

Higher Lower Bounds for Near-Neighbor and Further *Rich* Problems

Mihai Pătrașcu Mikkel Thorup



FOCS 2006

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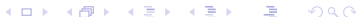
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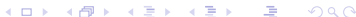
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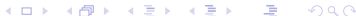
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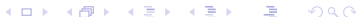
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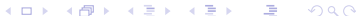
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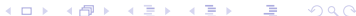
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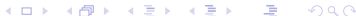
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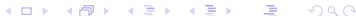
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Data Structures and Communication



<i>Cell-probe data structure</i>	<i>Communication game</i>
space S	querier sends $\lg S$ bits
d -bit cells	memory sends d bits
query time T	T rounds

Lost in Translation

Memory can remember past communication

Data Structures and Communication



<i>Cell-probe data structure</i>	<i>Communication game</i>
space S	querier sends $\lg S$ bits
d -bit cells	memory sends d bits
query time T	T rounds

Lost in Translation

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What can this prove?

Richness technique \Rightarrow one of $\begin{cases} \text{querier sends } A \text{ bits} \\ \text{memory sends } B \text{ bits} \end{cases}$

$$\Rightarrow \text{one of } \begin{cases} T \lg S \geq A \\ Td \geq B \end{cases} \Rightarrow T \geq \min\left\{\frac{A}{\lg S}, \frac{B}{d}\right\}$$

Best lower bound possible

$$A \sim d, B \sim n \Rightarrow T = \Omega\left(\min\left\{\frac{d}{\lg S}, \frac{n}{d}\right\}\right) \quad \dots \text{typically, } T = \Omega\left(\frac{d}{\lg S}\right)$$

Our result: A “black-box” improvement to $T = \Omega(d / \lg \frac{Sd}{n})$.

	$S = n^{O(1)}$	$T = \Omega\left(\frac{d}{\lg n}\right)$	
	$S = n(d \lg n)^{O(1)}$	$T = \Omega\left(\frac{d}{\lg n}\right)$	$T = \Omega\left(\frac{d}{\lg d}\right)$
$d = O(\lg n)$	$S = n^{O(1)}$	$T = O(1)$	
$d = O(\lg n)$	$S = n \lg^{O(1)} n$	—	$T = \Omega\left(\frac{\lg n}{\lg \lg n}\right)$

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Known Richness Results

- partial match — *curse?*
- near neighbor search:

	Deterministic	Randomized
Exact	<i>curse?</i>	<i>curse?</i>
Approximate	<i>curse?</i>	<i>no curse</i>

MSNW STOC'95: p.m. $\Omega(\sqrt{\lg d})$ rand

Borodin, Ostrovsky, Rabani STOC'99: p.m., ENN $\Omega(\lg d)$ rand

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Liu IPL'04: ANN $\Omega(d)$ det

Andoni, Indyk, P. FOCS'06: ANN $\Omega(\frac{1}{\epsilon^2} \lg n)$ rand

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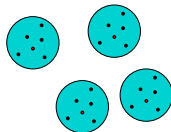
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Proof Technique

- 1 break input into k subproblems
e.g.: for NN, k far clusters, $\frac{n}{k}$ points each



- 2 simulate k queries in parallel:

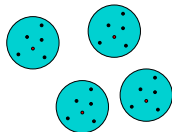
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NB: $k \lg \frac{S}{k} < k \lg S$. This step outside communication paradigm.

- 3 prove direct-sum law for richness:
one problem has $[A, B]$ lower bound **by richness**
 $\Rightarrow k$ problems have $[\Omega(kA), \Omega(kB)]$ lower bound

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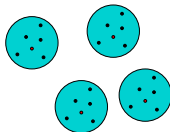
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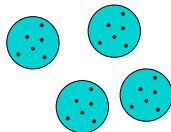
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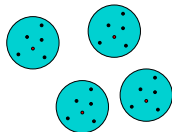
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deterministic richness: simple combinatorics

randomized richness: complicated combinatorics

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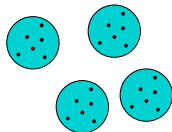
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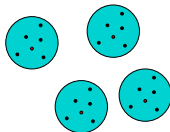
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Conclusion: $\begin{cases} Tk \lg \frac{S}{k} = \Omega(kA) = \Omega(kd) \\ Tkd = \Omega(kB) = \Omega(k \frac{n}{k}) \end{cases} \Rightarrow T = \Omega(d / \lg \frac{Sd}{n}).$

Conclusions

GOOD: simulating k queries at the same time

BAD: breaking into k subproblems

works if hardness comes from dimension (not for ANN)

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