# Algorithmic Aspects of Reinforcement Learning I 

## Ankur Moitra (MIT)

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## TUTORIAL GOALS

(1) Overview of theoretical foundations of RL

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(2) Gaps in our algorithmic understanding

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(1) Overview of theoretical foundations of RL
(2) Gaps in our algorithmic understanding
(3) Deep dive into some success stories, emphasizing connections to other areas

## INTRODUCTION

Success stories of reinforcement learning:


## robotic manipulation

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Success stories of reinforcement learning:
playing strategic games

## robotic manipulation



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Goal: Agent learns by interacting with the environment


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The basic model is a Markov Decision Process

## MARKOV DECISION PROCESSES



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Goal: Find a strategy that maximizes expected reward

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Goal: Find a straegy that maximizes expected reward policy $\pi: \mathcal{S} \rightarrow \mathcal{A}$ i.e. choice of action doesn't depend on past

## OUTLINE

Part I: Tabular Markov Decision Processes

- Planning vs. Learning

Interlude: Can We Make RL Algorithmically Tractable?
Part II: Partial Observations and the Curse of History

- Beyond Worst-Case Analysis

Part III: Planning and Belief Contraction

- Polynomial vs. Exponential Rates

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## MAIN PROBLEMS

## Planning <br> (computational)

Given full description of the MDP, compute an optimal policy

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Uses planning as a subroutine

## VALUE FUNCTIONS

Key: Keep track of how well you will do in the future

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Definition: The value function of a policy is

$$
V_{h}^{\pi}(s)=\mathbb{E}\left[R_{h}\left(s_{h}, a_{h}\right)+\cdots+R_{H}\left(s_{H}, a_{H}\right) \mid s_{h}=s\right]
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i.e. it gives the expected future reward starting from state $s$ at timestep h

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Bellman Optimality: An optimal policy $\pi^{*}$ must satisfy

$$
V_{h}^{\pi^{*}}(s)=\max _{a} R_{h}(s, a)+\mathbb{E}_{s^{\prime}}\left[V_{h+1}^{\pi^{*}}\left(s^{\prime}\right)\right]
$$

for every state s, i.e. value function must be consistent

## VALUE ITERATION

This gives an efficient algorithm for planning:

Initialize V = 0 (assuming no rewards at step H)
Repeat until convergence
Scan through states, update any violated V constraint

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Of course, this is just dynamic programming
Moreover can find the optimal policy from the $V^{\pi^{*}}$ values

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Given full description of the MDP, compute an optimal policy
e.g. value iteration, policy iteration,
linear programming

Given budget of interactions with the environment, learn an optimal policy

## EXPLORATION VS EXPLOITATION

First non-asymptotic result:
Theorem [Kearns, Singh '02]: There is an algorithm that has polynomial running time and sample complexity that outputs an $\epsilon$-suboptimal policy in tabular MDPs

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(2) Trade off playing the optimal policy in current model vs. discovering new states

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Tight regret bounds given by [Azar, Osband, Munos "17]

## POLICY GRADIENTS

Suppose we parameterize the class of policies by $\theta$--- i.e. we want to maximize

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random trajectory under $\pi_{\theta}$

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random trajectory under $\pi_{\theta}$
How can we compute the gradient without full knowledge of the environment?

Policy Gradient Theorem: In fact

$$
\nabla J(\theta)=\mathbb{E}_{\pi_{\theta}}\left[R(\tau) \nabla \log \pi_{\theta}(\tau)\right]
$$

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Many challenges both in theory and practice, e.g. delayed feedback can cause gradients to be extremely small
e.g. see [Agarwal, Kakade, Lee, Mahajan '19]

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Given budget of interactions with the environment, learn an optimal policy
e.g. model based, Q-learning, actor-critic policy gradient

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Partially observable MDPs (POMDPs)

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Existing theory is built around (1) and (2) but what do we miss out on by ignoring (3)?

## WHAT ABOUT COMPUTATIONAL COMPLEXITY?

Returning to our earlier picture


Are there computationally efficient algorithms with strong end-to-end provable guarantees?

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## MARKOV DECISION PROCESSES



## PARTIALLY OBSERVABLE MDPS (POMDPS)



## PLANNING IS HARD

Classic lower bound:

Theorem [Papadimitriou, Tsitsiklis]: Optimal planning in a POMDP is PSPACE hard

## THE CURSE OF HISTORY

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| MDPs | POMDPs |
| :---: | :---: |
| Optimal action only <br> depends on current state | Optimal action depends on <br> action/observation history |
| $\pi: \mathcal{S} \rightarrow \mathcal{A}$ | $\pi: \mathcal{A} \times \mathcal{O} \cdots \times \mathcal{O} \rightarrow \mathcal{A}$ |
| Alternatively, it depends |  |
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Natural approaches use exponential space $(|\mathcal{A} \| \mathcal{O}|)^{H}$ or $C^{|\mathcal{S}|}$

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Even worse news:
Theorem [Golowich, Moitra, Rohatgi '23]: Unless the exponential time hierarchy collapses, there is no polynomial sized description of an approximately optimal policy

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Why should real-world POMDPs have succinct descriptions of good policies?

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## "The observations don't tell you anything about the state"

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Could this enable tractable planning/learning?

## BEYOND WORST-CASE ANALYSIS

Definition: We say the POMDP is $\gamma$-observable if for all $h$ and all distributions $b, b^{\prime}$ on states we have

$$
\left\|\mathbb{O}_{h} b-\mathbb{O}_{h} b^{\prime}\right\|_{1} \geq \gamma\left\|b-b^{\prime}\right\|_{1}
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i.e. well-separated distributions on states lead to well-separated distributions on observations

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Introduced by [Even-Dar, Kakade, Mansour] for understanding stability of beliefs in HMMs under misspecification

Key Point: No assumption on transition dynamics like e.g. deterministic transitions or mixing (under every possible policy)

## PLANNING VIA STABILITY

There is a quasi-polynomial time algorithm for planning under observability:

Theorem [Golowich, Moitra, Rohatgi '23]: Given description of a $\gamma$-observable POMDP there is an algorithm running in time

$$
H(|\mathcal{O}||\mathcal{A}|)^{C \log (|\mathcal{S}| H / \epsilon) / \gamma^{4}}
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that outputs an $\epsilon$-suboptimal policy

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Key Idea: The Bayes filter is exponentially stable compute posterior on states, given actions/observations

Parallels well-known stability results for Kalman filtering

## LOWER BOUNDS

Moreover these results are tight

Theorem [Golowich, Moitra, Rohatgi '23]: Under the Exponential Time Hypothesis, there is no algorithm running in time

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(|\mathcal{S}||\mathcal{A}| H|\mathcal{O}|)^{o(\log (|\mathcal{S}||\mathcal{A}| H|\mathcal{O}| / \epsilon) / \gamma)}
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for finding an $\epsilon$-suboptimal policy in a $\gamma$-observable POMDP

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It's hard even in the lossy case, where you observe the state with probability $\gamma$ independently at each step

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## BELIEF CONTRACTION

Theorem: Fix any $\gamma$-observable POMDP and policy $\pi$. Then

$$
\mathbb{E}_{\tau}\left[\left\|b_{t}-b_{t}^{\prime}\right\|_{1}\right] \leq\left(1-\gamma^{4}\right)^{t}|\mathcal{S}|
$$

posterior, starting from arbitrary belief state
posterior, starting from uniform belief state
where $\tau$ is the trajectory from the POMDP by playing $\pi$

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posterior, starting from arbitrary belief state
posterior, starting from uniform belief state
where $\tau$ is the trajectory from the POMDP by playing $\pi$
Thus, we could ignore all but the most recent history

## BELIEF UPDATES

Definition: The Bayes operator, given an observation

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updates the posterior on states, and is defined as

$$
B_{h}(b, y)(x)=\frac{\mathbb{O}_{h}(y \mid x) b(x)}{\sum_{z \in \mathcal{S} \mathbb{O}_{h}(y \mid z) b(z)}}
$$

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Definition: And the update operator, given both an action and observation

based on the chosen action, takes a step
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U_{h}(b, a, y)=B_{h}\left(\mathbb{T}_{h}(a) b, y\right)
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## TOWARDS A WEAKER BOUND

From the data processing inequality, we have that for any action

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But for some observations, the Bayes operator can increase the KL-divergence

Do we make progress in expectation?

Lemma: Given beliefs $b, b^{\prime}$

$$
\underset{y \sim \mathbb{O}_{h} b}{\mathbb{E}}\left[K L\left(B_{h}(b, y) \| B_{h}\left(b^{\prime}, y\right)\right)\right]=K L\left(b \| b^{\prime}\right)-K L\left(\mathbb{O}_{h} b \| \mathbb{O}_{h} b^{\prime}\right)
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Lemma: Given beliefs $b, b^{\prime}$

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Notation: Let $P_{X}=b, Q_{X}=b^{\prime}$ and $P_{Y \mid X=x}=\mathbb{O}_{h}(\cdot \mid x)$.

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Proof: $\quad K L\left(b \| b^{\prime}\right)=K L\left(P_{X} \| Q_{X}\right)$

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$$
\left.=K L\left(P_{X, Y} \| Q_{X, Y}\right) \underset{x}{+\mathbb{E}\left[K P_{X}\right.}\left[P_{Y \mid X=x} \| Q_{Y \mid X=x}\right)\right]
$$

..using the chain rule

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\left.=K L\left(P_{X, Y} \| Q_{X, Y}\right) \underset{x \sim P_{X}}{\substack{\mathbb{E}}} \underset{P_{Y}}{ } L\left(P_{Y \mid X=x} \mid Q_{Y \mid X=x}\right)\right]
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..using the chain rule, and the fact that $Q_{Y \mid X=x}=P_{Y \mid X=x}$

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$$
\begin{aligned}
& =K L\left(P_{X, Y} \| Q_{X, Y}\right) \\
& =K L\left(P_{Y} \| Q_{Y}\right)+\mathbb{E}\left[K L\left(P_{X \mid Y=y} \| Q_{X \mid Y=y}\right)\right] \\
& y
\end{aligned}
$$

... this time using the chain rule in opposite order

Lemma: Given beliefs $b, b^{\prime}$
$\underset{y \sim \mathbb{O}_{h} b}{\mathbb{E}}\left[K L\left(B_{h}(b, y) \| B_{h}\left(b^{\prime}, y\right)\right)\right]=K L\left(b \| b^{\prime}\right)-K L\left(\mathbb{O}_{h} b \| \mathbb{O}_{h} b^{\prime}\right)$
Notation: Let $P_{X}=b, Q_{X}=b^{\prime}$ and $P_{Y \mid X=x}=\mathbb{O}_{h}(\cdot \mid x)$.
$P_{X, Y}=P_{Y \mid X} P_{X}$ and $Q_{X, Y}=P_{Y \mid X} Q_{X}$
Proof: $\quad K L\left(b \mid b^{\prime}\right)=K L\left(P_{X} \| Q_{X}\right)$

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& y \sim P_{Y} \\
& K L\left(B_{h}(b, y) \| B_{h}\left(b^{\prime}, y\right)\right)
\end{aligned}
$$

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$$

Does this imply fast enough convergence?

Using Pinsker's inequality (1) and observability (2), we have

$$
\underset{y \sim \mathbb{O}_{h} b}{\mathbb{E}}\left[K L\left(B_{h}(b, y) \| B_{h}\left(b^{\prime}, y\right)\right)\right] \stackrel{(1)}{\leq} K L\left(b \| b^{\prime}\right)-\frac{1}{2}\left\|\mathbb{O}_{h} b-\mathbb{O}_{h} b^{\prime}\right\|_{1}^{2}
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(1) & \leq K L\left(b \| b^{\prime}\right)-\frac{1}{2}\left\|\mathbb{O}_{h} b-\mathbb{O}_{h} b^{\prime}\right\|_{1}^{2} \\
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& (2) \\
& \leq K L\left(b \| b^{\prime}\right)-\frac{\gamma^{2}}{2}\left\|b-b^{\prime}\right\|_{1}^{2}
\end{aligned}
$$

Using reverse Pinsker's inequality, we get

$$
\underset{y \sim \mathbb{O}_{h} b}{\mathbb{E}}\left[K L\left(B_{h}(b, y) \| B_{h}\left(b^{\prime}, y\right)\right)\right] \leq K L\left(b \| b^{\prime}\right)-c \gamma^{2} K L\left(b \| b^{\prime}\right)^{2}
$$

provided that $\left\|b / b^{\prime}\right\|_{\infty}$ is bounded

Theorem [Even-Dar et al.]: Fix any $\gamma$-observable POMDP and policy $\pi$. Then we have

$$
\mathbb{E}_{\tau}\left[K L\left(b_{t} \| b_{t}^{\prime}\right)\right] \leq \epsilon
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provided that $t \geq 1 /\left(\gamma^{2} \epsilon\right)$

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Inverse polynomial, rather than exponential convergence :(

Unfortunately there are cases where progress can be slow, but...

## A WIN-WIN ARGUMENT

We show that either
(1) A stronger reverse Pinsker holds, i.e.
$\underset{y \sim \mathbb{O}_{h} b}{\mathbb{E}}\left[K\left(B_{h}(b, y) \| B_{h}\left(b^{\prime}, y\right)\right)\right] \leq K L\left(b \| b^{\prime}\right)-\frac{\gamma^{2}}{32} \min \left(K L\left(b \| b^{\prime}\right), 1\right)$

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or instead
(2) Progress is anti-concentrated, i.e. for some event $\mathcal{E} \subset \mathcal{O}$

$$
\underset{y \sim \mathbb{O}_{h} b}{\mathbb{E}}\left[\left(K L\left(B_{h}(b, y) \| B_{h}\left(b^{\prime}, y\right)\right)-K L\left(b \| b^{\prime}\right)\right) \mathbf{1}_{y \in \mathcal{E}}\right] \leq-\frac{\gamma}{8} K L\left(b \| b^{\prime}\right)
$$

## A WIN-WIN ARGUMENT

As a result, we get:
Corollary: For any $\gamma$-observable POMDP

$$
\begin{aligned}
& \underset{y \sim \mathbb{O}_{h} b}{\mathbb{E}\left[\sqrt{K L\left(B_{h}(b, y) \| B_{h}\left(b^{\prime}, y\right)\right)}\right]} \\
& \quad \leq\left(1-\Omega\left(\frac{\gamma^{2}}{\left.\max \left(1, K L\left(b \| b^{\prime}\right)\right)\right)}\right)\right) \sqrt{K L\left(b \| \mid b^{\prime}\right)}
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Variations on this argument lead to different rates of contraction

Open: Prove sharp rates that match Chernoff bounds

## BELLMAN UPDATES

How does this lead to better algorithms for planning?

## $\operatorname{Value}(x)=\underset{\text { actions a }}{\operatorname{Max}} E\left[\operatorname{Reward}(a)+\operatorname{Value}\left(x^{\prime}\right)\right]$ <br> current action/obs. sequence <br> new action/obs. sequence latent state sampled from current belief, stochastic transition based on chosen action

## BELLMAN UPDATES

Belief contraction allows us to truncate

## TRUNCATED BELLMAN UPDATES

Belief contraction allows us to truncate

latent state sampled from truncated belief, with uniform prior

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Belief contraction allows us to truncate

latent state sampled from truncated belief, with uniform prior

We only need a quasi-polynomial number of belief states

## OUTLINE

Part I: Tabular Markov Decision Processes

- Planning vs. Learning

Interlude: Can We Make RL Algorithmically Tractable?
Part II: Partial Observations and the Curse of History

- Beyond Worst-Case Analysis

Part III: Planning and Belief Contraction

- Polynomial vs. Exponential Rates

Part IV: Learning

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## WHAT ABOUT LEARNING?

In observable POMDPs:


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In observable POMDPs:
PSPACE Quasi-polynomial time planning

Can belief contraction be used for learning too?

## SAMPLE EFFICIENT LEARNING

Assumption 1: The POMDP is undercomplete, i.e. $|\mathcal{S}| \leq|\mathcal{O}|$ And moreover $\sigma_{\min }\left(\mathbb{O}_{h}\right) \geq \alpha$ for all $h$

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Theorem [Jin, Kakade, Krishnamurthy, Liu '20]: Given access to an optimistic planning oracle, there is an algorithm that uses

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i.e. given a constrained, non-convex set of POMDPs, find the maximum value achievable by any policy in the set

But optimism is very hard!

## Alternatively:

[Lin, Chung, Szepesvari, Jin '23] gave a framework based on optimistic maximum likelihood estimation

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Can we circumvent optimism?

## COMPUTATIONALLY EFFICIENT LEARNING

We show:
Theorem [Golowich, Moitra, Rohatgi '23]: There is an algorithm with running time and sample complexity

$$
(|\mathcal{O} \| \mathcal{A}|)^{C \log (H|\mathcal{S} \| \mathcal{O}| / \epsilon \gamma) / \gamma^{4}}
$$

that outputs an $\epsilon$-suboptimal policy in a $\gamma$-observable POMDP

## APPROXIMATION BY MDPS

Corollary: Any $\gamma$-observable POMDP P can be approximated by an MDP M with a quasi-polynomial number of states

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(1) P can be thought of as an MDP on belief states
(2) Construct M as follows:
states $=$ length $L$ sequences of actions/observations
transitions = shift in/out the newest/oldest actions/obs.
(3) States in $M$ can mapped to beliefs (using a uniform prior).

By belief contraction, $M$ and $P$ approximate each other

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Can we learn M efficiently?

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Can we learn M efficiently?

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How can we find a mixture of policies that visits all latent states?

## BARYCENTRIC SPANNERS

Definition: Given a set $\mathcal{X} \subseteq \mathbb{R}^{d}$, a $\lambda$-approximate barycentric spanner is a set $\mathcal{C} \subseteq \mathcal{X}$ of size $d$ such that every point in $\mathcal{X}$ can be expressed as a linear combination of points in $\mathcal{C}$ with coefficients in the range $[-\lambda, \lambda]$

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Theorem [Awerbuch, Kleinberg '04]: Given an oracle for optimizing linear functions over $\mathcal{X}$, there is a polynomial time algorithm for constructing a $\lambda$-approximate barycentric spanner with

$$
O\left(d^{2} \log _{\lambda} d\right)
$$

calls to the optimization oracle (assuming $\mathcal{X}$ is compact)

## POLICY COVERS

Now let
$\mathcal{\chi}=\begin{aligned} & \text { set of all distributions on observations } \\ & \text { at step } h \text { that can be obtained by a policy }\end{aligned}$

## POLICY COVERS

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Claim: By observability, if we can construct policies

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\pi_{1}, \pi_{2}, \ldots, \pi_{|\mathcal{O}|}
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whose induced distributions on observations at step $h$ are an approximate barycentric spanner

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whose induced distributions on observations at step $h$ are an approximate barycentric spanner, we must visit each latent state with nonnegligible probability

## ITERATIVE EXPLORATION

Our approach is:

$M_{h}$<br>MDP that<br>approximates<br>P up to step h

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MDP that
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## ITERATIVE EXPLORATION

Our approach is:


Reaches all
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## ITERATIVE EXPLORATION

Our approach is:


MDP that
Barycentric spanner approximates for observation
Pup to step h distributions at step h

Estimate next


Reaches all
latent states

## ITERATIVE EXPLORATION

Our approach is:


Estimate next layer of transitions

Reaches all
latent states

## ITERATIVE EXPLORATION

Our approach is:

| $\mathbf{M}_{\mathrm{h}+1}$ |  |
| :---: | :---: |
| MDP that <br> approximates <br> P up to step $\mathrm{h}+1$ | Barycentric spanner <br> for observation |
| $\mathbf{X}_{\mathrm{h}}$ |  |
| distributions at step h |  |



Estimate next


Reaches all
latent states

Without explorability, need more complex measure of progress

## LOOKING FORWARD

To get end-to-end algorithmic guarantees, we need to explore new assumptions and frameworks

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e.g. in [Golowich, Moitra '22], we took a learning-augmented algorithms approach:
"Can you improve Q-learning with advice?"

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"Can you improve Q-learning with advice?"

Takeaway: Improved regret bounds, where you only need to explore state-action pairs with substantially inaccurate predictions, even without knowing which ones are accurate in advance

## Summary:

- Modern RL is built on computationally intractable oracles. Are there end-to-end guarantees?
- Quasi-polynomial time algorithm for planning in observable POMDPs, no assumption on dynamics
- New framework for learning without optimism


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## Thanks! Any Questions?

