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Identification

Segment Loading Table Manager D. H. Johnson

Purpose

The segment loading table (SLT) is accessed directly by Multics Initializer procedures to obtain information about segments. However, there are several frequently used examinations of the SLT which involve more than a simple one item reference. The SLT manager includes utility procedures which provide these queries. Also included is a procedure to add an entry to the SLT.

Segment loading table manager procedures

The segment loading table manager consists of the following five procedures.

- build entry Given a logical header record describing a 1. segment, this procedure assigns a segment number and constructs an SLT entry.
- get seg ptr Given a segment name, this procedure returns a 2. pointer to the base of the segment.
- get seg name Given a pointer to a segment, this procedure 3. returns the name of the segment.
- 4. get text seg ptr - Given a pointer to a linkage section, this procedure returns a pointer to the base of the associated text segment.
- 5. get link seg ptr - Given a pointer to a text segment, this procedure returns a pointer to the associated linkage section.

These procedures are described in more detail below. The following PL/I statement defines the parameters to all of the SLT manager procedures.

dcl seg ptr ptr, /* pointers to segments */ text seg ptr ptr, link seg ptr ptr, seg name char(*), /* segment name */ er ret label, /* error return */ header ptr; /* pointer to logical header array */

The logical header array is similar to an SLT entry (see MSPM BL.2.01). It consists of:

SECTION BL.2.02

entry words 1 - 5 number of segment names segment names (8 words for each name) size of path name in characters path name

1. call slt manager\$build_entry(header,seg_ptr,er_ret) ;

This procedure builds an entry in the SLT for the segment described by the data at argument header. The information at header is in the format of a logical header record on the Multics System Tape (MSPM BL.1.01). This information indicates whether the segment it describes is an initialization or hard-core Either the SLT item last sup seg or supervisor segment. last init seq value is incremented by one and used as the segment number (and SLT index) for the segment. The information at header is copied into the SLT entry exactly as it appears except for the segment names and path name. They are placed in the auxiliary segment name seg and pointers to the names are stored in the SLT entry. The segment number is used to construct and return a pointer to the base of the segment in seg ptr. If the SLT entry can not be built, control returns to the label er ret.

2. call slt manager\$get seg ptr(seg name, seg ptr,er ret) ;

This procedure attempts to return a pointer to the base of a segment identified by the name given in the argument seg_name. It searches the SLT entries in turn looking for the specified segment name. If the name is found, the SLT entry index is used to construct a pointer to the base of the segment. The pointer is placed in seg_ptr. If the name could not be found in the SLT, control transfers to the label er ret.

3. call slt manager\$get seg name(seg ptr, seg name, er ret) ;

This procedure checks whether the segment number given in seg ptr has an SLT entry. If it doesn't, control returns to the label er ret. Otherwise, the SLT entry is examined and the <u>first</u> name of the segment is placed in argument seg name.

This procedure is given a pointer to a linkage section and is requested to return a pointer to the associated text segment. The SLT entry specified by the linkage section segment pointer is examined.

If the entry item link sect sw is OFF, the segment is not a linkage section. Control transfers to the error return, er ret.

If the linkage section segment is not one of the special combined linkage segments as indicated by item link sect status, a pointer to the text segment is constructed from the SLT entry item text link segno. Control then returns to the caller.

If the linkage section segment is a combined linkage segment for active, loaded, or wired down hard-core supervisor segments, the definition pointer is obtained from the header in the linkage section pointed to by link seg ptr. The definition pointer is assumed to point to either the text segment or the original linkage section segment that was loaded from the system tape. The SLT entry for the segment referenced by the definition pointer is examined. If the entry indicates that this segment has combined linkage, i.e., linkage sw is set ON, and its link section and link offset entry items match link seg ptr, then the definition pointer is used to construct a pointer to the base of the text segment. If, however, the SLT entry pointed at by the definition pointer indicates that the segment is a normal section, i.e., link sect sw is set linkage ON, the text link segno item is used to construct a pointer to the text sequent. If the SLT entry pointed to by the definition pointer fails these tests, control is transferred to the error return, er ret.

This procedure attempts to return a pointer to a linkage section given a pointer to its text segment. The SLT entry referenced by text seg_ptr is examined. If it doesn't have a linkage section associated with it, i.e., linkage sw is set OFF, the error return, cr_ret, is used. Otherwise, the combine linkage switch, combine sw, is tested to see if the segment's linkage is part of a combined linkage segment. If it is, the link section and link offset items are used to construct a pointer to the linkage section. If the text segment only has the original linkage segment, a pointer to it is constructed using the text link segno item.