Published: 09/28/66

Identification

"Entry variables" in PL/I. Fake_entry\$create, fake_entry\$call. D. B. Wagner

Purpose

A problem with the PL/I language is that there is no such thing as an "entry variable," that is, for example, there is no equivalent to the MAD sequence.

FUNCTION NAME A

A = SIN.

• • •

B = A. (C)

This causes trouble in a number of places: in particular in the Request Dispatcher (described in BY.6.01). The most reasonable design for the Request Dispatcher involves a calling argument which is an array of entries.

The library routines fake_entry\$create and fake_entry\$call provide a way of getting around this difficulty.

<u>Usage</u>

The calling forms are:

a = fake_entry\$create (q);
call fake_entry\$call (a,...);

The ellipsis represents an arbitrary sequence of arguments. \underline{A} and \underline{g} are declared as follows:

dcl q entry,

a bit (216); /* 216=6*36 */

fake_entry\$create bears a family resemblance to <u>unspec</u>: it takes an entry and returns a bit-string representation of it. Fake_entry\$call is used to call the entry represented by such a bit-string, giving any arguments the user desires.

<u>Implementation</u>

The bit-string \underline{a} above is 6 words long, and these words are used as follows. The first and second contain an ITS

pair pointing to the entry, the third and fourth contain an ITS pair representing a stack level (always present, but not used unless **g** is an internal procedure), and the fifth and sixth contain validating information (not yet specified). This is precisely the same form as a PL/I label variable.

It should be noted that since a fake entry contains a stack pointer, one should be chary of putting it into static or controlled storage.