Algorithmic Aspects of Reinforcement Learning I

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February 24th, ITA+ALT Tutorial

TUTORIAL GOALS

(1) Overview of theoretical foundations of RL

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- (2) Gaps in our *algorithmic* understanding

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- (1) Overview of theoretical foundations of RL
- (2) Gaps in our *algorithmic* understanding
- (3) Deep dive into some success stories, emphasizing connections to other areas

Success stories of reinforcement learning:



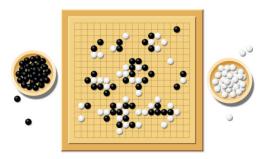
robotic manipulation

Success stories of reinforcement learning:



robotic manipulation

playing strategic games

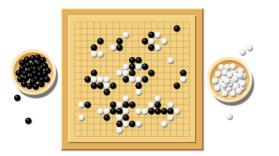


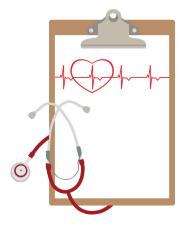
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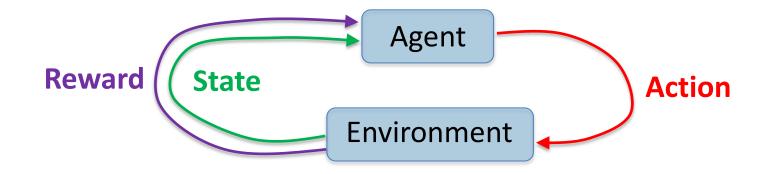
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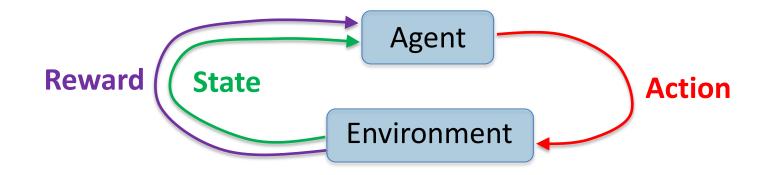


personalized treatment in medicine

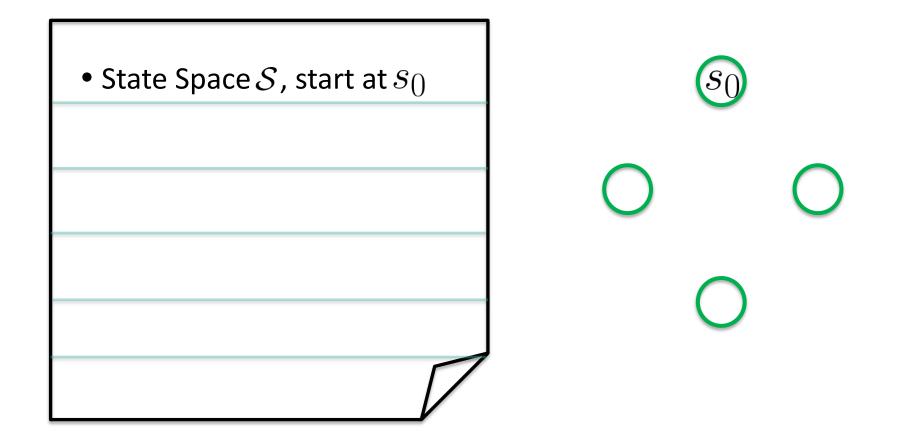
Goal: Agent learns by interacting with the environment

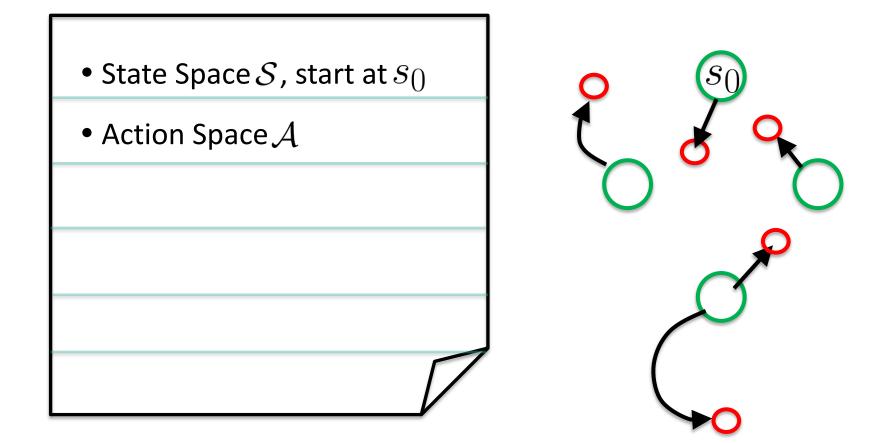


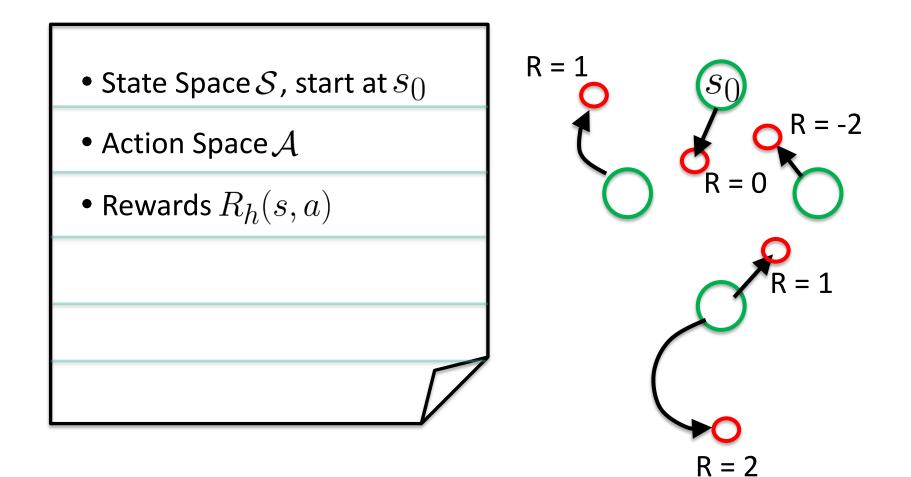
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The basic model is a Markov Decision Process

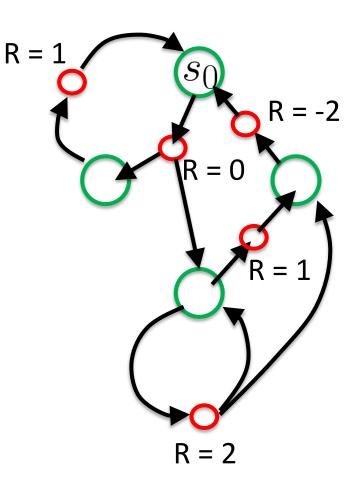






- State Space ${\cal S}$, start at s_0
- Action Space ${\cal A}$
- Rewards $R_h(s,a)$
- Transition Probabilities

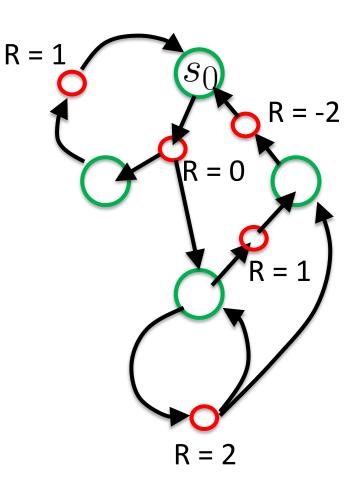
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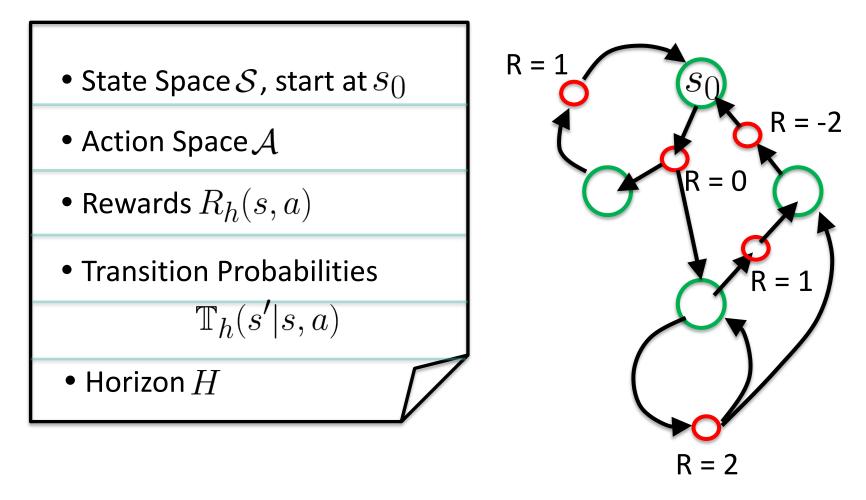


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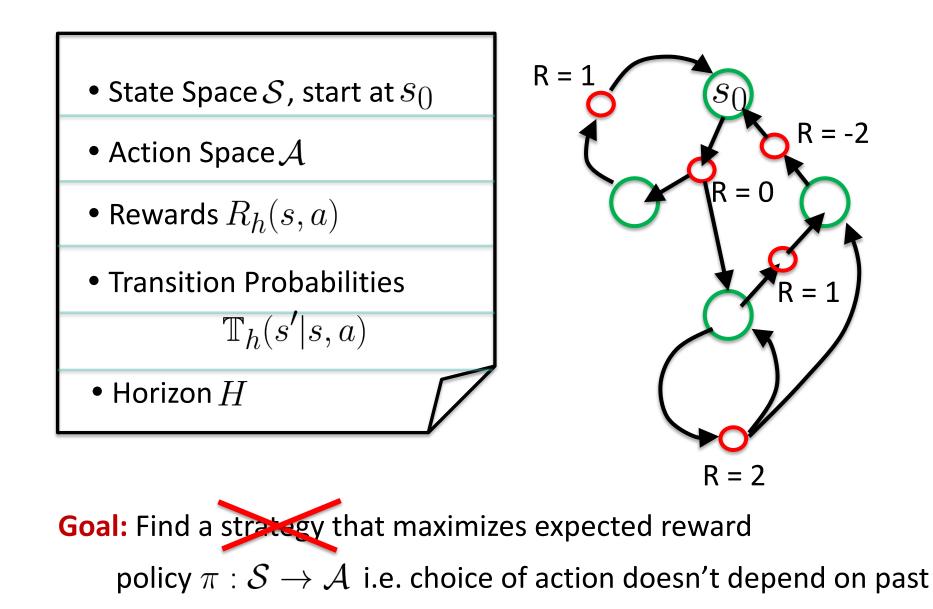
 $\mathbb{T}_h(s'|s,a)$

 $\bullet \operatorname{Horizon} H$





Goal: Find a strategy that maximizes expected reward



OUTLINE

Part I: Tabular Markov Decision Processes

• Planning vs. Learning

Interlude: Can We Make RL Algorithmically Tractable?

Part II: Partial Observations and the Curse of History

• Beyond Worst-Case Analysis

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Planning (computational)	
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(computational)	(statistical)
Given full description of the MDP, compute an optimal policy	Given budget of interactions with the environment, learn an optimal policy Uses planning as a subroutine

VALUE FUNCTIONS

Key: Keep track of how well you will do in the future

VALUE FUNCTIONS

Definition: The **value function** of a policy is

$$V_h^{\pi}(s) = \mathbb{E}[R_h(s_h, a_h) + \dots + R_H(s_H, a_H)|s_h = s]$$

i.e. it gives the expected future reward starting from state s at timestep h

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Bellman Optimality: An optimal policy π^* must satisfy

$$V_{h}^{\pi^{*}}(s) = \max_{a} R_{h}(s, a) + \mathbb{E}_{s'}[V_{h+1}^{\pi^{*}}(s')]$$

for every state s, i.e. value function must be consistent

VALUE ITERATION

This gives an efficient algorithm for planning:

```
Initialize V = 0 (assuming no rewards at step H)

Repeat until convergence

Scan through states, update any

violated V constraint
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Moreover can find the optimal policy from the V^{π^*} values

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EXPLORATION VS EXPLOITATION

First non-asymptotic result:

Theorem [Kearns, Singh '02]: There is an algorithm that has polynomial running time and sample complexity that outputs an ϵ -suboptimal policy in tabular MDPs

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- (2) Trade off playing the optimal policy in current model vs. discovering new states

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Tight regret bounds given by [Azar, Osband, Munos '17]

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random trajectory under π_{θ}

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How can we compute the gradient without full knowledge of the environment?

Policy Gradient Theorem: In fact

$$\nabla J(\theta) = \mathbb{E}_{\pi_{\theta}}[R(\tau)\nabla \log \pi_{\theta}(\tau)]$$
probability of trajectory under π_{θ}

Thus we can approximate the gradient through samples

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Many challenges both in theory and practice, e.g. delayed feedback can cause gradients to be extremely small

e.g. see [Agarwal, Kakade, Lee, Mahajan '19]

MAIN PROBLEMS

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MAIN PROBLEMS

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an optimal policy	learn an optimal policy
e.g. value iteration,	e.g. model based,
policy iteration,	Q-learning, actor-critic
linear programming	policy gradient

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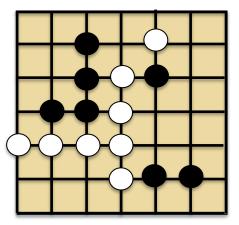
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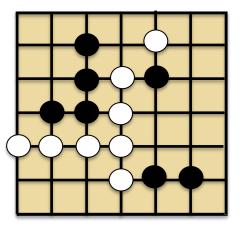
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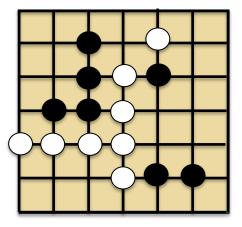
function approximation, block MDPs, etc

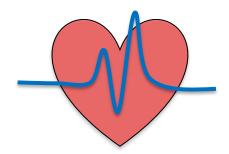


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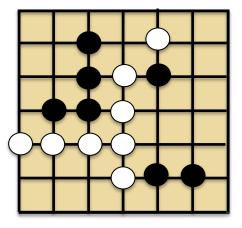


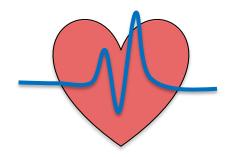
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Partially observable MDPs (POMDPs)

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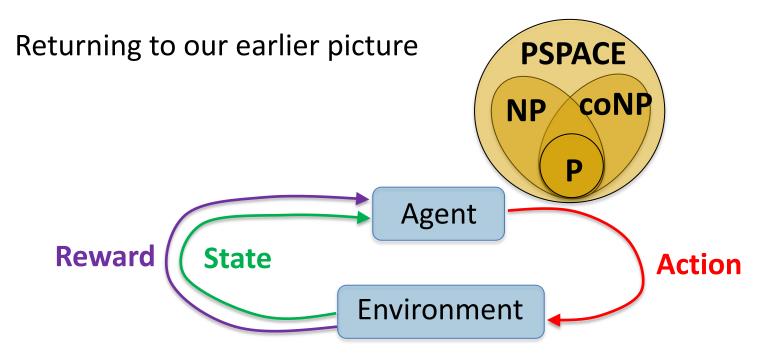
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Existing theory is built around (1) and (2) but what do we miss out on by ignoring (3)?

WHAT ABOUT COMPUTATIONAL COMPLEXITY?



Are there computationally efficient algorithms with strong end-to-end provable guarantees?

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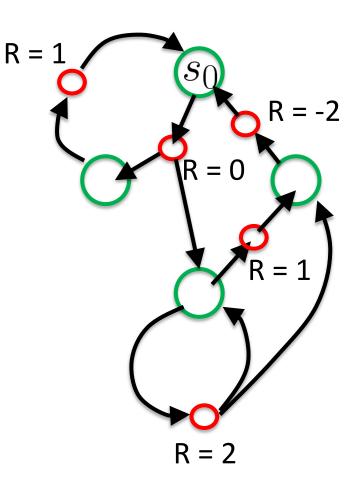
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MARKOV DECISION PROCESSES

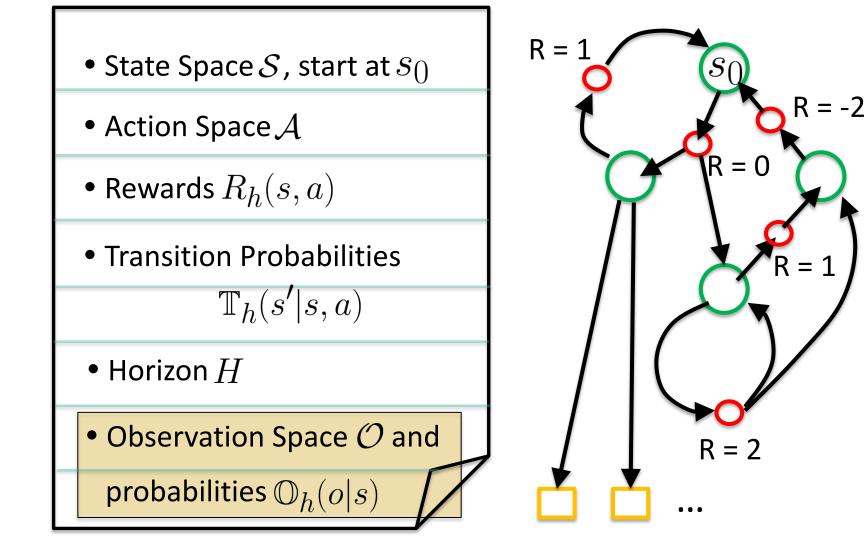
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• Horizon H



PARTIALLY OBSERVABLE MDPS (POMDPS)



PLANNING IS HARD

Classic lower bound:

Theorem [Papadimitriou, Tsitsiklis]: Optimal planning in a POMDP is PSPACE hard

MDPs	POMDPs
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MDPs	POMDPs
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Can you succinctly represent an optimal policy?

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Natural approaches use exponential space $(|\mathcal{A}||\mathcal{O}|)^H$ or $C^{|\mathcal{S}|}$

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Even worse news:

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Why should real-world POMDPs have succinct descriptions of good policies?

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Could this enable tractable planning/learning?

Definition: We say the POMDP is γ -observable if for all h and all distributions b,b' on states we have

$$\|\mathbb{O}_h b - \mathbb{O}_h b'\|_1 \ge \gamma \|b - b'\|_1$$

i.e. well-separated distributions on states lead to well-separated distributions on observations

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BEYOND WORST-CASE ANALYSIS

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Key Point: No assumption on transition dynamics like e.g. deterministic transitions or mixing (under every possible policy)

PLANNING VIA STABILITY

There is a quasi-polynomial time algorithm for planning under observability:

Theorem [Golowich, Moitra, Rohatgi '23]: Given description of a γ -observable POMDP there is an algorithm running in time $H(|\mathcal{O}||\mathcal{A}|)^{C\log(|\mathcal{S}|H/\epsilon)/\gamma^4}$

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Parallels well-known stability results for Kalman filtering

LOWER BOUNDS

Moreover these results are tight

Theorem [Golowich, Moitra, Rohatgi '23]: Under the Exponential Time Hypothesis, there is no algorithm running in time $(|\mathcal{S}||\mathcal{A}|H|\mathcal{O}|)^{o(\log(|\mathcal{S}||\mathcal{A}|H|\mathcal{O}|/\epsilon)/\gamma)}$

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It's hard even in the lossy case, where you observe the state with probability γ independently at each step

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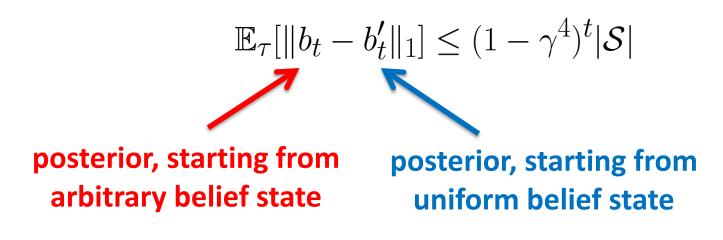
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BELIEF CONTRACTION

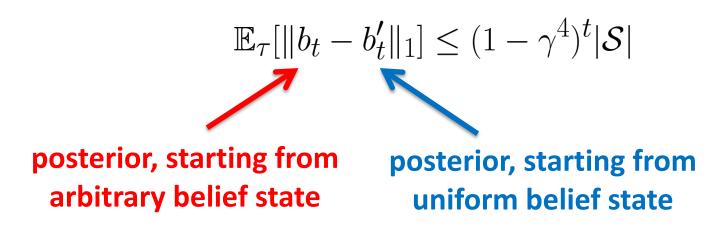
Theorem: Fix any γ -observable POMDP and policy π . Then



where au is the trajectory from the POMDP by playing π

BELIEF CONTRACTION

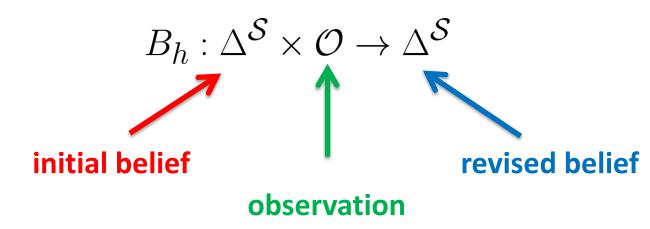
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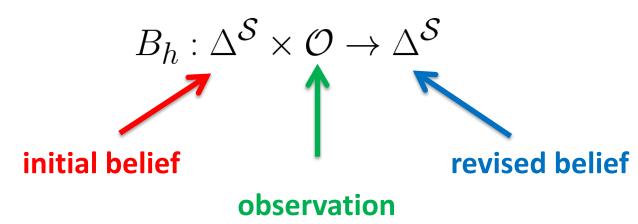
Thus, we could ignore all but the most recent history

Definition: The **Bayes operator**, given an observation



updates the posterior on states

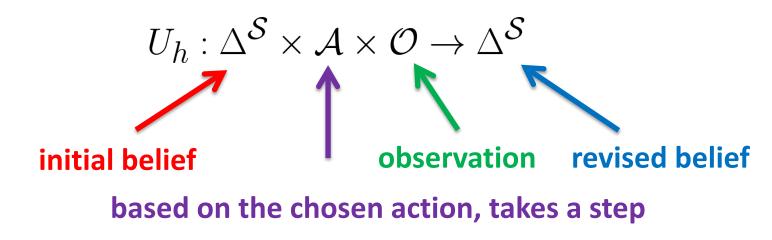
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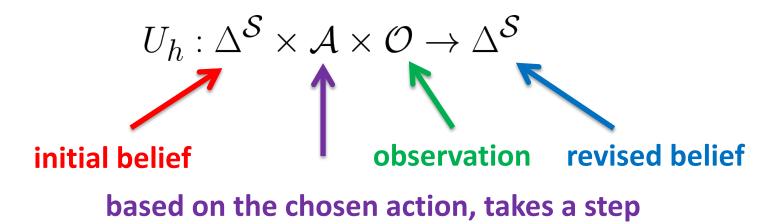
$$B_h(b,y)(x) = \frac{\mathbb{O}_h(y|x)b(x)}{\sum_{z \in \mathcal{SO}_h(y|z)b(z)}}$$

Definition: And the **update operator**, given both an action and observation



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TOWARDS A WEAKER BOUND

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Do we make progress in expectation?

 $\mathbb{E}[KL(B_h(b,y)||B_h(b',y))] = KL(b||b') - KL(\mathbb{O}_h b||\mathbb{O}_h b')$

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Proof: $KL(b||b') = KL(P_X||Q_X)$

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...using the chain rule, and the fact that $Q_{Y|X=x} = P_{Y|X=x}$

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... this time using the chain rule in opposite order

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Does this imply fast enough convergence?

Using Pinsker's inequality (1) and observability (2), we have $\mathbb{E}[KL(B_{h}(b,y)||B_{h}(b',y))] \stackrel{(1)}{\leq} KL(b||b') - \frac{1}{2}||\mathbb{O}_{h}b - \mathbb{O}_{h}b'||_{1}^{2}$

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Using reverse Pinsker's inequality, we get

 $\mathbb{E}[KL(B_h(b,y)||B_h(b',y))] \le KL(b||b') - c\gamma^2 KL(b||b')^2$

provided that $\|b/b'\|_{\infty}$ is bounded

Theorem [Even-Dar et al.]: Fix any γ -observable POMDP and policy π . Then we have

 $\mathbb{E}_{\tau}[KL(b_t||b_t')] \le \epsilon$

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Inverse polynomial, rather than exponential convergence :(

Unfortunately there are cases where progress can be slow, but...

We show that either

(1) A stronger reverse Pinsker holds, i.e.

 $\mathbb{E}[K_{u} [K_{h}(b, y) || B_{h}(b', y))] \le KL(b || b') - \frac{\gamma^{2}}{32} \min(KL(b || b'), 1)$

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$$\mathbb{E}[KL_{y \sim \mathbb{O}_{h}b}(B_{h}(b, y)||B_{h}(b', y))] \le KL(b||b') - \frac{\gamma^{2}}{32}\min(KL(b||b'), 1)$$

or instead

(2) Progress is **anti-concentrated**, i.e. for some event $\mathcal{E} \subset \mathcal{O}$

$$\mathbb{E}\left[\left(KL(B_{h}(b,y)||B_{h}(b',y)) - KL(b||b')\right)\mathbf{1}_{y\in\mathcal{E}}\right] \leq -\frac{\gamma}{8}KL(b||b')$$

As a result, we get:

Corollary: For any γ -observable POMDP

$$\mathbb{E}\left[\sqrt{KL(B_{h}(b, y)||B_{h}(b', y))}\right] \leq \left(1 - \Omega\left(\frac{\gamma^{2}}{\max(1, KL(b||b')))}\right) \sqrt{KL(b||b')}$$

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Variations on this argument lead to different rates of contraction

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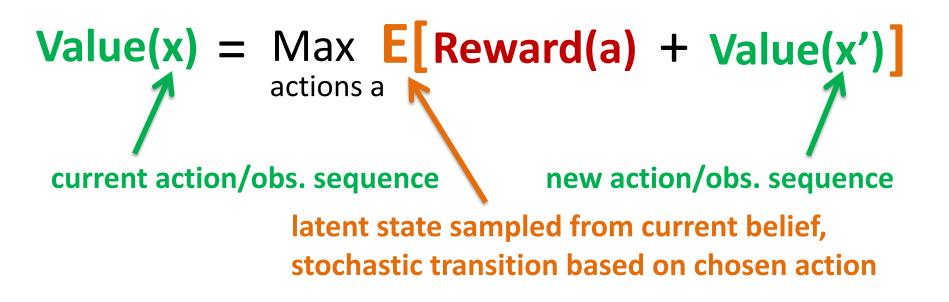
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Variations on this argument lead to different rates of contraction

Open: Prove sharp rates that match Chernoff bounds

BELLMAN UPDATES

How does this lead to better algorithms for planning?



BELLMAN UPDATES

Belief contraction allows us to truncate

TRUNCATED BELLMAN UPDATES

Belief contraction allows us to truncate



latent state sampled from truncated belief, with uniform prior

TRUNCATED BELLMAN UPDATES

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We only need a quasi-polynomial number of belief states

OUTLINE

Part I: Tabular Markov Decision Processes

• Planning vs. Learning

Interlude: Can We Make RL Algorithmically Tractable?

Part II: Partial Observations and the Curse of History

• Beyond Worst-Case Analysis

Part III: Planning and Belief Contraction

• Polynomial vs. Exponential Rates

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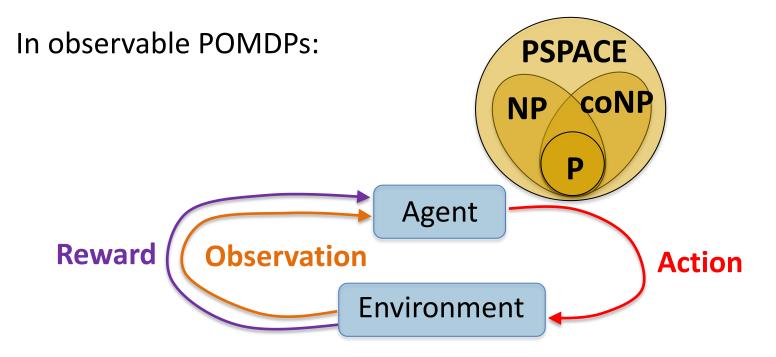
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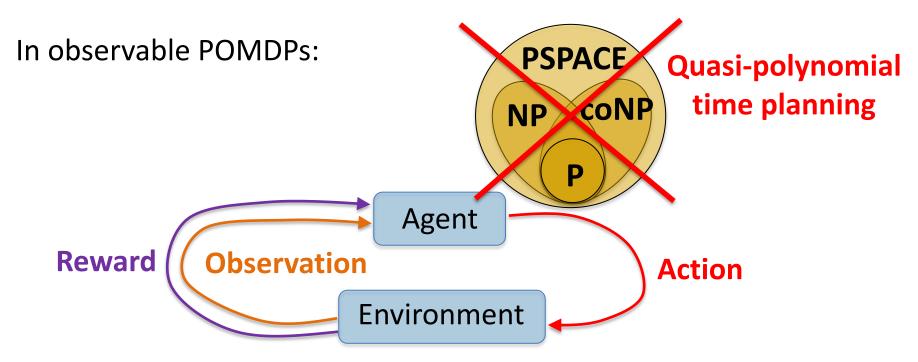
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Part IV: Learning

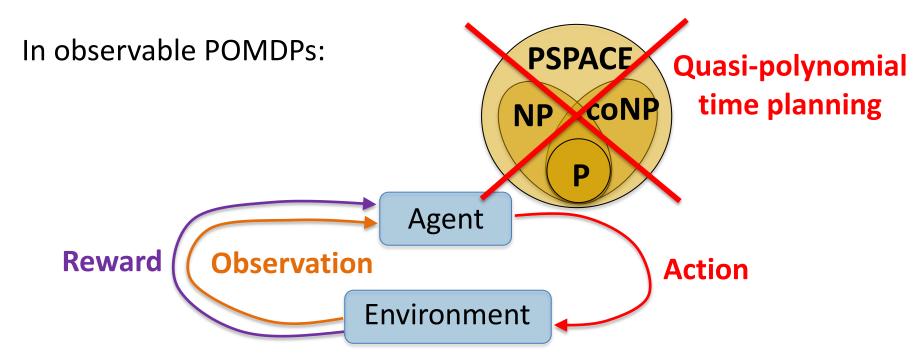
WHAT ABOUT LEARNING?



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Can belief contraction be used for learning too?

Assumption 1: The POMDP is undercomplete, i.e. $|\mathcal{S}| \leq |\mathcal{O}|$ And moreover $\sigma_{min}(\mathbb{O}_h) \geq \alpha$ for all h

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Theorem [Jin, Kakade, Krishnamurthy, Liu '20]: Given access to an **optimistic planning oracle**, there is an algorithm that uses $poly(|\mathcal{S}|, |\mathcal{A}|, H, |\mathcal{O}|, 1/\alpha)$

samples and finds an ϵ -suboptimal policy under Assumption 1

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But optimism is very hard!

Alternatively:

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Can we circumvent optimism?

COMPUTATIONALLY EFFICIENT LEARNING

We show:

Theorem [Golowich, Moitra, Rohatgi '23]: There is an algorithm with running time and sample complexity

$$(|\mathcal{O}||\mathcal{A}|)^{C\log(H|\mathcal{S}||\mathcal{O}|/\epsilon\gamma)/\gamma^4}$$

that outputs an ϵ -suboptimal policy in a γ -observable POMDP

Corollary: Any γ -observable POMDP P can be approximated by an MDP M with a quasi-polynomial number of states

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(3) States in M can mapped to beliefs (using a uniform prior).

By belief contraction, M and P approximate each other

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How can we find a mixture of policies that visits all latent states?

BARYCENTRIC SPANNERS

Definition: Given a set $\mathcal{X} \subseteq \mathbb{R}^d$, a λ -approximate barycentric spanner is a set $\mathcal{C} \subseteq \mathcal{X}$ of size d such that every point in \mathcal{X} can be expressed as a linear combination of points in \mathcal{C} with coefficients in the range $[-\lambda, \lambda]$

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Theorem [Awerbuch, Kleinberg '04]: Given an oracle for optimizing linear functions over \mathcal{X} , there is a polynomial time algorithm for constructing a λ -approximate barycentric spanner with

$$O(d^2 \log_\lambda d)$$

calls to the optimization oracle (assuming ${\mathcal X}$ is compact)

POLICY COVERS

Now let

$\mathcal{X}=rac{\mathrm{set}\ \mathrm{of}\ \mathrm{all}\ \mathrm{distributions}\ \mathrm{on}\ \mathrm{observations}}{\mathrm{at}\ \mathrm{step}\ \mathrm{h}\ \mathrm{that}\ \mathrm{can}\ \mathrm{be}\ \mathrm{obtained}\ \mathrm{by}\ \mathrm{a}\ \mathrm{policy}}$

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Claim: By observability, if we can construct policies

$$\pi_1, \pi_2, \ldots, \pi_{|\mathcal{O}|}$$

whose induced distributions on observations at step h are an approximate barycentric spanner

POLICY COVERS

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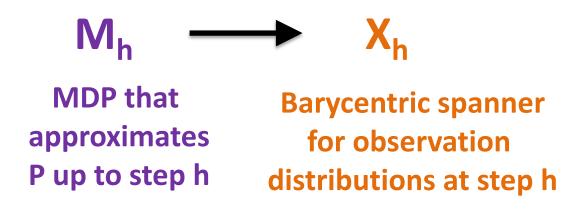
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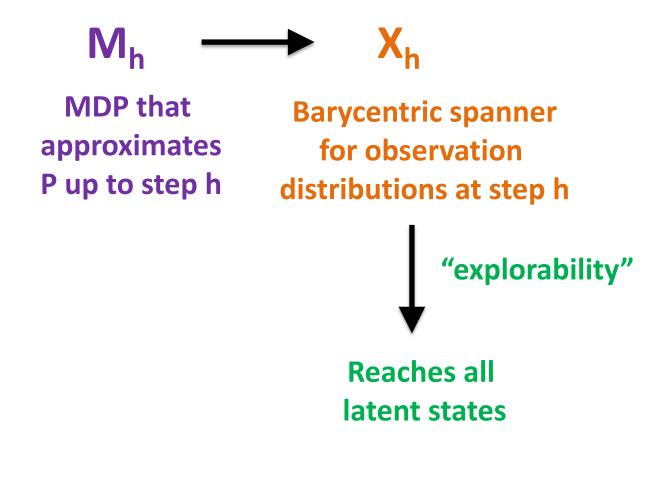
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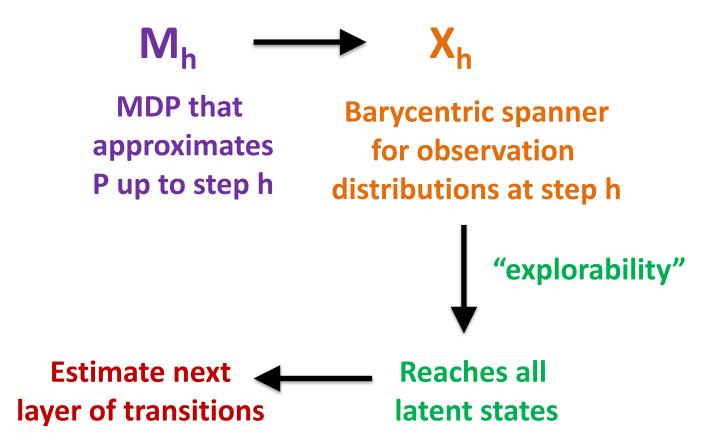
whose induced distributions on observations at step h are an approximate barycentric spanner, we must visit each latent state with nonnegligible probability

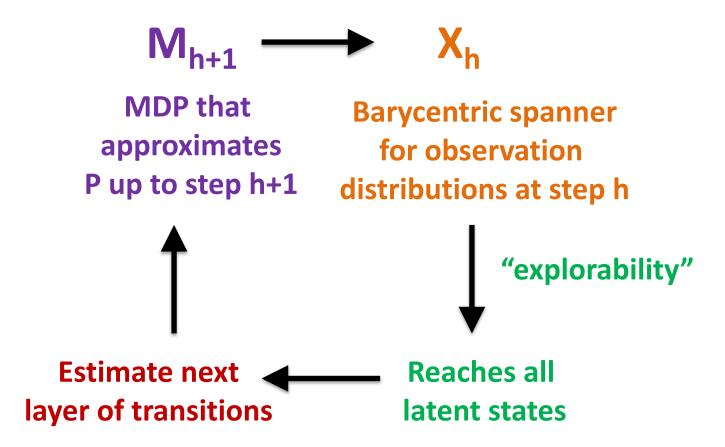
Our approach is:

M_h MDP that approximates P up to step h

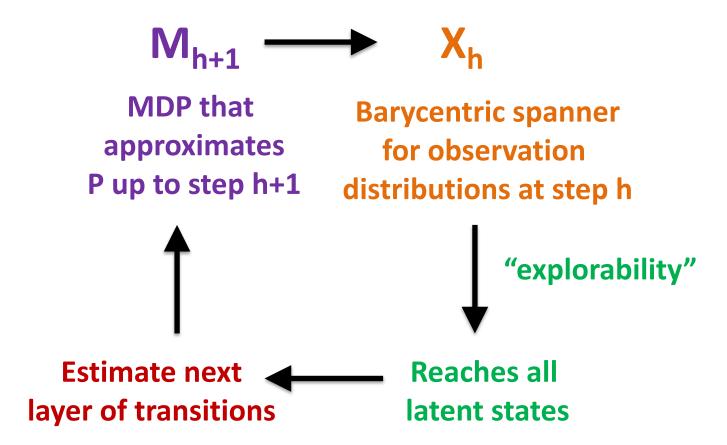








Our approach is:



Without explorability, need more complex measure of progress

LOOKING FORWARD

To get end-to-end algorithmic guarantees, we need to explore new assumptions and frameworks

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"Can you improve Q-learning with advice?"

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Takeaway: Improved regret bounds, where you only need to explore state-action pairs with substantially inaccurate predictions, even without knowing which ones are accurate in advance

Summary:

- Modern RL is built on computationally intractable oracles. Are there end-to-end guarantees?
- Quasi-polynomial time algorithm for planning in observable POMDPs, no assumption on dynamics
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Thanks! Any Questions?