

Abelian Square-Free Dithering and Recoding for Iterated Hash Functions

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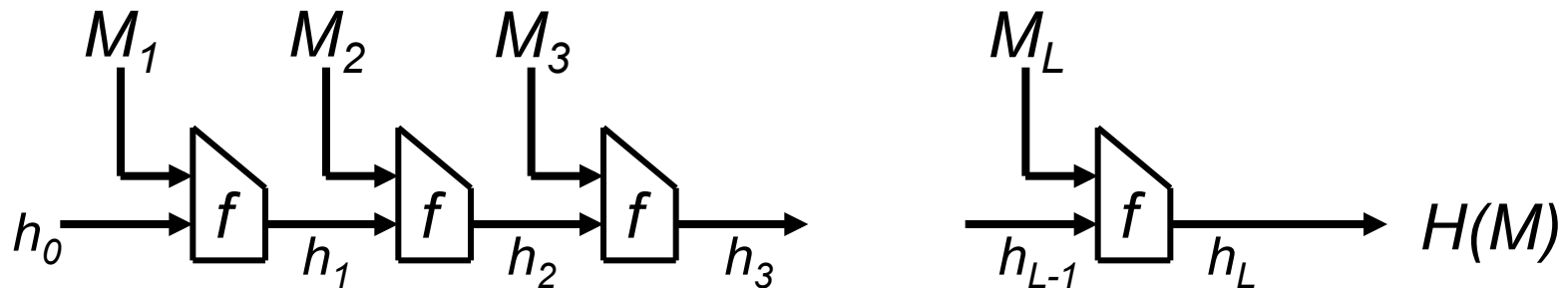
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Outline

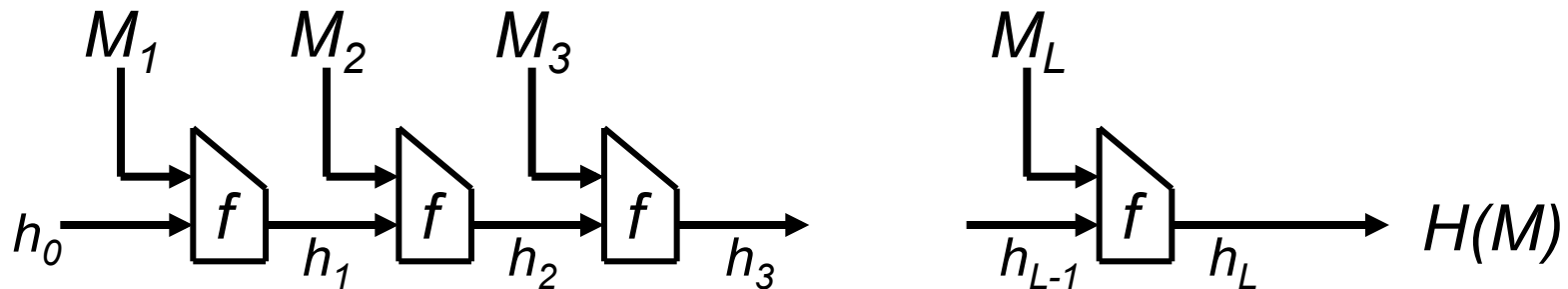
- ◆ Dean/Kelsey/Schneier Attacks
- ◆ Square-Free Sequences
 - Prouhet-Thue-Morse Sequences
 - Towers of Hanoi
- ◆ Abelian Square-Free Sequences
 - Keränen's Sequence
- ◆ Dithering and Recoding
- ◆ Open Questions
- ◆ Conclusions

Typical Iterated hashing



- ◆ Message extended with 10^* & length (MD)
- ◆ f is compression function.
- ◆ h_0 is initialization vector (IV)
- ◆ h_i is i -th chaining variable
- ◆ Last chaining variable h_L is hash output $H(M)$

Dean/Kelsey/Schneier Attacks



- ◆ Assumes one can find *fixpoint* h for f, M_0 :
$$h = f(h, M_0)$$
- ◆ Can then have *message expansion attacks* that find *second preimage* by
 - Finding many fixpoint pairs (h, M)
 - Finding a fixpoint h in actual chain for given message
 - Finding another shorter path from h_0 to some chaining variable
 - Creating second preimage with this new starting path using message expansion to handle Merkle-Damgard strengthening

Dithering and Recoding

- ◆ Make hash function round *dependent on round index i* as well as h_{i-1} and M_i

- ◆ *Dithering*: include dither input d_i to compression function:

$$h_i = f(h_{i-1}, M_i, d_i)$$

- ◆ *Recoding*: Include dither input as part of i -th message block

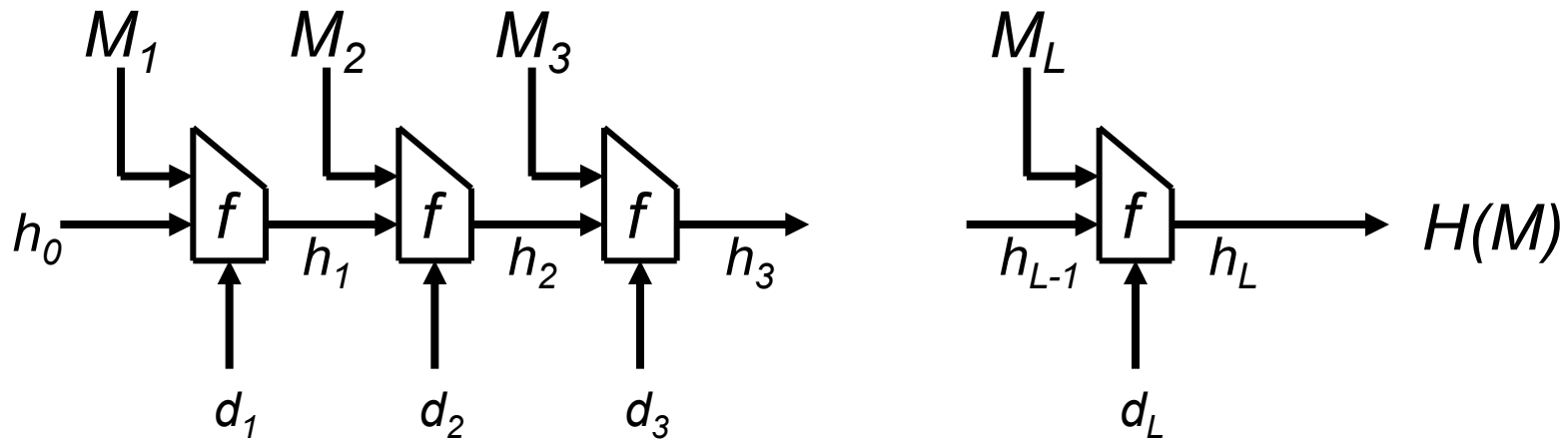
$$h_i = f(h_{i-1}, M'_i)$$

where

$$M'_i = (M_i, d_i)$$

- ◆ (These are equivalent, of course...)

Iterated hashing with dithering



- ◆ How to choose dither input d_i ?
 - Could choose $d_i = i$
 - Could choose $d_i = r_i$ (pseudo-random)
 - Use *square-free sequence* d_i
(repetition-free sequence; no repeated symbols or subwords.)

Square-Free Sequence

- ◆ A sequence is *square-free* if it contains no two equal adjacent subwords.
- ◆ Examples:
 - abracadabra is square-free
 - hobbit is not (repeated "b")
 - banana is not (repeated "an")
- ◆ Dithering with a square-free sequence prevents message expansion attacks. (Would need fixpoint that works for all dither inputs.)

Infinite square-free sequences

- ◆ There exists infinite square-free sequences over 3-letter alphabet.
- ◆ Start with parity sequence:
0110100110010110...
 i -th element is parity of integer i .
This (Prouhet-Thue-Morse, or PTM) sequence is only *cube-free*, but...
- ◆ Sequence of inter-zero gap lengths in PTM is square-free:
2102012101202102012021...

Generating infinite sf sequences

◆ Or:

- Take two copies of PTM sequence; shift second one over by one, then code vertical pairs:

A = 00, B = 01, C = 10, D = 11:

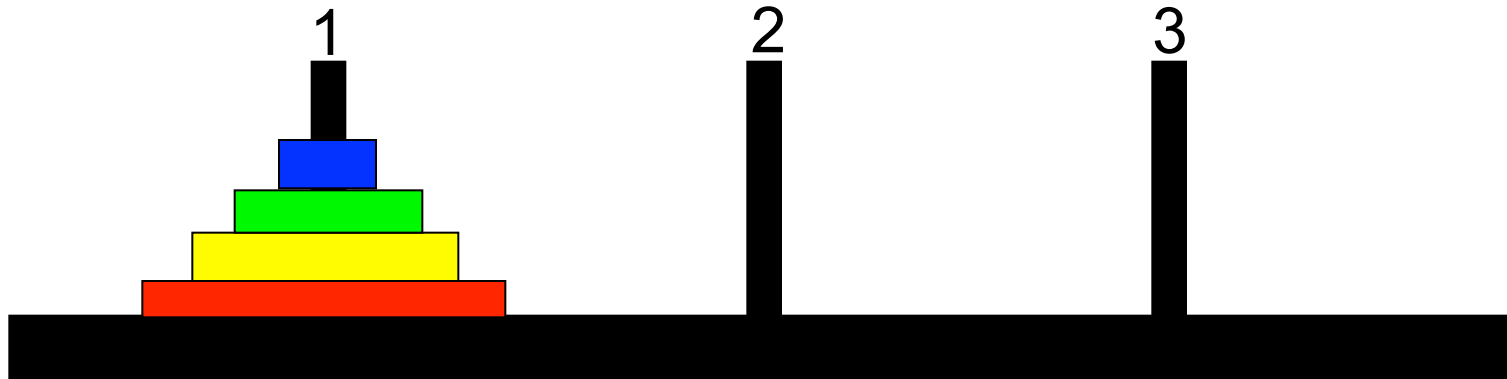
0 1 1 0 1 0 0 1 1 0 0 1 0 1 ...

- 0 1 1 0 1 0 0 1 1 0 0 1 0 ...

- C D B C B A C D B A C B C ...

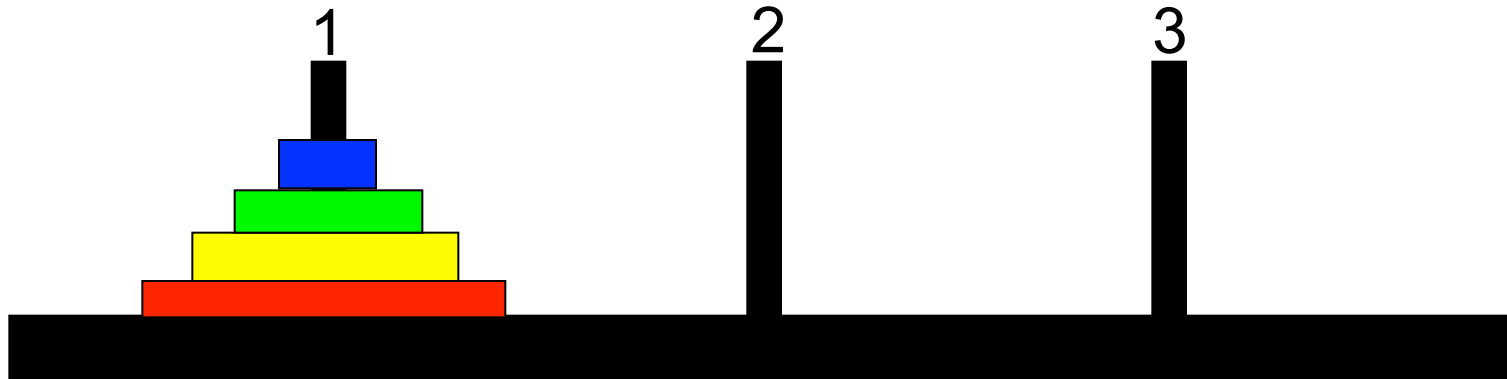
◆ Result is also square-free.

Towers of Hanoi Sequence



- ◆ Optimal play moves small disk on odd moves cyclically $1 \rightarrow 2 \rightarrow 3 \rightarrow 1 \rightarrow 2 \rightarrow 3 \dots$; even moves are then forced.
- ◆ Code moves with six letters as $A[1 \rightarrow 2], B[1 \rightarrow 3], C[2 \rightarrow 1], D[2 \rightarrow 3], E[3 \rightarrow 1], F[3 \rightarrow 2]$
- ◆ Optimal sequence is square-free! (Shallit &c)

Towers of Hanoi Sequence



- ◆ Code moves with six letters as
A[1-→2], B[1-→3], C[2-→1], D[2-→3], E[3-→1], F[3-→2]
- ◆ Optimal play:

A B D A E F A B D C...

- ◆ Easy to generate sequence for infinitely many disks...

Abelian square-free sequences

- ◆ An even stronger notion of "repetition-free" than (ordinary) square-free.
- ◆ A sequence is *abelian square-free* if it contains no two adjacent subwords yy' where y' is a permutation of y (possibly identity permutation).
- ◆ Example:
 `abelianalien`
is square-free but not abelian square-free,
since "alien" is a permutation of "elian".

Infinite ASF sequences exist

- ◆ Thm (Keränen). There exists infinite ASF sequences on four letters.
- ◆ Keränen's sequence based on "magic sequence" S of length 85:
abcacdcbcadcdbdabacabadbabc
bdbcbacbcdcacbabdabacadcdbcda
cdbcbacbcdcacdcdbdcdadbdcba
- ◆ Let $\sigma(w)$ denote word w with all letters shifted one letter cyclically:
 $\sigma(\text{abcacd}) = \text{bcdbda}$

Generating infinite asf sequence(I)

- ◆ Start with Keränen's magic sequence
 $S = \text{abcac...dcbca}$ (length 85)

- ◆ Apply morphism:

$$a \rightarrow S = \text{abcac...dcbca}$$

$$b \rightarrow \sigma(S) = \text{bcdbd...adcdb}$$

$$c \rightarrow \sigma^2(S) = \text{cdaca...badac}$$

$$d \rightarrow \sigma^3(S) = \text{dabdb...cbabd}$$

simultaneously to all letters.

- ◆ Repeat to taste (each sequence is prefix of next, and of infinite limit sequence).

Generating infinite asf sequence(II)

- ◆ Count $i = 0$ to infinity in base 85
- ◆ Apply simple four-state machine to base-85 representation of i (high-order digit processed first).
- ◆ Output $a/b/c/d$ is last state.
- ◆ Requires constant (amortized) time per output symbol.

Dithering with ASF sequence

- ◆ Since Keränen's ASF sequence on four letters is so easy to generate efficiently, we propose using it to dither an iterated hash function.
- ◆ This add negligible computational overhead, and only two new bits of input to compression function.

Recoding with ASF sequence

- ◆ Can also recode message using given ASF sequence. (This is essentially equivalent to dithering, just viewed another way...)

Open Questions

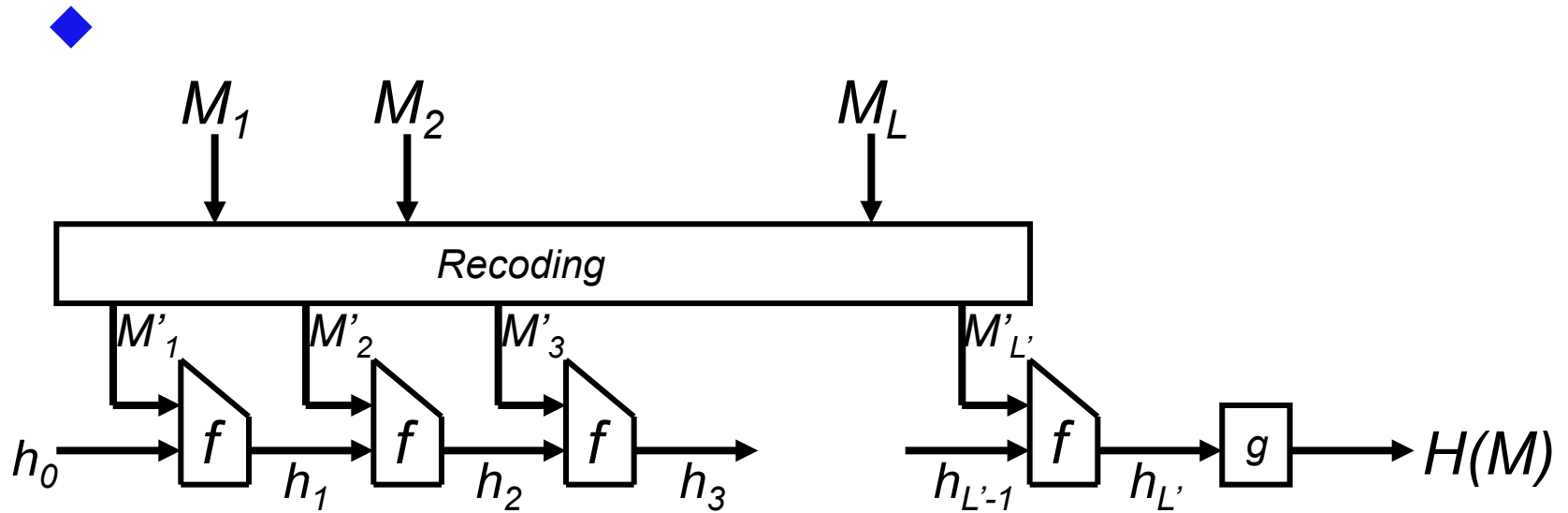
- ◆ Can Dean/Kelsey/Schneier attacks be adapted to defeat use of ASF sequences in hash function?
- ◆ Does ASF really add anything over SF?
- ◆ Are there generalizations of ASF that could be used? ("Even more" pattern-free?)
- ◆ Where else in cryptography can ASF sequences be used?

Conclusions

- ◆ Abelian square-free sequences seem to be a very inexpensive way to prevent repetitive inputs from causing vulnerabilities in hash functions.
- ◆ (Thanks to Jeff Shallit and Veikko Keränen for teaching me about square-free and abelian square-free sequences.)

(The End)

Iterated hashing



Iterated hashing with dithering

