

Lecture 14

Udlearn.17
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Learning halfspaces (linear threshold fctns)

Def. $h(x) = \text{sign}(w \cdot x - \theta)$ is a "halfspace function"

$$\begin{array}{c} \uparrow \\ \text{sign}(x) \begin{cases} +1 & \text{if } x \geq 0 \\ -1 & \text{o.w.} \end{cases} \end{array}$$

Thm Let h be a halfspace over $\{\pm 1\}^n$
then h has Fourier concentration $\alpha(\epsilon) = \frac{c}{\epsilon^2}$

$$\text{(ie. } \sum_{|S| \geq \frac{c}{\epsilon^2}} \hat{h}(S)^2 \leq \epsilon)$$

(Will prove soon)

Corr low degree algorithm learns halfspaces under
uniform distribution with $n^{O(1/\epsilon^2)}$ uniform samples.

(Actually can learn in $O(n^5)$ but we'll get a "big win"
from this approach soon...)

Key idea: Noise sensitivity

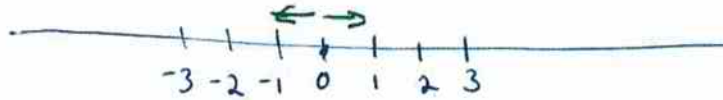
3) $f(x) = \text{Maj}(x_1, \dots, x_n)$

$n S_\epsilon(f) = O(\sqrt{\epsilon})$

Sketch

• $\text{Maj}(x) \sim$ random walk on line starting at 0

$P(L) = \frac{1}{2} \quad P(R) = \frac{1}{2}$



$X_i = 1 \Rightarrow R \times 2$

$X_i = -1 \Rightarrow L \times 2$

$\sum X_i =$ end pt of walk (if start at 0)

eg. $x = (11 \rightarrow 11 \rightarrow 11)$



ends at 0

Equivalent process:

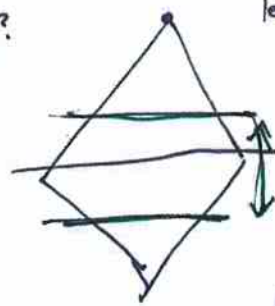
pick random pt on hypercube + output the level

Well Known fact:

expected distance from startpt $(|X_1 + X_2 + \dots + X_n|)$ after n steps is \sqrt{n}

+ likely to be close to expectation

why?



level n

hypercube

level 0

most nodes w/in \sqrt{n} of level 0

level $-n$

• $N_\epsilon(x) \sim$ random walk on ϵn bits each has twice value later

expected displacement $2\sqrt{\epsilon n}$

↑ since change
+1 \rightarrow -1
-1 \rightarrow +1

Another view: • take walk according to x
• then take walk according to $N_\epsilon(x)$

6) any f :

Thm $f: \{\pm 1\}^n \rightarrow \{\pm 1\}$

$$NS_\epsilon(f) = \frac{1}{2} - \frac{1}{2} \sum_s (1-2\epsilon)^{|s|} \hat{f}(s)^2$$

note for parity fctns, this gives

$$NS_\epsilon(\chi_s) = \frac{1}{2} - \frac{1}{2} \cdot (1-2\epsilon)^{|s|} \quad \checkmark$$

Pf.

$$NS_\epsilon(f) = \Pr_{x, y \in N_\epsilon(x)} [f(x) \neq f(y)]$$

$$= E_{x, y \in N_\epsilon(x)} \left[\frac{1 - f(x)f(y)}{2} \right]$$

$$= E_{x, y} \left[\frac{(f(x) - f(y))^2}{4} \right] = \frac{1}{4} E \left[\underbrace{f(x)^2 + f(y)^2}_{\text{always 1}} - 2f(x)f(y) \right]$$

$$= \frac{1}{2} - \frac{1}{2} E[f(x)f(y)]$$

$$= \frac{1}{2} - \frac{1}{2} E \left[\sum_s \hat{f}(s) \chi_s(x) \sum_T \hat{f}(T) \chi_T(y) \right]$$

$$= \frac{1}{2} - \frac{1}{2} \sum_{s, T} \hat{f}(s) \hat{f}(T) \underbrace{E[\chi_s(x) \chi_T(y)]}_{=0 \text{ if } s \neq T}$$

show using standard techniques using pairing on $i \in SAT$

but, if $s = T$?

$$E_{x, y} [\chi_s(x) \chi_T(y)] = 1 \cdot \Pr[\chi_s(x) = \chi_T(y)] + (-1) \cdot \Pr[\chi_s(x) \neq \chi_T(y)]$$

$$= 1 - 2 \Pr[\chi_s(x) \neq \chi_T(y)]$$

$$= \frac{1}{2} - \frac{1}{2} \sum_s (1-2\epsilon)^{|s|} \hat{f}(s)^2$$

$$NS_\epsilon(\chi_s) = \frac{1}{2} - \frac{1}{2} (1-2\epsilon)^{|s|}$$

$$= (1-2\epsilon)^{|s|}$$

□

Back to $\frac{1}{2}$ spaces:

Corr 2 for $\frac{1}{2}$ space $h: \{\pm 1\}^n \rightarrow \{\pm 1\}$

$$\sum_{|s| \geq \alpha(\frac{1}{\epsilon})} \hat{f}(s)^2 \leq \epsilon$$

PF

for $\frac{1}{2}$ space $ns_\epsilon(h) \leq 8.8\sqrt{\epsilon}$ so use

$$\beta(\epsilon) = 8.8\sqrt{\epsilon}$$

$$\beta^{-1}(x) = \left(\frac{x}{8.8}\right)^2$$

so for $m \equiv \frac{1}{\left(\frac{\epsilon}{20}\right)^2}$

$$\beta^{-1}\left(\frac{\epsilon}{2.32}\right) \leq \left(\frac{\frac{\epsilon}{2.32}}{8.8}\right)^2$$

$$\leq \left(\frac{\epsilon}{20}\right)^2 \quad (20.416)$$

□

⇒ Can learn any $\frac{1}{2}$ space over uniform
 with $n = O\left(\frac{1}{\epsilon^2}\right)$ random examples
 but, can do a lot better!
 still, technique extends...

Learn any fctn of k $\frac{1}{2}$ spaces:

h_1, \dots, h_k are k half spaces

$g: \{\pm 1\}^k \rightarrow \{\pm 1\}$ Boolean fctn

$f: \{\pm 1\}^n \rightarrow \{\pm 1\}$ st. $f(x) = g(h_1(x), \dots, h_k(x))$

Thm $ns_\epsilon(f) = 8.8k\sqrt{\epsilon}$