

# Semantic Goal-Oriented Communication

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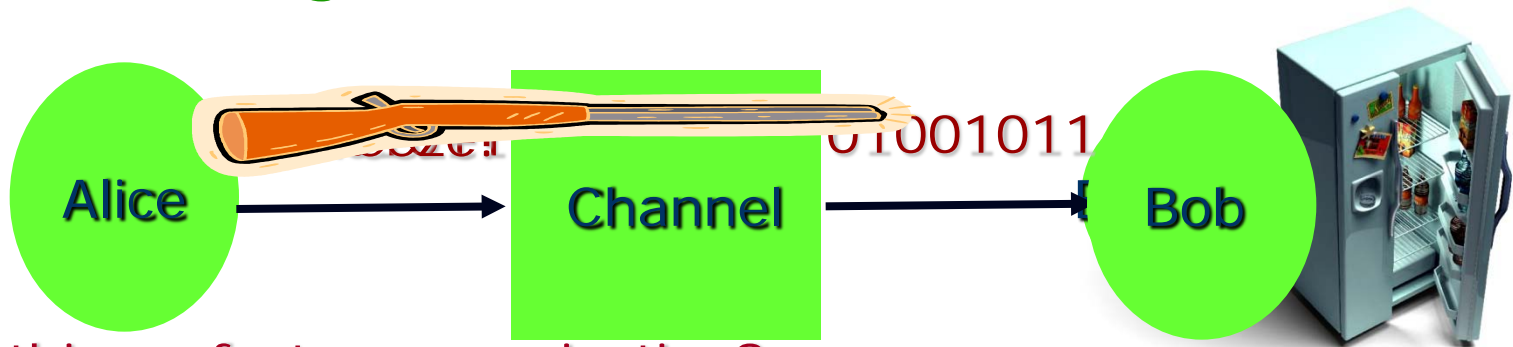
Microsoft Research + MIT

Joint with **Oded Goldreich** (Weizmann) and **Brendan Juba** (MIT).

# Disclaimer

- Work in progress (for ever) ...
- Comments/Criticisms/Collaboration/Competition welcome.

# The Meaning of Bits



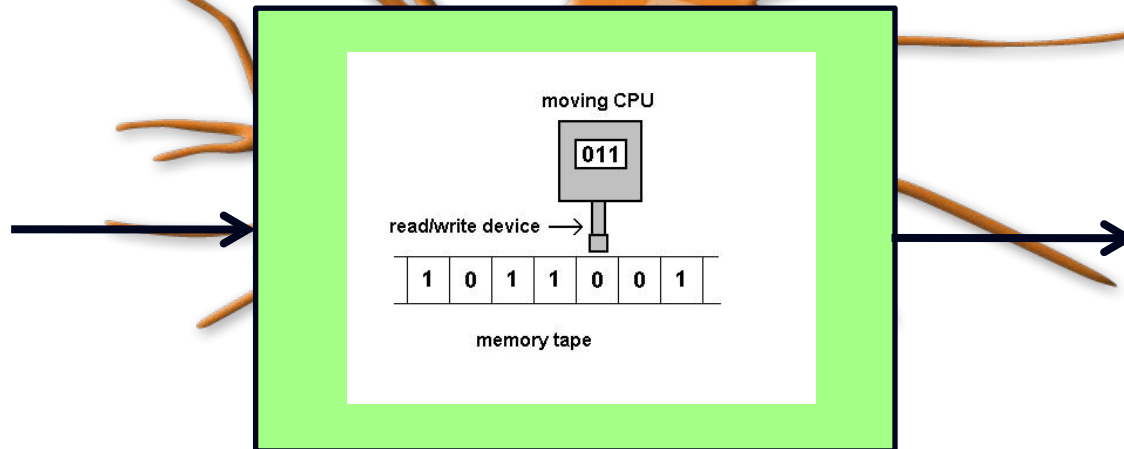
- Is this perfect communication?
- What if Alice is trying to send instructions?
  - Aka, an algorithm
  - Does Bob understand the correct algorithm?
  - What if Alice and Bob speak in different (programming) languages?

# Miscommunication (in practice)

- Exchanging (powerpoint) slides.
  - Don't render identically on different laptops.
- Printing on new printer.
  - User needs to "learn" the new printer, even though printer is quite "intelligent".
- Many such examples ...
  - In all cases, sending bits is insufficient.
  - Notion of meaning ... intuitively clear.
  - But can it be formalized?
    - Specifically? Generically?
    - While conforming to our intuition

# Modelling Computing

- Classically: Turing Machine/(von Neumann) RAM.
  - Described most computers being built?



- Modern computers: more into communication than computing.
  - What is the mathematical model?
  - Do we still have universality?

# Bits vs. their meaning

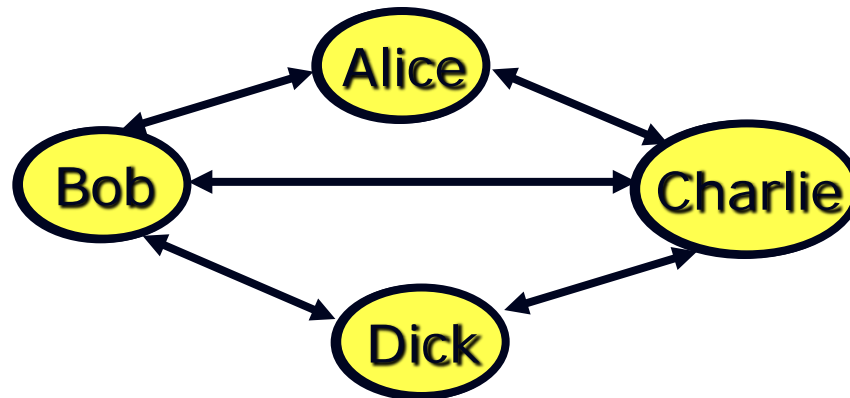
- Say, User and Server know different programming languages. Server wishes to send an algorithm  $A$  to User.
  - $A$  = sequence of bits ... (relative to prog. language)
- Bad News: Can't be done
  - For every User, there exist algorithms  $A$  and  $A'$ , and Servers  $S$  and  $S'$  such that  $S$  sending  $A$  is indistinguishable (to User) from  $S'$  sending  $A'$
- Good News: Need not be done.
  - From Bob's perspective, if  $A$  and  $A'$  are indistinguishable, then they are equally useful to him.
- What should be communicated? Why?

## Aside: Why communicate?

- Classical "Theory of Computing"



- Issues: Time/Space on DFA? Turing machines?
- Modern theory:



- Issues: Reliability, Security, Privacy, Agreement?
- If communication is so problematic, then why not "Just say NO!"?

# Motivations for Communication

- Communicating is painful. There must be some compensating gain.
- What is User's Goal?
  - "Control": Wants to alter the state of the environment.
  - "Intellectual": Wants to glean knowledge (about universe/environment).
- Thesis: By studying the goals, can enable User to overcome linguistic differences (and achieve goal).



# Part II: Computational Motivation

# Computational Goal for Bob

- Why does User want to learn algorithm?
  - Presumably to compute some function  $f$   
( $A$  is expected to compute this function.)
  - Lets focus on the function  $f$ .
- Setting:
  - User is prob. poly time bounded.
  - Server is computationally unbounded, does not speak same language as User, but is "helpful".
  - What kind of functions  $f$ ?
    - E.g., uncomputable, PSPACE, NP, P?

# Setup

User

$f(x) = 0/1?$

$R \leftarrow \text{\$}\$ \$$

Server

$q_1$

Different from interactions in  
cryptography/security:

There, User does not **trust** Server,  
while here he does not  
**understand** her.

Computes  $P(x, R, a_1, \dots, a_k)$

Hopefully  $P(x, \dots) = f(x)$ !

# Intelligence & Cooperation?

- For User to have a non-trivial interaction, Server must be:
  - Intelligent: Capable of computing  $f(x)$ .
  - Cooperative: Must communicate this to User.
- Formally:
  - Server  $S$  is helpful (for  $f$ ) if
    - $\exists$  some (other) user  $U'$  s.t.
      - $\forall x$ , starting states  $\sigma$  of the server  
( $U'(x) \leftrightarrow S(\sigma)$ ) outputs  $f(x)$

# Successful universal communication

- Universality: **Universal User U** should be able to talk to any (every) helpful server **S** to compute **f**.
- Formally:
  - **U** is **universal**, if
$$\forall \text{ helpful } S, \forall \sigma, \forall x$$
$$(U(x) \leftrightarrow S(\sigma)) = f(x) \text{ (w.h.p.)}$$
- What happens if **S** is not helpful?
  - Paranoid view  $\Rightarrow$  output "**f(x)**" or "?"
  - Benign view  $\Rightarrow$  Don't care (everyone is helpful)

# Main Theorems [Juba & S. '08]

- If  $f$  is PSPACE-complete, then there exists a  $f$ -universal user who runs in probabilistic polynomial time.
  - Extends to checkable problems
    - $(NP \cap \text{co-NP})$ , breaking cryptosystems)
    - $S$  not helpful  $\Rightarrow$  output is safe
- Conversely, if there exists a  $f$ -universal user, then  $f$  is PSPACE-computable.
  - Scope of computation by communication is limited by misunderstanding (alone).

# Proofs?

- Positive result:
  - $f \in \text{PSPACE} \Rightarrow$  membership is verifiable.
  - User can make hypothesis about what the Server is saying, and use membership proof to be convinced answer is right, or hypothesis is wrong.
- Negative result:
  - In the absence of proofs, sufficiently rich class of users allow arbitrary initial behavior, including erroneous ones.
  - (Only leads to finitely many errors ...)

# Implications

- No universal communication protocol ☹
  - If there were, should have been able to solve every problem (not just (PSPACE) computable ones).
- But there is gain in communication:
  - Can solve more complex problems than on one's own, but not every such problem.
- Resolving misunderstanding? Learning Language?
  - Formally **No!** No such guarantee.
  - Functionally **Yes!** If not, how can user solve such hard problems?



# Principal Criticisms

- Solution is no good.
  - Enumerating interpreters is too slow.
    - Approach distinguishes **right/wrong**; does not solve search problem.
    - Search problem needs new definitions to allow better efficiency.
- Problem is not the right one.
  - Computation is not the goal of communication. Who wants to talk to a PSPACE-complete server?



Next part of talk

# Part III: Generic Goals

# Generic Communication [Goldreich, J., S.]

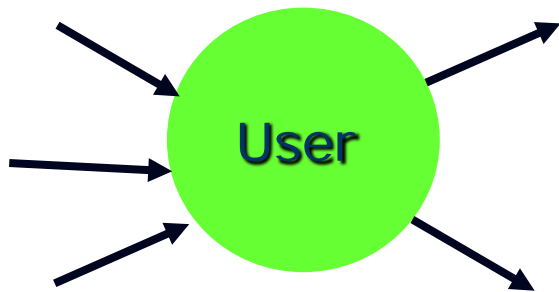
- Not every goal is computational!
  - Even if it is, is Semantic Communication only possible if Server is "much better" than User?
- What are generic goals?
  - Why do we send emails?
  - Why do I browse on Amazon?
  - Why do we listen to boring lectures?
    - (or give inspirational ones 😊)
- Seemingly rich diversity

# Universal Semantics for such Goals?

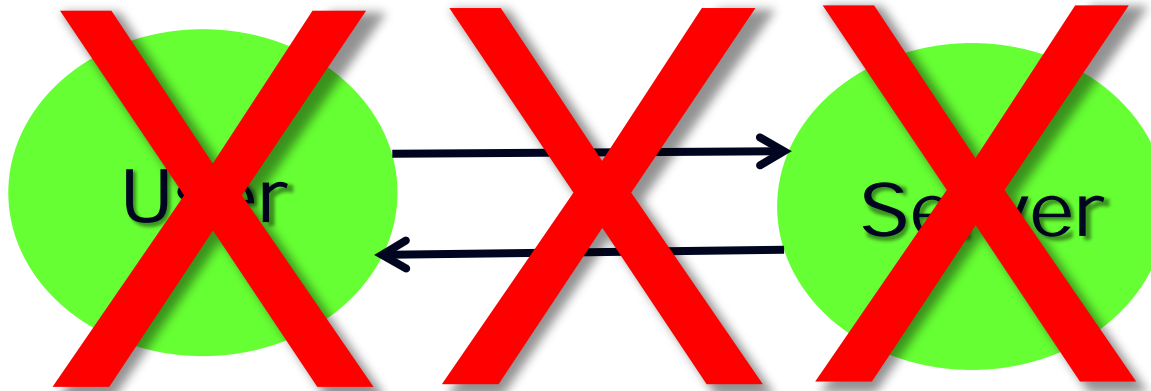
- Can we still achieve goal without knowing common language?
  - Seems feasible ...
    - If user can detect whether goal is being achieved (or progress is being made).
  - Just need to define
    - Sensing Progress?
    - Helpful + Universal?
    - ...
    - Goal?
    - User?

# Modelling User/Interacting agents

- (standard AI model)
- User has state and input/output wires.
  - Defined by the map from current state and input signals to new state and output signals.



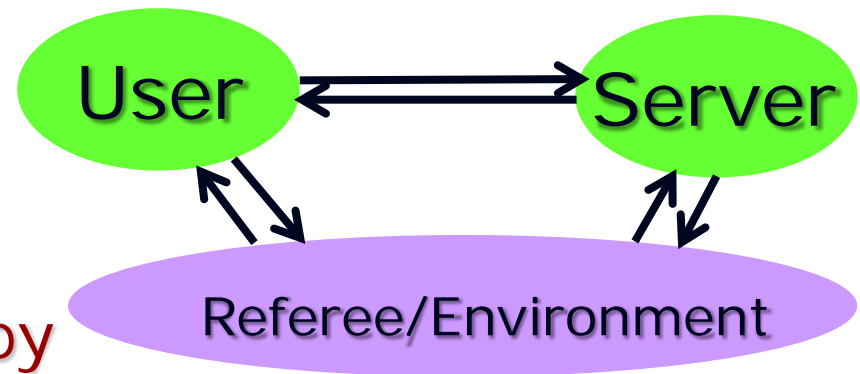
# Generic Goal?



- Goal = function of ?
  - User? – But user wishes to change actions to achieve universality!
  - Server? – But server also may change behaviour to be helpful!
  - Transcript of interaction? – How do we account for the many different languages?

# Generic Goals

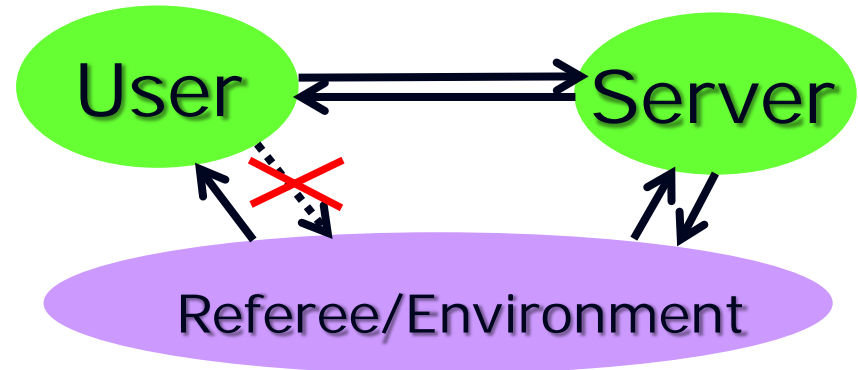
- Key Idea: Introduce 3rd entity: Referee
  - Poses tasks to user.
  - Judges success.



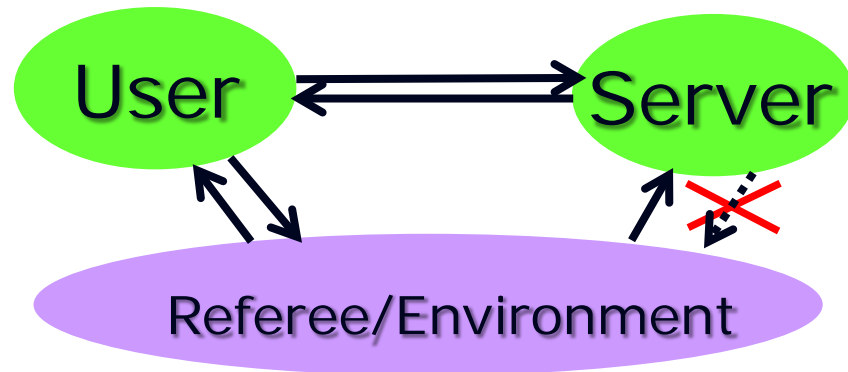
- Generic Goal specified by
  - Referee (just another agent)
  - Boolean Function determining if the state evolution of the referee reflects successful achievement of goal.
  - Class of users/servers.

# Generic Goals

- Pure Control



- Pure Informational





# Sensing & Universality (Theorems)

- To achieve goal, User should be able to sense progress.
  - I.e., user should be able to compute a function that (possibly with some delay, errors) reflects achievement of goals.
  - "User simulates Referee"
- Generalization of positive result:
  - Generic goals (with technical conditions) universally achievable if  $\exists$  sensing function.
- Generalization of negative result:
  - If non-trivial generic goal is achieved with sufficiently rich class of helpful servers, then it is safely achieved with every server.

# Why is the paper so long?

- To capture fully general models ...
  - User/Server live for ever and Goal achievement is a function of infinite sequence of states.
  - User/Server should be allowed to be probabilistic.
  - Referee should be allowed to be non-deterministic.
  - (Getting quantifiers right non-trivial.)

# When Server is less powerful than User

- What should the goal be?
  - Can't expect server to solve (computational) problems user can't.
  - So what can user try to do? Why?
    - Ask Server: Repeat after me ...
    - Test if Server has some computational power ... solves simple (linear/quadratic time) problems.
    - Has memory ... can store/recall.
    - Can act (pseudo-)independently – is not deterministic, produces incompressible strings.
    - Can challenge me with puzzles.
- Each Goal/combo can be cast in our framework.
  - (Sensing functions do exist; communication is essential to achieving Goals; problems are more about control ...)

# (Generalized) Turing tests

- Distinguish between "Intelligent"/"Not"
- Distinguish between "Humans"/"Bots"
  - Generically:
    - Class of servers = H union N:
      - $H = \{ (1,i) \mid i \}$
      - $N = \{ (0,i) \mid i \}$
      - Referee: gets identity of server from server  $(b,s),;$  and after interaction between user and server, gets User's guess  $b'$ . Accepts if  $b = b'$ .
- Sensing function exists? Depends on H vs. N.

# Conclusions - 1

- Goals of communication.
  - Should be studied more.
  - Suggests good heuristics for protocol design:
    - What is your goal?
    - Server = Helpful?
    - User = Sensing?

# References

- Juba & S. (computational)
  - ECCC TR07-084: <http://eccc.uni-trier.de/report/2007/084/>
- Goldreich, Juba & S. (generic)
  - ECCC TR09-075: <http://eccc.uni-trier.de/report/2009/075/>
- Juba & S. – 2. (examples)
  - ECCC TR08-095: <http://eccc.uni-trier.de/report/2008/095/>

**Thank You!**