_1(1 - u_0^2) \pmod{I^{2t}} \quad [\text{as } u \in I^t] \end{aligned}$$

Since g_1 and h_1 are not elements in I , for the last equation to hold, it must be the case that $u_0 \in I^t$. Therefore, we conclude the lemma. ■

2 Outline of Factoring, revisited

We now give a more complete outline for factoring bivariate polynomials.

Given a monic $f(x, y) \in \mathbb{F}_q[x, y]$, with total degree d , the factoring algorithm *SPLIT* proceeds as follows:

1. If $g = \gcd(f, \frac{\partial f}{\partial x}) \neq 1$, then output $(g, f/g)$ and stop. Otherwise, continue the following steps.
2. Find $y_0 \in \mathbb{F}$ such that $f(x, y_0)$ has no repeated factors. This can be done by computing $\text{Res}\left(f, \frac{\partial f}{\partial x}\right)$, and plugging in $y_0 = 1, 2, \dots$ until we find one that makes the resultant non-zero.

We claim that this will terminate in at most d^2 iterations, as $\text{Res}\left(f, \frac{\partial f}{\partial x}\right)$ is a polynomial in y with degree at most d^2 . Furthermore, the first step ensures that f does not have repeated roots; therefore, $\text{Res}\left(f, \frac{\partial f}{\partial x}\right)$ is not a zero polynomial. Hence it has at most d^2 roots.

3. Put $f_{y_0}(x) = f(x, y) \pmod{(y - y_0)} = f(x, y_0)$ and factor it. This can be done by using the factoring algorithm for univariate polynomial over \mathbb{F} . Let g be an irreducible factor of $f_{y_0}(x)$, and h such that $f = gh \pmod{(y - y_0)}$.
4. Now we apply Hensel's Lifting to obtain $f = g_1h_1 \pmod{(y - y_0)^t}$ for a $t \approx d^2$.
5. Next, from g_1 we ask if we can find a nontrivial factor \tilde{g} of f . This is done through the "Jump" step, which tries to find polynomials \tilde{g} and \tilde{h} such that $\tilde{g} = g_1\tilde{h} \pmod{(y - y_0)^t}$, and \tilde{g} has small degrees in y (smaller than d) and minimal degree in x .
6. Finally, return \tilde{g} and f/\tilde{g} .