TO: MSPM Distribution
FROM: P. G. Neumann
SUBJECT: BF.2.22, The Registry File Maintainer
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The attached BF.2.22 is a major revision of the former issue, 08/14/67. Many of the calls have been redesigned, although the module is functionally the same.

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# <u>Identification</u>

The Registry File Maintainer S. I. Feldman

#### Purpose

Every I/O device known to a system has an associated file called its Registry File (RF). The Registry File Maintainer (RFM) is called by outer modules and other interested users to get information from or to store information into a Registry File. This section describes the functions of Registry Files, describes the calls to the RFM and the implementation thereof.

## <u>Registry Files</u>

Protection of all I/O devices is implemented by file system protection of their Registry Files. Non-volatile information about the device is stored in this file for use by the I/O System. The following is a brief discussion of the form, content, and use of these files.

The RFs are organized by type of device into "type" directories. These type directories are immediately inferior to the Registry File Directory (path name ">rfd"). A Registry File is identified by two 32-character strings, called the "type" and the "name". The type is the entry name of the type directory in the Registry File Directory; the name is the entry name of the RF in its type In addition to the regular RFs, these directory. tvpe directories also contain certain other files: an 1/0 Assignment Table (IOAT), normally accessible from the hardcore ring only by the 1/0 Assignment Module (10AM); see BF.2.26). This table lists names of the users who possess and control the file and its associated devices at any time. There may be several read-only files in the directory for use in storing certain critical RF-related information that cannot be left in the normal RFs for security reasons. Finally, there are usually one or more normal Registry Files called prototype RFs. These files are used to created new RFs of the given type.

The normal RFs are made up of several distinct parts. The first is of standard form and content. This part of the RF is used by the Attachment Module (see BF.2.23) and the Registry File Maintainer (RFM; see below). This fixed part has information on the RF, and has "pointers" to certain RFs. (These pointers consist of the type and name of the target RF). Each RF has a "level" number associated with it. The smaller the level number, the more directly related the device is to the GIOC. Thus, a typewriter channel has level equal to 1, while the typewriter has level equal to 2. "Down" is defined to be the direction of decreasing level numbers. The up and down chaining represents the connections between devices. A file may point to more than one up and one down RF, although this is not the usual case.

Some of these links between RFs are permanent, representing physical wiring. Others are transient, and change with the configuration or for other reasons. See below for a discussion of logical channels. As an example of a temporary link, a typewriter is associated with a channel only while it is dialed in. The association is broken when it hangs up. The following is a plausible RF arrangment for tapes: The level 4 RF represents a tape reel. There is a temporary association (down link) with the level 3 RF for the tape drive upon which the reel is presently mounted. There is also a temporary up link from the drive to the reel file. There is a permanent link down from the level 3 drive RF to the level 2 controller RF. There are many permanent up links from the controller, one for each of the drives connected to it. There are also some temporary down links from the level 2 controller file to the level 1 files for the High Performance Channel RFs. There is a temporary up link from each of these channel files to the appropriate controller file. The association between controller and channel can be changed by using the peripheral switch; any modification will be reflected in the Device Configuration Table. (See the discussion of logical channels.)

The fixed part of the RF contains, in addition to names of RFs in the chain, information to guide the Attachment Module in following the list. There is also information on the type of DCM to be used, and part of the mode string to be passed to that DCM. A RF has at least one device associated with it. If several devices are closely related, they may be considered as a single resource and may therefore share a single RF. For example, a full-duplex typewriter channel can be implemented by connecting two half-duplex typewriter GIOC channels to a single data set. There would be a single RF for the pair of channels, but there is information on both devices separately in the RF.

Each device has a "device profile" in the RF. Once the size is set (either when the RF is created or when the first <u>rfm\$set profile</u> call is made), it does not change. This profile is meant to contain relatively constant physical device information. For example, the tab settings of typewriters will be stored in the profile. However, data as changeable as the line and column number are not stored in the profile. A tape profile might include information on density and the names of other reels that make up a single multi-reel file.

In addition to the profile, there is a behavior log for each device. This log usually contains information on faulty performance by the device. Whenever such a thing happens, the appropriate I/O System module calls <u>rfm\$add log</u> to store this the information at the end of the chain of log entries. log Each entry contains the identification of the process and user in control, the time, and an arbitrary bit string supplied by the caller. The log can be read and entries can be deleted bv appropriate calls to the RFM. The log itself is a chained list bing a relative pointer to the next in of structures a

the list. The fixed part of the RF contains a relative pointer to the first and last blocks on the chain.

At some installations, the assocation between certain devcies and certain GIOC channels may depend upon the settings of the peripheral switch. Registry Files have the ability to handle such a connection. When the Attachment Module finds that the "logchans" bit is on in the RF, it calls a hardcore supervisor procedure, the Device Configuration Table Manager, which will translate a "logical channel" name stored in the RF to the real RF name. The Attachment Module then stores that information in the RF and proceeds as if that were the down name stored in the RF normally.

Above, we mentioned the existence of certain special read-only files in the type directories. These files contain information relating to Universal Device Manager processes. Such a process is a system process that can handle several devices of a given type. The Attachment Module needs to know the process group id of the relevant UDMP and also the name of a certain data base used by the process, called its PDT. This information is placed in the special read-only file. Only certain types of RFs require these files. These are the last RFs examined by the Attachment Module while it traces through a chain of files. Typically, that file represents the channel or controller. For the RF with name "X", there is a file with name "X\_ro" in the same type directory. Since several devices can use the same read-only file, a single file will most likely have several names.

# Registry File Declaration

The following is the EPL declaration for a Registry File. The first two declarations, rf and rfx, together form the fixed part of the RF discussed above; there are two parts for implementation reasons. The third declaration, rf\_ro, is the special read-only file which exists for reasons of system security.

dc1 2	<pre>1 rf based(rfp), level fixed bin(35), """"""""""""""""""""""""""""""""""""</pre>	<pre>/*level=1 for a GIOC channel, 2 for a device connected directly to a GIOC channel, etc.*/</pre>
2	<pre>force_udmp bit(1),     ""</pre>	<pre>/*if 1, force the use of a universal     device manager process*/</pre>
2	<pre>in_use_switch bit(36),</pre>	<pre>/*set ON at attach time and OFF    at detach time*/</pre>
2	hangupable bit(1),	/*if ON, device can hang up*/
2	logchans bit(1), "	<pre>/*if ON, the down_names for this     device are to be filled in by a call</pre>
2	allocate bit(1),	to the hardcore ring to get the present RF name corresponding to the logical channel names. If this bit is UN, no more RFs are to be searched.*/ /*if ON, Reserver should be called

11 with each resource\_name as argument.\*/ 2 temp\_link bit(1), /\*connection with next file is 11 temporary. Blank out down name entries upon detachment\*/ /\*number of entries in up array\*/
/\*number of entries in down array\*/ 2 nup fixed bin(35), 2 ndown fixed bin(35), 2 ndev fixed bin(35), /\*number of entries in devices array\*/ 2 ntypes fixed bin(35), /\*number of entries in att\_types array\*/ 2 present\_type\_index fixed bin(35), /\*index in att\_types array of 11 type with which device was last attached\*/ 2 down\_slot fixed bin(35), /\*position of upname for this file in up array of next registry file\*/ 2 alloc\_type char(32), /\*use this type in calls to the Reserver alloc\$resource 11 entry\*/ 2 lock bit(144), /\*for locking RF when threading 11 or deleting behavior log entries or 11 modifying the profile.\*/ /\*registry files pointing to this one\*/ 2 up(rfp->rf.nup), 3 uptype char(32), 3 upname char(32), 2 devices(rfp->rf.ndev), /\*entries for devices associated with this registry file\*/ 3 resource\_name char(32), /\*name used in calls to the Reserver and the Device Assignment Module\*/ 3 profile\_relp bit(18), /\*relp to device profile for this ... device\*/ 3 profile\_length fixed bin, /\*number of bits in this profile\*/ 3 oldest\_log\_relp bit(18), /\*relp to oldest entry in behavior log\*/ 3 newest\_log\_relp bit(18), /\*relp to most recent entry in behavior log\*/ /\*number of entries in behavior log\*/ 3 nlog fixed bin, 3 device\_type fixed bin(35), 2 rfxrelp bit(18), /\*relative pointer to RF extension\*/ 2 free\_storage area((15000)); /\* \*/ dcl 1 rfx based(rfxp), 2 att\_types(rfp->rf.ntypes),/\*special information for each type by which this device may be known\*/ 3 type\_name char(32),

- 3 ccm\_type char(32),
- 5 trace\_down bit(1),
- 3 alloc\_down bit(1),

11

11

- /\*name of code conversion driving
   table to be used\*/
  /\*if ON, trace down to next registry
   file. Otherwise, stop here\*/
- /\*if ON, must call Reserver to allocate a device of type down\_type, and use returned resource\_name as down\_name(1).

11 In either case, find next RF by ... using down type and down\_name(1)\*/ /\*keep tracing down to other KFs 3 look\_only bit(1), 11 under trace\_down control, but 11 only to compute code conversion .. driving table name\*/ /\*used as described above\*/ 3 down\_type char(32), 3 down name(rfp->rf.ndown) char(32), /\*used as described above\*/ 3 logical\_channel (rfp->rf.ndown) char(32), /\*array of 11 names to be used in call to get present 11 equivalent RF name from 11 info in DCT. Used only if the 11 logchans bit is on\*/ /\*character string to be 3 extra\_mode\_char(32), concatenated with mode to be 11 11 passed to DCM\*/ /\*used as type in attach call to 3 dcm\_type char(32), 11 DCM if trace\_down is OFF or 11 look only is ON\*/ 3 dcm\_name char(32); /\*used as ioname2 of attach call to 11 DCM if trace\_down is 11 OFF or look\_only is Ol'\*/

\*/ dcl 1 rf\_ro based(p), 11 11 11 .. 11 .. 11

/\*

/\*special Registry File. There is a file of this format associated with each regular RF, with name equal to the name of the normal RF concatenated with "\_ro". This file contains certain data that must be protected against tampering and is therefore read-only to most users.\*/ /\*name of PDT in DMP\*/ 2 udmp\_user\_id char(50); /\*user\_id of universal device manager

for this device, if any\*/

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Calls and Arguments
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2 pdt\_name char(32),

call rfm\$get\_devices(type,name,device\_types,resource\_names,cstatus); call rfm\$set\_profile(type,name,devnumber,dataptr,nbits,cstatus); call rfm%get profile(type,name,devnumber,dataptr,nbits,cstatus); call rfm\$get\_nlog(type,name,nlogs,cstatus); call rfm\$add\_log(type,name,devnumber,dataptr,nbits,cstatus); call rfm\$delete\_log(type,name,devnumber,first,number,cstatus); call rfm\$read\_log(type,name,devnumber,first,number,infoptr,cstatus); call rfm\$get\_ups(type,name,uptypes,upnames,nreturned,cstatus); call rfm%get\_down(type,name,down\_type,down\_names,cstatus); call rfm\$link(toptype,topname,att\_type\_index,down\_index, bottomtype,bottomname,upindex,cstatus);

declare

.

.

type char(*).	/*type of device*/		
name char(*),	/*name of device*/		
toptype char(*),	/*type of device with larger level number		
topname char(*).	<pre>/*name of that device*/</pre>		
bottomtype char(*).	/*type of device with smaller level		
II	number*/		
bottomname char(*).	/*name of that device*/		
devnumber fixed bin.	/*index in devices array for		
II	this device*/		
latantr otr	/*nointer to bit string of		
n n	length phits into which profile		
11	is to be stored or from which profile		
11	or log entry is to be conject*/		
phits fixed him	/toumber of bits in bit string See		
nores rixed offi,	above*/		
first fixed bin	/tindex of first behavior log entry		
	(starting from the oldest) which		
11	is to be read or delated*/		
number fixed bin	Athen many log entries are to be		
	read or doletodt/		
device $types(*)$ fixed bin	/tarray into which dovice types are		
uevice_cypes(") fixed offi;	to be stored +/		
resource names(*) char(32)	/tarray into which the resource names		
	are to be stored		
infontr ptr	/the inter to the following structuret/		
1 info(number) based(infontr)	/*pointer to the forfowing structure*/		
n	contain information on the behavior		
1	log ontring. The length of this prray		
II.	tells the DEM how many to read out #/		
$2^{\circ}$ time fixed bin(71)	/time when REM stored log ontrut/		
2 prop id bit(36)	/*cline when Kim Stored Tog entry*/		
	log ontrut/		
2  usor id char(50)	/tusor id of above usert/		
2 user_id Char(50),	/*user to of above user*/		
2 dotootr otr	/*number of bits in log entry*/		
2 using $(+)$ char(32)	/*pointer to the log entry*/		
uprypes(*) = char(32),	/*iist of uppers for the KF*/		
nlogs(*) fixed bin.	/*list of number of behavior log		
	entries for each device*/		
down type char(*)	/*down type for RE*/		
down name( $\star$ ) char(32).	/*list of down names*/		
down index fixed bin.	/*index in down name array where		
	down name is to be stored*/		
up index fixed bin	/*index in upnames array where up type		
	and up name are to be stored*/		
nreturned fixed hin.	/*number of upnames returned*/		
niecainea iixea biny	y manual of apranes recurrically		
Implementation of Calls			
THE CHARGE CALLS			
get devices			

This call is made to get a list of device types and resource names from a given RF. In response to the call, the following steps are taken:

1. If the RF with type <u>type</u> and name <u>name</u> is non-existent, set bit 1 of <u>cstatus</u> and return. If the file is not readable by this user, set bit 8 of <u>cstatus</u> and return.

2. Store as many device\_types as possible in the <u>device types</u> array. If there are more in the RF, set bit 18 of <u>cstatus</u>. If the array is larger than the array in the RF, fill in the remaining elements with zeroes.

3. Store as many resource names from the RF into the <u>resource names</u> array as possible. If there are more in the RF, set bit 18 of <u>cstatus</u>. If the array is larger than the array in the RF, fill in the remaining elements with blanks.

4. Return.

#### <u>set profile</u>

This call is used to store a device profile. In response to the call, the following steps are taken:

1. Find the RF with the given type and name. If the file does not exist, set bit 1 of <u>cstatus</u> and return. If the file is not readable by this user, set bit 8 of <u>cstatus</u> and return. If the file is accessible but not writable by this user from the caller's validation level, set bit 4 of <u>cstatus</u> and return.

2. Lock the RF using rf.lock as a lock structure.

3. If <u>devnumber</u> if less than one or greater than rf.ndev, set bit 2 of <u>cstatus</u>, unlock the RF, and return.

above arguments are all right, then if lf the 4. rf.devices(devnumber).profile\_relp is not zero, go to step 5. the Otherwise. nbits store in . rf.devices(devnumber).profile\_length, allocate a bit string of the appropriate size in the area, and store the relp to that bit string in rf.devices(devnumber).profile\_relp. If the area is not large enough, set bit 7 of cstatus, unlock the RF and return.

5. Store the bit string (length <u>nbits</u>) pointed to by <u>dataptr</u> in the appropriate device profile and return. The rules of bit assignment relating to padding and truncating apply.

# get profile

This call is made to read out a profile. The following steps are taken in response to the call:

1. If the file is non-existent ,set bit 1 of <u>cstatus</u> and return. If this user does not have read permission for the file, set bit 8 of <u>cstatus</u> and return. 2. If devnumber is less than one or greater than rf.ndev, set bit 2 of <u>cstatus</u> and return.

3. Lock the RF using rf.lock as a lock structure.

4. If the arguments are valid, then if rf.devices(devnumber).profile\_relp is zero, set the bit string pointed to by <u>dataptr</u> (length <u>nbits</u>) equal to the null string, unlock the RF, and return.

4. Set the hit string described above equal to the <u>devnumber</u>th device profile using bit string assignment rules. Unlock the RF and return.

<u>get nlog</u>

This call is made to find out how many behavior log entries for the devices associated with the RF. In response to the call, the following steps are taken:

1. If the RF with the given type and name is non-existent, set bit 1 of <u>cstatus</u> and return. If this user does not have read permission for the file, set bit 8 of <u>cstatus</u> and returrn.

2. Store as many of the nlog entries in rf.devices into the <u>nlogs</u> array as possible. If the argument array has more elements than there are devices, fill in the rest of the array with zeroes. If there are elements in the RF that have not been returned, set bit 18 of <u>cstatus</u>.

3. Return.

add log

The following call is made to add a behavior log entry for a device associated with the file. In response to the call, the following steps are taken:

1. If the RF with the given <u>type</u> and <u>name</u> is non-existent, set bit 1 of <u>cstatus</u> and return. If this user does not have read permission for he file, set bit 8 of <u>cstatus</u> and return. If the file is not writable from the caller's ring, set bit 4 of <u>cstatus</u> and return.

2. Lock the RF.

3. If <u>devnumber</u> is less than one or greater than rf.ndev, set bit 2 of <u>cstatus</u>, unlock the RF and return.

4. Otherwise, allocate a structure like the following in the area in the RF:

dcl l log\_entry based(p),

2 next\_relp bit(18), 2 time fixed bin(71), 2 proc\_id bit(36), 2 user\_id char(50), 2 nbits fixed bin, 2 data bit(nbits);

See the declaration of all arguments of RFM calls for the meaning of the elements of the structure. If the area is not large enough, set bit 7 of <u>cstatus</u>, unlock the RF, and return.

5. If the allocation succeeds, copy the bit string of length <u>nbits</u> pointed to by <u>dataptr</u> into the appropriate part of the structure. Store the present process id, user id, and timer value in the structure, as well as <u>nbits</u>. Set the next\_relp equal to zero. If rf.devices(devnumber).oldest\_log\_relp is zero, set it and the corresponding newest\_log\_relp equal to the offset of the newly allocated structure. If the oldest relp is nonzero, to bv then set the relp in the structure pointed rf.devices(devnumber).newest\_log\_relp equal to the offset of the new structure, and then store that same offset ⇒in newest\_log\_relp.

6. Increment rf.devices(devnumber).nlog by one, unlock the RF, and return.

### d<u>elete lo</u>g

This call is used to delete a set of behavior log entries for a particular device associated with a Registry File. In response to such a call, the following steps are taken:

1. If the RF with given <u>type</u> and <u>name</u> is non-existent, set bit 1 of <u>cstatus</u> and return. If the file is not readable by this user, set bit 8 of <u>cstatus</u> and return. If the file is not writable from the caller's ring, set bit 4 of <u>cstatus</u> and return.

2. Call the Locker and lock the RF using rf.lock structure.

3. If <u>devnumber</u>, <u>first</u>, or <u>number</u> is less than one or if <u>devnumber</u> is greater than rf.ndev or if the sum of <u>first</u> and <u>number</u> is greater than rf.devices(devnumber).nlog, then unlock the RF, set bit 2 of <u>cstatus</u> and return.

4. Otherwise, follow the chain of relps for the behavior log of the device. Starting at the <u>first</u> element (counting the oldest link on the chain as number 1), free <u>number</u> of them. Make the next\_relp of the last entry not freed point to the one after the gap, or set it to zero if it is now the last entry on the chain. Modify oldest\_log\_relp and newest\_log\_relp in the devices array of the RF as necessary.

5. Unlock the RF and return.

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# <u>read log</u>

This call is used to read out a set of log entries (non-destructively). In response to the call, the following steps are taken:

1. If the file is non-existent, set bit 1 of <u>cstatus</u> and return. If the file is not readable by this user, set bit 8 of <u>cstatus</u> and return.

2. If <u>devnumber</u>, <u>first</u>, <u>or number</u> is less than one, or if <u>devnumber</u> is greater than rf.ndev or if the sum of <u>first</u> and <u>number</u> is greater than rf.devices(devnumber).nlog then set bit 2 of <u>cstatus</u> and return.

3. Chase the chain of relps until the entry with index equal to <u>first</u> is found. Copy out the contents of the log entry into the corresponding elements of the info array. (The next\_log relp is not copied and a pointer to the data string is stored in dataptr). If there are more entry logs on the chain, set bit 13 of <u>cstatus</u>, Return.

### <u>get ubs</u>

This call is used to get the uptypes and upnames (identifications of the RFs that point down to this one). In response to the call, the following steps are taken:

1. If the file is non-existent, set bit 1 of <u>cstatus</u> and return. If the file is not readable by this user, set bit 8 of <u>cstatus</u> and return.

2. Copy as many elements of the uptype and upname arrays (elements of rf.up) into the <u>uptypes</u> and <u>upnames</u> character array arguments. Set <u>nreturned</u> equal to the number of complete pairs of RF names. If there are more than this, set bit 18 of <u>cstatus</u>.

3. Return.

## get down

This call is used to find the next RF in a chain, assuming they have been previously linked. In response to the call, the following steps are taken:

1. If the RF with type <u>type</u> and name <u>name</u> is non-existent or inaccessible, set bit 1 of <u>cstatus</u> and return.

2. If rf.present\_type\_index is not greater than zero and less than or equal to rf.ntypes, set bit 5 of <u>cstatus</u> and return.

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of the corresponding down\_names into the <u>down names</u> array. If there are more element of the argument array than down names, fill in the extra elements with blanks. If there are more down names, set bit 18 of <u>cstatus</u>. Return.

## link

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This call is used to link up two Registry Files, an upper one and a lower one. This call is meant for use by certain levice Control Modules. In response to the call, the following steps are taken:

1. Find the Registry File with type <u>toptype</u> and name <u>topname</u>. If the file does not exist, set bit 1 of <u>cstatus</u> and return. If the file is not readable by this user, set bit 8 of <u>cstatus</u> and return. If the file is not writable by this user from the caller's ring, set bit 4 of <u>cstatus</u> and return.

2. Do the same checking for the file with type <u>bottomtype</u> and name <u>bottomname</u>.

3. If any of the following conditions holds, set bit 2 of <u>cstatus</u> and return:

<u>upindex</u> is less than one or greater than rf.nup in the lower file.

<u>down index</u> is less than one or greater than rf.ndown in the upper file.

<u>att type index</u> is negative or greater than rf.ntypes in the upper file. If that argument is zero, then if rf.present\_type\_index is less than one or greater than rf.ntypes.

4. If <u>att type index</u> is non-zero, store it in rf.present\_type\_index for the upper file. Call the value of the present\_type\_index N.

5. Store <u>bottom name</u> in rfx.att\_types(N).down\_name(down\_index) in the upper file.

6. Store <u>bottomtype</u> and <u>bottomname</u> in the corresponding elements of rf.up(upindex) in the lower file.

7. Return.

# Summary of Cstatus Bits

- 1 File non-existent
- 2 Number out of range
- 3 Typename not found
- 4 File not writable (<u>set profile</u>, <u>add log</u>, <u>delete log</u>, and <u>link</u> calls only)
- 5 Unlinked Registry File (<u>get down</u> call only)
- 6 Non-RFM unexpected error

7 Area too small (<u>add log</u> and <u>set profile</u> calls only) 8 File not readable by this user 18 More data available (array not large enough)

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