# **Universal Semantic Communication**

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## The Meaning of Bits



- What if Alice is trying to send instructions?
  - Aka, an algorithm
  - Does Bob understand the correct algorithm?
  - What if Alice and Bob speak in different (programming) languages?
- Sales pitch: Crucial to the future of computing. More and more heterogenous computers, data, and software interact without human in the loop.

## Some modelling

- Say, Alice and Bob know different programming languages. Alice wishes to communicate an algorithm to Bob.
- Bad News: Can't be done
  - For every Bob, there exist algorithms A and A', and Alices, Alice and Alice', such that the two are indistinguishable to Bob.
- Good News: Need not be done.
  - From Bob's perspective, if A and A' are indistinguishable, then they are equally useful to him.

## What should be communicated? Why?

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## Modelling Bob and his perspective

Bob: 
$$\Omega \times \Sigma^k \to \Omega \times \Gamma^\ell$$
,  
where  $\Omega = \text{countable state space}$   
 $\Sigma^k = \text{input signals}$   
 $\Gamma^\ell = \text{output signals.}$ 



- Bob speaks to some environment (a multitude of entities).
- Why? Has some goal!
  - "Control": Wants to alter the state of the environment.
  - "Intellectual": Wants to glean knowledge (about universe/environment).
- Claim: By studying the goals, can enable Bob to overcome linguistic differences (and achieve goal).

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**Rest of the talk** 



- Part I: Bob wishes to solve hard problem, is computationally limited, and Alice can solve the problem.
- Part II: Bob is a teacher and wants to test student's ability.
- Part III: Generic goals.

# Part I: A Computational Goal

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## **Computational Goal for Bob**

- Bob is prob. poly time bounded. Wants to decide membership in set S.
- Alice is computationally unbounded, does not speak same language as Bob, but is "helpful".
- What kind of sets S?
  - E.g., undecidable?, decidable? PSPACE, NP, BPP?

## Setup



## **Intelligence & Cooperation?**

- For Bob to have a non-trivial interaction, Alice must be:
  - Intelligent: Capable of deciding if x in S.
  - Cooperative: Must communicate this to Bob.
- Formally:
  - Alice is **S-helpful**

if  $\exists$  probabilistic poly time (ppt) Bob B' s.t.  $\forall$  initial state of mind  $\sigma$ ,  $A(\sigma) \leftrightarrow B'(x)$  accept w.h.p. iff  $x \in S$ .

## Successful universal communication

- Bob should be able to talk to any S-helpful Alice and decide S.
- Formally,
  - Ppt B is S-universal if for every  $x \in \{0, 1\}^*$ - A is S-helpful  $\Rightarrow [A \leftrightarrow B(x)] = 1$  iff  $x \in S$  (whp). A is not S-helpful  $\Rightarrow$  Nothing!!

Or should it be ...

A is not S-helpful  $\Rightarrow [A \leftrightarrow B(x)] = 1$  implies  $x \in S$ .

## Main Theorem

- If S is PSPACE-complete, then there exists a Suniversal Bob (generalizes to other checkable sets S).
- Conversely, if there exists a S-universal Bob, then S is in PSPACE.

#### In other words:

- If S is moderately stronger than what Bob can do on his own, then attempting to solve S leads to non-trivial (useful) conversation.
- If S too strong, then leads to ambiguity.
- Uses IP=PSPACE

### Few words about the proof

Positive result: Enumeration + Interactive Proofs Guess: Interpreter;  $x \in S$ ?



Proof works  $\Rightarrow x \in S$ ; Doesnt work  $\Rightarrow$  Guess wrong. Alice S-helpful  $\Rightarrow$  Interpreter exists!

## **Proof of Negative Result**

L not in PSPACE implies Bob makes mistakes.

- Suppose Alice answers every question so as to minimize the conversation length.
  - (Reasonable effect of misunderstanding).
- Conversation comes to end quickly.
- Bob has to decide.
- Conversation + Decision simulatable in PSPACE (since Alice's strategy can be computed in PSPACE).
- Bob must be wrong if S is not in PSPACE.
- Warning: Only leads to finitely many mistakes.

# Part II: Intellectual Curiosity

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## Setting: Bob more powerful than Alice

#### What should Bob's Goal be?

- Can't use Alice to solve problems that are hard for him.
- Can pose problems and see if she can solve them. E.g., Teacher-student interactions.
- But how does he verify "non-triviality"?
- What is "non-trivial"? Must distinguish ...



## Setting: Bob more powerful than Alice

- Concretely:
  - Bob capable of TIME(n<sup>10</sup>).
  - Alice capable of TIME(n<sup>3</sup>) or nothing.
  - Can Bob distinguish the two settings?
- Definition:
- Alice is  $n^{3-\epsilon}$ -helpful if  $\neg$  Dob D'  $\neg$  TIME $(m^{3-\epsilon})$ 
  - if  $\exists$  Bob  $B' \in \text{TIME}(n^{3-\epsilon})$  s.t.  $\forall S \in \text{TIME}(n^3)$ , and  $\forall$  initial state of mind  $\sigma$ ,  $A(\sigma) \leftrightarrow B'(x_1, \dots, x_n)$  computes  $S(x_1), \dots, S(x_n)$ .
- Theorem: There exists a universal Bob that distinguishes helpful Alices from trivial ones.
- Moral: Language (translation) should be simpler than problems being discussed.

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## Part III: Generic Goals

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Semantic Communication

17

## Generically

Bob interacts with an environment (collection of entities).



- Function of transcript of interaction
- States of environment!
- But not Bob itself!
- Should forgive finite prefixes.

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## **Generic Helpfulness, Universality**

- Consider a class of Alices  $\mathcal{A}_{\text{r}}$  a class of Bobs  $\mathcal{B}$  , and some goal G
  - (G,  $\mathcal{B}$ )-Helpful: Helpful for some Bob in  $\mathcal{B}$ .
  - (G,  $\mathcal{A}$ )-Universal: Works with all Alices in  $\mathcal{A}$ .
- A -Verifiable protocol: For every A in A, Protocol accepts iff goal is achieved.



Theorem: Verifiable Goals can be achieved universally.

## Conclusions

- Communication of "meaning/context" is feasible; provided goals are explicit.
- Verifying "goal achievement" for non-trivial goals is the (only?) way to learn languages.
- Currently the learning is slow ... is this inherent?
  Better class of Alices?
- What are interesting goals, and how can they be verified?

# **Thank You!**

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