

## The Harshman-Jennrich Algorithm

this algorithm has been rediscovered many times, originates in psychometrics

Setup: We are given  $T$ , assumed to be of the form

$$T = \sum_{l=1}^r u^{(l)} \otimes v^{(l)} \otimes w^{(l)}$$

definition. We say two sets of factors

$$\{ (u^{(l)}, v^{(l)}, w^{(l)}) \} \text{ and } \{ (\hat{u}^{(l)}, \hat{v}^{(l)}, \hat{w}^{(l)}) \}$$

are equivalent if there is a permutation  $\pi$  s.t.

$$u^{(l)} \otimes v^{(l)} \otimes w^{(l)} = \hat{u}^{\pi(l)} \otimes \hat{v}^{\pi(l)} \otimes \hat{w}^{\pi(l)}$$

i.e. they produce equivalent low rank decompositions

Main Question When are the factors of  $T$  determined up to equivalence?

Theorem [Harshman, Jennrich] suppose the following conditions hold

- (1) The vectors  $\{u^{(l)}\}$  are linearly indep.
- (2) same for  $\{v^{(l)}\}$
- (3) every pair of vectors in  $\{w^{(l)}\}$  are linearly indep.

Then the factors are uniquely determined up to equivalence, and there is a polynomial time algorithm to find them

The algorithm is simple

- Choose  $a, b \in S^{n_3}$  uniformly at random, set

$$T_a = \sum_{i=1}^{n_3} a_i T(\cdot, \cdot, i) ; T_b = \sum_{c=1}^{n_3} b_c T(\cdot, \cdot, c)$$

- Compute the eigen decompositions of

$$T_a (T_b)^T \text{ and } (T_a^T T_b)^T$$

Let  $\hat{u}$  and  $\hat{v}$  be eigenvectors with non-zero eigenvalue

Pair up  $\hat{u}^{(i)}$  and  $\hat{v}^{(j)}$  iff their eigenvalues are reciprocals

- Solve for  $\hat{w}^{(i)}$  in the linear system

$$T = \sum_{c=1}^r \hat{u}^{(i)} \otimes \hat{v}^{(i)} \otimes \hat{w}^{(i)}$$

The analysis follows by tracking the structure of  $T$  through the algorithm, and using standard uniqueness facts about eigen decomp.

Let  $D_a = \text{Diag}(a^T w^{(i)})$ ;  $D_b = \text{Diag}(b^T w^{(i)})$

Lemma: we have that

$$T_a = U D_a V^T \text{ and } T_b = U D_b V^T$$

Proof: Since the operation of computing  $T_a$  from  $T$  is linear, we can do it just for a rank one term:

If  $T = u \otimes v \otimes w$  then  $T_a = (a^T w) u \otimes v$

Thus, in general, we have

$$T_a = \sum_{i=1}^r (a^T w^{(i)}) u^{(i)} \otimes v^{(i)} = U D_a V^T \quad \square$$

For simplicity, let's assume  $T_a$  and  $T_b$  are invertible. Then

$$\begin{aligned} T_a T_b^{-1} & \stackrel{\text{lemma}}{=} U D_a V^T (V^T)^{-1} D_b^{-1} U^{-1} \\ & = U D_a D_b^{-1} U^{-1} \end{aligned}$$

From property (3), almost surely the diag. entries of

$$D_a D_b^{-1}$$

will be distinct. Thus  $T_a T_b^{-1}$  has distinct eigenvalues  $\Rightarrow$  its eigen decomposition

is unique  $\Rightarrow$  we can find  $U$ , up to a permutation/  
of its columns rescaling

Similarly we have

$$\begin{aligned}(T_a^{-1} T_b)^T &= (V^T)^{-1} D_a^{-1} U^{-1} U D_b V^T)^T \\ &= ((V^T)^{-1} D_a^{-1} D_b V^T)^T \\ &= V D_a^{-1} D_b^{-1} V^{-1}\end{aligned}$$

Thus we can determine  $V$ , up to a permutation/rescaling of its columns, again from uniqueness of the eigendecomp.

Moreover pairing succeeds, again because the diagonal entries of  $D_a^{-1} D_b$  are distinct

Finally, we show:

Lemma: The matrices  $u^{(i)} \otimes v^{(i)}$  are linearly independent

Proof: Suppose not. Then  $\exists \alpha_i$ 's s.t.

$$\sum_{i=1}^r \alpha_i u^{(i)} \otimes v^{(i)} = 0 \quad \text{Suppose WLOG } \alpha_1 \neq 0$$

By condition (1), we know  $\exists x$  s.t.

$$x^T u^{(1)} \neq 0, \quad x^T u^{(i)} = 0 \quad \forall i \neq 1$$

Now using the identity above, we get

$$(\alpha, a^T u^{(i)}) v^{(i)} = 0 \Rightarrow v^{(i)} = 0$$

which contradicts condition (2).  $\square$

Why was this algorithm forgotten?

Psychometrics generally cared about uniqueness, and there are better non-algorithmic uniqueness theorems known

Now returning to factor analysis

Given:  $T = \sum_{i=1}^r u^{(i)} \otimes v^{(i)} \otimes w^{(i)}$ , when are the

factors determined up to equivalence?

Harshman-Jennrich: when  $\{u^{(i)}\}$  and  $\{v^{(i)}\}$  are linearly indep., and no pair in  $\{w^{(i)}\}$  are scalar multiples of each other