

Notes for talk

8/12/69

History of Multics, reason for doing, w/r CTSS problems  
Length of Project, current status. (brief)

Follow hardware lines; talk on hardware

DSU-10 → DSU-220; capacity as you need it  
plan (Temporary) to bring in 2314-style (DSU-120)  
to get out of DSU-10 mode, if necessary.

Development work in earnest and for testing; vegetation  
earliest underway as to its future status after 12/71.

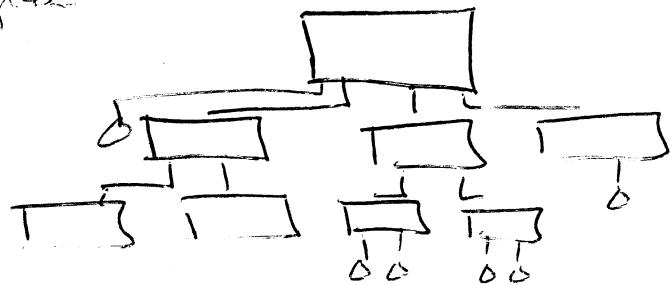
Block configured for M37 TTY, 1050/2141, AKA's  
possible address coin: TN 300, M35/33. Dura Corp.)  
(only thing going for the 33 is that everyone seems taken on)

### System Organization

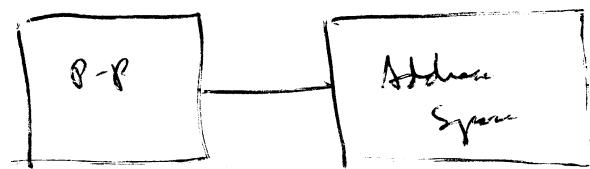
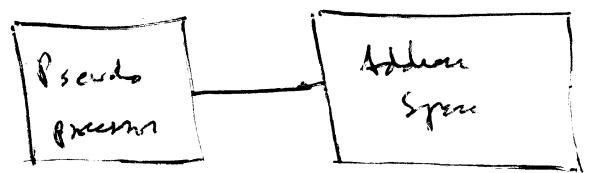
Hardware can grow to 3-4 times present capacity; organization  
allows for a module to be plugged in one at a time.

## System Organization

### 1. Hierarchically Organized System



### 2. Multi process environment



a. Address space can overlap.

b. One gets all files by mapping them into his own address space

c. Supervisor is a set of processes appearing in your address space; it is shared.

d. Interprocess communication (control) is provided

3. Reliability is not quite up to CTSS, but it is still very usable. A "stable" system will have 1-2 crashes/day (probably software bugs). Occasional hardware instability will keep system down for a day. New software systems sometimes are buggy. (Virtue: many bugs only kill one user.)

Our  
Plans.

IFC has committed it self to run Multics for at least 3 years; GE is solidly behind it with probable new hardware, etc., in future.

CSK Gray will continue to upgrade it; not a on-site expert in system.

ANPA Net, Display Support, Upgrade hardware, new functions (e.g., alescarf)

Documentation

~~User Manual~~

Programmers Manual (In Draft) (Good Draft until Oct. 1)

Condensed Guide

Subsystem Writers Guide (Available Drafts of 1-9; 1 coming, publication this year)

Papers (5 coming this fall)

Internal - MS(M), etc.

FC/I : Mission : - I/O, File system  
(One Multics call to achieve  
much more sophisticated semantics)  
- Shared fixed point (boundary)  
- Turing

Most everything else is there, implemented, and working.  
e.g., structures, pointers, ~~—~~, and some  
type conversion, string functions (quite  
good implementation) etc.

This is our 2<sup>nd</sup> computer, and it benefitted a lot —  
It performs very good code.

Performance, etc.

Depends on class of cows.  
 Today. { P D F . 8 semidwarf wheat that grows well on  
 Daffy range 20 very heavy ears.  
 $20 \xrightarrow{\text{P}} 30$

I expect to see the 30 and climb to the 50-60  
area during next year. 20 might go to 25 at most, since  
limitation is outside of experience; aspect of tools.

Cancor scheduler is biased toward heavy user, and  
~~is profit oriented~~.

give away care like it was Kleenex. Result: miserable

response on small jobs. This will be fixed by out.

### Disadvantages

1. GE is a different path than IBM.

Community of users consists of

GE 625, 635 users

GE 605 users

Future customers of 645

(Several have shown interest in port, but

have adopted a "wait for the software" attitude.)

PL/I communication possible  $\in$  IBM

ASCTI character set maps into EBCDIC with some care. (Data)

FORTNIXON communication possible

2. System is new; there are lots of rough edges to be polished and human engineering to be worked out before we are satisfied - (But system is worth using)

One process;

Registration rings;

Ring-zero contains supervisor

ring-1 contains command language library, library-

User can invoke additional rings if he wants.

All commands are independent of previous programs.

Type ~~"x"~~; ~~it~~ describes my directions,

then library, for

something named X, then

use it in, then call it.

{ call "x", taken } ~~done done~~.

Sharing / See bottom

Languages:

SYNTRAN ~~is~~ complete, working, has been used  
only a small amount, will probably show up with minor  
bugs. No NAMELIST; some simple, short  
~~short~~ ~~free~~ free-format I/O will be provided so to  
supplement the usual Fortran-like I/O.