Data-Centric Execution of Speculative Parallel Programs

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Executive summary

Many-cores must exploit cache locality to scale

Current speculative systems, e.g. TLS or TM, do not exploit locality

Spatial Hints: run tasks likely to access the same data in the same place

- A software-given hint denotes the data a new task is likely to access
- Hardware maps tasks with the same hint to the same place
- Hardware uses hints to perform locality-aware load balancing

Our techniques make speculative parallelism practical at large scale

- It is easy to modify programs to convey locality through hints
- Performance improves by 3.3x at 256 cores
- We reduce network traffic by 6.4x and wasted work by 3.5x

Prior speculative systems scale poorly

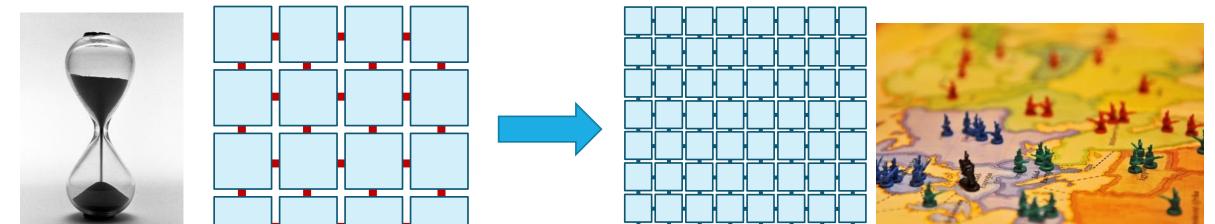
TRANSACTIONAL MEMORY (TM) SCHEDULERS SPATIAL HINTS

Reduce wasted work of coarse-grain txns

Limit concurrency: **When** to run a task?

Make accesses local for fine-grain tasks

Less data movement: Where to run a task?



Spatially map tasks for improved locality and less waste

Prior non-speculative locality techniques do not work for speculation

STATIC TASK MAPPING

DYNAMIC TASK MAPPING

Data dependences known a priori

• Linear algebra, Anton 2 [ASPLOS '13]

Graph partitioning

- Localizes communication and scheduling
- Slow preprocessing step
- Cannot adapt to imbalance

Work stealing

- Cheap, local enqueues
- Steals to adapt to imbalance
- Limited application types
- Stealing interferes with speculation

Baseline Architecture: Swarm [MICRO '15]

Baseline Swarm execution model

Programs consist of timestamped tasks

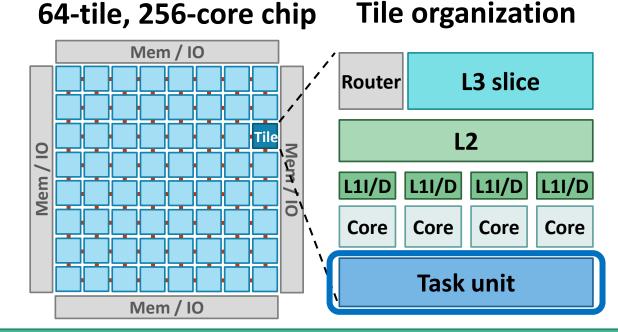
- Tasks can create children tasks with >= timestamp
- Tasks appear to execute in timestamp order

General execution model supports ordered and unordered parallelism

Baseline Swarm architecture

Speculatively executes tasks out of order

Large hardware task queues Scalable ordered speculation Scalable ordered commits

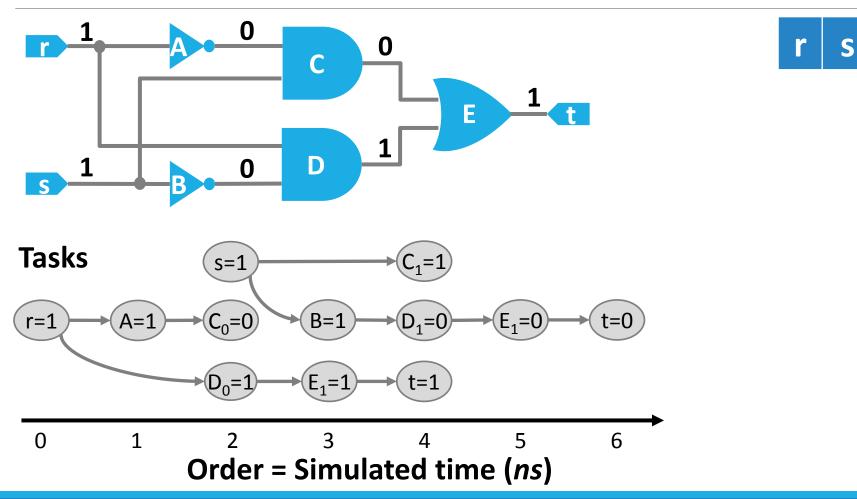


Efficiently supports tiny speculative tasks

Spatial Hints in Action

COMBINING SPECULATION AND LOCALITY

Example: Discrete event simulation (DES)

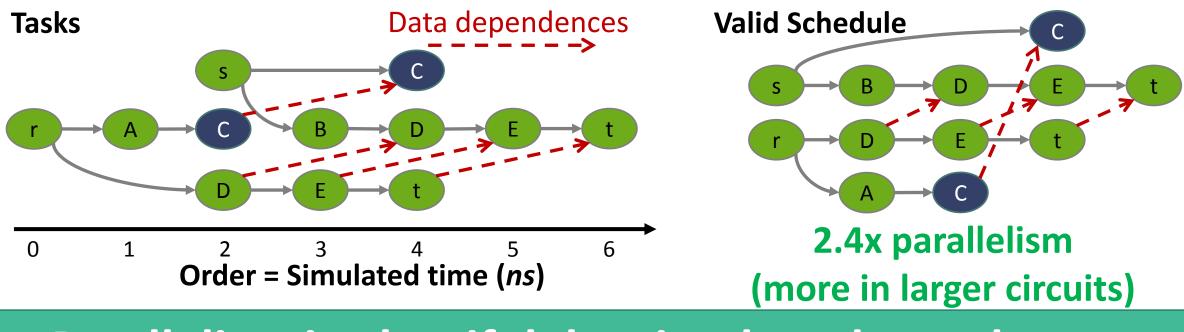


DATA-CENTRIC EXECUTION OF SPECULATIVE PARALLEL PROGRAMS

t = r XOR s

Extracting parallelism in DES

Execute independent tasks out of order



Parallelism is plentiful despite data dependences

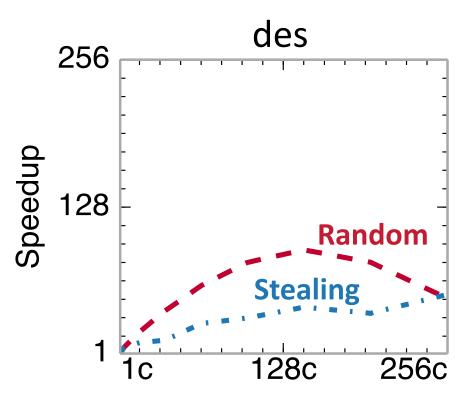
Speculation scales poorly without locality

Swarm sends new tasks to random tiles

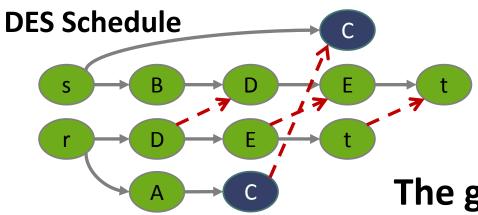
- Good for **load balance**
- **Poor locality** hurts scalability beyond 100 cores

Work stealing: a non-speculative scheduler

- Enqueue new tasks locally
- Steal from the most-loaded tile
- Not a good strategy for DES



Where is the locality?

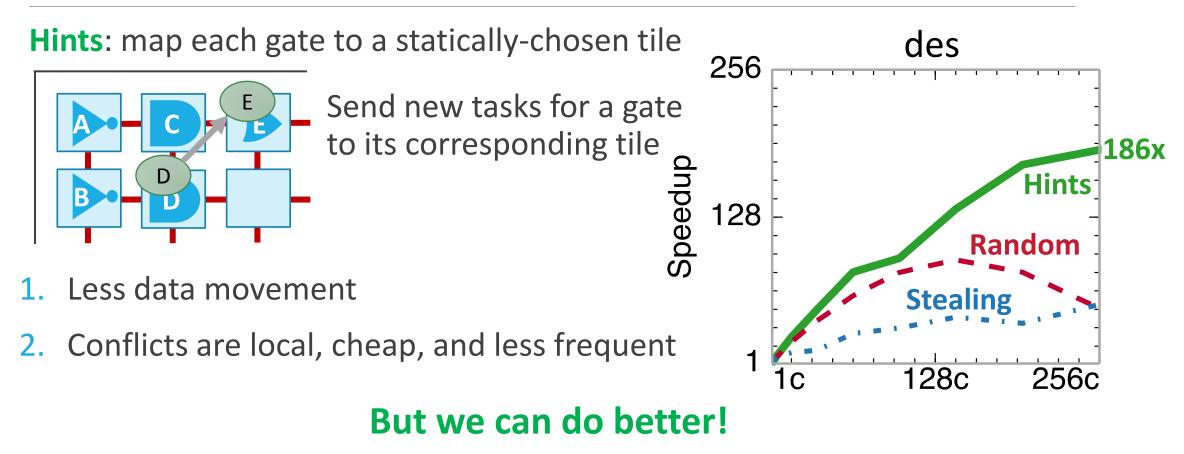


Each task operates on a single gate

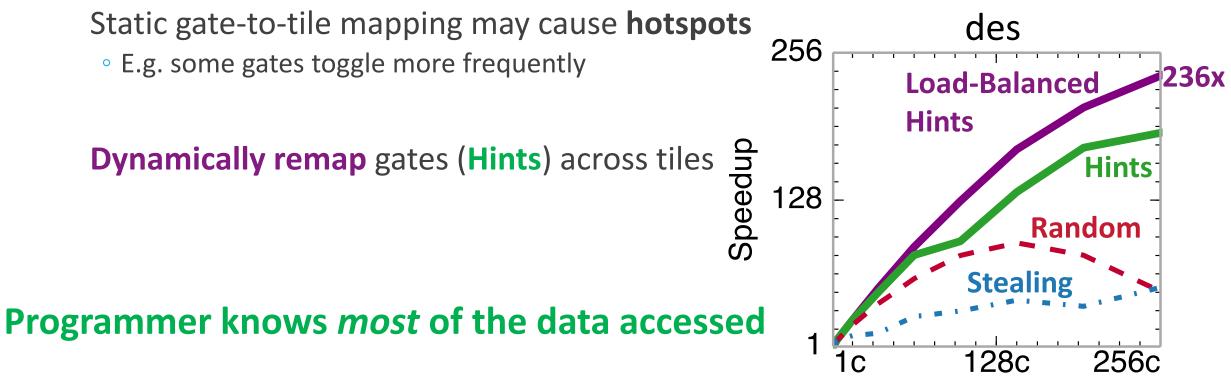
The gate is known when the task is created

With fine-grain tasks, most data accessed is known at creation time

Data-centric speculation scales well



Load-balanced speculation scales best



Spatial Hints convey program-level knowledge to exploit locality

Spatial Hints Implementation

Hint mechanisms are straightforward

SOFTWARE

HARDWARE

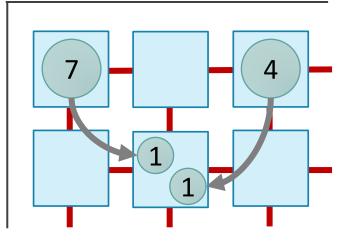
A **Spatial Hint** is an integer value

- Given at task creation time
- Denotes data likely to be accessed by the task
- E.g. the gate ID in DES

Localize most data accesses within a tile Serialize tasks likely to conflict

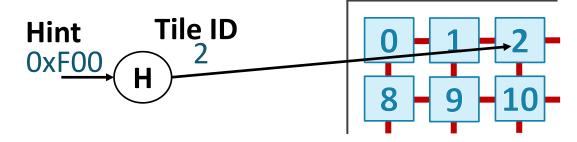
Hashes each new task's **Hint** to a tile ID

Serializes same-Hint tasks

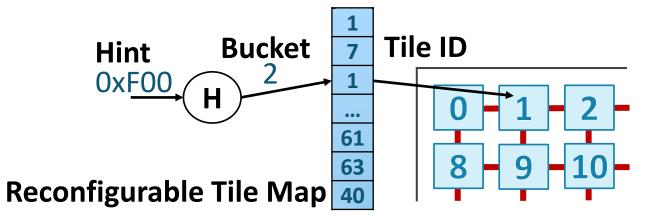


Load balance with a level of indirection

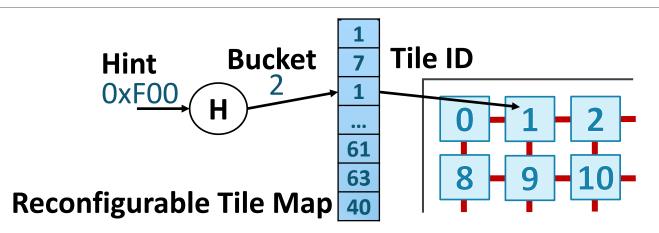
Static hint-to-tile mapping may cause imbalance



Instead, periodically **remap hints** across tiles to equalize load



"Load" is different for speculation



Non-speculative systems use # queued tasks as a proxy for load

When imbalanced, speculative systems often

- Don't run out of work
- Abort more work or strain speculation resources

Remap hints to tiles to balance # of committed cycles per tile

Adding hints to applications is easy

One line of code to express the Gate ID as a Hint

Adding hints to applications is easy

Benchmark	Hint	Why?
des	Gate ID	Map tasks for same gate to same tile
nocsim	Router ID	Frequent intra-router communication
bfs, sssp, astar, color	Cache-line address	Several vertices reside on the same line
silo	(Table ID, primary key)	Each task accesses one database tuple
genome, kmeans	Multiple	

See the paper for more details!

Load balance reconfiguration algorithm

Choice of application hints

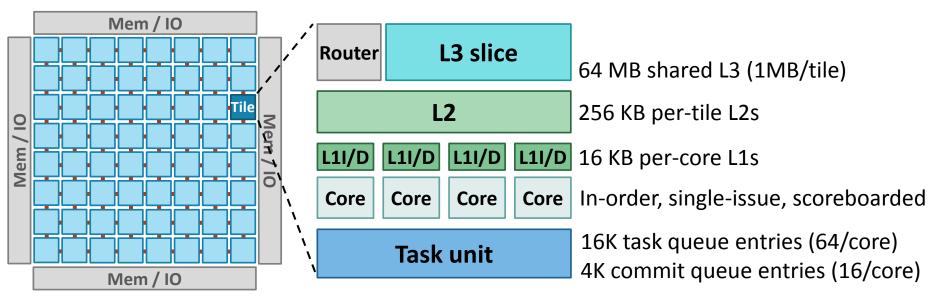
Relationship between task size and hint effectiveness

Evaluation

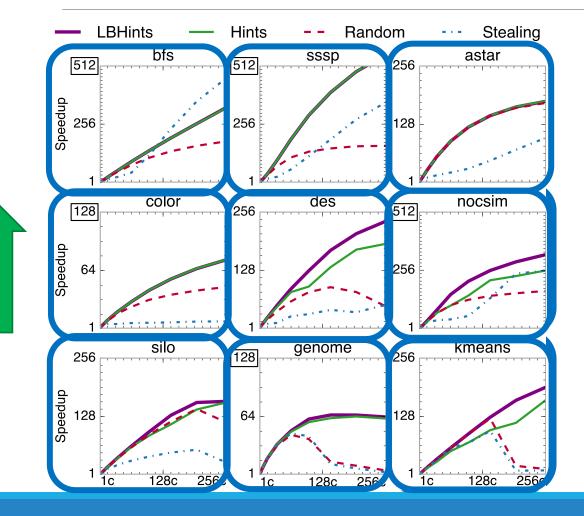
Methodology

Event-driven, Pin-based simulator Scalability experiments from 1–256 cores • Scaled-down systems have fewer tiles

Target system: 256-core, 64-tile chip



Hints make speculation practical on large-scale systems

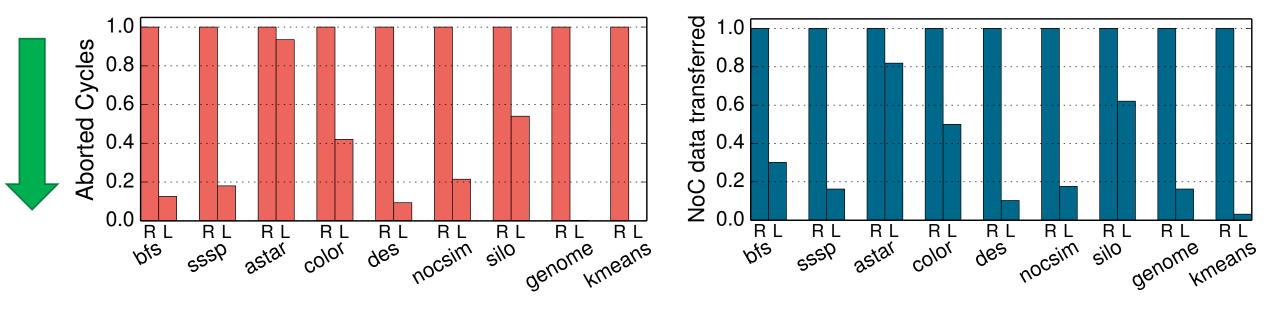


Load-Balanced Hints 3.3x faster than Random (193x gmean vs 58x)

Load-Balanced Hints 17% – 27% faster than Hints

Stealing is inconsistent across benchmarks

Hints make speculation more efficient



Reduce wasted work by 6.4x

Reduce network traffic by 3.5x

Conclusion

Speculative architectures must exploit locality to scale to 100s of cores

Important to simplify parallel programming

Spatial Hints convey app-level knowledge to exploit cache locality

Hardware leverages hints by:

- Sending tasks likely to access the same data to the same tile
- Serializing tasks likely to conflict
- Balancing work in a locality-aware and speculation-friendly way

Our techniques make speculation practical on large-scale systems

Thank you! Questions?

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