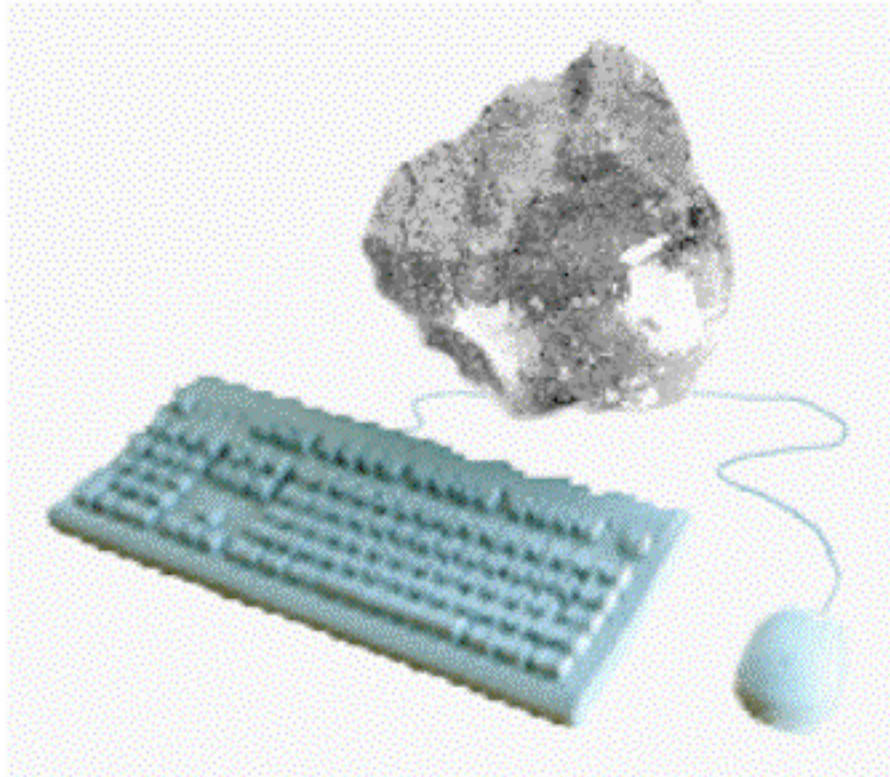




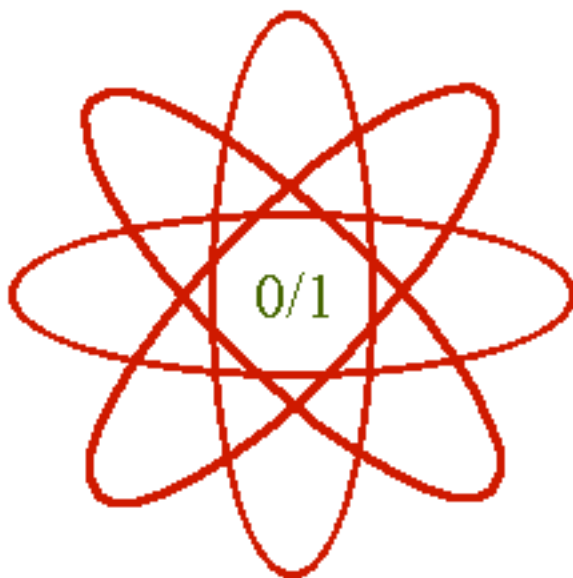
Computing Beyond Silicon Summer School

Physics becomes the computer



Norm Margolus

Physics becomes the computer



Emulating Physics

- » *Finite-state, locality, invertibility, and conservation laws*

Physical Worlds

- » *Incorporating comp-universality at small and large scales*

Spatial Computers

- » *Architectures and algorithms for large-scale spatial computations*

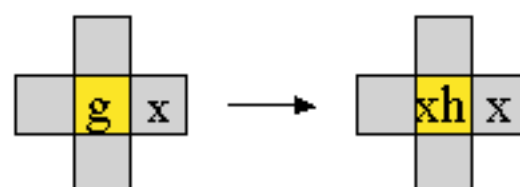
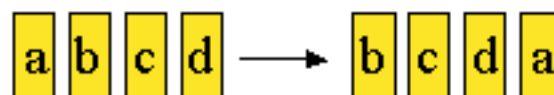
Nature as Computer

- » *Physical concepts enter CS and computer concepts enter Physics*

Review: CA's with conservations

To make reversibility and other conservations manifest, we employ a multi-step update, in each step of which either

1. *The data are rearranged without any interaction, or*
2. *The data are partitioned into disjoint groups of bits that change as a unit. Data that affect more than one such group don't change.*



We would like to be able to study large-scale computations with this kind of conservation-friendly format, which allows them to map efficiently onto microscopic physics, and also allow them to have interesting macroscopic behavior.

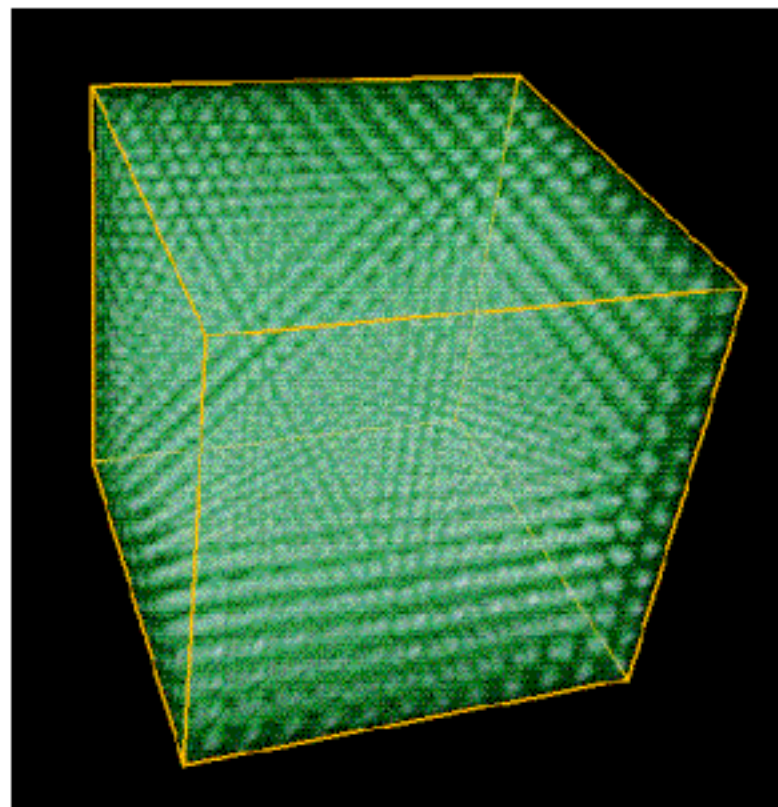
Spatial Computers



(ISSN: 4/5)

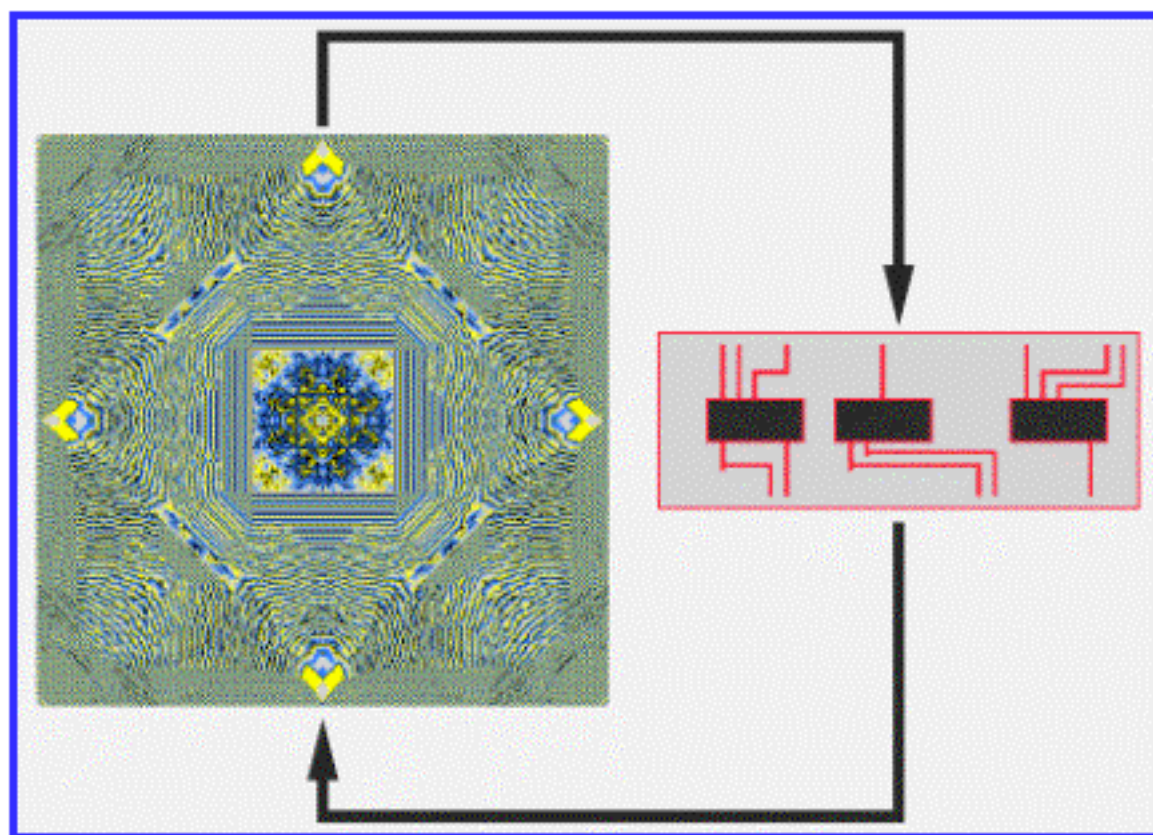
Spatial Computers

- Finite-state computations on lattices are interesting
- Crystalline computers will one day have highest performance
- Today's ordinary computers are terrible at these computations



We may one day compute with regular crystals of atoms.

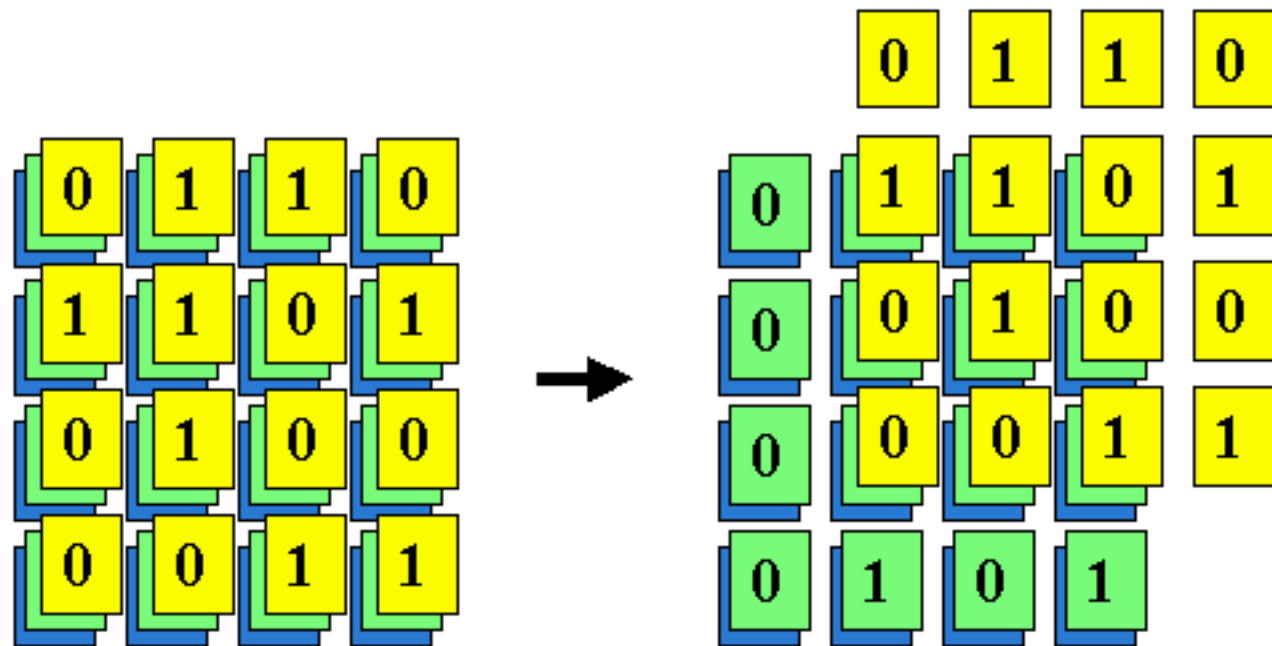
CAMs 1 to 6



- Tom Toffoli and I wanted to see CA-TV
- Stimulated by new rev rules
- Rule = a few TTL chips!
- Later, rule = SRAM LUT
- LGA's needed more "serious" machine!

CAM-8 Programming Model

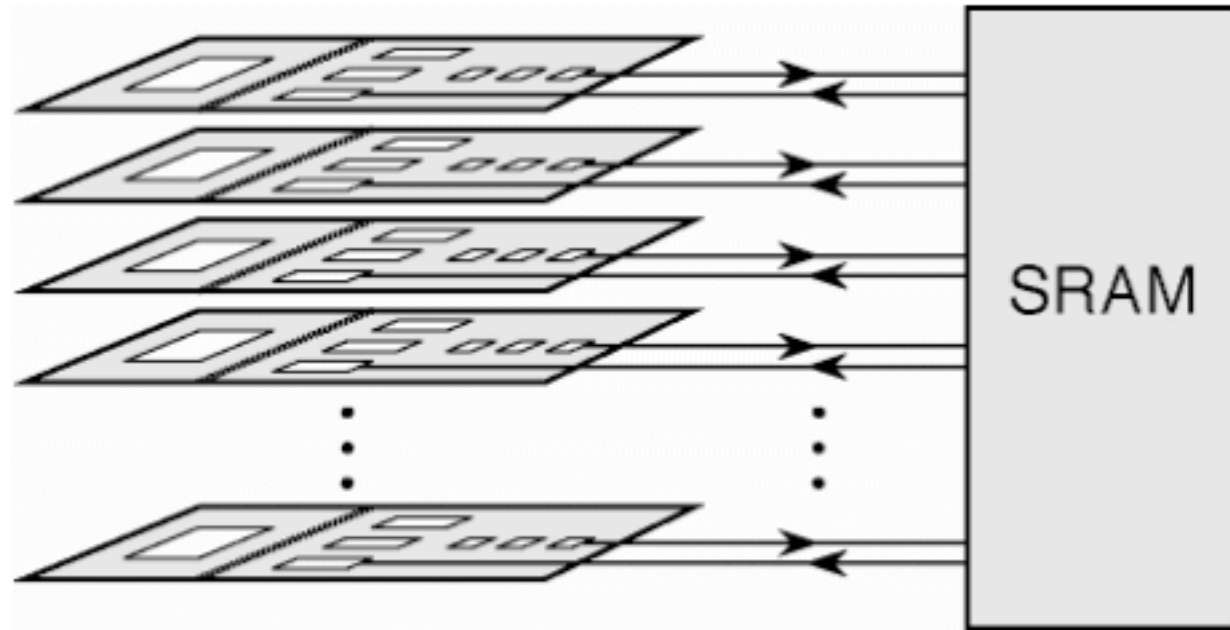
*Data
movement
step:*



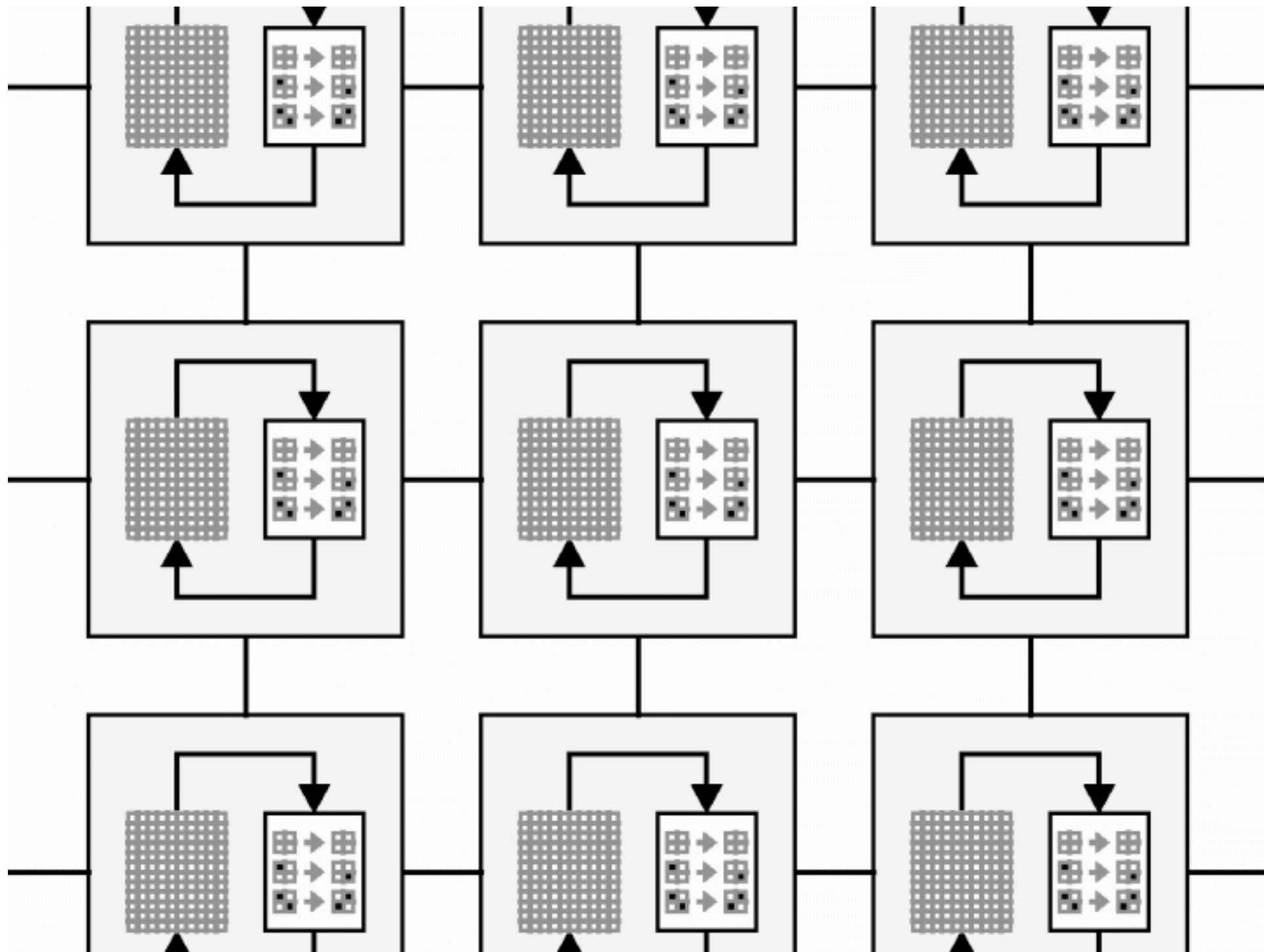
*Site
update
step:*



CAM-8 node

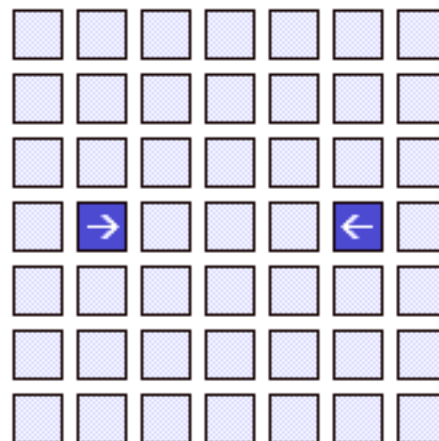


- maps directly onto programming model (*bit fields*)
- 16-bit lattice sites (use internal dimension to get more)
- each node handles a chunk of a larger space



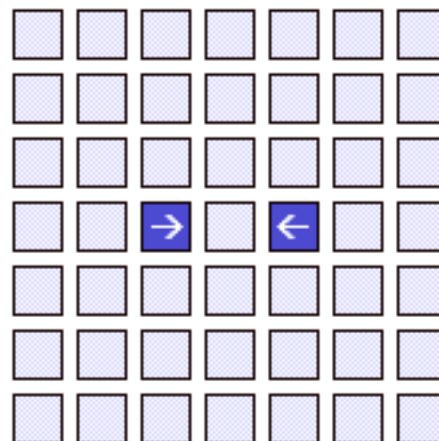
Lattice example

```
/** Four direction lattice gas */  
  
2D-square-lattice {512 512}  
create-bit-planes {N S E W}  
define-step  
  shift-planes {(N,0,-1) (S,0,1)  
              (E,1,0) (W,-1,0)}  
  for-each-site  
    if {N==S && E==W}  
      {xchng(N,E); xchng(S,W)}  
    endif  
  end-site  
end-step
```



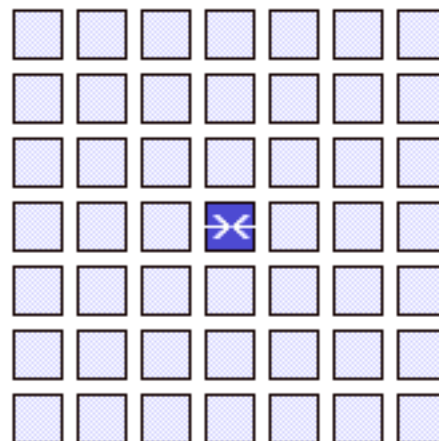
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```



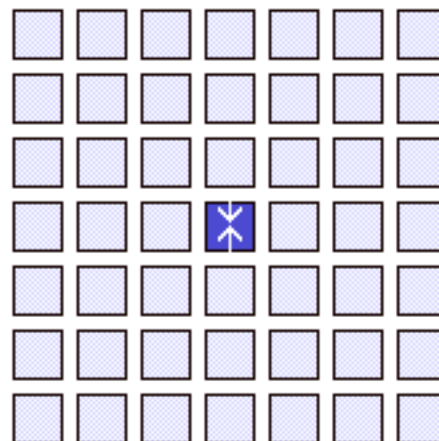
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end-step
```



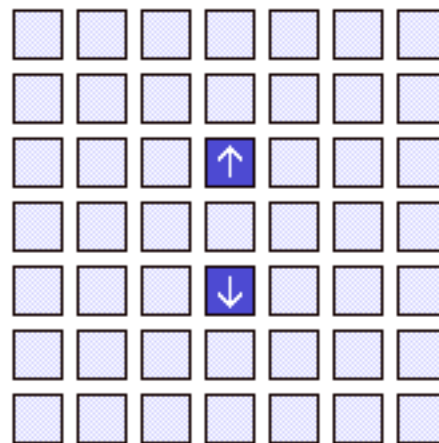
Lattice example

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```



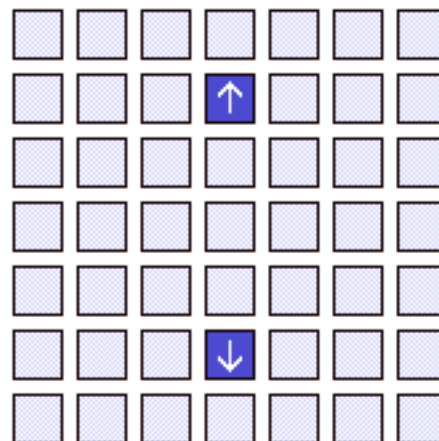
Lattice example

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end-step
```



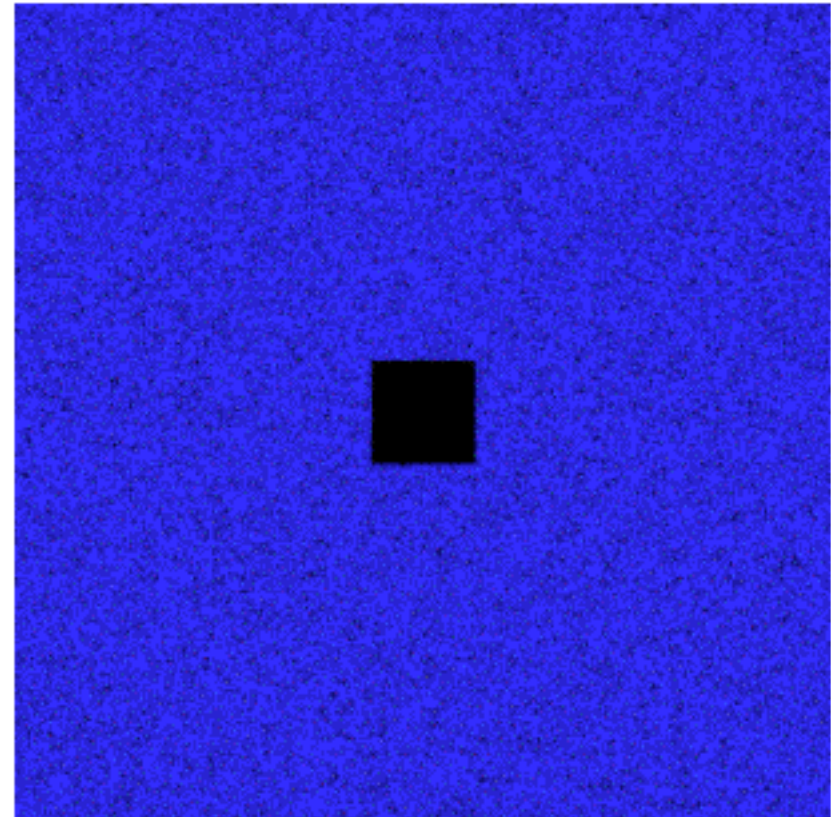
Lattice example

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```

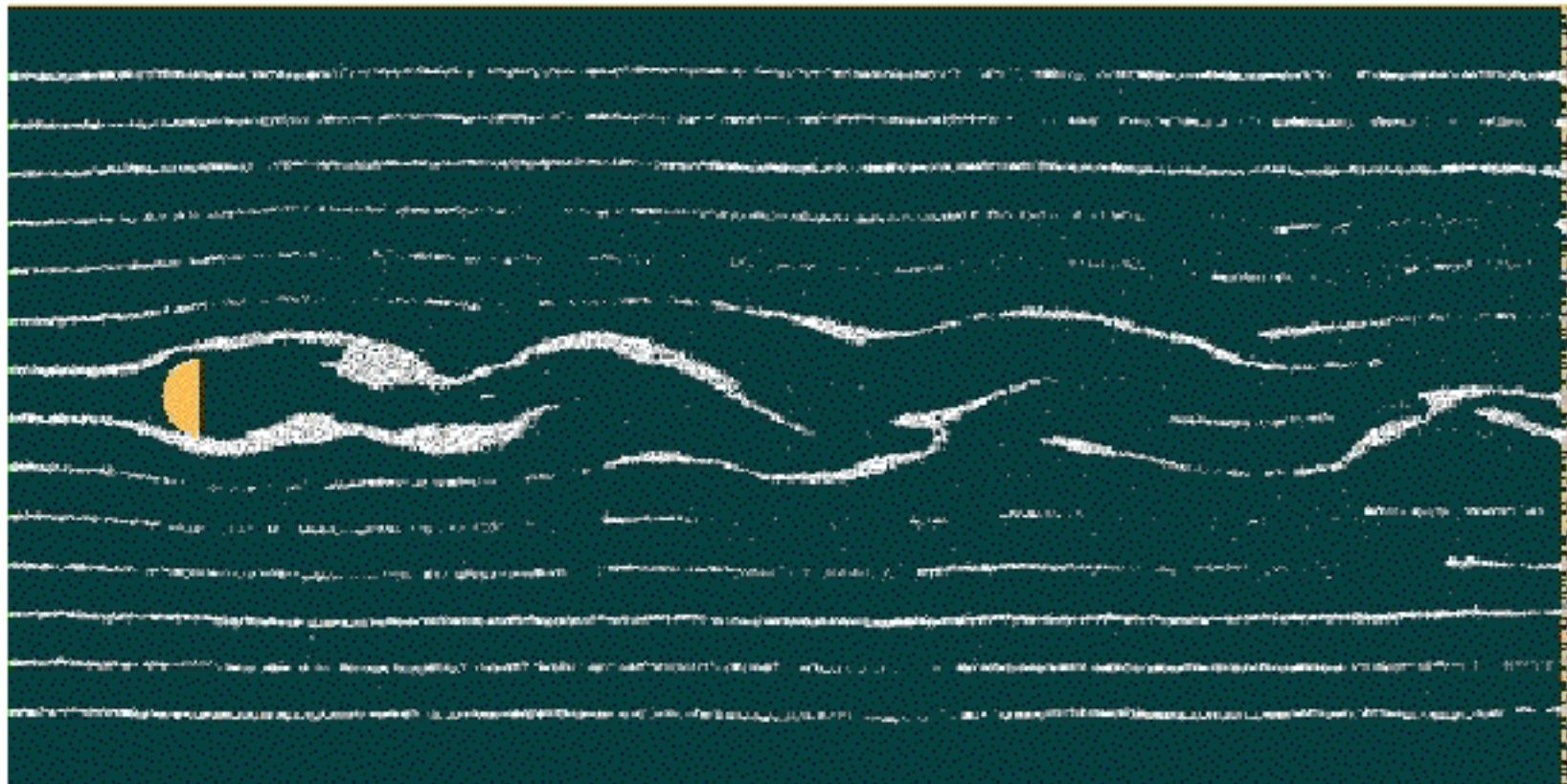


Lattice example

```
/** Four direction lattice gas */  
  
2D-square-lattice {512 512}  
create-bit-planes {N S E W}  
define-step  
  shift-planes {(N,0,-1) (S,0,1)  
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```

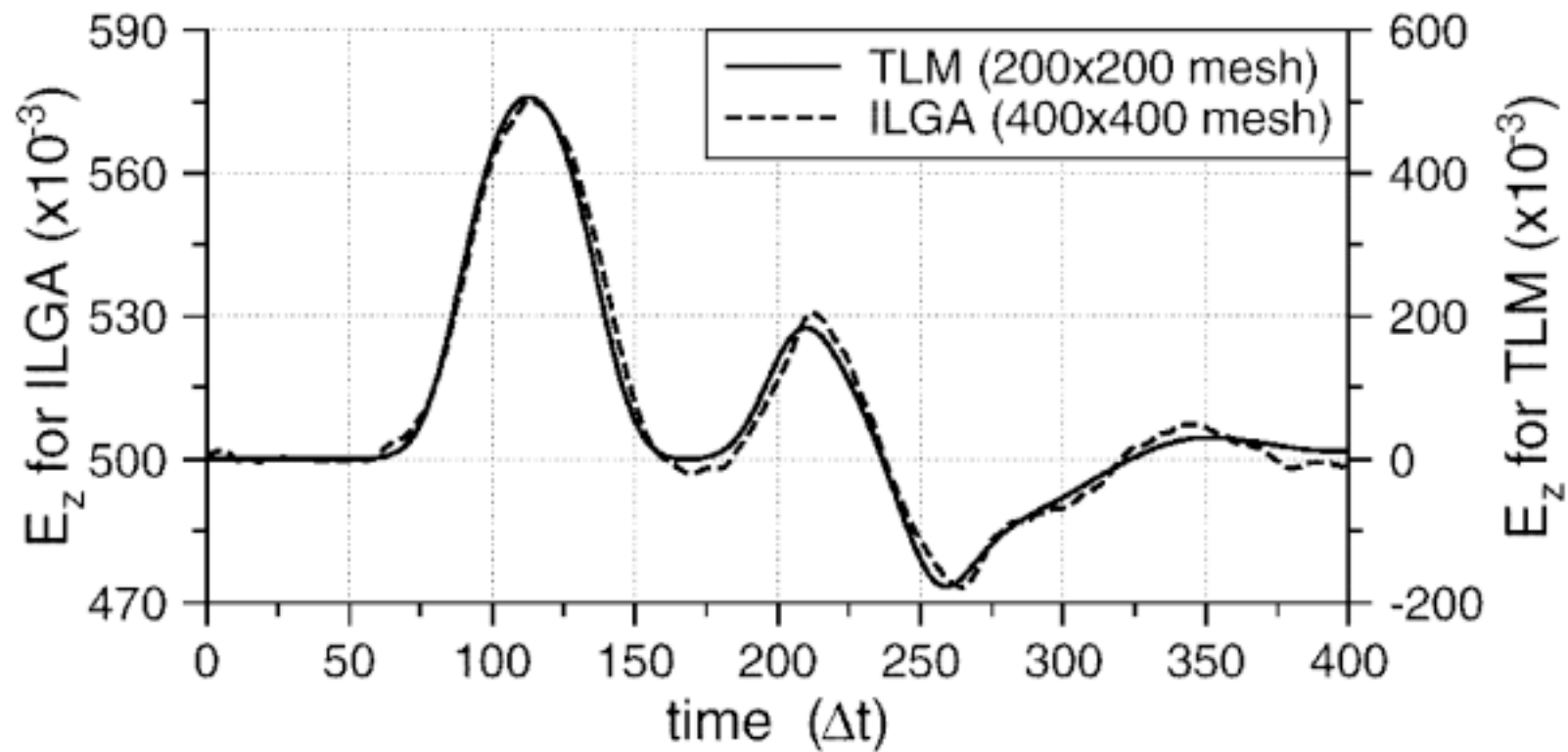


CAM-8: discrete hydrodynamics

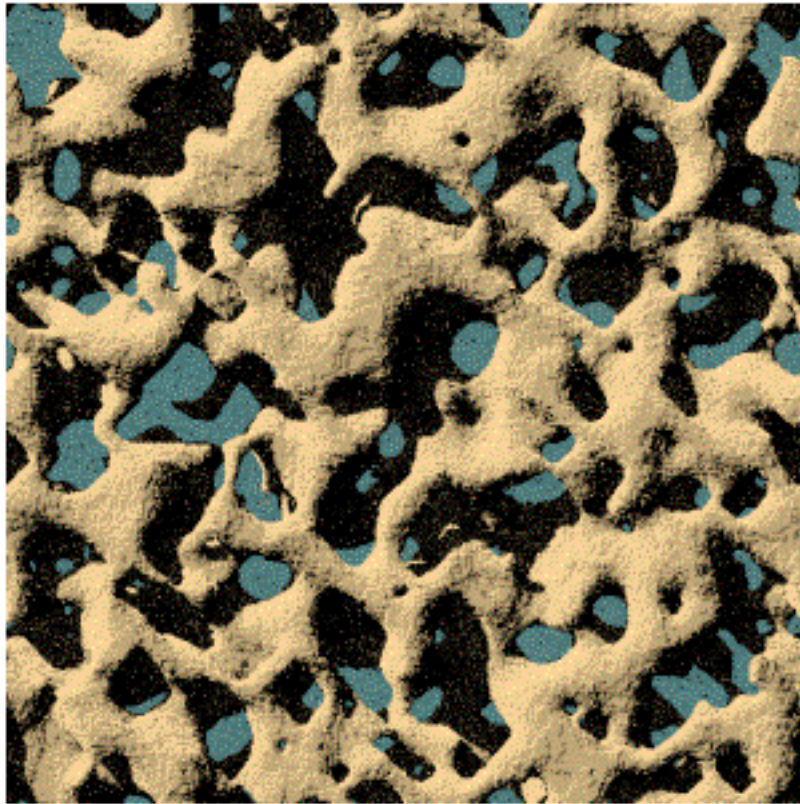


Six direction LGA flow past a half-cylinder. System is 2Kx1K.

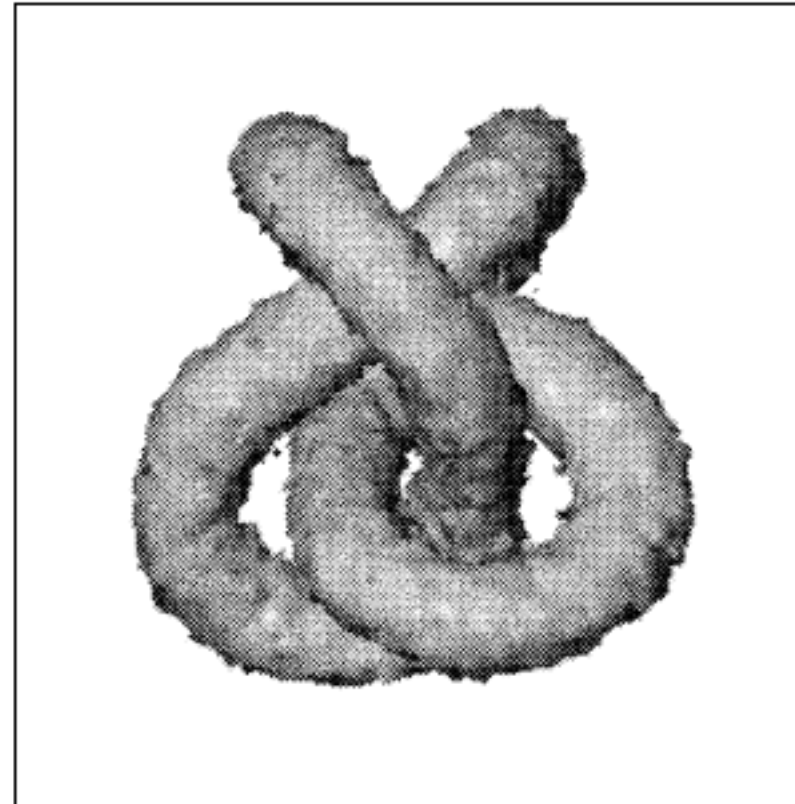
CAM-8: EM scattering



CAM-8: materials simulation

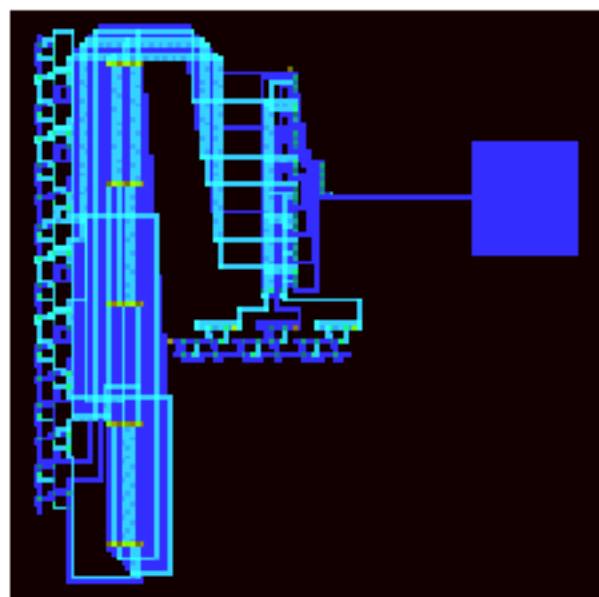


512x512x64 spin system, 6up/sec w rendering



256x256x256 reaction-diffusion (Kapral)

CAM-8: logic simulation

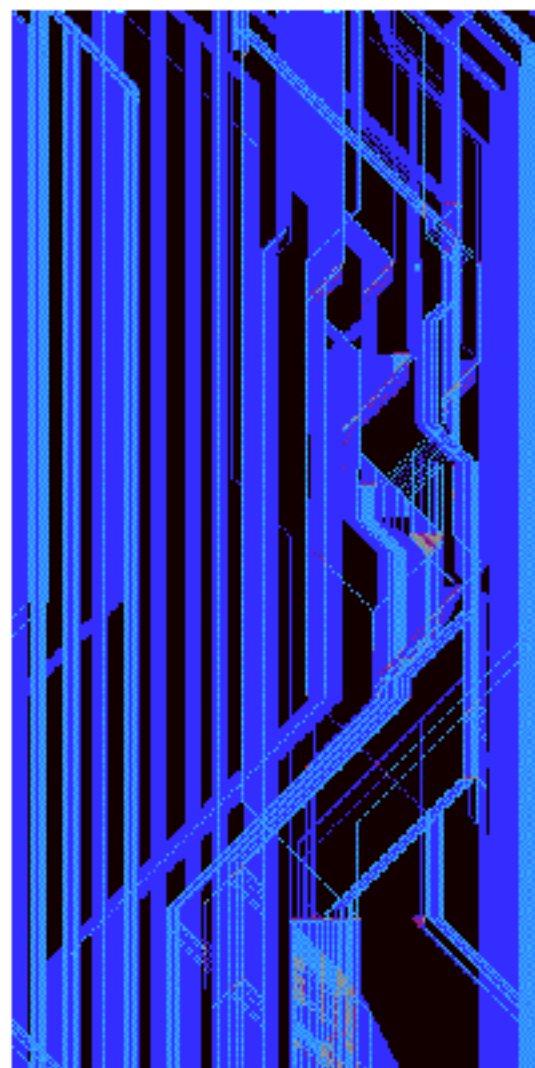


A simple logic simulation in a gate-array-like lattice dynamics.

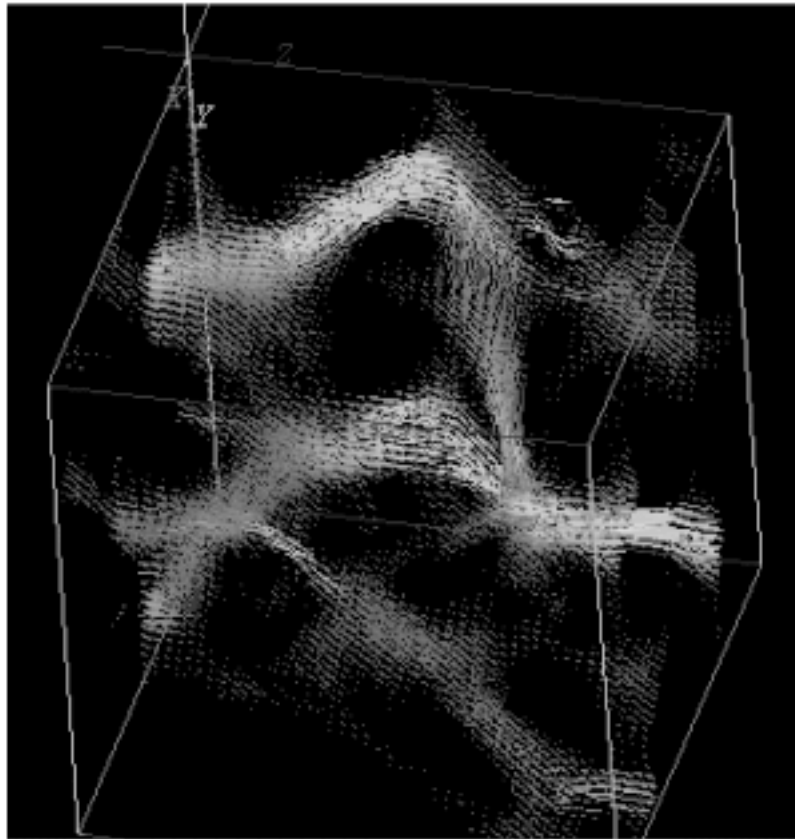
- Physical sim
- Logical sim
- Wavefront sim
- Microprocessor at right runs at 1KHz

CBSSS 6/26/02

*Wavefront simulation
(R. Agin)*



CAM-8: porous flow



64x64x64 simulation (.5mm/side)

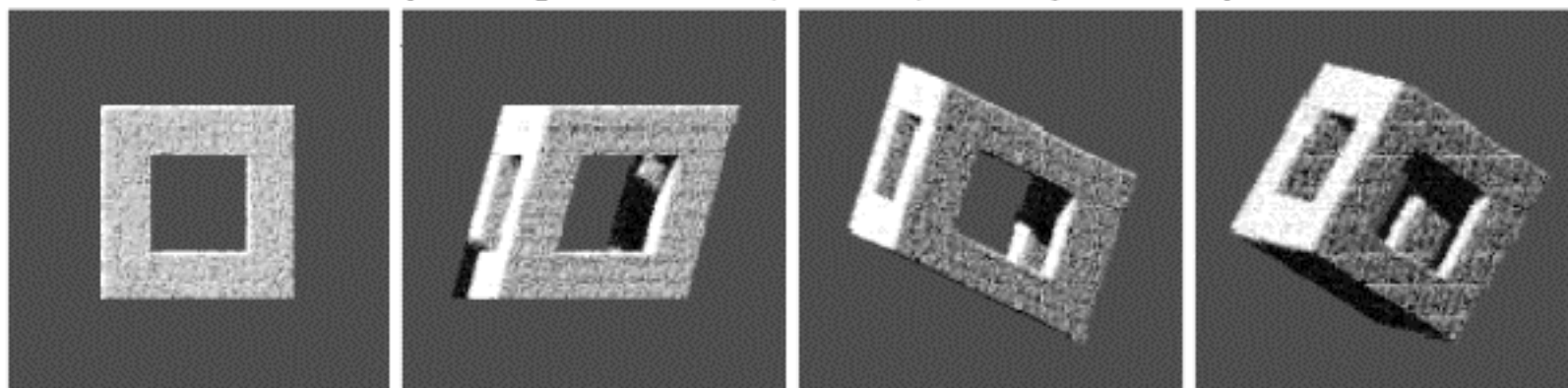
- Used MRI data from a rock (Fontainebleau sandstone)
- Simulation of flow of oil through wet sandstone agrees with experiment within 10%
- 3D LGA's used for chemical reactions, other complex fluids

CAM-8: image processing



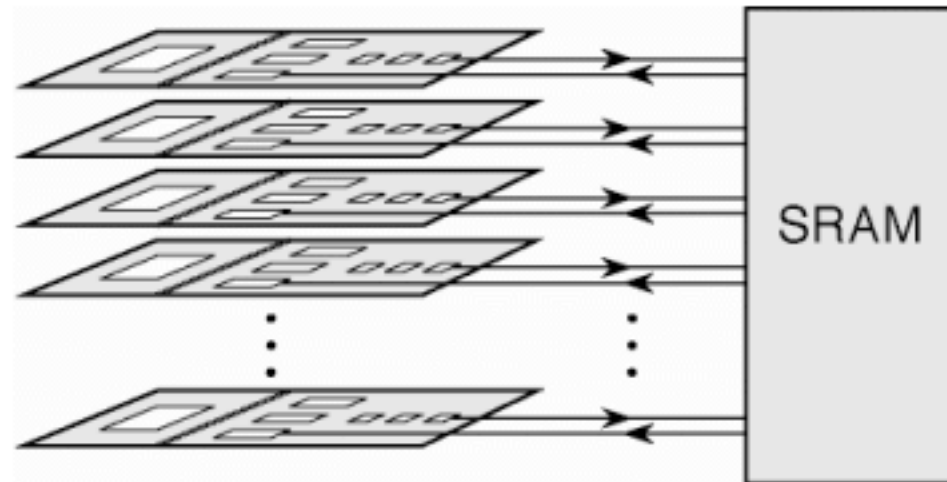
- Local dynamics can be used to find and mark features
- Bit maps can be “continuously” rotated using 3 shears

Three-dimensional rotation (Euler angles: $\alpha = -10^\circ$, $\beta = 30^\circ$, $\gamma = -15^\circ$) achieved by three shears.



CAM-8 node (1988 technology)

- 25 Million 16-bit site-updates/sec
- 8 node prototype
- arbitrary shifts of all bit fields at no extra cost
- limited by mem bandwidth
- still 100x PC!

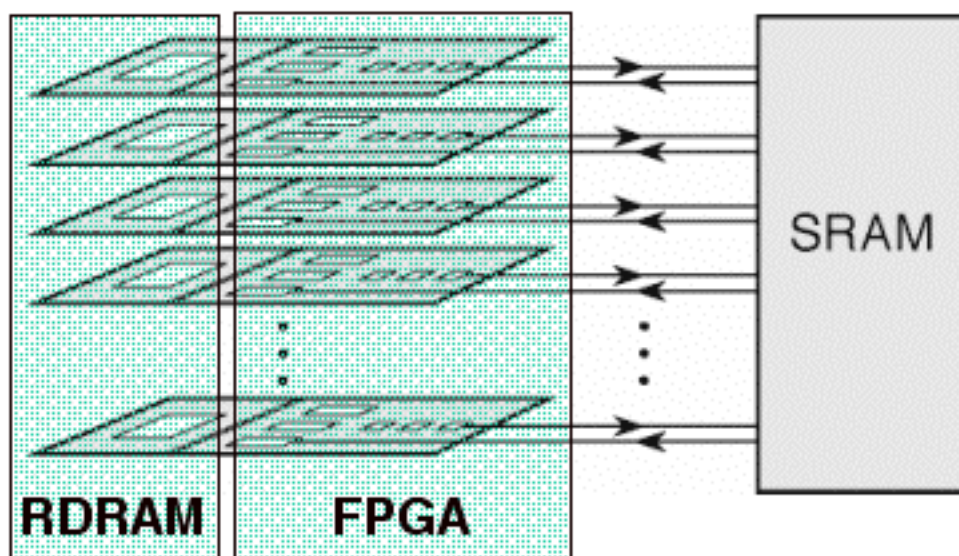


**6MB/sec x
16 DRAMs
+ 16 ASICs**

**25 MHz
64Kx16**

FPGA node design (1997)

- 100 x faster per DRAM chip
- No custom hardware
- Needed FPGA interface to RDRAM



600MB/sec
1 RDRAM
+ 1 FPGA

150 MHz
64Kx16
pipelined

Embedded DRAM (2000)

*4 Mbit Block
of DRAM:*

2K

2K



2K bits / 40 ns

*20 Blocks of DRAM
(80 Mbits total):*

40K

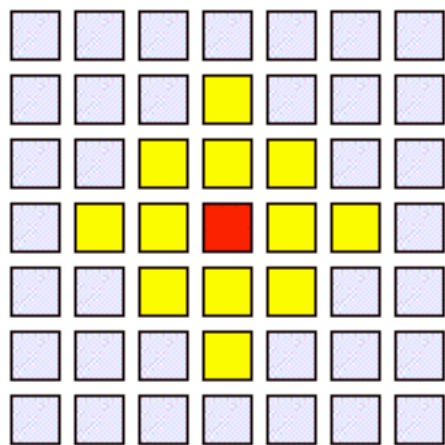
2K



40K bits / 40 ns = **1 Tbit/sec per chip!**

Each *chip* would be 1000x faster than 128-DRAM CAM-8!

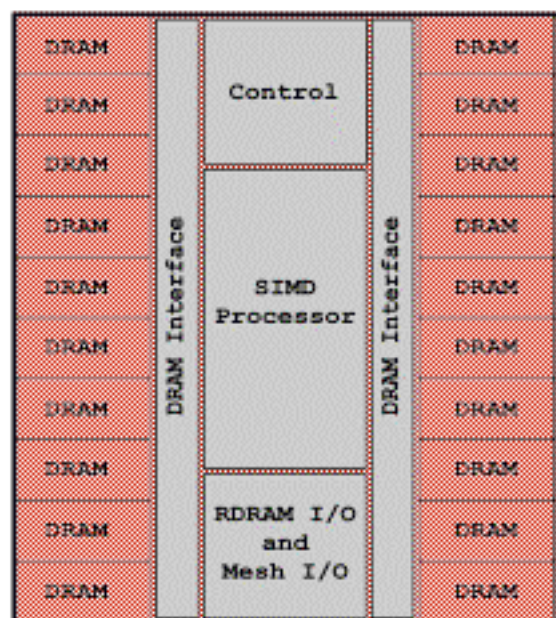
SPACERAM: A general purpose crystalline lattice computer



$$\begin{aligned} f_{x,y,t+1} = & \alpha f_{x+1,y,t} + \beta f_{x-1,y,t} + \gamma f_{x,y+1,t} \\ & + \delta f_{x,y-1,t} + \epsilon f_{x+1,y+1,t} + \zeta f_{x+1,y-1,t} \\ & + \eta f_{x-1,y+1,t} + \iota f_{x-1,y-1,t} + \kappa f_{x+2,y,t} \\ & + \mu f_{x-2,y,t} + \nu f_{x,y+2,t} + \sigma f_{x,y-2,t} \\ & + \tau f_{x,y,t} \end{aligned}$$

- Large scale prob with spatial regularity
- 1 Tbit/sec $\approx 10^{10}$ 32-bit mult-accum / sec (sustained) for finite difference
- Many *image processing* and *rendering* algorithms have spatial regularity.
- Brute-force physical simulations (complex MD, fields, etc.)
- Bit-mapped 3D games and virtual reality (simpler and more direct)
- Logic simulation (physical, wavefront)
- Pyramid, multigrid and multiresolution computations

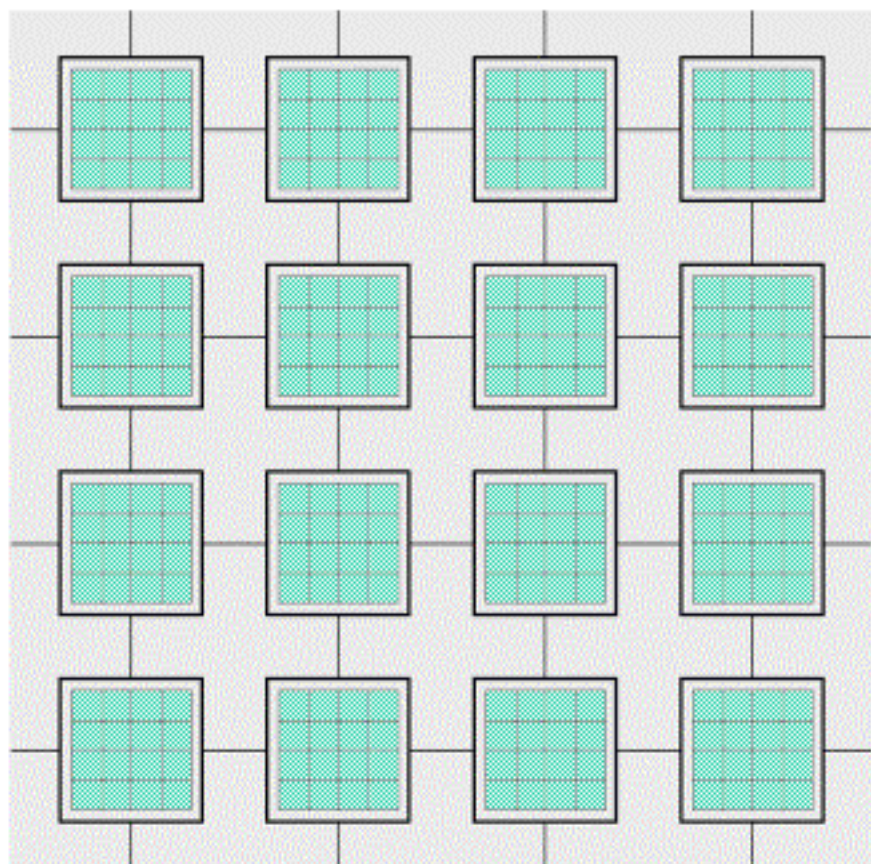
SPACERAM



4Mbits/DRAM, 256 bit I/O \times 20
@200 MHz \rightarrow 1 Tbit/sec, 10 MB,
130 mm² (.18 μ m) and 8 W (mem)

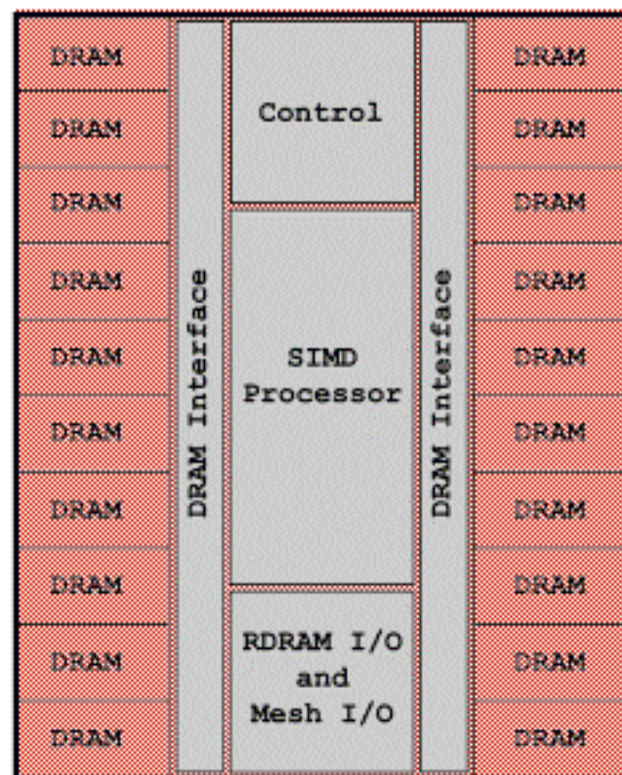
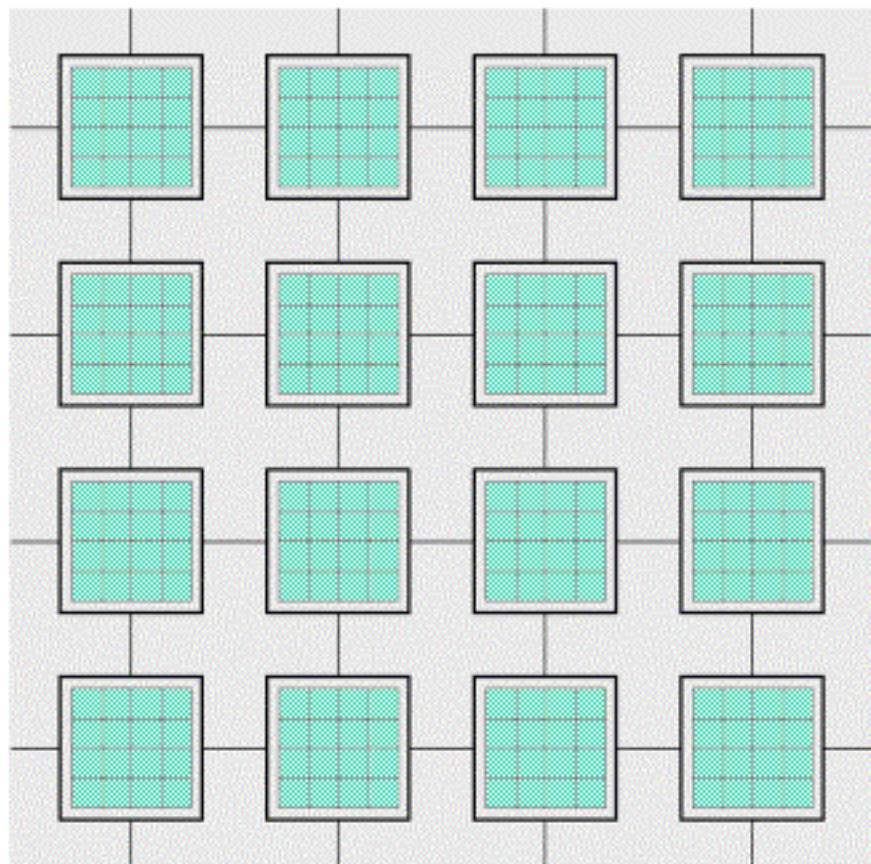
- The Task: Large-scale brute force computations with a crystalline-lattice structure
- The Opportunity: Exploit row-at-a-time access in DRAM to achieve speed *and* size.
- The Challenge: Dealing efficiently with memory granularity, commun, proc
- The Trick: Data movement by layout and addressing

Mapping a lattice into hardware



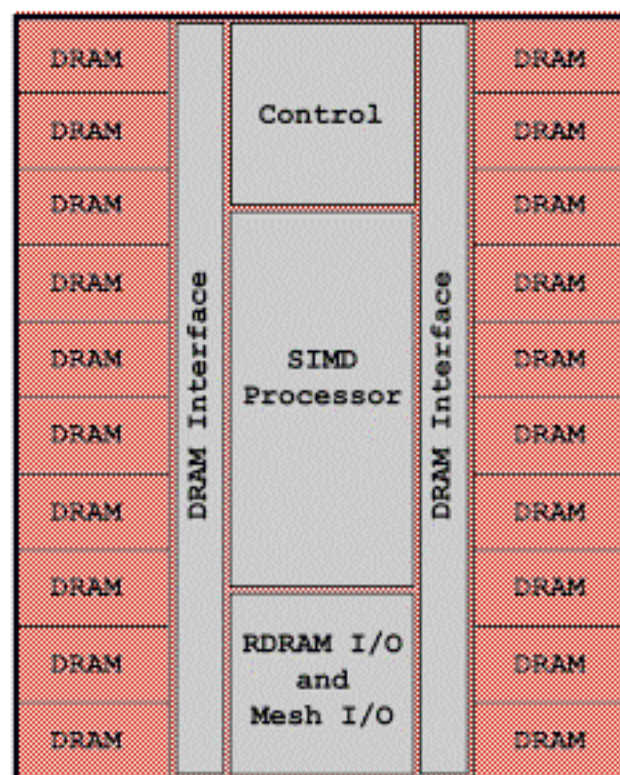
- Divide the lattice up evenly among mesh array of chips
- Each chip handles an equal sized *sector* of the emulated lattice

Mapping a lattice into hardware

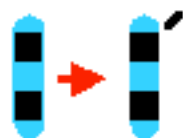
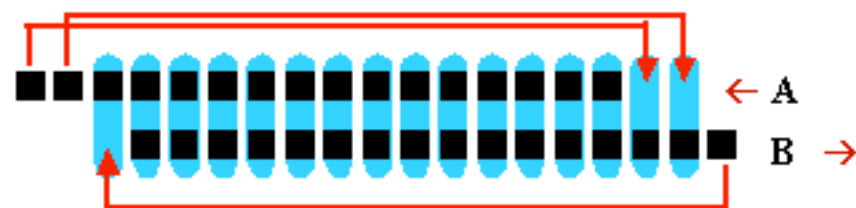


Mapping a lattice into hardware

- Use blocks of DRAM to hold the lattice data (Tbits/sec)
- Use virtual processors for large comps (depth-first, wavefront, skew)
- Main challenge: mapping lattice computations onto granular memory



Lattice computation model

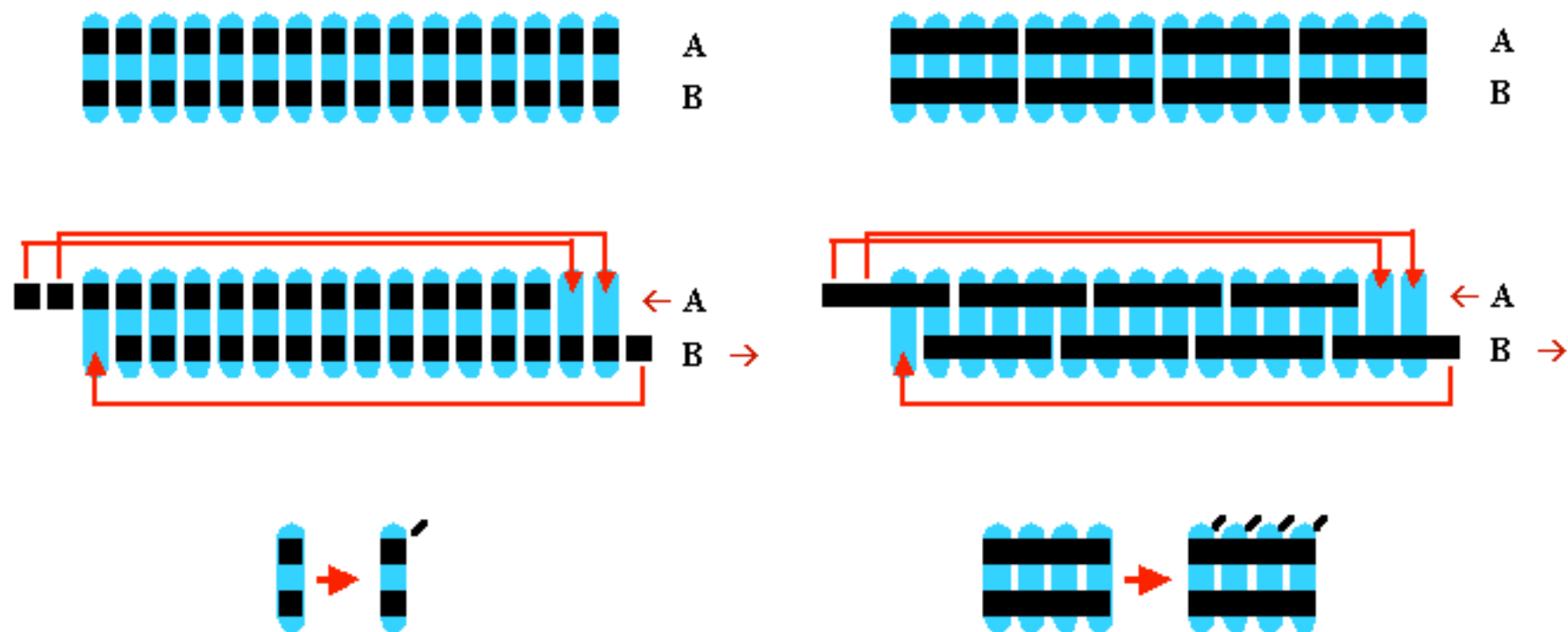


1D example: the rows are bit-fields and columns sites

Shift bit-fields periodically and uniformly (may be large)

Process sites independently and identically (SIMD)

Mapping the lattice into memory



Mapping the lattice into memory

CAM-8 approach: group adjacent into memory words



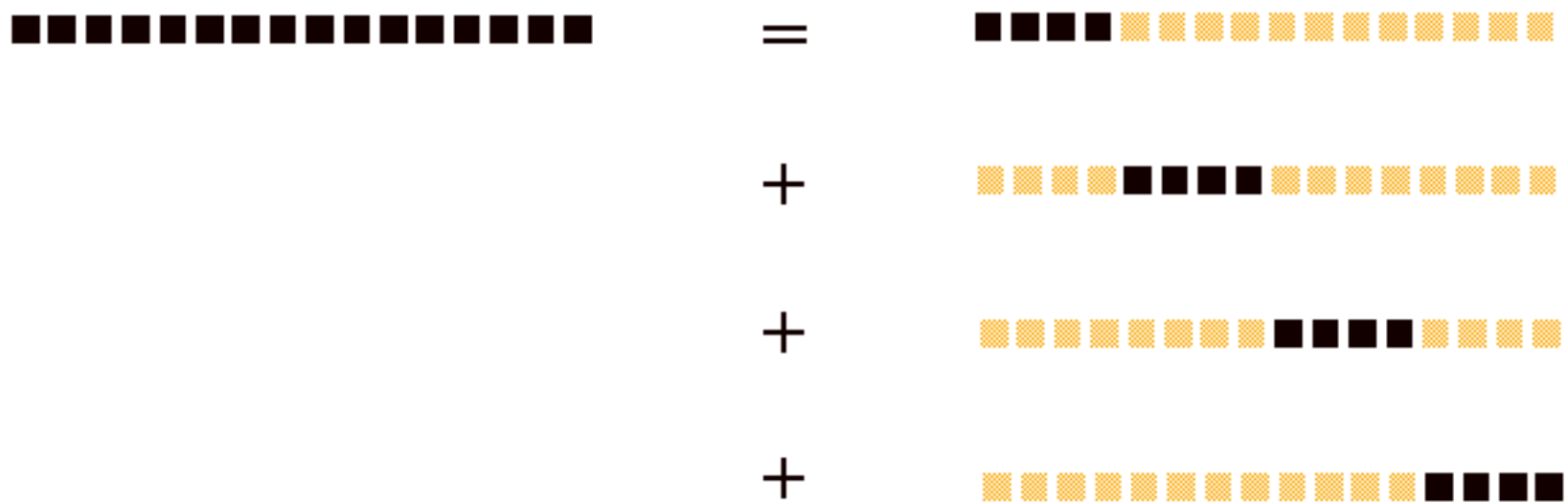
Shifts can be long but re-aligning messy → chaining



*Can process word-wide
(CAM-8 did site-at-a-time)*



Mapping the lattice into memory



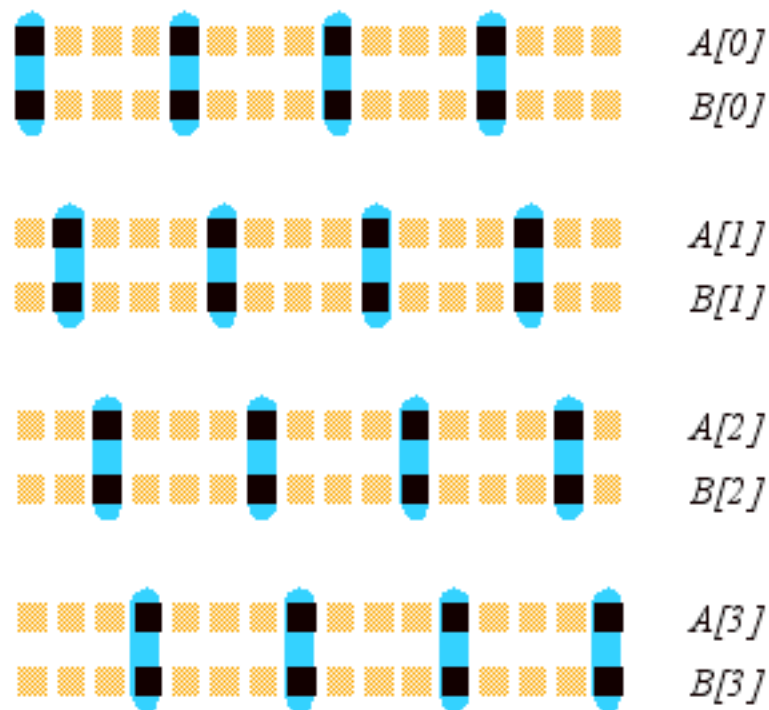
CAM-8: Adjacent bits of bit-field are stored together

Mapping the lattice into memory



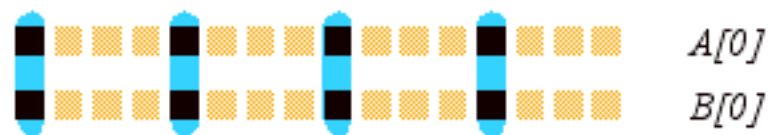
SPACERAM: bits stored as evenly spread skip samples.

Processing skip-samples

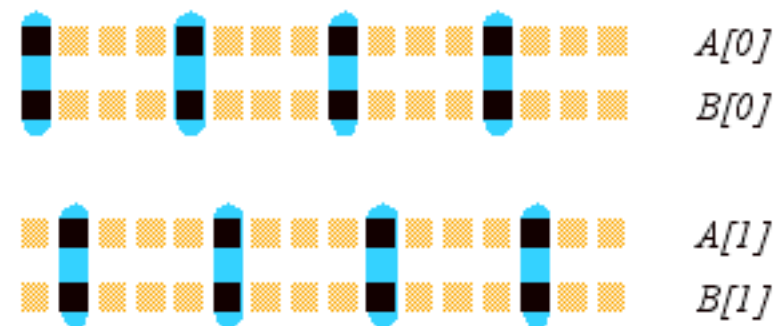
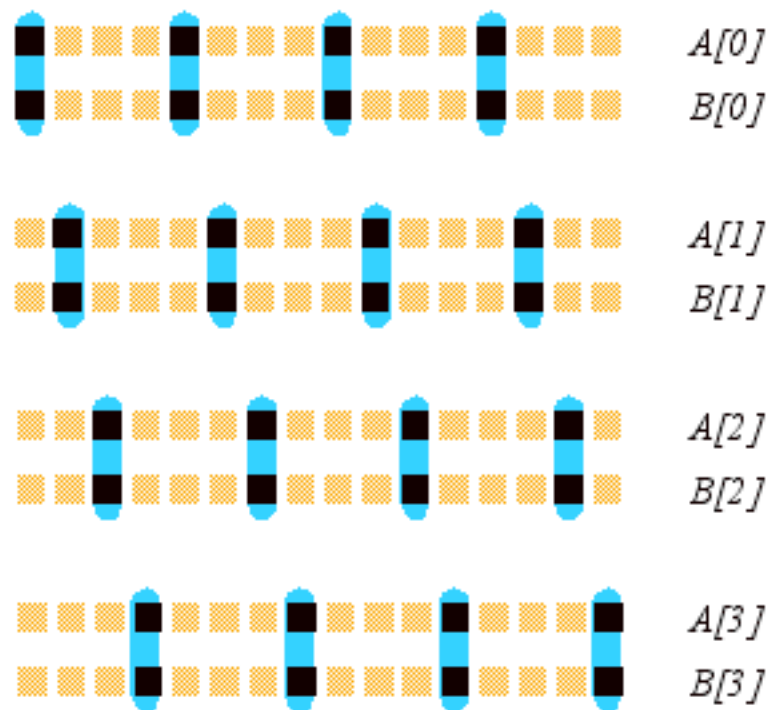


- Process sets of evenly spaced lattice sites
(site groups)
- Data movement is done by reading shifted data

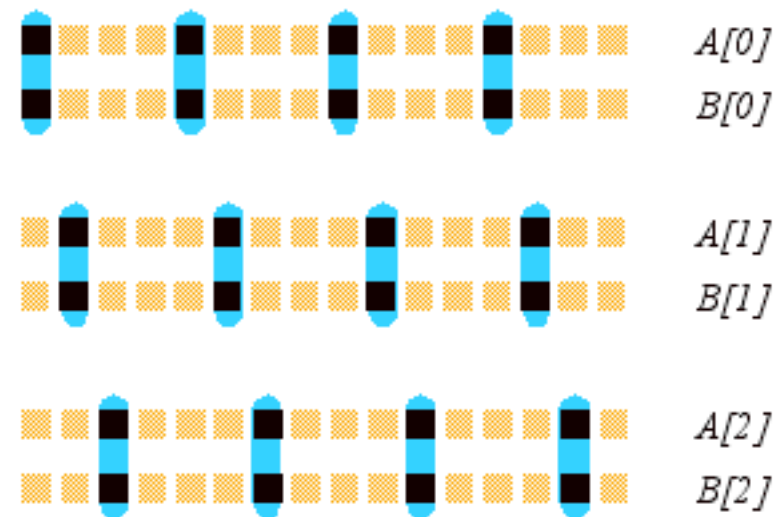
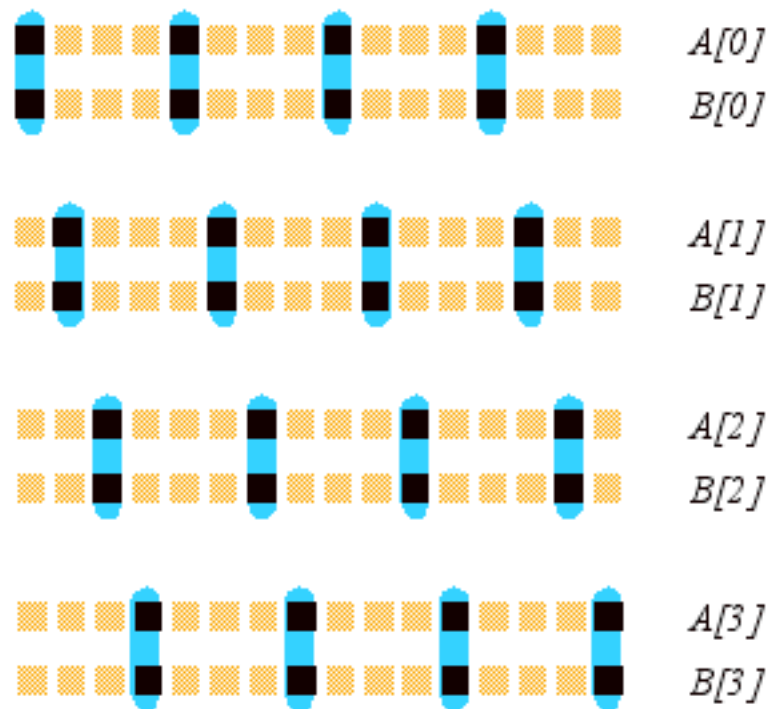
Processing skip-samples



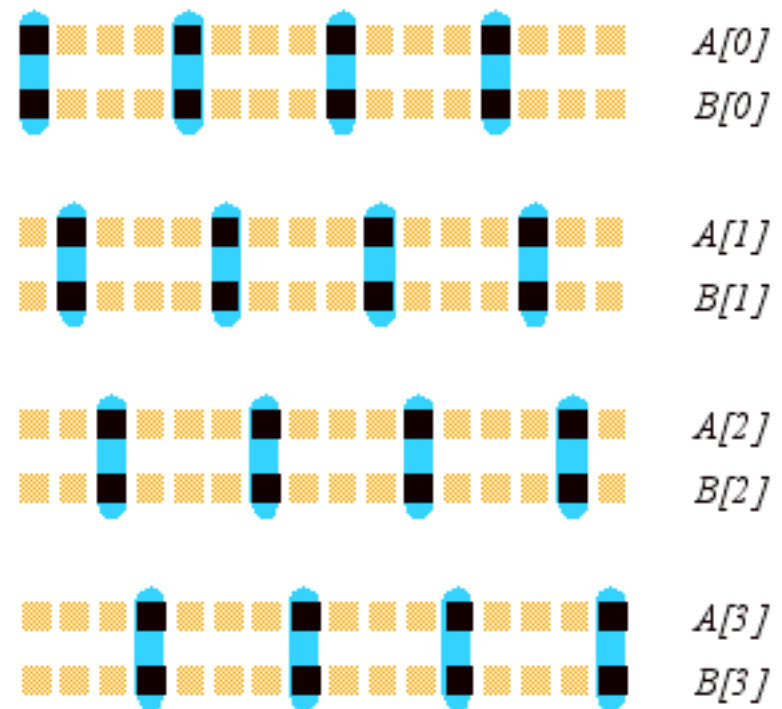
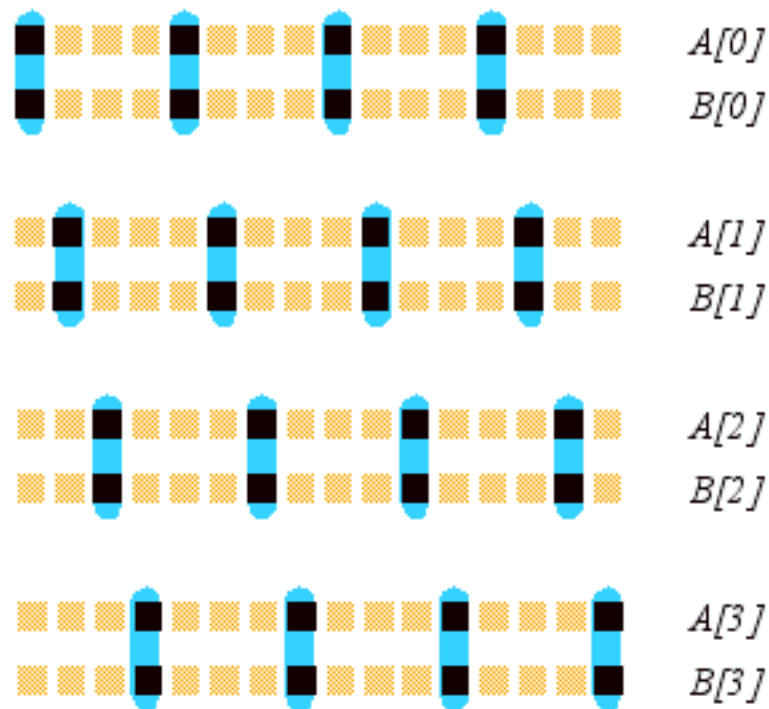
Processing skip-samples



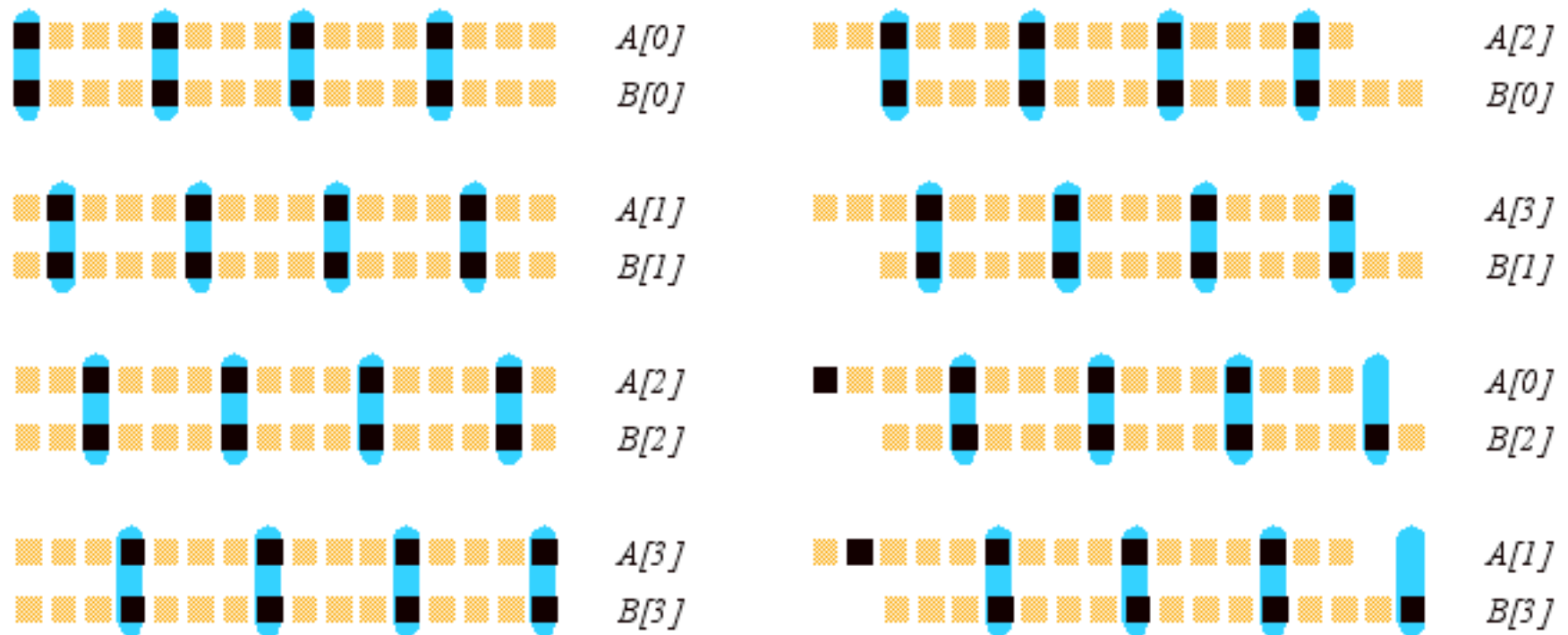
Processing skip-samples



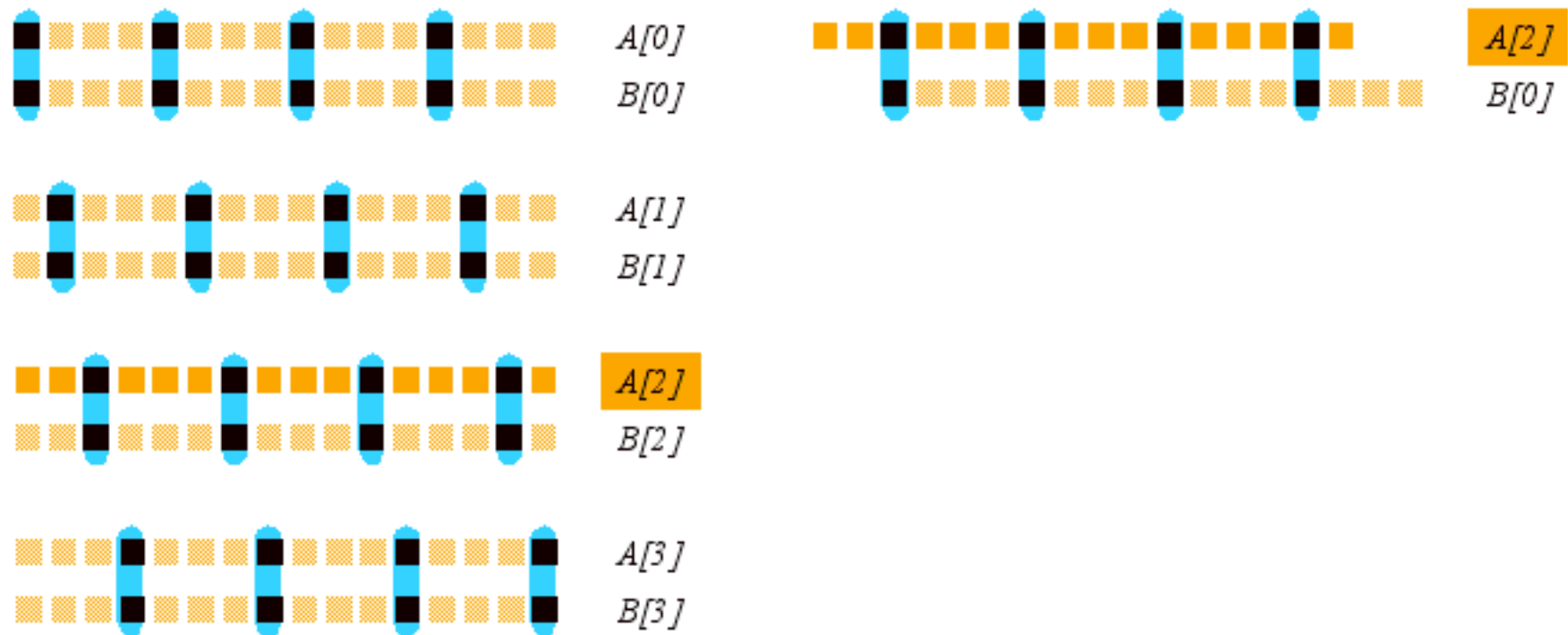
Processing skip-samples



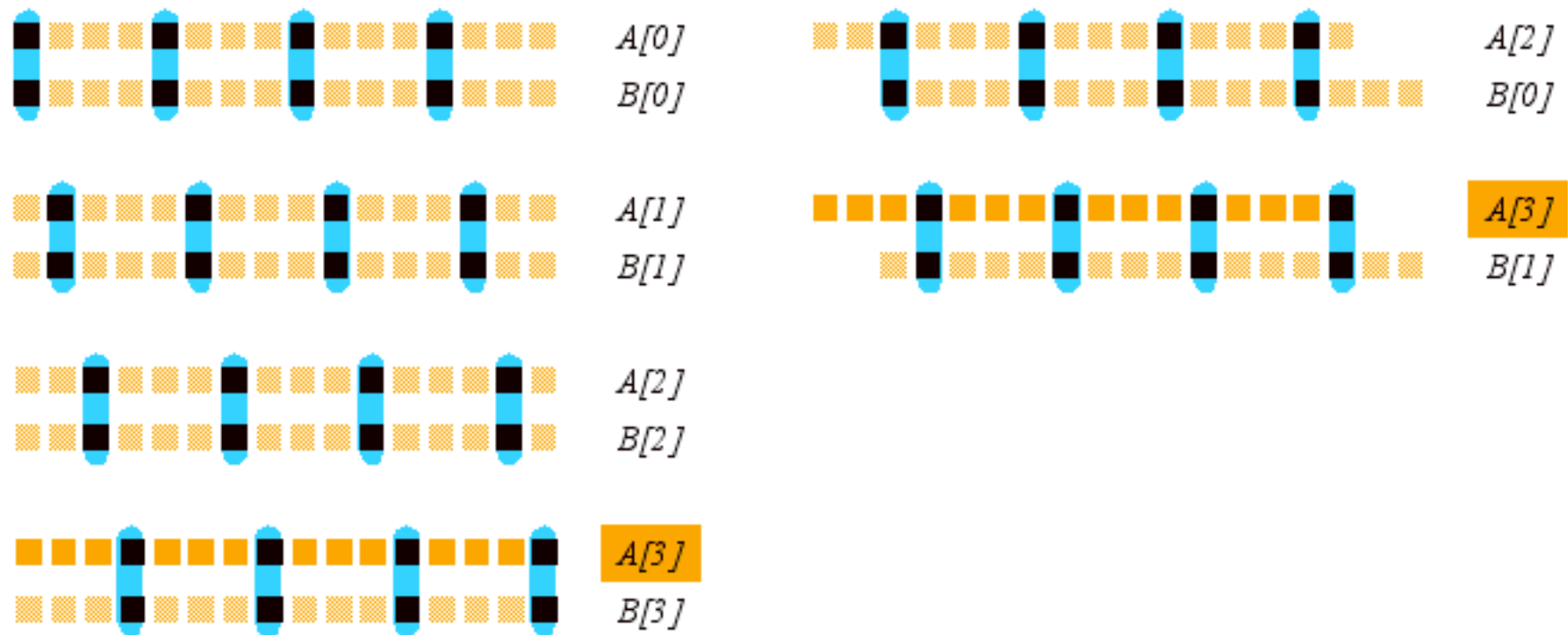
Processing skip-samples



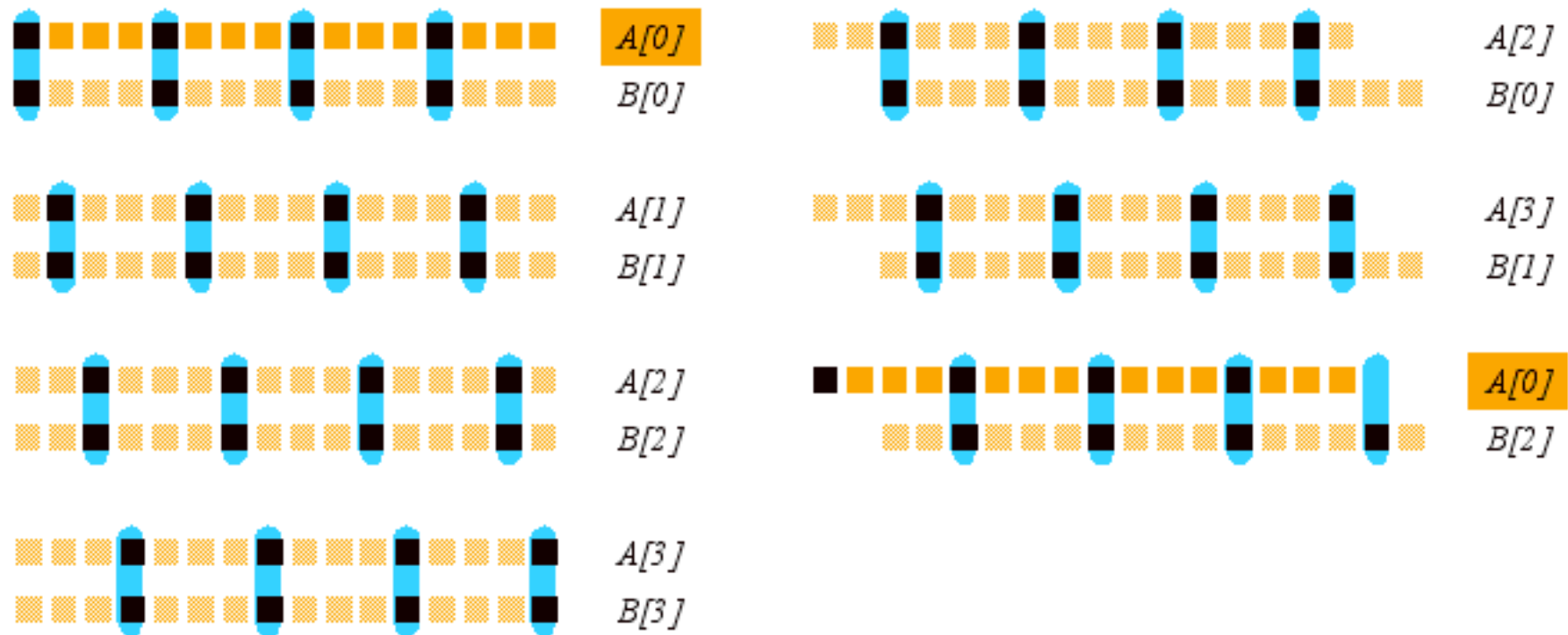
Processing skip-samples



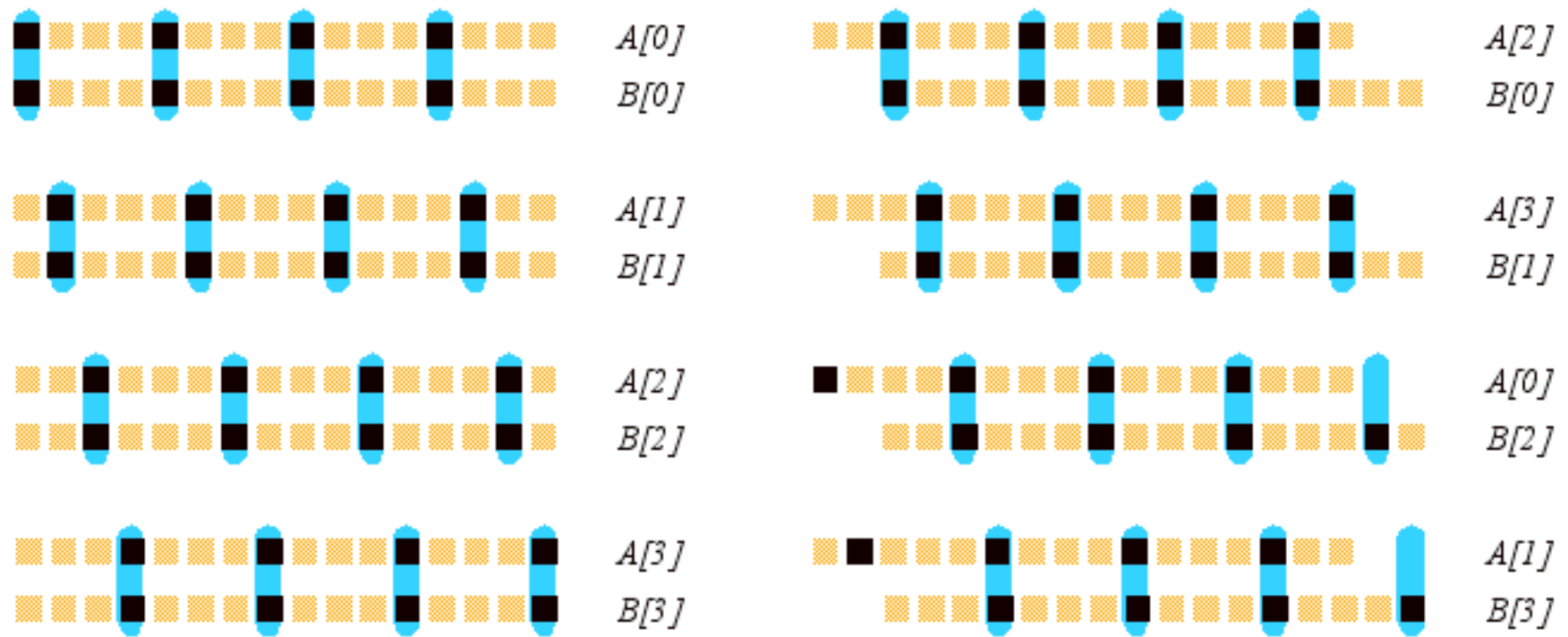
Processing skip-samples



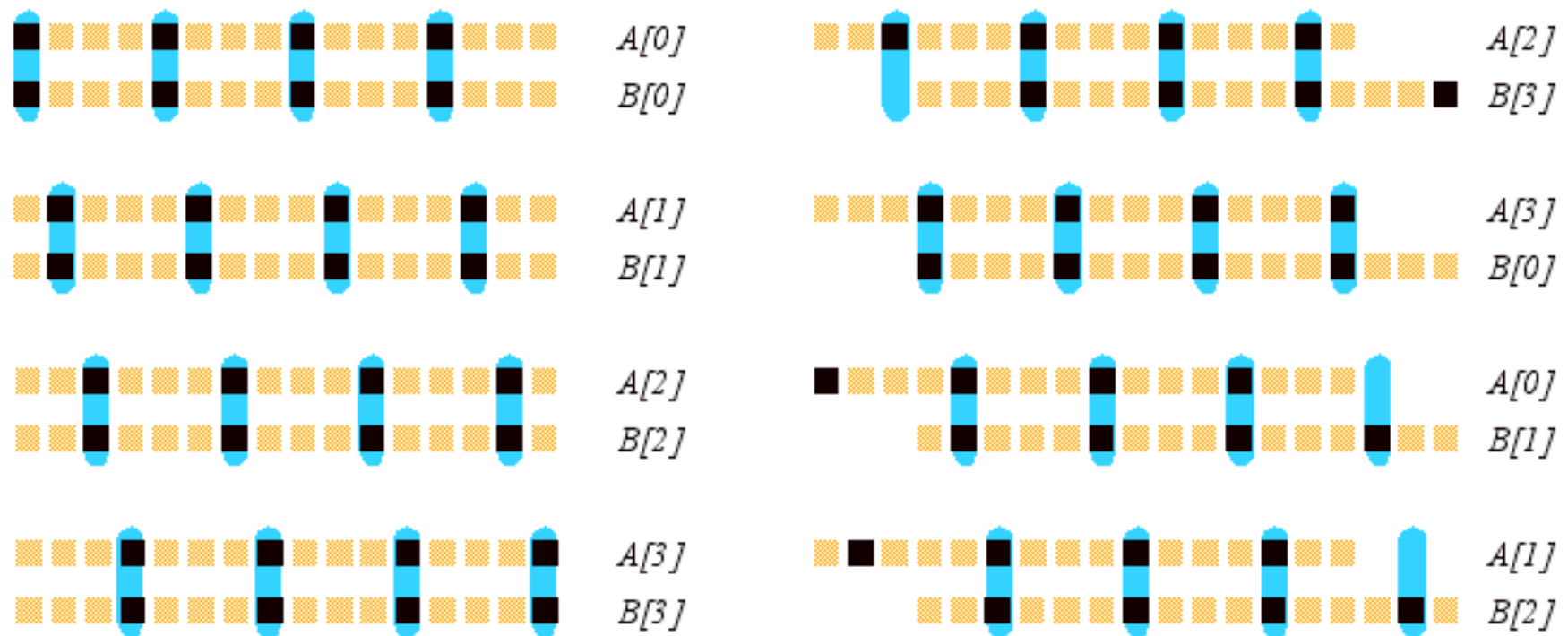
Processing skip-samples



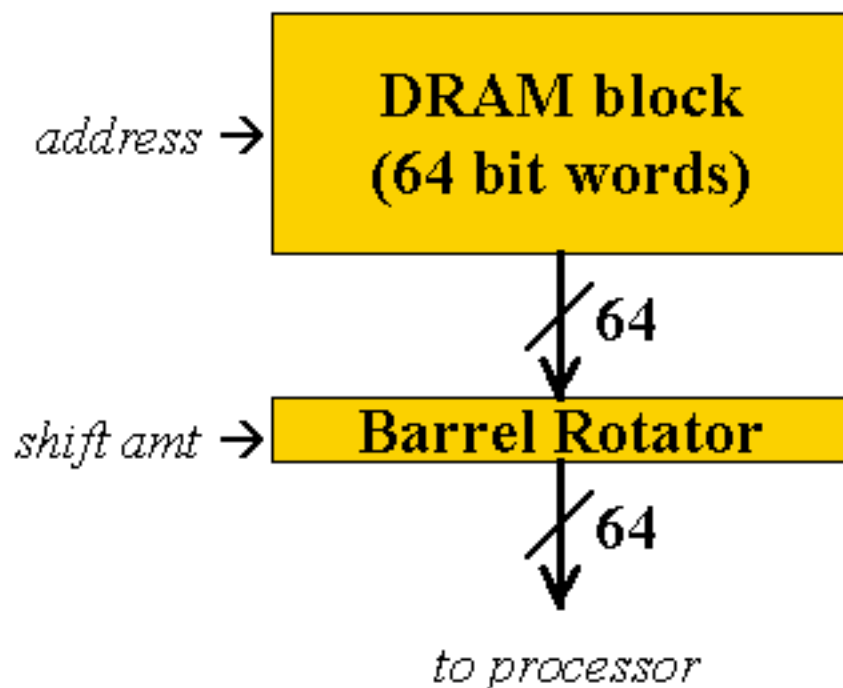
Processing skip-samples



Processing skip-samples



DRAM module



- Assume hardware word size is 64-bits
- Assume we can address any word
- For each site group: address next word we need; rot if necessary
- Each bit field resides in a different module

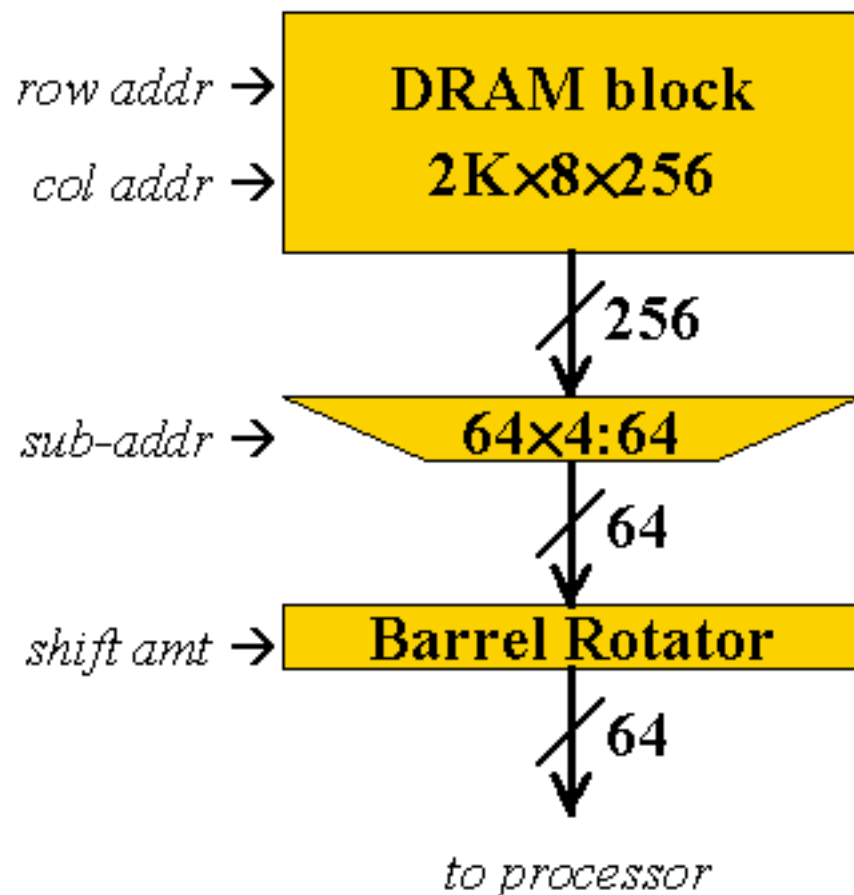
Problem: granularity

- DRAM block has structure: eg., all words in a row must be used before switching rows

Solution: grouping

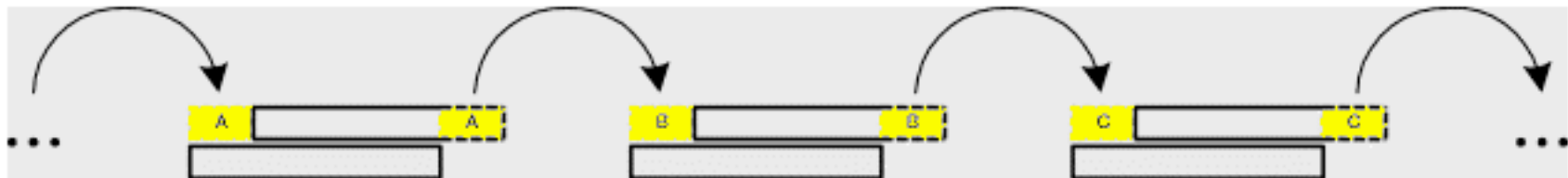


DRAM module



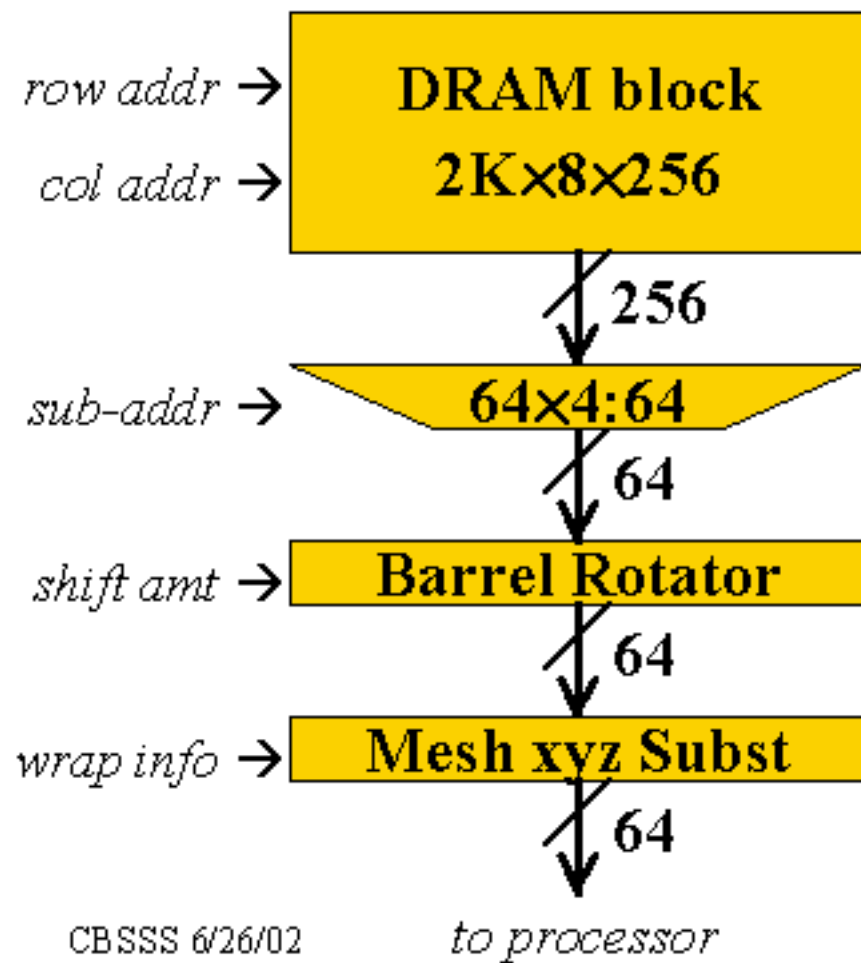
- 3 levels of granularity: 2K-bit row, 256-bit word, 64-bit word
- All handled by data layout, addressing, and rot within 64-word

Gluing chips together



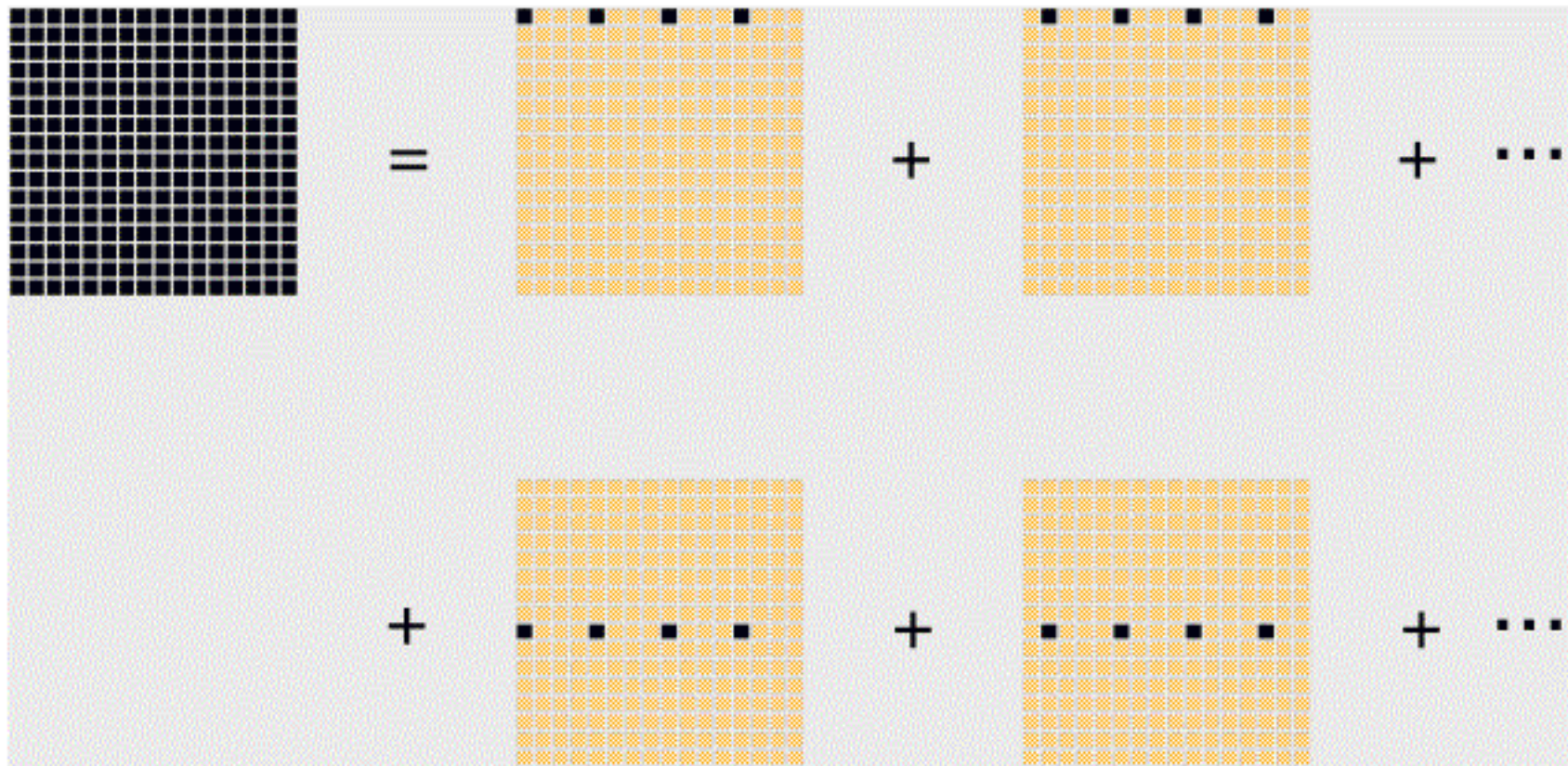
- All chips operate in lockstep
- Shift is periodic within each sector
- To glue sectors, bits that stick out replace corresponding bits in the next sector

DRAM module

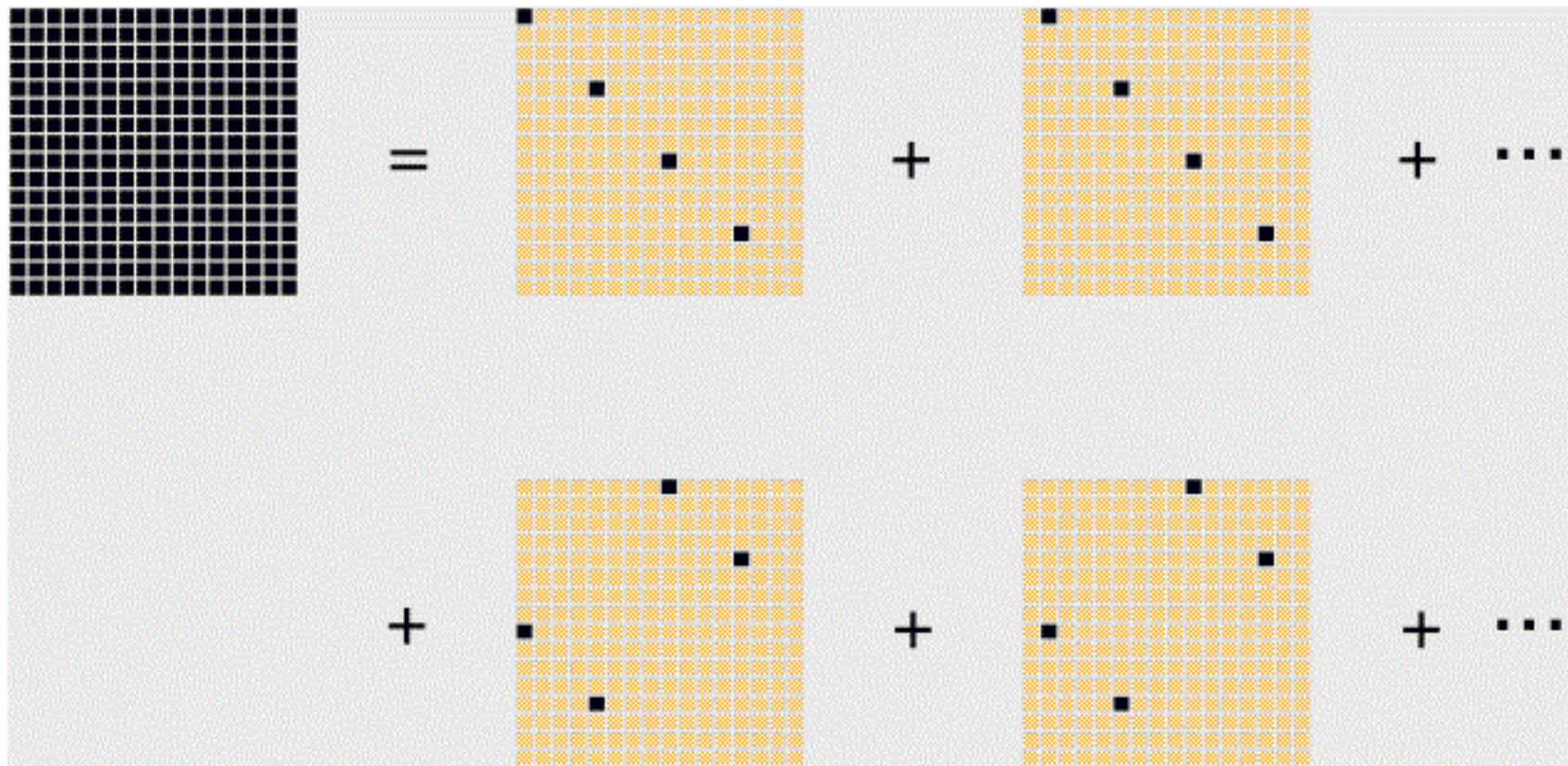


- Any data that wraps around is transmitted over the mesh
- Corresponding bit on corresponding wire is substituted for it
- This is done sequentially in each of the three directions

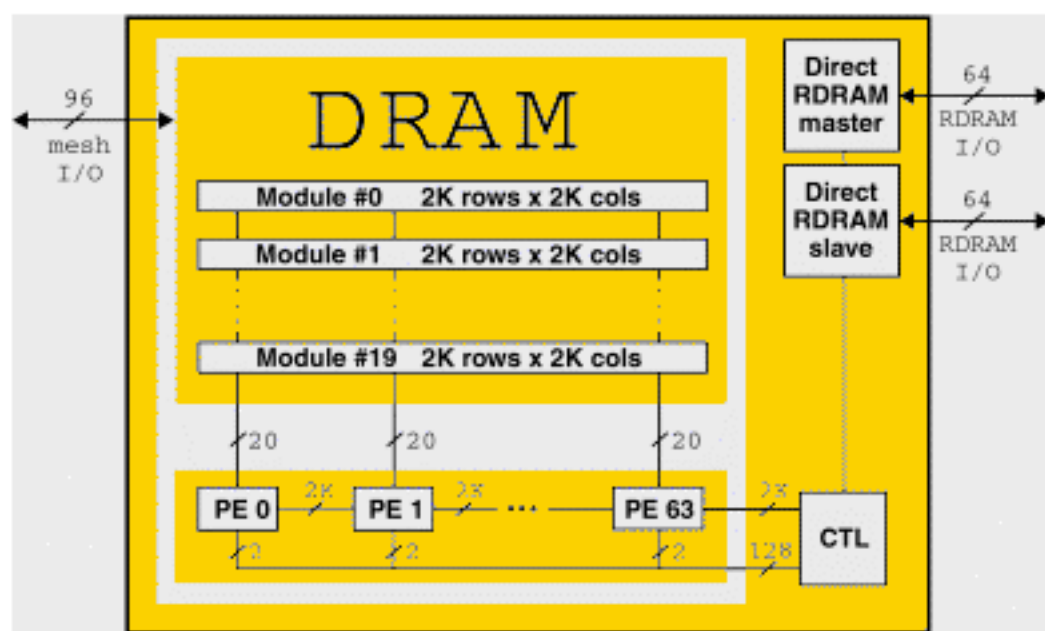
Partitioning a 2D sector



Partitioning a 2D sector

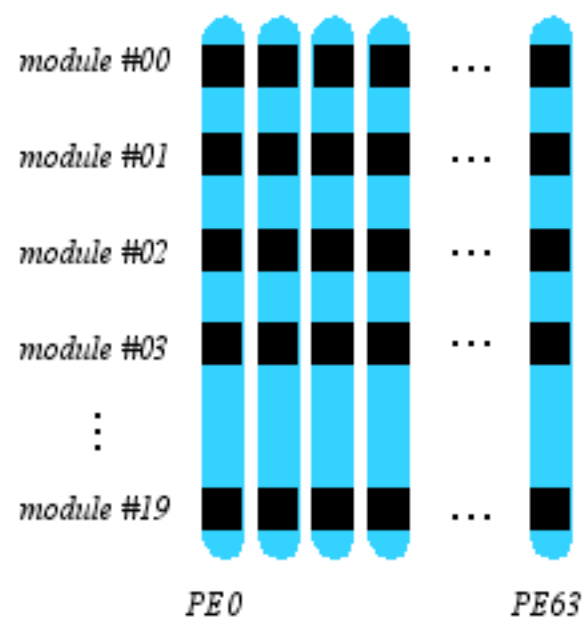
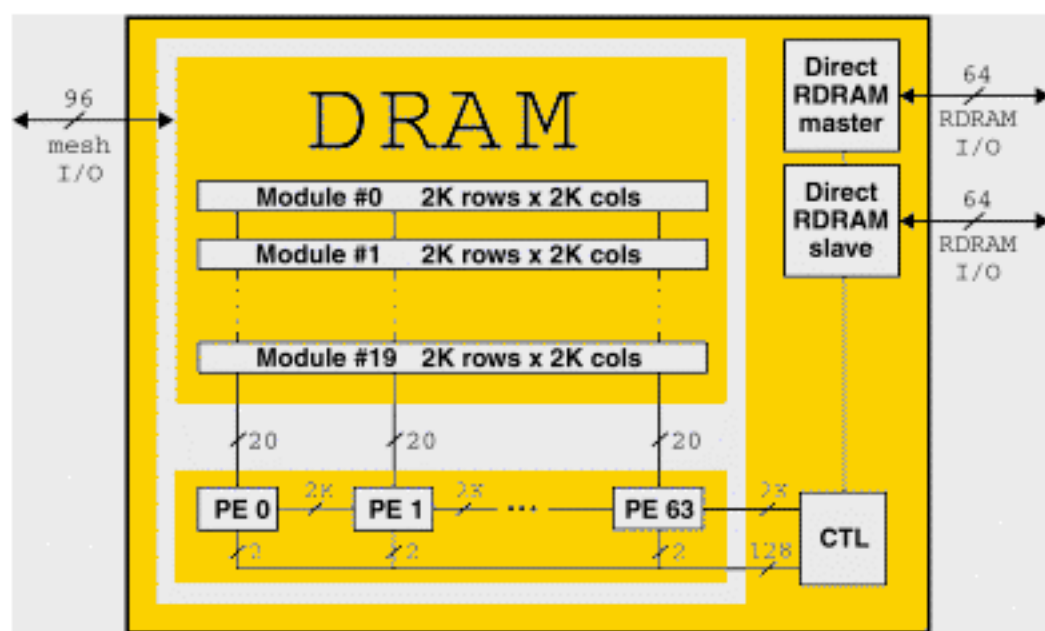


The SPACERAM chip

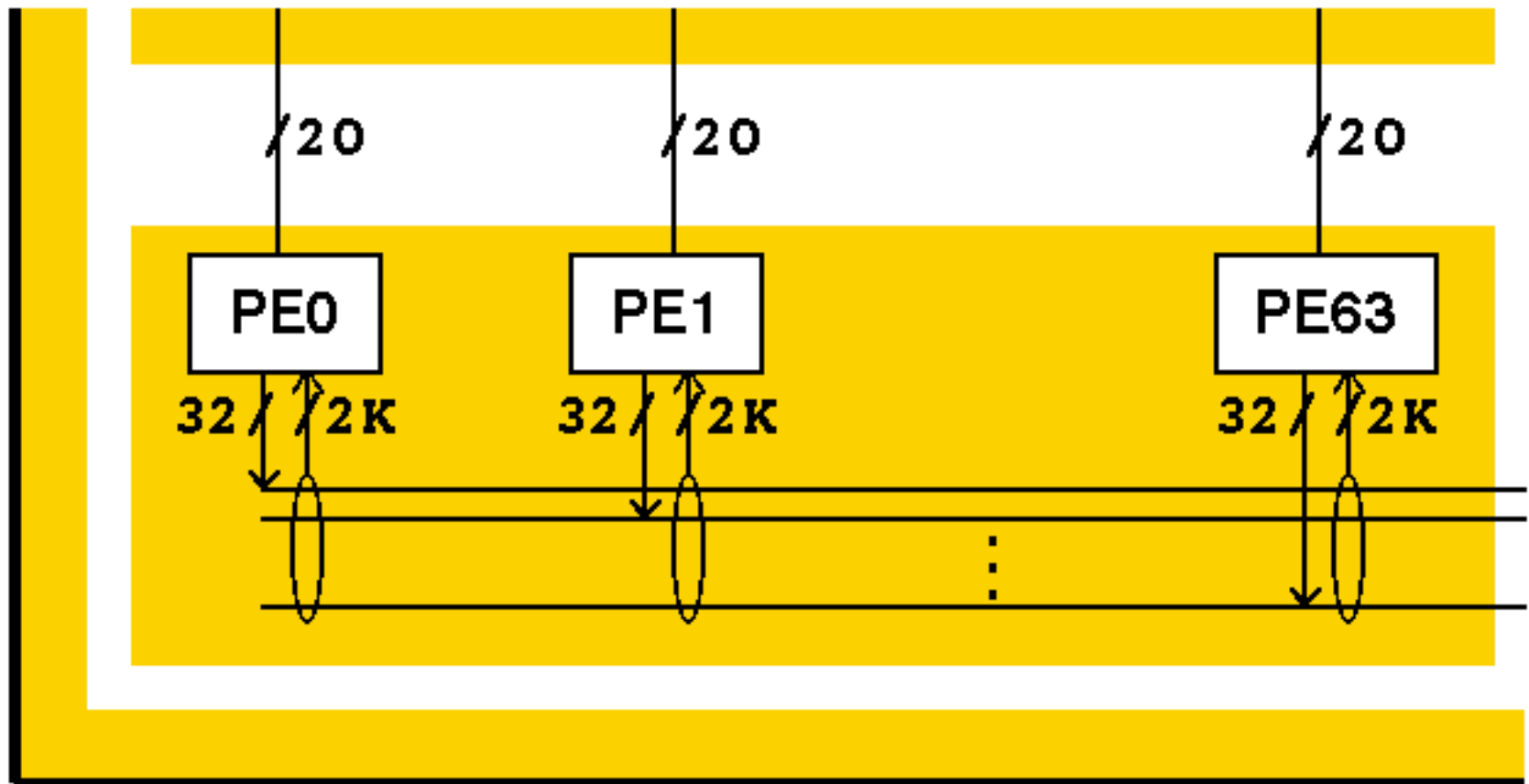


- .18 μm CMOS DRAM proc. (20W, 215 mm²)
- 10% of memory bandwidth \rightarrow control
- Memory hierarchy (external memory)
- Single DRAM address space

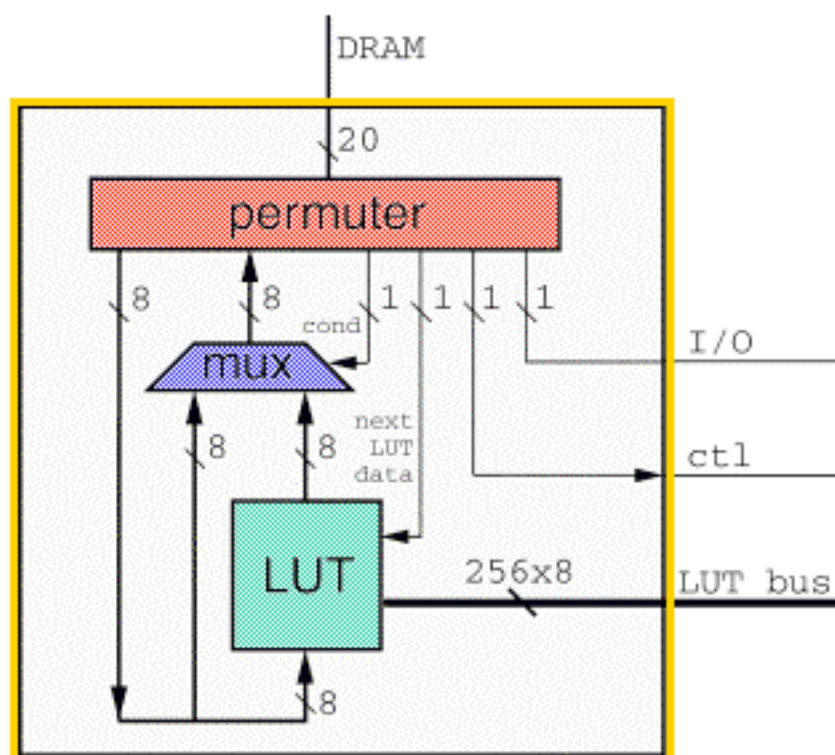
The SPACERAM chip



The SPACERAM chip

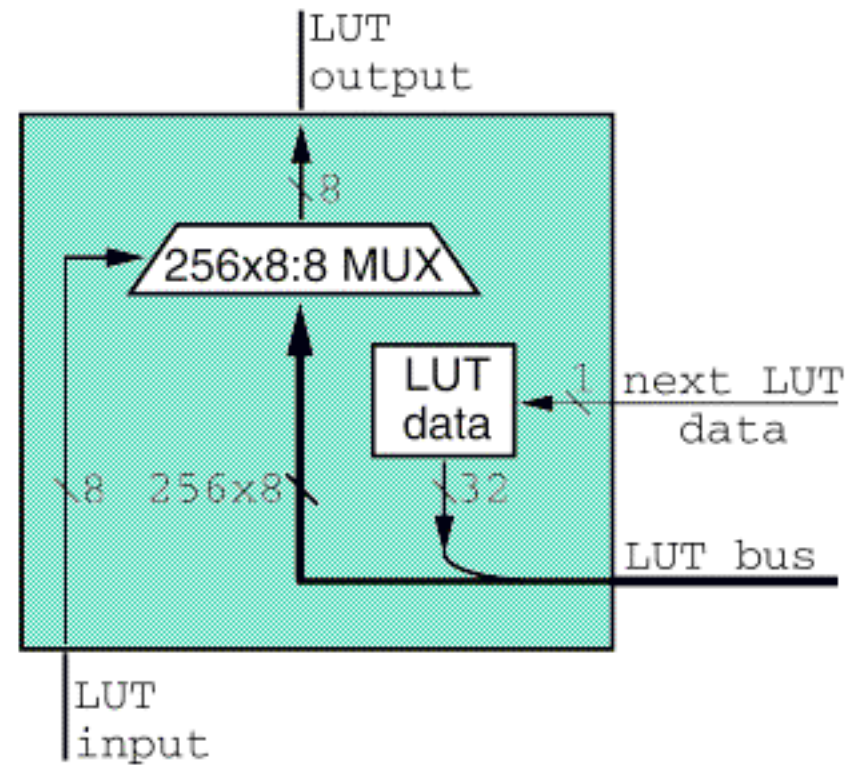
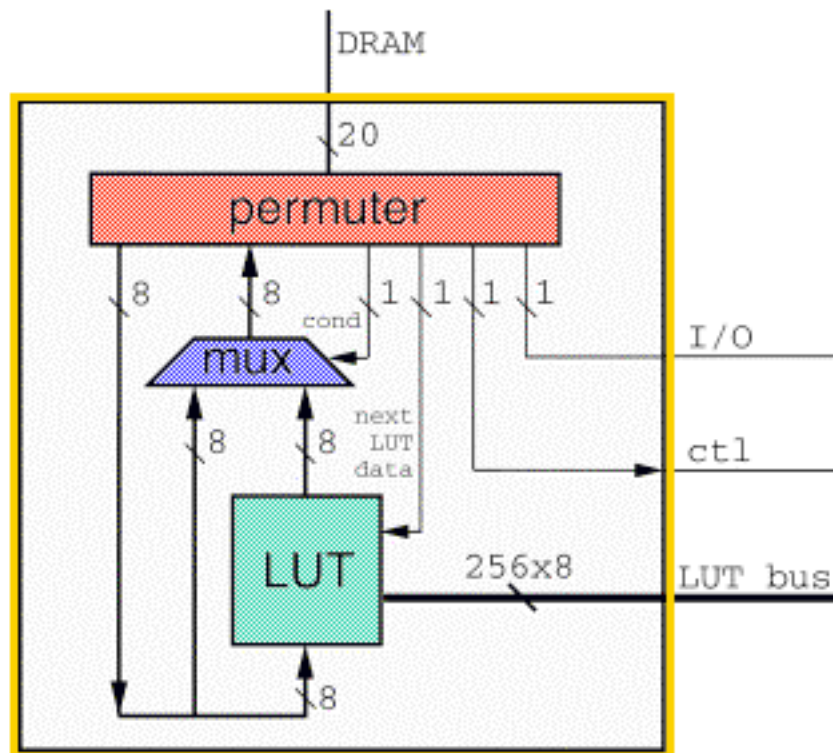


SPACERAM: symbolic PE

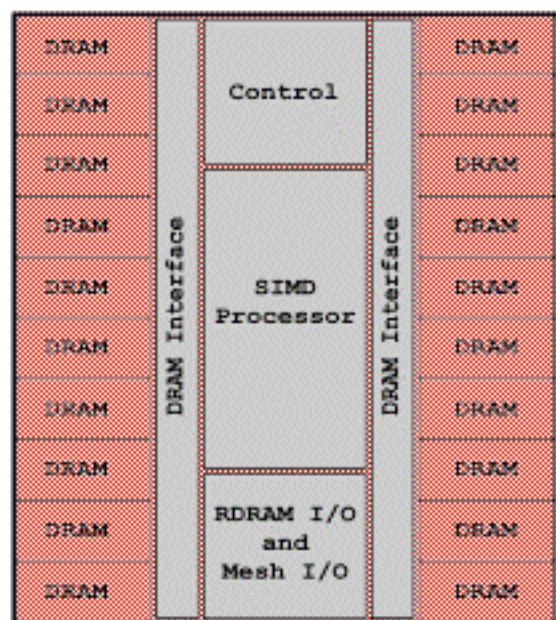


- Computation by LUT
- Permuter lets any DRAM wire play any role
- MUX lets LUT output be used conditionally
- LUT is really a MUX, with data bussed

SPACERAM: symbolic PE



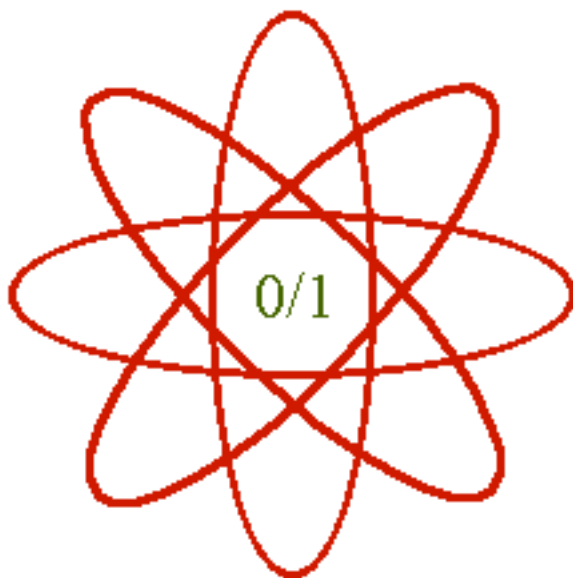
Conclusions



*. 18 μ m CMOS, 10 Mbytes DRAM,
1 Tbit/sec to mem, 215 mm², 20W*

- Addressing is powerful
- Simple hardware provides a general mechanism for lattice data movement
- We can exploit the parallelism of DRAM access without extra overhead

Physics becomes the computer



Emulating Physics

- » *Finite-state, locality, invertibility, and conservation laws*

Physical Worlds

- » *Incorporating comp-universality at small and large scales*

Spatial Computers

- » *Architectures and algorithms for large-scale spatial computations*

Nature as Computer

- » *Physical concepts enter CS and computer concepts enter Physics*