Universal Semantic Communication

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An fantasy setting (SETI)



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No common language! Is meaningful communication possible?



Bob

What should Bob's response be?

If there are further messages, are they reacting to him? Is there an intelligent Alien (Alice) out there?

Classical Paradigm for interaction



New paradigm



Goal of this talk

Definitional issues and a definition: What is successful communication? What is intelligence? cooperation? Theorem: "If Alice and Bob are intelligent and cooperative, then communication is feasible" (in one setting) Proof ideas: Suggest: Protocols, Phenomena ... Methods for proving/verifying intelligence

A first attempt at a definition

Alice and Bob are "universal computers" (aka programming languages)
Have no idea what the other's language is!
Can they learn each other's language?

Good News: Language learning is finite. Can enumerate to find translator.

Bad News: No third party to give finite string!
 Enumerate? Can't tell right/wrong ⁽³⁾

Communication & Goals

Indistinguishability of Right/Wrong: Consequence of "communication without goal".

 Communication (with/without common language) ought to have a "Goal".

Before we ask how to improve communication, we should ask why we communicate?

> *"Communication is not an end in itself, but a means to achieving a Goal*"



Contrast with Interactive Proofs

 Similarity: Interaction between Alice and Bob.
 Difference: In IP, Bob does not trust Alice. (In our case Bob does not understand Alice).

Famed Theorem: IP = PSPACE [LFKN, Shamir].
 Membership in PSPACE solvable S can be proved interactively to a probabilistic Bob.
 Needs a PSPACE-complete prover Alice.

Intelligence & Cooperation?

For Bob to have a non-trivial interaction, Alice must be:

Intelligent: Capable of deciding if x in S.

Cooperative: Must communicate this to Bob.

• Formally: Alice is S-helpful if \exists probabilistic poly time (ppt) Bob B' s.t. For all states σ of Alice $A(\sigma) \leftrightarrow B'(x)$ accept w.h.p. iff $x \in S$.

Successful universal communication

Formally,

Bob should be able to talk to any S-helpful Alice and decide S.

Ppt B is S-universal if for every $x \in \{0,1\}^*$ $-x \in S$ and A is S-helpful $\Rightarrow [A \leftrightarrow B(x)] = 1$ (whp). - (For S-helpful A) $[A \leftrightarrow B(x)] = 1$ w.h.p. $\Rightarrow x \in S$.

Main Theorem

 If S is PSPACE-complete (aka Chess), then there exists an S-universal Bob. (Generalizes to any checkable set S.)
 If there exists an S-universal Bob then S is in PSPACE.

In English:

If S is moderately stronger than what Bob can do on his own, then attempting to solve S leads to non-trivial (useful) conversation.
If S too strong, then leads to ambiguity.
Uses IP=PSPACE

Few words about the proof

• Positive result: Enumeration + Interactive Proofs Guess: Interpreter; $x \in S$?



Proof works $\Rightarrow x \in S$; Doesnt work \Rightarrow Guess wrong. Alice S-helpful \Rightarrow Interpreter exists!

Few words about the proof

Positive result: Enumeration + Interactive Proofs

Negative result:

- Suppose Alice answers every question so as to minimize the conversation length.
 - (Reasonable effect of misunderstanding).
- Conversation comes to end quickly.
- Bob has to decide.
- Decision can be computed in PSPACE (since Alice's strategy can be computed in PSPACE).
- Bob must be wrong if L is not in PSPACE.
- Warning: Only leads to finitely many mistakes.

Is this language learning?

End result promises no language learning: Merely that Bob solves his problem.

In the process, however, Bob learns Interpreter!

But this may not be the right Interpreter.

All this is Good!
 No need to distinguish indistinguishables!