

Disintegrating Manycores: Which Applications Lose and Why?

Isidor R. Brkić, Mark C. Jeffrey

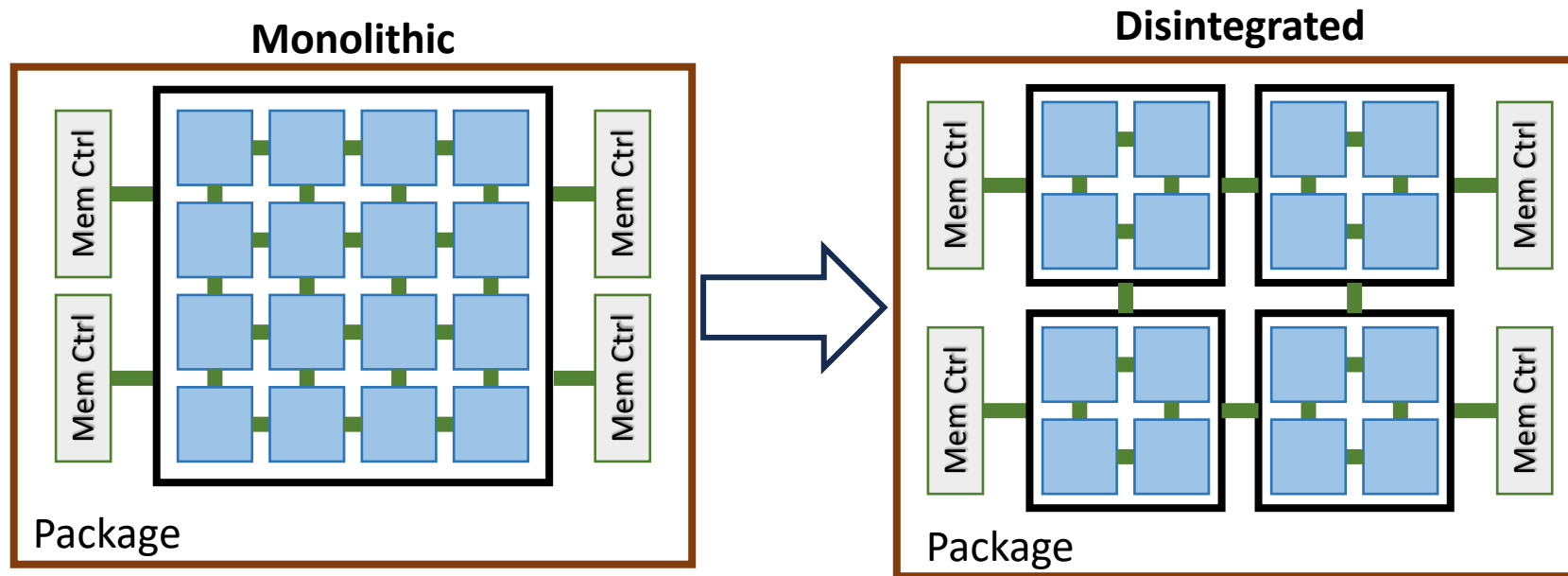
28 October 2023

NoCArc 2023



UNIVERSITY OF
TORONTO

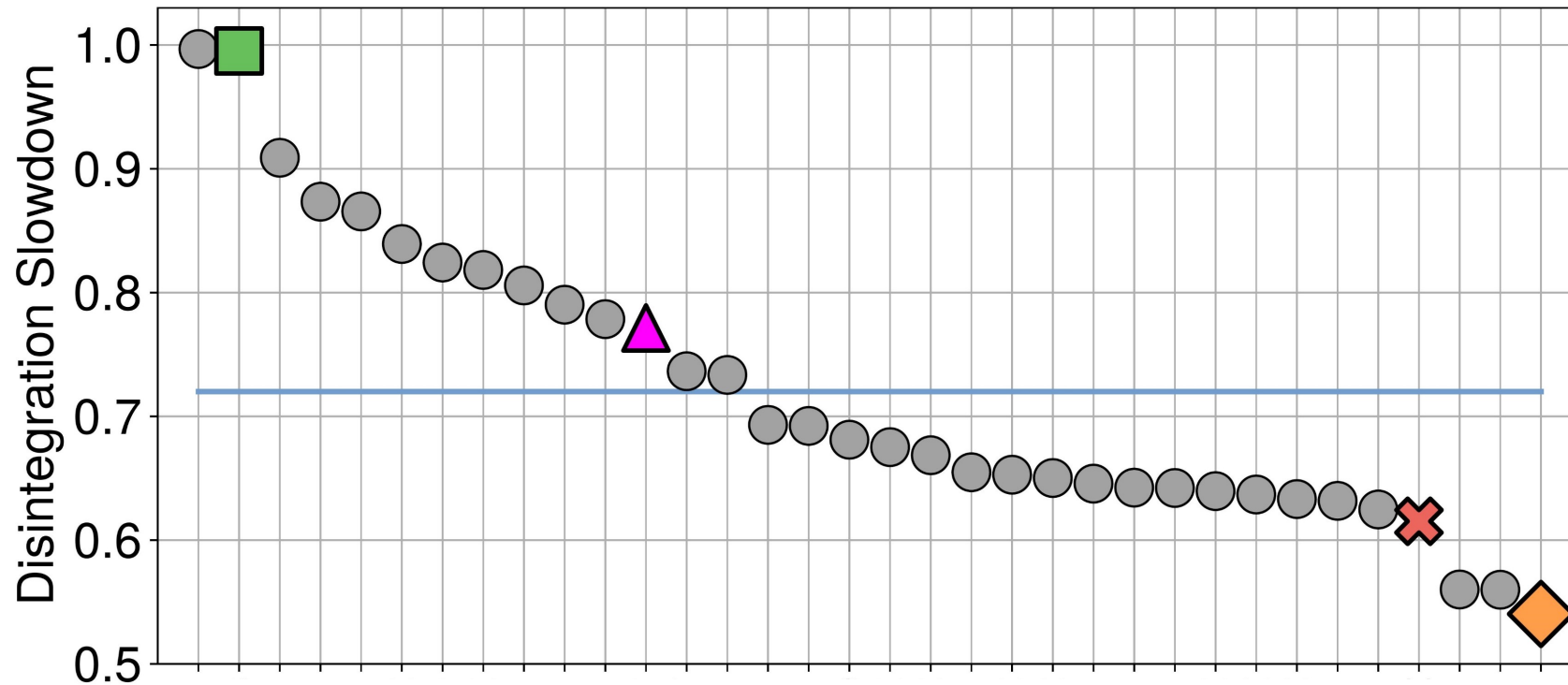
Disintegration Saves Cost (\$)



Company	Architecture	Chiptlets	"Cores"
Intel	CPU	4	60
AMD	CPU	9-13	64
Apple	CPU/GPU/Accelerator	2	20/64/32
NVIDIA	GPU	4	256
NVIDIA	Accelerator	36	576
Xilinx	FPGA	?	N/A

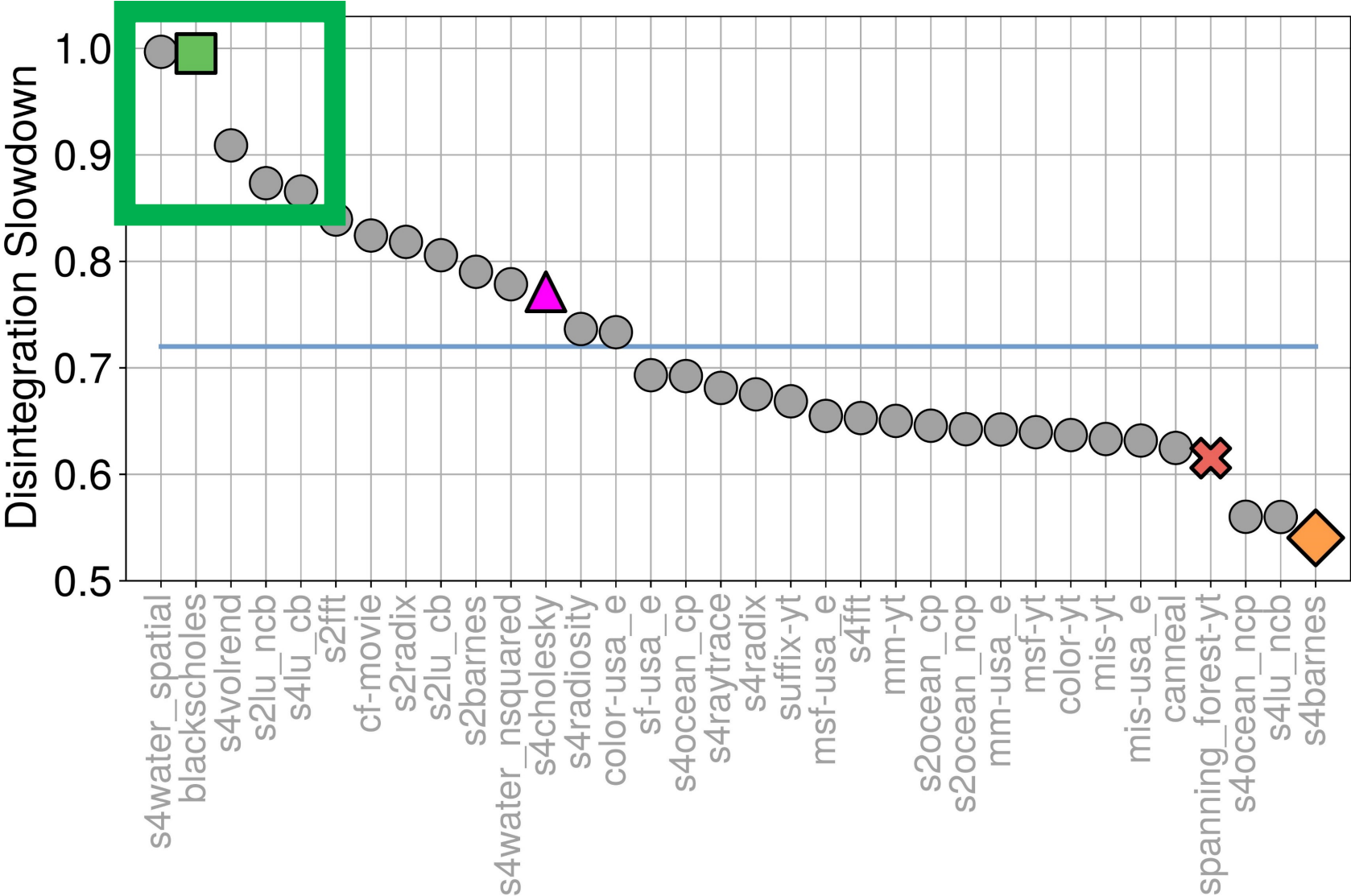
Disintegration Harms Performance

There has been no broad study of how disintegration impacts performance



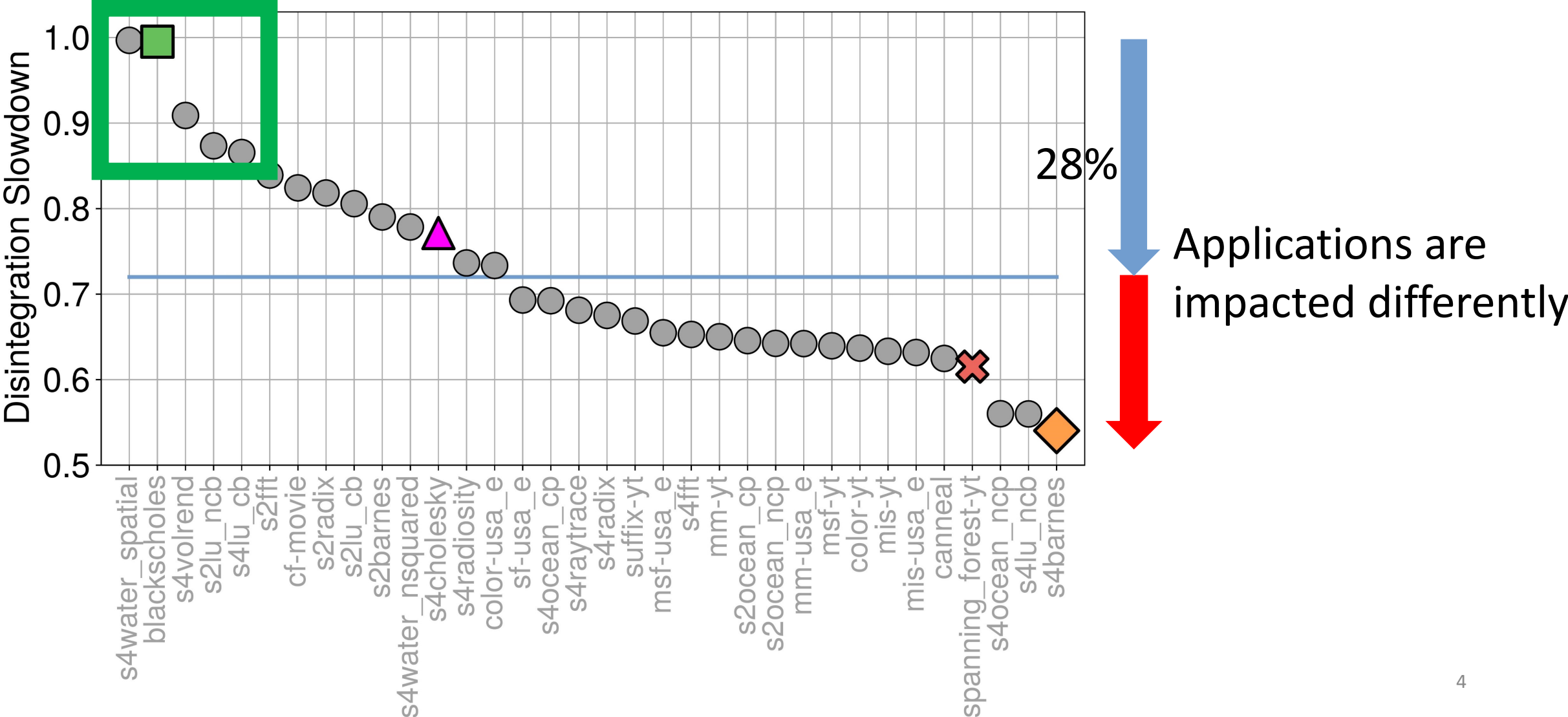
Disintegration Harms Performance

There has been no broad study of how disintegration impacts performance



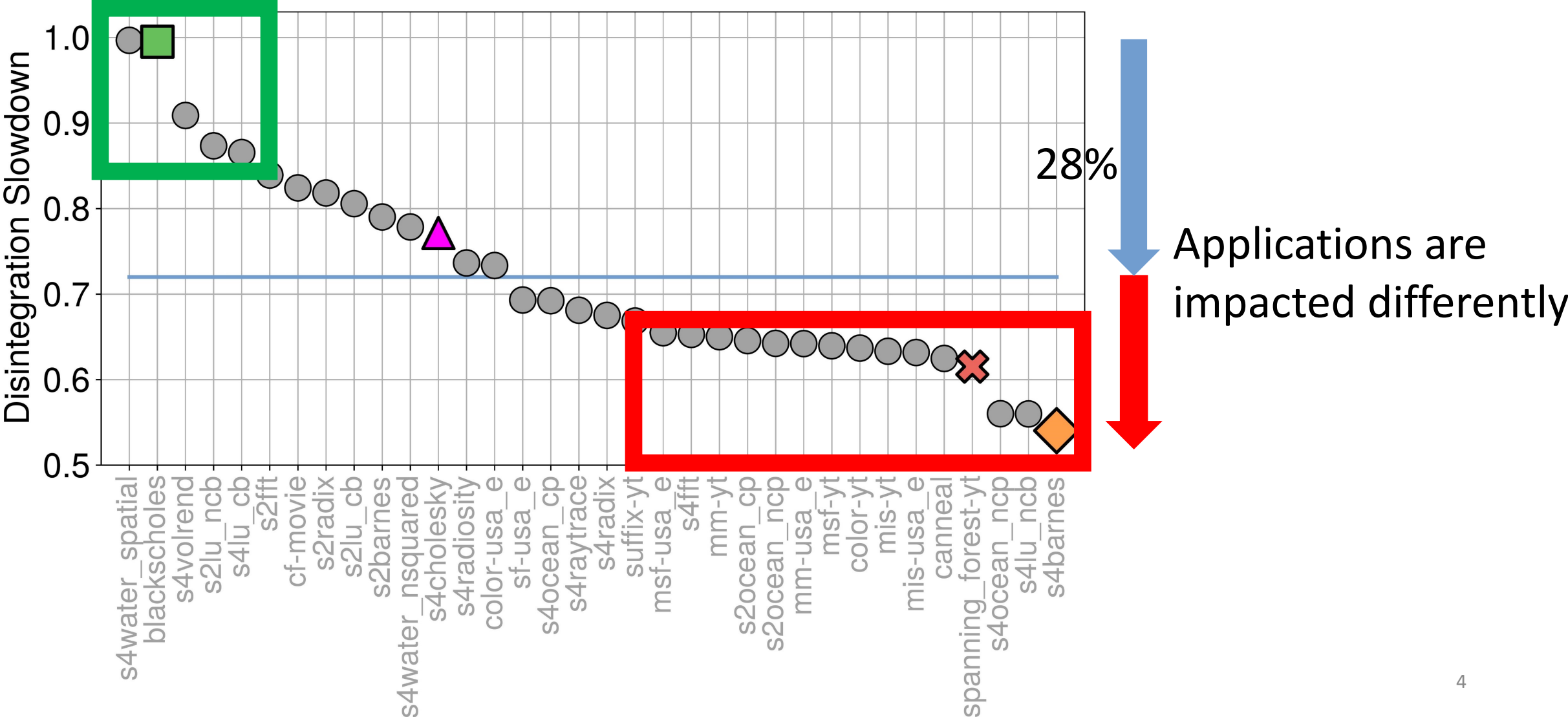
Disintegration Harms Performance

There has been no broad study of how disintegration impacts performance



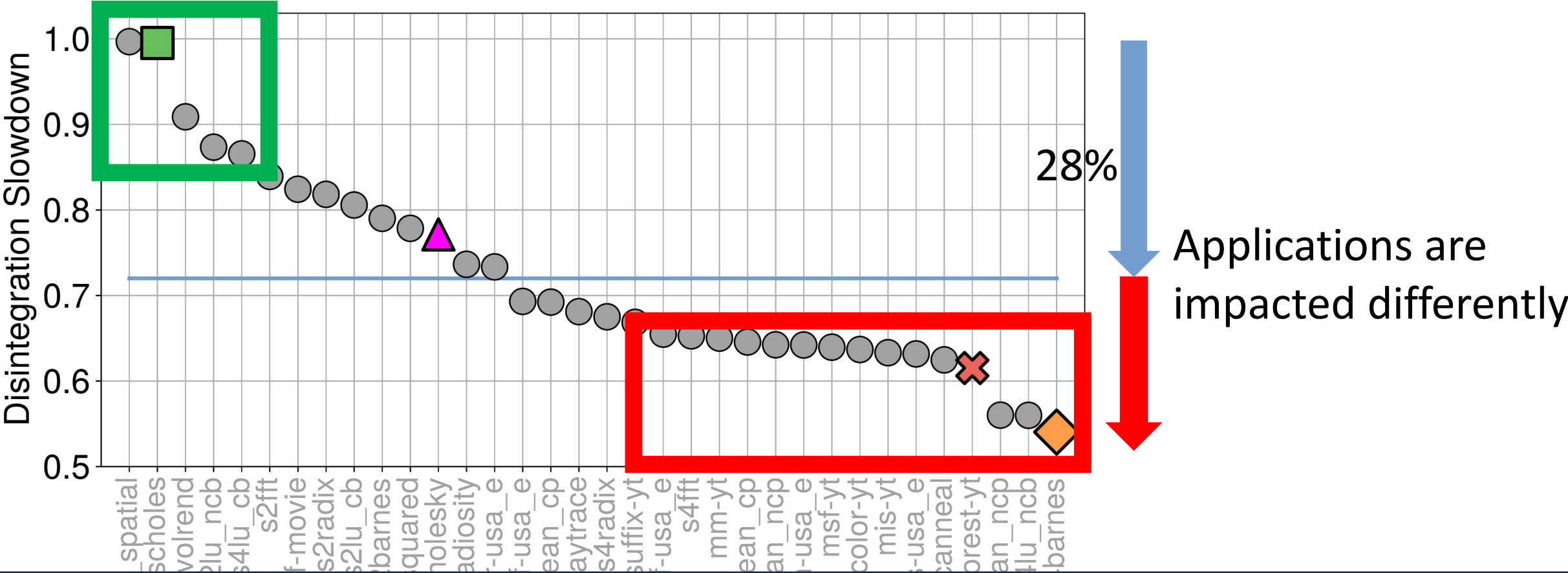
Disintegration Harms Performance

There has been no broad study of how disintegration impacts performance



Disintegration Harms Performance

There has been no broad study of how disintegration impacts performance



What application characteristics matter?

s4

sp

Executive Summary

We demonstrate that:

Unsurprisingly, scalable applications still scale on disintegrated systems

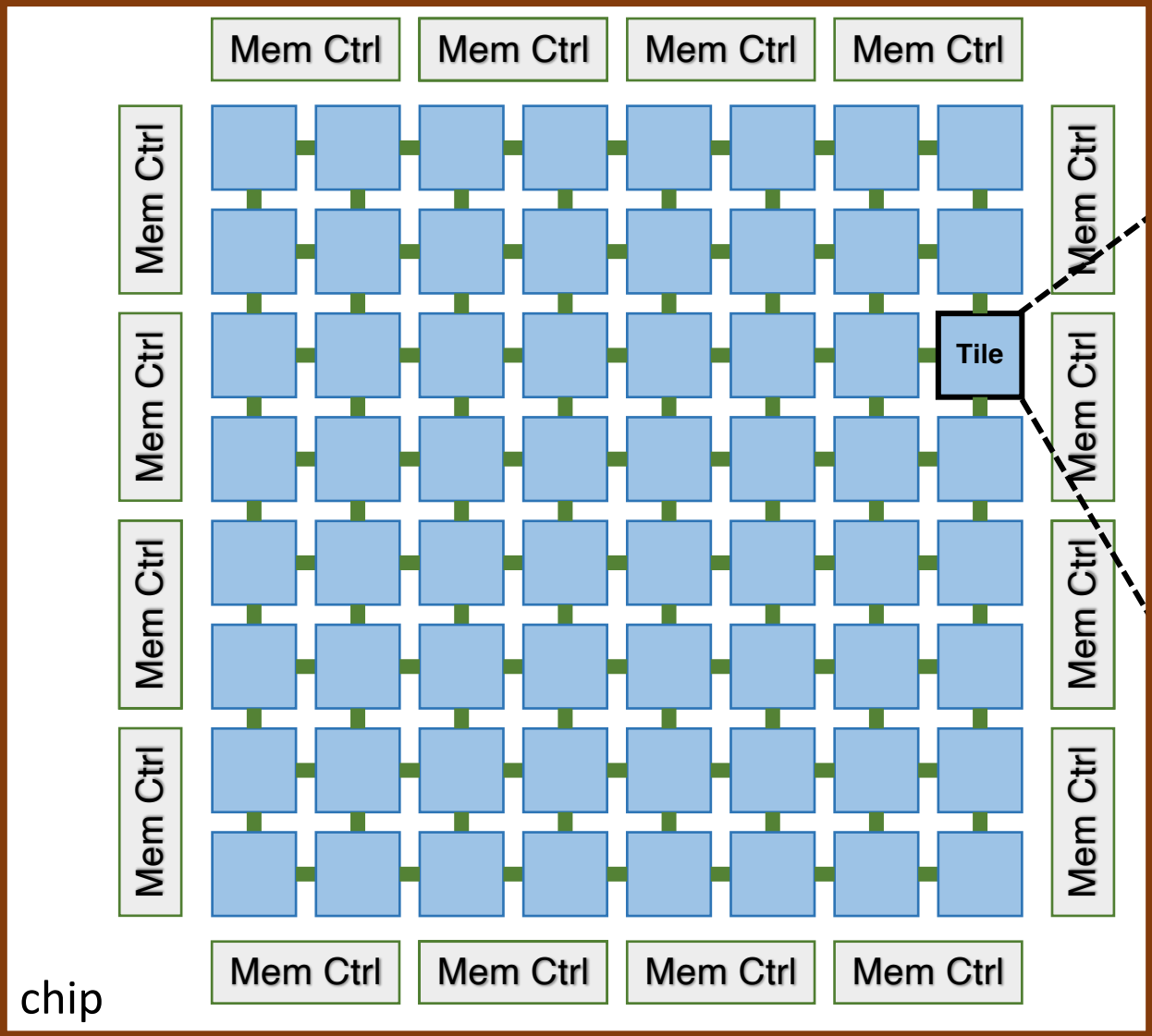
Disintegration penalty varies significantly across applications and is not correlated with monolithic speedup

Data sharing and network injection bandwidth lead to larger disintegration penalty

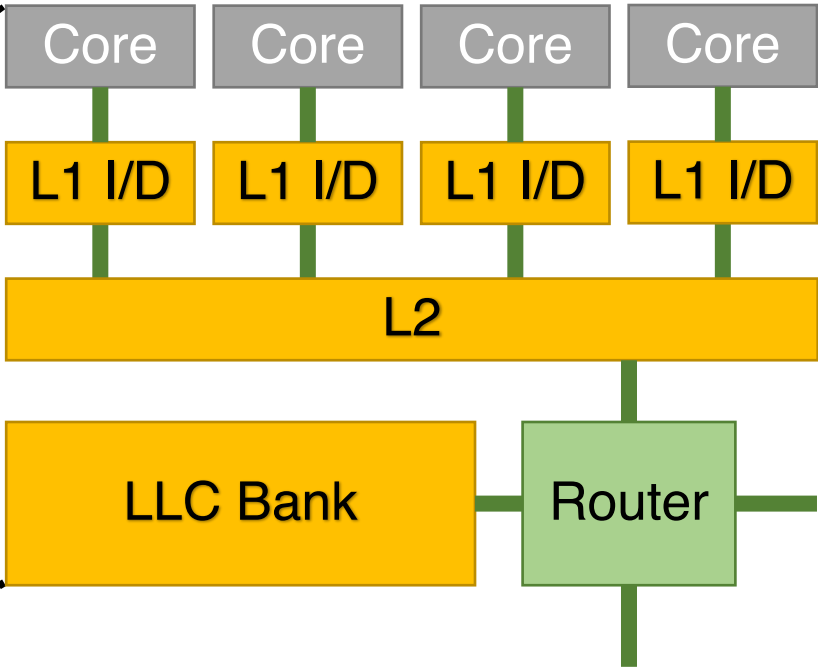
Methodology

Baseline Manycore Monolith

256 core, 64 tile chip



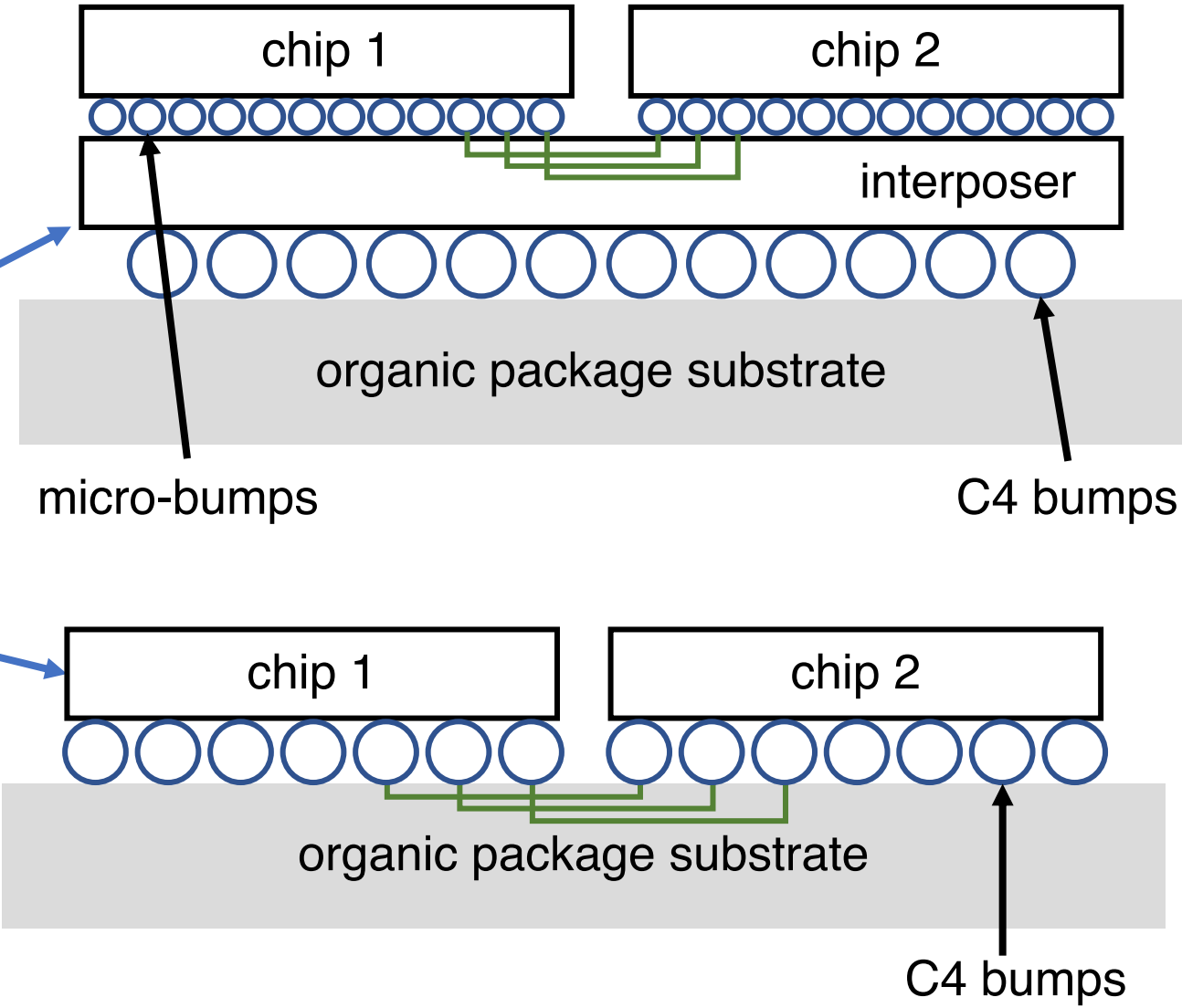
Logical Tile Organization



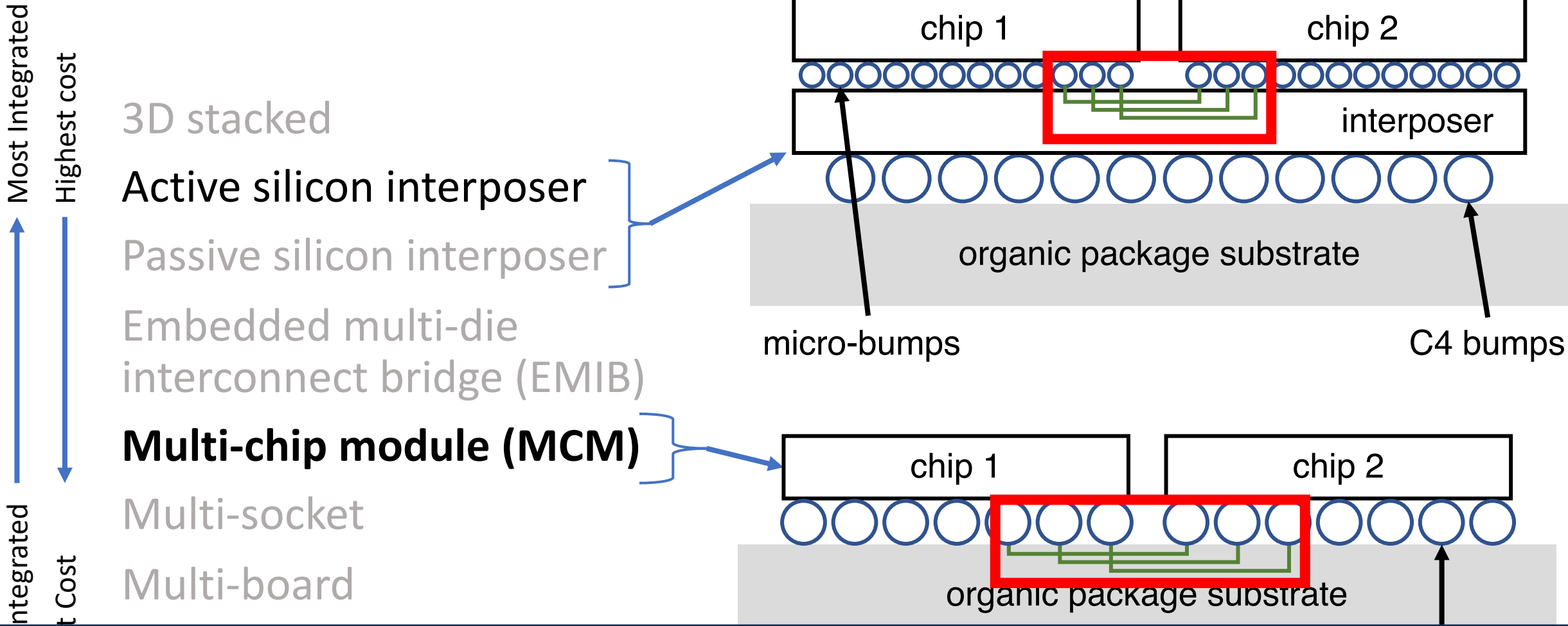
Many Ways to Connect Chiplets

Most Integrated
Highest cost
↑
↓
Least Integrated
Lowest Cost

- 3D stacked
- Active silicon interposer
- Passive silicon interposer
- Embedded multi-die interconnect bridge (EMIB)
- Multi-chip module (MCM)**
- Multi-socket
- Multi-board



Many Ways to Connect Chiplets



Disintegration constrains the inter-chip bandwidth and latency

Methodology

Diverse suite of 29 multithreaded benchmarks

- Splash-2x, Splash-4, Parsec, PBBS, etc.
- Regular and irregular parallel algorithms
- 10 runs/benchmark

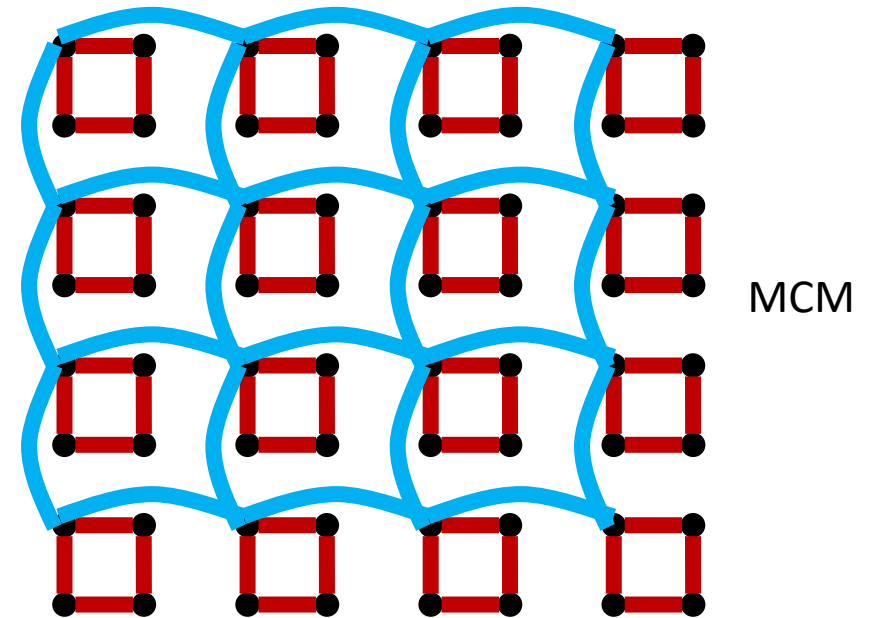
Cycle level simulator¹

- 256 cores, 2-wide out-of-order
- 4 cores/tile, 1 on-chip router/tile
- 16 chips, 16 cores/chip

Implement 3 network topologies

- Monolith (mesh)
- MCM: Hierarchical Mesh

1: <https://github.com/SwarmArch/sim>



Motivation

The disintegration slowdown varies across applications

Why Disintegrate? → To Increase Performance/\$

$$\frac{Perf_{Disintegrated}(A)}{\$Disintegrated} > \frac{Perf_{Monolith}(A)}{\$Monolith}$$

Why Disintegrate? → To Increase Performance/\$

$$\frac{Perf_{Disintegrated}(A)}{Perf_{Monolith}(A)} > \frac{\$_{Disintegrated}}{\$_{Monolith}}$$

Why Disintegrate? → To Increase Performance/\$

$$\frac{Perf_{Disintegrated}(A)}{Perf_{Monolith}(A)} > \frac{\$_{Disintegrated}}{\$_{Monolith}}$$

Disintegration slowdown

- Varies per Application

Cost savings

- Fixed for a system
- e.g., 0.59x for AMD EPYC¹

Why Disintegrate? → To Increase Performance/\$

This work →

$$\frac{Perf_{Disintegrated}(A)}{Perf_{Monolith}(A)} > \frac{\$_{Disintegrated}}{\$_{Monolith}}$$

Disintegration slowdown

- Varies per Application

Cost savings

- Fixed for a system
- e.g., 0.59x for AMD EPYC¹

What is the performance penalty of disintegration across applications?

Why Disintegrate? → To Increase Performance/\$

This work →
$$\frac{Perf_{Disintegrated}(A)}{Perf_{Monolith}(A)} > \frac{\$_{Disintegrated}}{\$_{Monolith}}$$

Disintegration slowdown

- Varies per Application

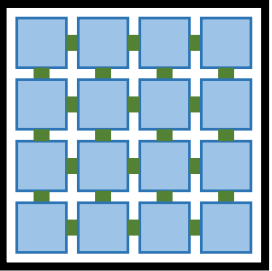
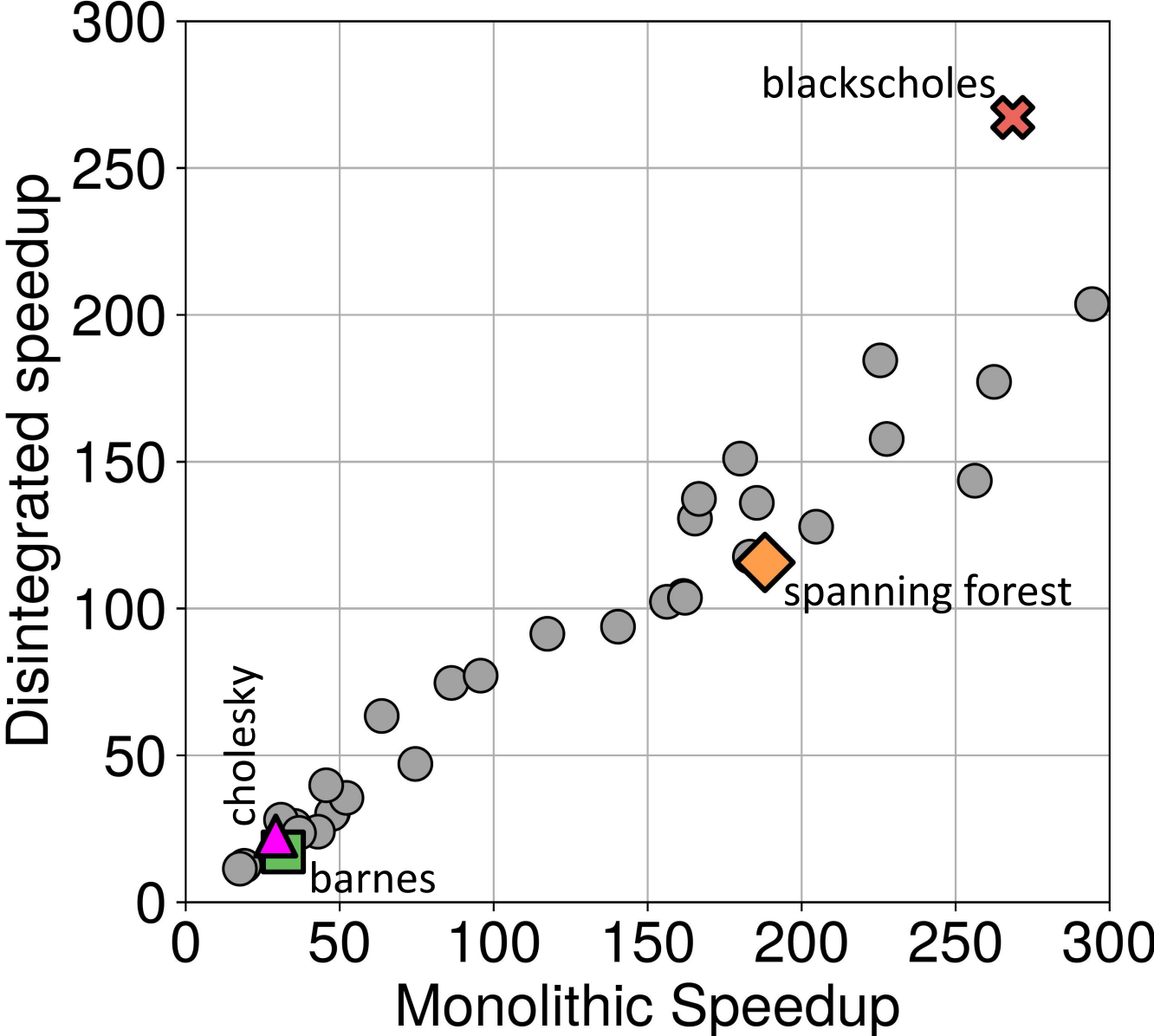
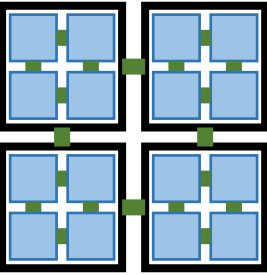
Cost savings

- Fixed for a system
- e.g., 0.59x for AMD EPYC¹

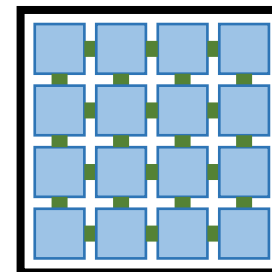
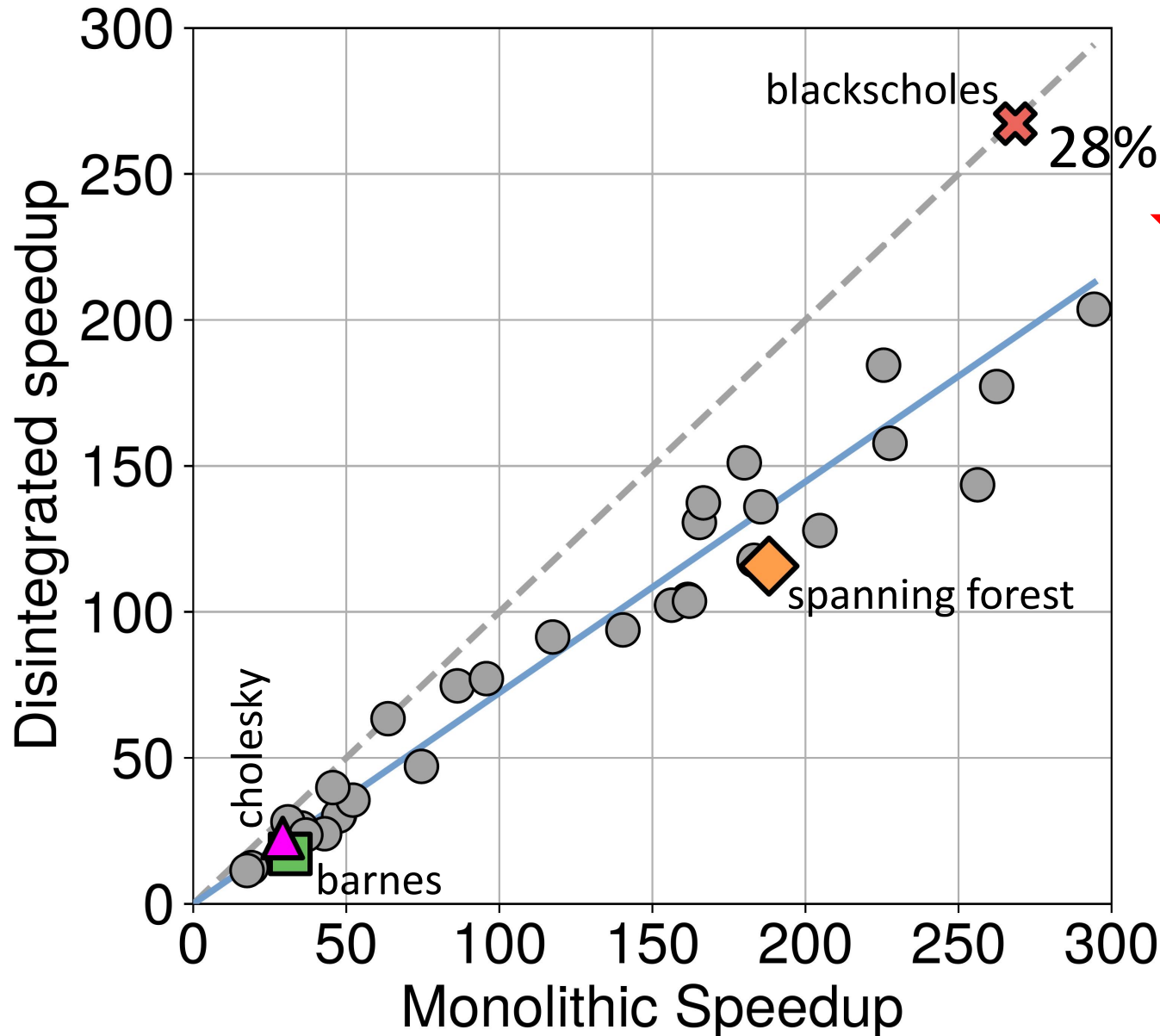
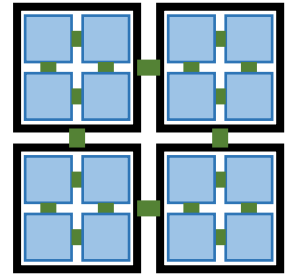
What is the performance penalty of disintegration across applications?

What application characteristics lead to slowdown on disintegrated systems?

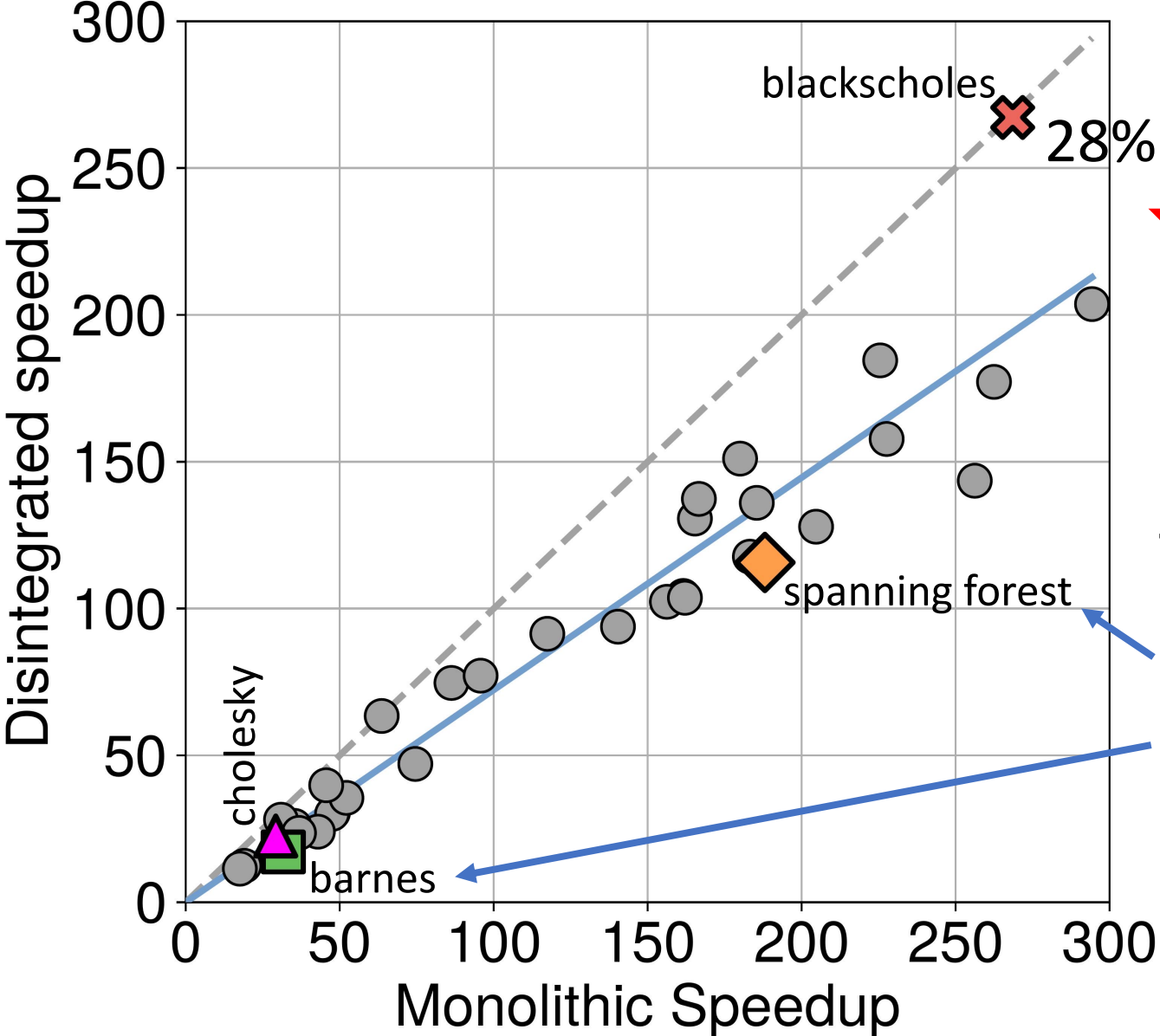
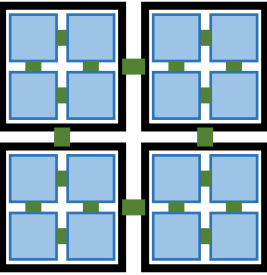
Unsurprisingly, Applications Continue to Scale



Unsurprisingly, Applications Continue to Scale



Unsurprisingly, Applications Continue to Scale



Equal performance

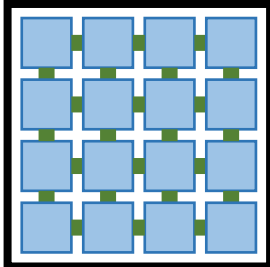


Line of best fit performance

$$\frac{Perf_{Disintegrated}}{Perf_{Monolith}}$$

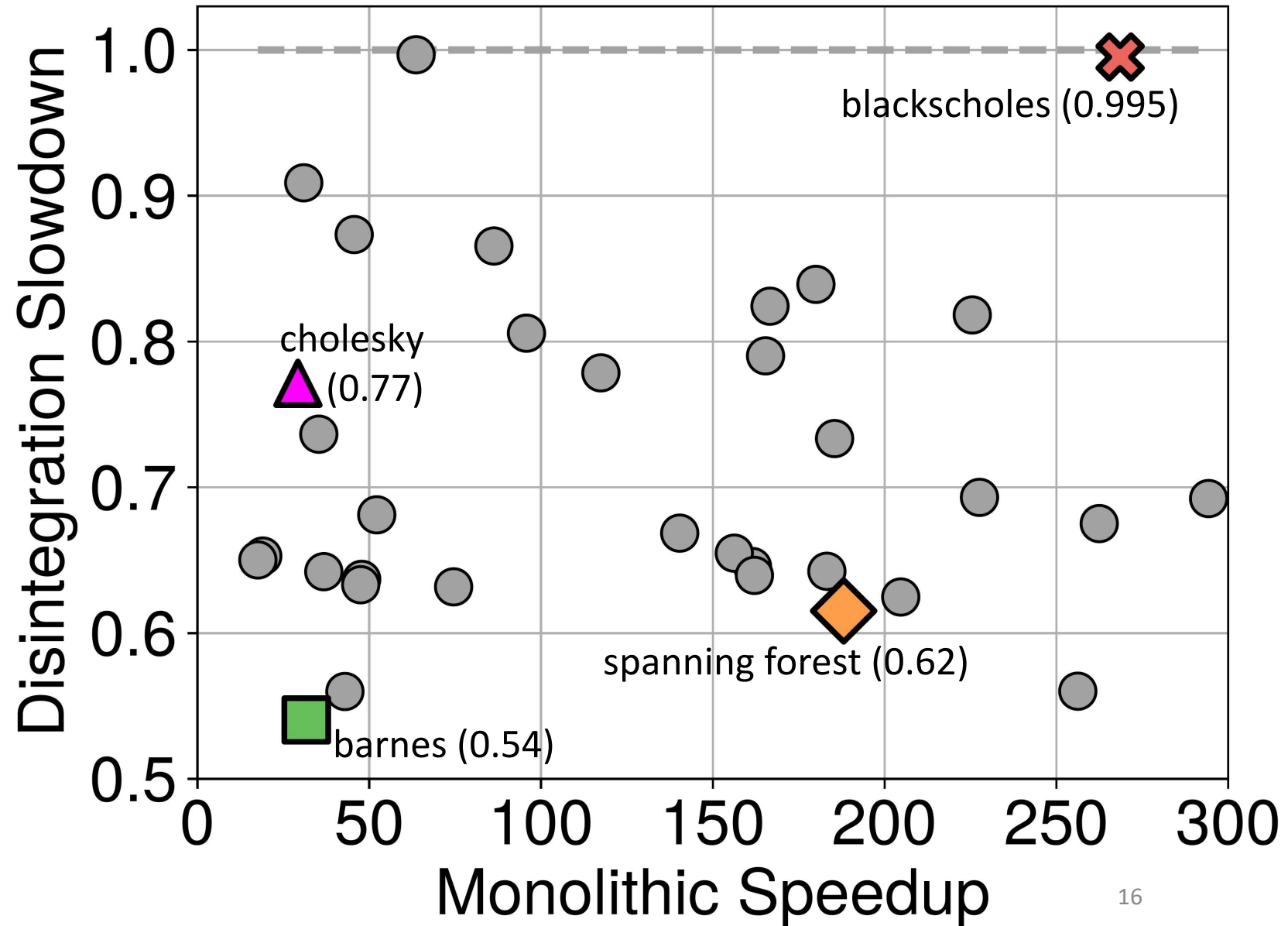
0.62x (spanning forest)

0.54x (barnes)



Large Variance in Application Slowdown

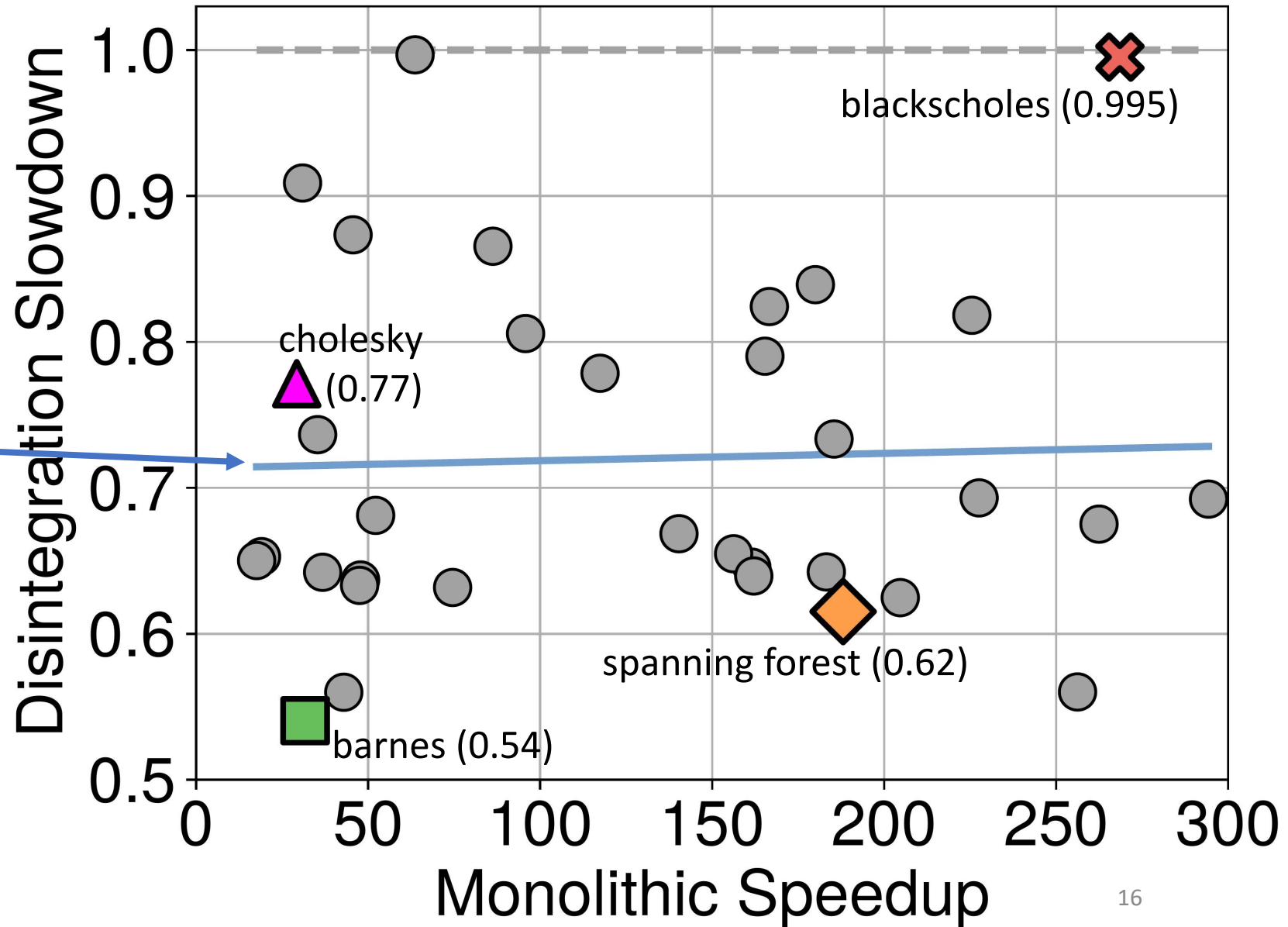
$$\text{Disintegration Slowdown} = \frac{\text{Perf}_{\text{Disintegrated}}(A)}{\text{Perf}_{\text{Monolith}}(A)}$$



Large Variance in Application Slowdown

$$\text{Disintegration Slowdown} = \frac{\text{Perf}_{\text{Disintegrated}}(A)}{\text{Perf}_{\text{Monolith}}(A)}$$

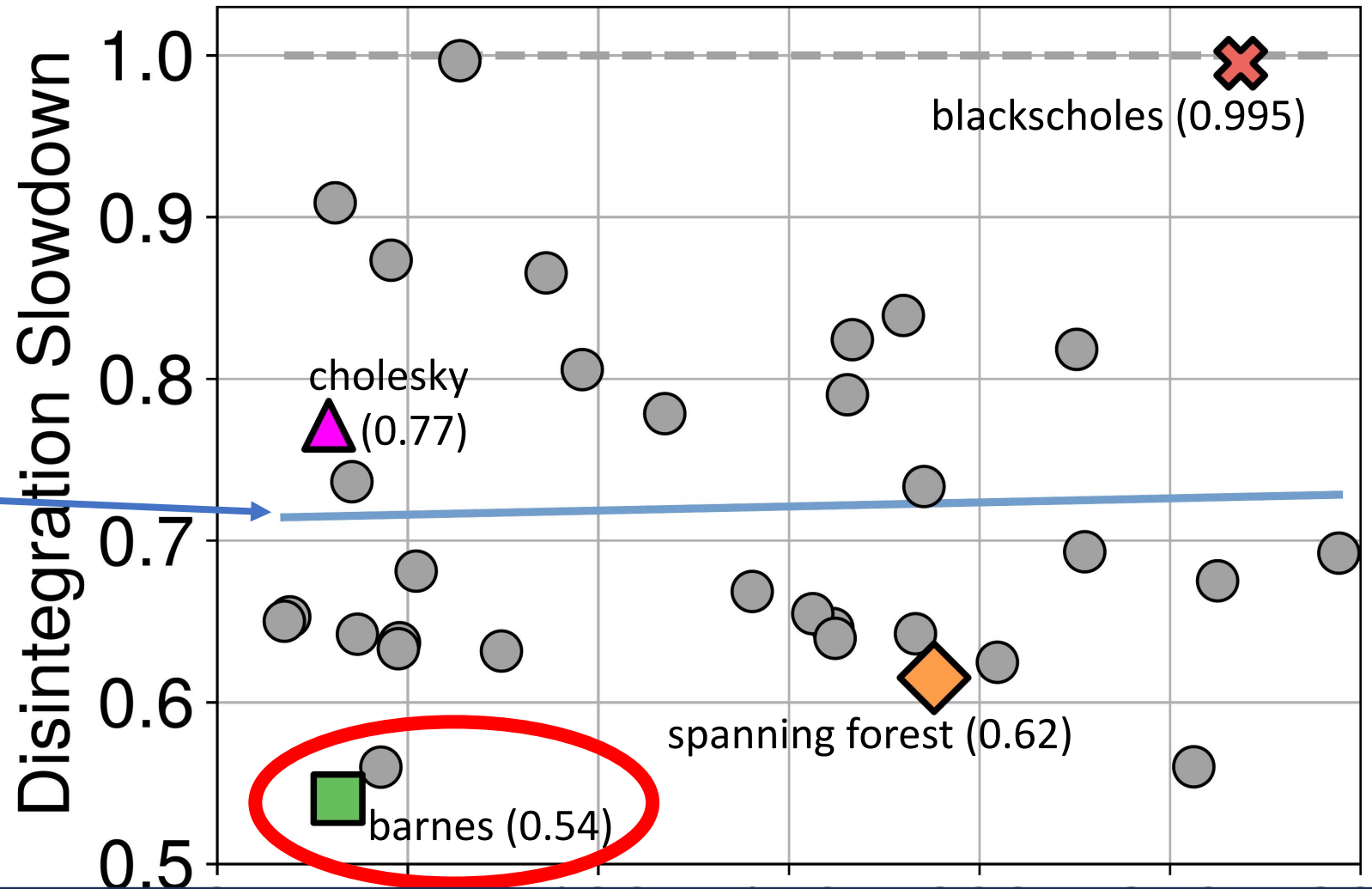
0.72x – average disintegration slowdown



Large Variance in Application Slowdown

$$\text{Disintegration Slowdown} = \frac{\text{Perf}_{\text{Disintegrated}}(A)}{\text{Perf}_{\text{Monolith}}(A)}$$

0.72x – average disintegration slowdown

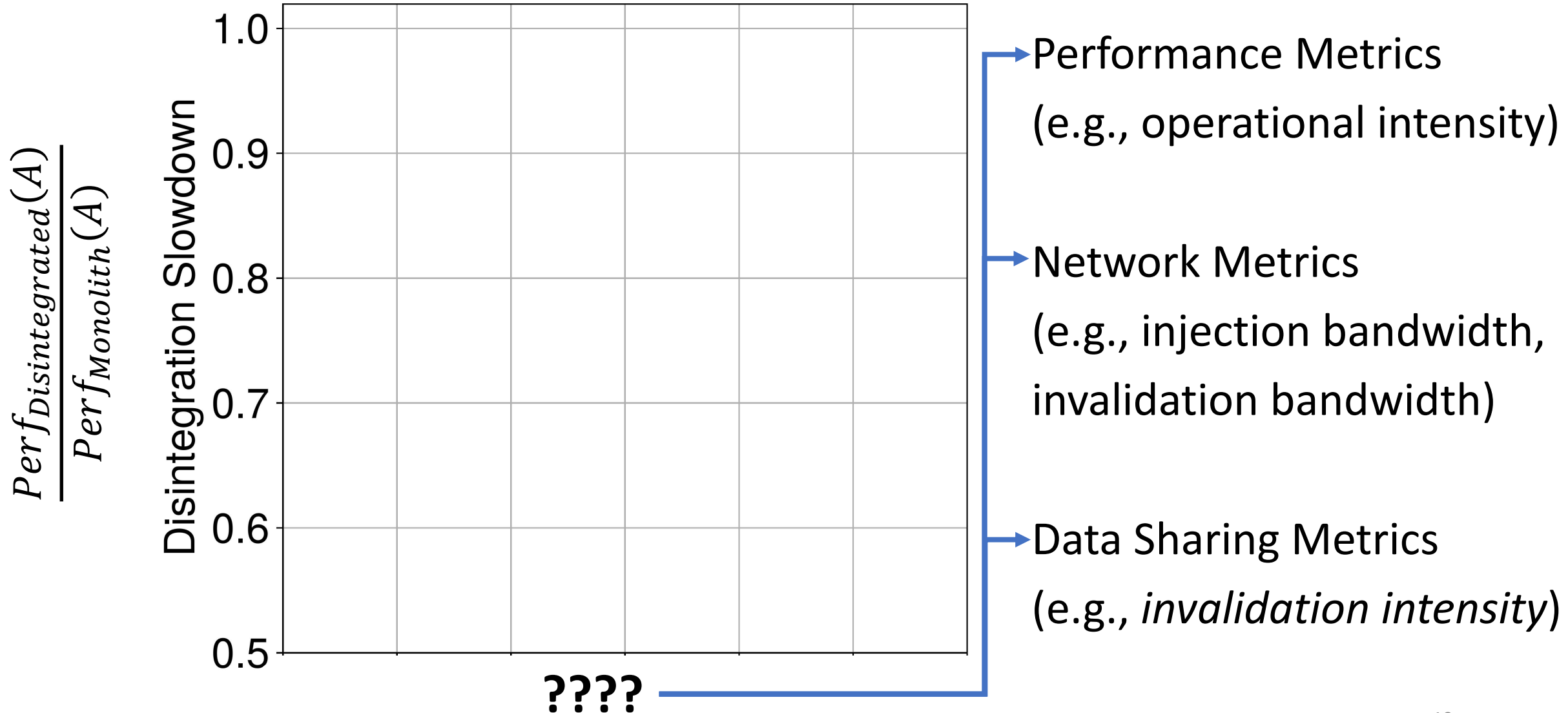


Monolithic scalability does not predict disintegration slowdown

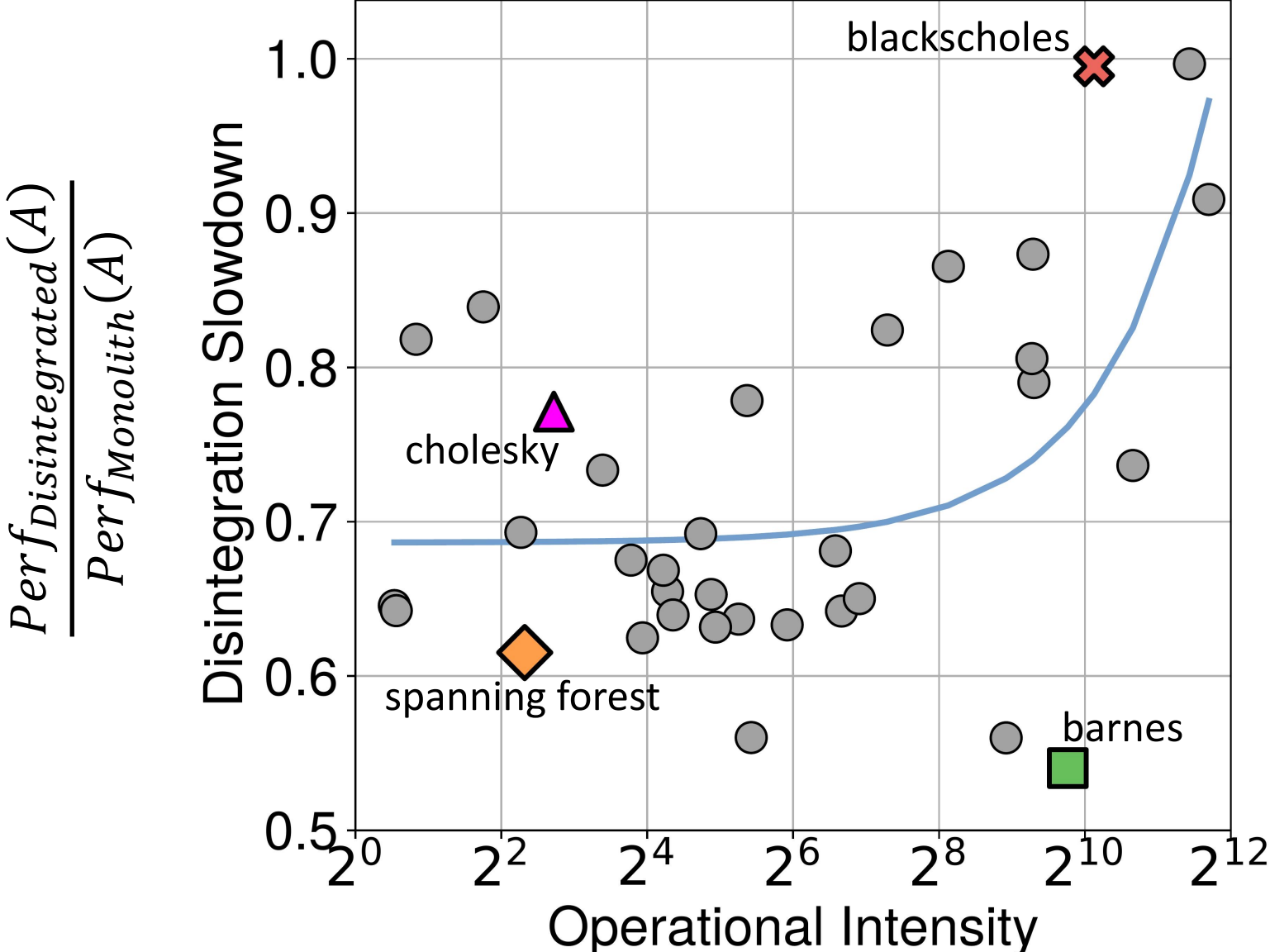
Results

What application characteristics predict slowdown on disintegrated systems?

What Metrics Correlate with Slowdown?

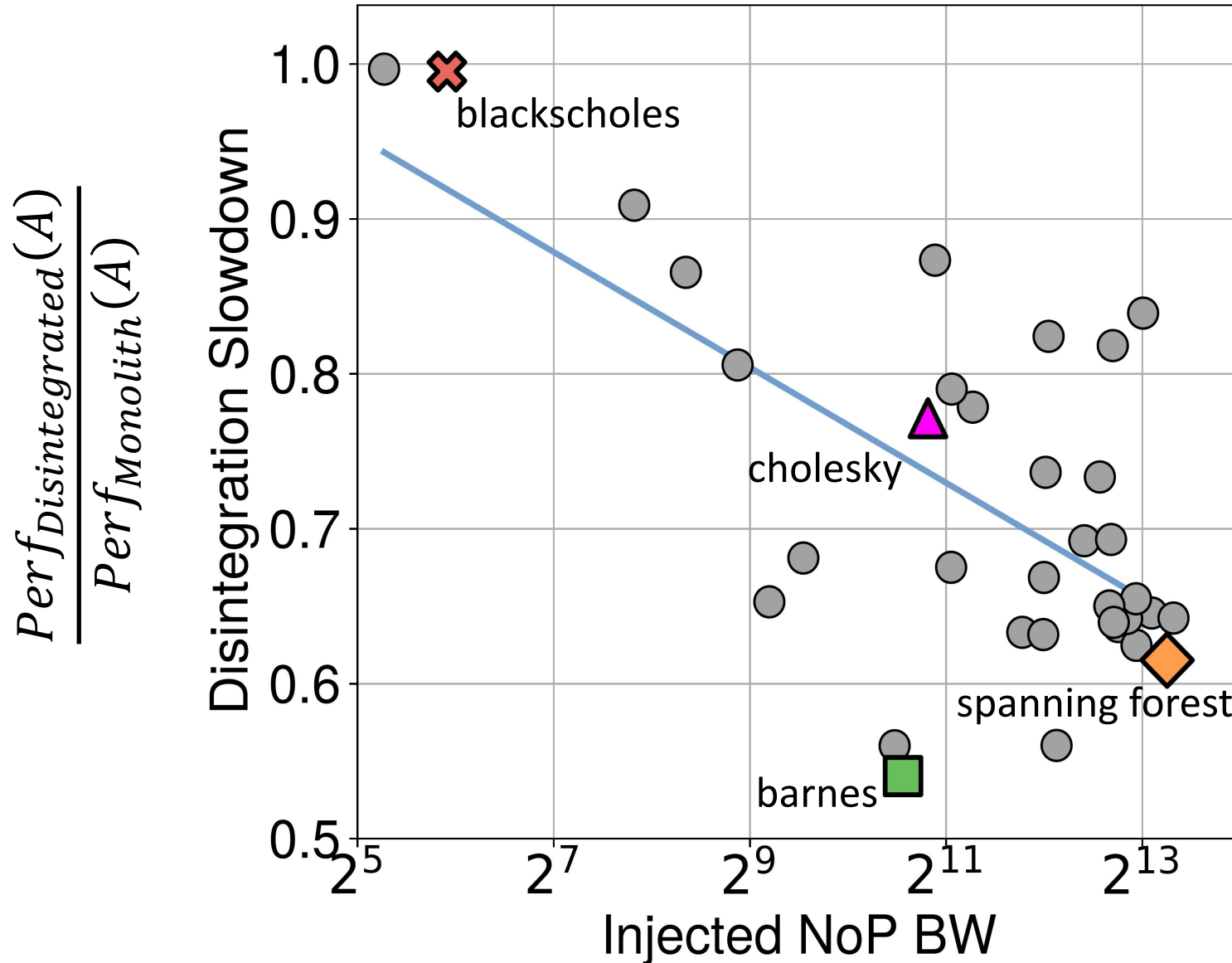


Performance Metrics: Do Not Explain Disintegration Slowdown



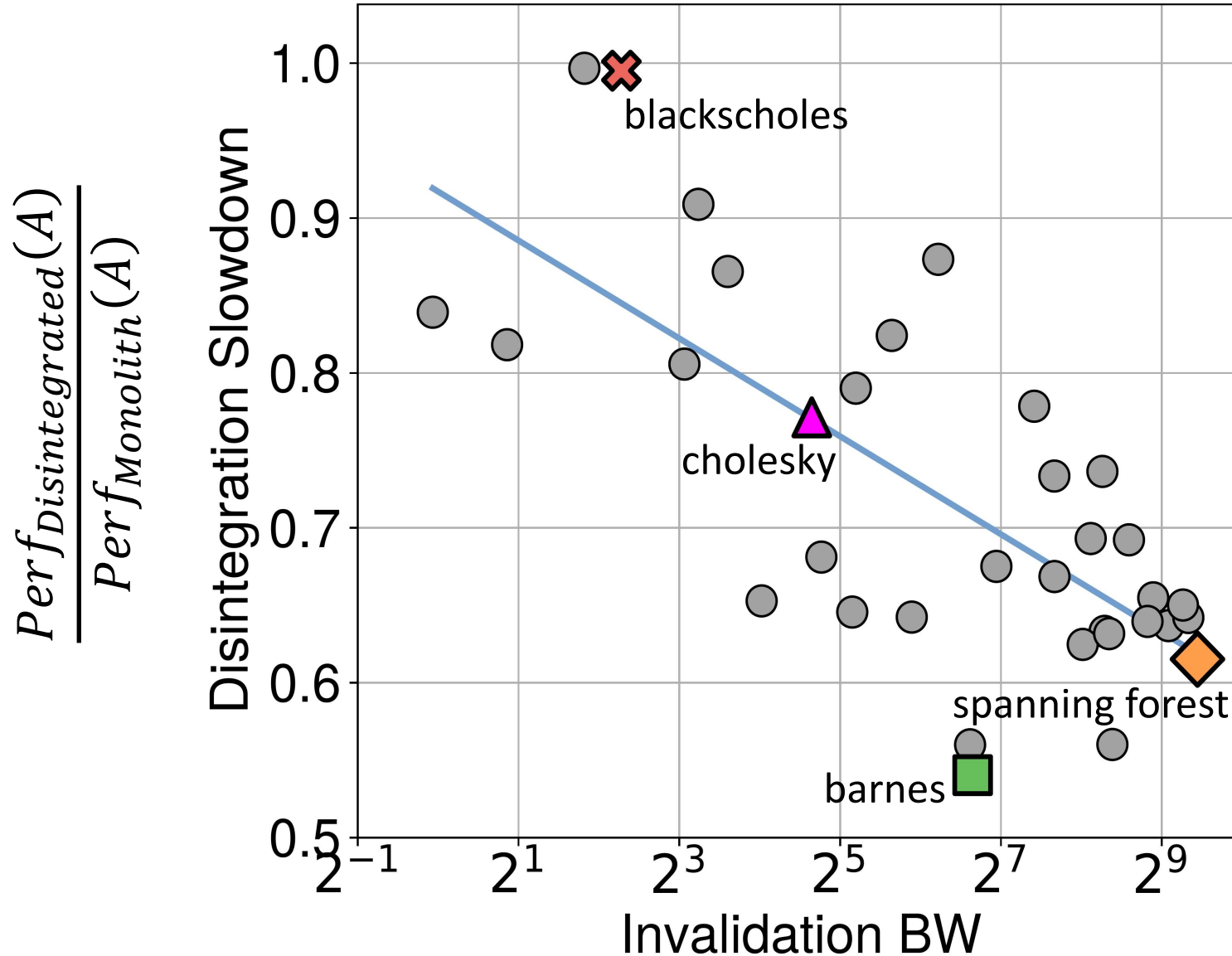
Operational Intensity:
instrs/byte
correlation up to 0.32

Network Bandwidth: Correlates Better With Disintegration Slowdown



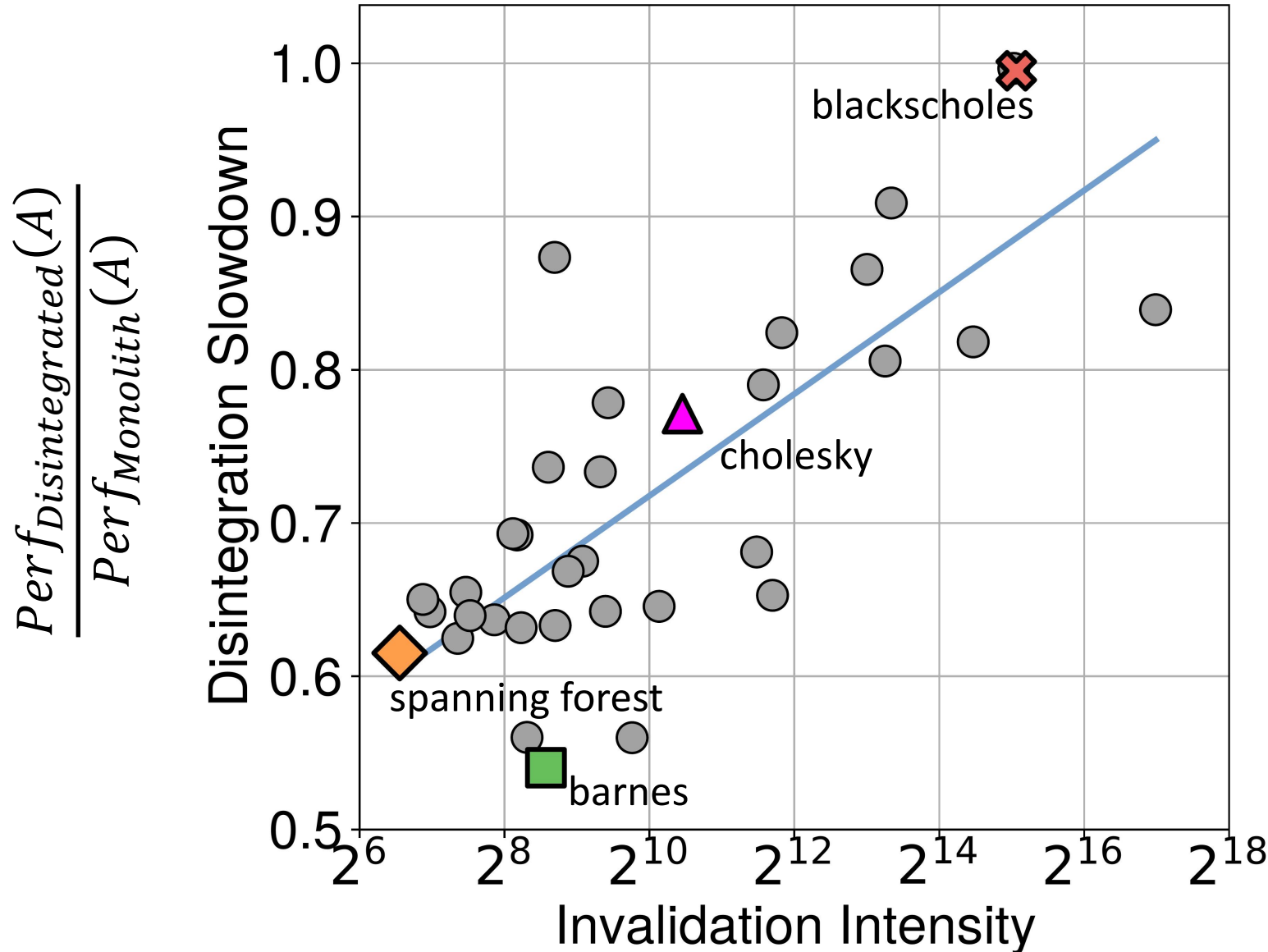
Network bandwidth:
MB/s/tile
correlation up to 0.42

Invalidation Bandwidth: Predict Slowdown Better Than Total Bandwidth



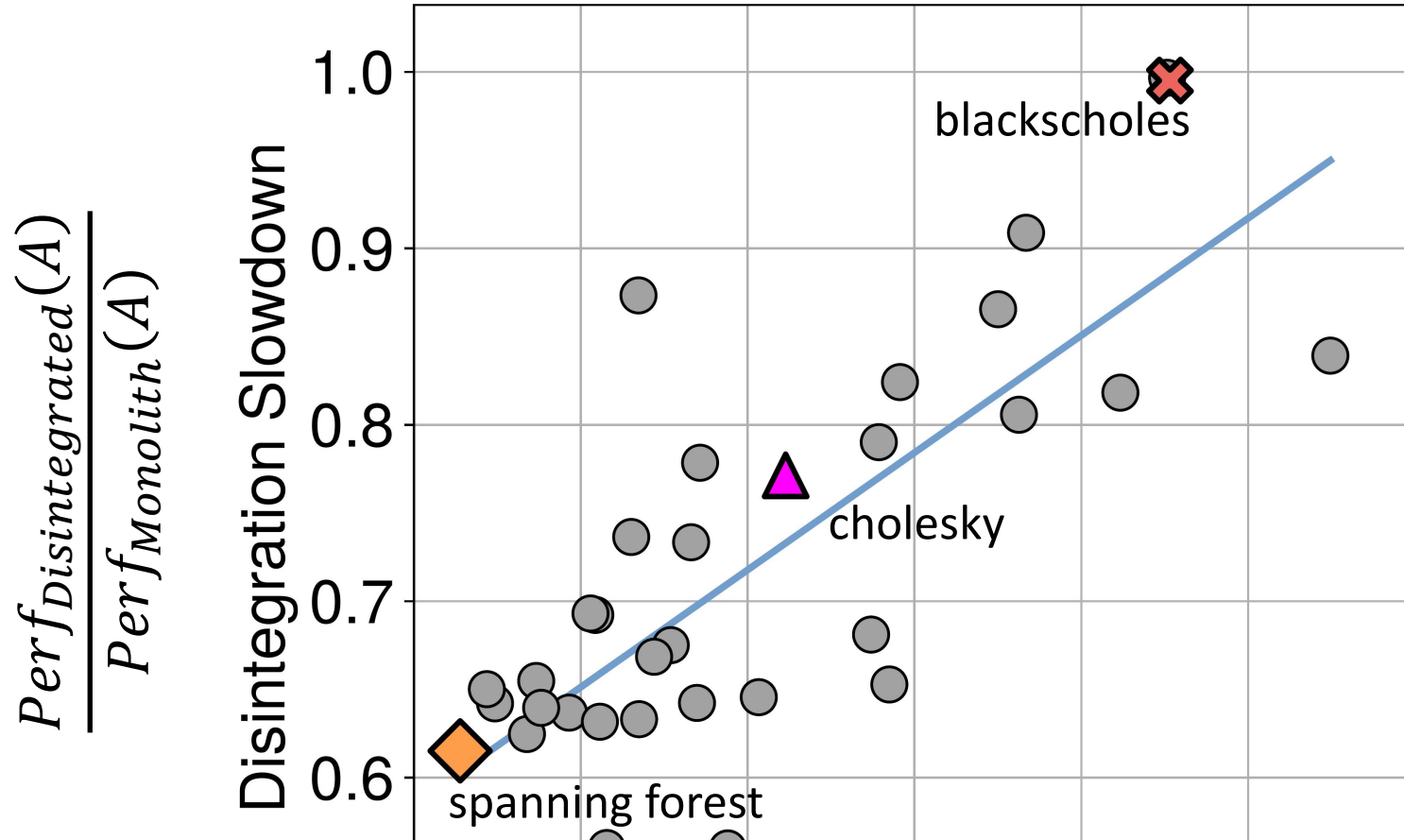
Invalidation bandwidth:
MB of Invalidations /s/tile
correlation up to 0.50

Our Data Sharing Metrics: Have The Best Observed Correlation



Invalidation intensity:
Instructions / Invalidation
correlation up to 0.58

Our Data Sharing Metrics: Have The Best Observed Correlation



Invalidation intensity:
Instructions / Invalidation
correlation up to 0.58

Invalidation and downgrade intensity are the most predictive known metrics

Invalidation Intensity

Metric Group	Metric	Correlation (ideal is 1.0)
Performance	IPC	0.14
	Operational Intensity	0.32
	Consumed Memory Bandwidth	0.13
Base Network	Network Injection Bandwidth	0.42
	Average Network Latency	0.06
Network Injection Bandwidth	GetS	0.41
	GetX	0.49
	Inv (Invalidate)	0.50
	InvX (Downgrade)	0.57
	Data Response	0.46
	PutS (Clean Eviction)	0.10
	PutX (Dirty Eviction)	0.05
Data Sharing	Read Sharers	0.07
	Invalidation Intensity	0.58
	Downgrade Intensity	0.59
	Sharing Fraction [Ferdman+, 2012]	0.31

Metric Group	Metric	Correlation (ideal is 1.0)
Performance	IPC	0.14

What application characteristics correlate with disintegration slowdown?

Base Network	Network Injection Bandwidth	0.42
	Average Network Latency	0.06
Network Injection Bandwidth	GetS	0.41
	GetX	0.49
	Inv (Invalidate)	0.50
	InvX (Downgrade)	0.57
	Data Response	0.46
	PutS (Clean Eviction)	0.10
	PutX (Dirty Eviction)	0.05
Data Sharing	Read Sharers	0.07
	Invalidation Intensity	0.58
	Downgrade Intensity	0.59
	Sharing Fraction [Ferdman+, 2012]	0.31

Metric Group	Metric	Correlation (ideal is 1.0)
Performance	IPC	0.14

What application characteristics correlate with disintegration slowdown?

Base Network	Network Injection Bandwidth	0.42
	Average Network Latency	0.06
Network Injection Bandwidth	GetS	0.41
	GetX	0.49
	Inv (Invalidate)	0.50
	InvX (Downgrade)	0.57
	Data Response	0.46
Data Sharing	PutS (Clean Eviction)	0.10
	PutX (Dirty Eviction)	0.05
	Read Sharers	0.07
	Invalidation Intensity	0.58
	Downgrade Intensity	0.59
	Sharing Fraction [Ferdman+, 2012]	0.31

Conclusion

Industry has adopted disintegrated systems to reduce cost

The penalty of disintegration varies across applications

Disintegration constrains inter-chiplet links

Sharing intensity is correlated with worse disintegration slowdown

Future work: Support performance for **all** applications on disintegrated systems

Q & A

Industry has adopted disintegrated systems to reduce cost

The penalty of disintegration varies across applications

Disintegration constrains inter-chiplet links

Sharing intensity is correlated with worse disintegration slowdown

Future work: Support performance for **all** applications on disintegrated systems

Isidor R. Brkić, Mark C. Jeffrey



UNIVERSITY OF
TORONTO