# Lagrangian Relaxation Algorithms for Inference in Natural Language Processing

Alexander M. Rush and Michael Collins

(based on joint work with Yin-Wen Chang, Tommi Jaakkola, Terry Koo, Roi Reichart, David Sontag)

## Decoding in NLP

**focus:** structured prediction for natural language processing decoding as a combinatorial optimization problem

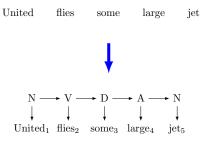
$$y^* = \arg\max_{y \in \mathcal{Y}} f(y)$$

where f is a scoring function and  $\mathcal{Y}$  is a set of structures

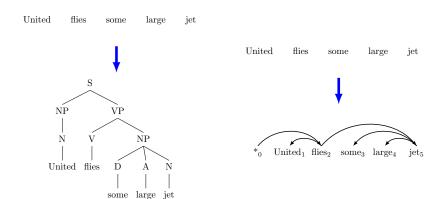
for some problems, use simple combinatorial algorithms

- dynamic programming
- minimum spanning tree
- min cut

## Structured prediction: Tagging



## Structured prediction: Parsing



## Decoding complexity

**issue:** simple combinatorial algorithms do not scale to richer models

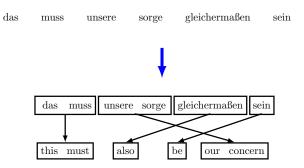
$$y^* = \arg\max_{y \in \mathcal{Y}} f(y)$$

need decoding algorithms for complex natural language tasks

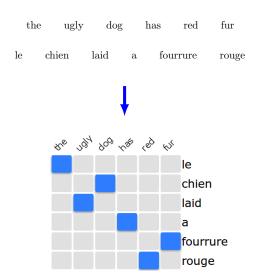
#### motivation:

- richer model structure often leads to improved accuracy
- exact decoding for complex models tends to be intractable

## Structured prediction: Phrase-based translation



## Structured prediction: Word alignment



## Decoding tasks

#### high complexity

- combined parsing and part-of-speech tagging (Rush et al., 2010)
- "loopy" HMM part-of-speech tagging
- syntactic machine translation (Rush and Collins, 2011)

#### NP-Hard

- symmetric HMM alignment (DeNero and Macherey, 2011)
- phrase-based translation (Chang and Collins, 2011)
- higher-order non-projective dependency parsing (Koo et al., 2010)

#### in practice:

- approximate decoding methods (coarse-to-fine, beam search, cube pruning, gibbs sampling, belief propagation)
- approximate models (mean field, variational models)

## Lagrangian relaxation

a general technique for constructing decoding algorithms solve complicated models

$$y^* = \arg\max_{y} f(y)$$

by decomposing into smaller problems.

upshot: can utilize a toolbox of combinatorial algorithms.

- dynamic programming
- minimum spanning tree
- shortest path
- · min cut
- •

## Lagrangian relaxation algorithms

Simple - uses basic combinatorial algorithms

Efficient - faster than solving exact decoding problems

#### Strong guarantees

- gives a certificate of optimality when exact
- direct connections to linear programming relaxations

## MAP problem in Markov random fields



**given:** binary variables  $x_1 \dots x_n$ 

goal: MAP problem

$$\arg\max_{x_1...x_n} \sum_{(i,j)\in E} f_{i,j}(x_i,x_j)$$

where each  $f_{i,j}(x_i, x_j)$  is a local potential for variables  $x_i$ ,  $x_j$ 

## Dual decomposition for MRFs (Komodakis et al., 2010)



goal:

$$\arg\max_{x_1...x_n} \sum_{(i,j)\in E} f_{i,j}(x_i,x_j)$$

#### equivalent formulation:

$$\arg\max_{x_1...x_n,y_1...y_n} \sum_{(i,j)\in\mathcal{T}_1} f'_{i,j}(x_i,x_j) + \sum_{(i,j)\in\mathcal{T}_2} f'_{i,j}(y_i,y_j)$$

such that for  $i = 1 \dots n$ ,

$$x_i = y_i$$

#### Lagrangian:

$$L(u,x,y) = \sum_{(i,j)\in T_1} f'_{i,j}(x_i,x_j) + \sum_{(i,j)\in T_2} f'_{i,j}(y_i,y_j) + \sum_i u_i(x_i-y_i)$$

#### Related work

- belief propagation using combinatorial algorithms (Duchi et al., 2007; Smith and Eisner, 2008)
- factored A\* search (Klein and Manning, 2003)

#### Tutorial outline

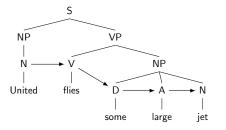
- 1. worked algorithm for combined parsing and tagging
- 2. important theorems and formal derivation
- 3. more examples from parsing and alignment
- 4. relationship to linear programming relaxations
- 5. practical considerations for implemention
- 6. further example from machine translation

## 1. Worked example

**aim:** walk through a Lagrangian relaxation algorithm for combined parsing and part-of-speech tagging

- · introduce formal notation for parsing and tagging
- · give assumptions necessary for decoding
- step through a run of the Lagrangian relaxation algorithm

## Combined parsing and part-of-speech tagging



goal: find parse tree that optimizes

$$score(S \rightarrow NP \ VP) + score(VP \rightarrow V \ NP) + ... + score(N \rightarrow V) + score(N \rightarrow United) + ...$$

## Constituency parsing

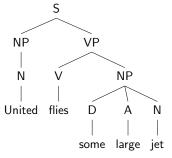
#### notation:

- $m{\cdot}$   $\mathcal{Y}$  is set of constituency parses for input
- $y \in \mathcal{Y}$  is a valid parse
- f(y) scores a parse tree

#### goal:

$$\arg\max_{y\in\mathcal{Y}}f(y)$$

example: a context-free grammar for constituency parsing



## Part-of-speech tagging

#### notation:

- ullet  ${\mathcal Z}$  is set of tag sequences for input
- $z \in \mathcal{Z}$  is a valid tag sequence
- g(z) scores of a tag sequence

#### goal:

$$arg \max_{z \in \mathcal{Z}} g(z)$$

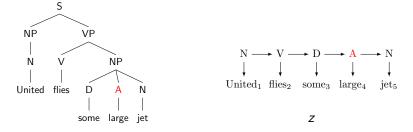
example: an HMM for part-of speech tagging

## Identifying tags

notation: identify the tag labels selected by each model

- y(i, t) = 1 when parse y selects tag t at position i
- z(i, t) = 1 when tag sequence z selects tag t at position i

**example:** a parse and tagging with y(4, A) = 1 and z(4, A) = 1



У

## Combined optimization

goal:

$$arg \max_{y \in \mathcal{Y}, z \in \mathcal{Z}} f(y) + g(z)$$

such that for all  $i = 1 \dots n$ ,  $t \in \mathcal{T}$ ,

$$y(i,t)=z(i,t)$$

i.e. find the best parse and tagging pair that agree on tag labels equivalent formulation:

$$arg \max_{y \in \mathcal{Y}} f(y) + g(I(y))$$

where  $\mathit{I}:\mathcal{Y}\to\mathcal{Z}$  extracts the tag sequence from a parse tree

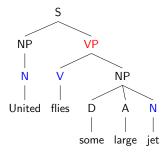
## Exact method: Dynamic programming intersection

can solve by solving the product of the two models

#### example:

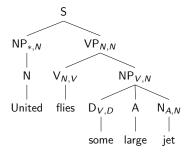
- parsing model is a context-free grammar
- tagging model is a first-order HMM
- can solve as CFG and finite-state automata intersection

replace  $VP \to V\ NP$  with  $VP_{N,V} \to V_{N,V}\ NP_{V,N}$ 



## Intersected parsing and tagging complexity

let G be the number of grammar non-terminals parsing CFG require  $O(G^3n^3)$  time with rules  $\operatorname{VP} \to \operatorname{V} \operatorname{NP}$ 



with intersection  $O(G^3n^3|\mathcal{T}|^3)$  with rules  $\mathrm{VP}_{\mathrm{N},\mathrm{V}} \to \mathrm{V}_{\mathrm{N},\mathrm{V}}$   $\mathrm{NP}_{\mathrm{V},\mathrm{N}}$  becomes  $O(G^3n^3|\mathcal{T}|^6)$  time for second-order HMM

### Parsing assumption

**assumption:** optimization with u can be solved efficiently

$$arg \max_{y \in \mathcal{Y}} f(y) + \sum_{i,t} u(i,t)y(i,t)$$

**example:** CFG with rule scoring function *h* 

$$f(y) = \sum_{X \to Y \ Z \in y} h(X \to Y \ Z) + \sum_{(i,X) \in y} h(X \to w_i)$$

where

$$\arg\max_{y\in\mathcal{Y}} f(y) + \sum_{i,t} u(i,t)y(i,t) =$$
 
$$\arg\max_{y\in\mathcal{Y}} \sum_{X\to Y} \sum_{Z\in\mathcal{Y}} h(X\to Y|Z) + \sum_{(i,X)\in\mathcal{Y}} (h(X\to w_i) + u(i,X))$$

## Tagging assumption

**assumption:** optimization with u can be solved efficiently

$$\arg \max_{z \in \mathcal{Z}} g(z) - \sum_{i,t} u(i,t)z(i,t)$$

**example:** HMM with scores for transitions T and observations O

$$g(z) = \sum_{t \to t' \in z} T(t \to t') + \sum_{(i,t) \in z} O(t \to w_i)$$

where

$$rg \max_{z \in \mathcal{Z}} \ g(z) - \sum_{i,t} u(i,t) z(i,t) =$$
  $rg \max_{z \in \mathcal{Z}} \ \sum_{t o t' \in z} T(t o t') + \sum_{(i,t) \in z} (O(t o w_i) - u(i,t))$ 

## Lagrangian relaxation algorithm

Set 
$$u^{(1)}(i,t) = 0$$
 for all  $i, t \in \mathcal{T}$ 

For 
$$k = 1$$
 to  $K$ 

$$y^{(k)} \leftarrow \arg\max_{y \in \mathcal{Y}} f(y) + \sum_{i,t} u^{(k)}(i,t)y(i,t)$$
 [Parsing]

$$z^{(k)} \leftarrow \arg\max_{z \in \mathcal{Z}} g(z) - \sum_{i,t} u^{(k)}(i,t)z(i,t)$$
 [Tagging]

If 
$$y^{(k)}(i,t) = z^{(k)}(i,t)$$
 for all  $i, t$  Return  $(y^{(k)}, z^{(k)})$ 

Else 
$$u^{(k+1)}(i,t) \leftarrow u^{(k)}(i,t) - \alpha_k(y^{(k)}(i,t) - z^{(k)}(i,t))$$

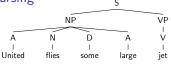
$$u(i,t) = 0$$
 for all  $i,t$ 

$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,t} u(i,t)y(i,t))$$

#### Viterbi Decoding

 $\mathsf{United}_1 \ \mathsf{flies}_2 \ \mathsf{some}_3 \ \mathsf{large}_4 \ \mathsf{jet}_5$ 

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$



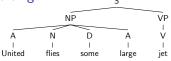
$$u(i,t) = 0$$
 for all  $i,t$ 

$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,t} u(i,t)y(i,t))$$

#### Viterbi Decoding

 $\mathsf{United}_1 \ \mathsf{flies}_2 \ \mathsf{some}_3 \ \mathsf{large}_4 \ \mathsf{jet}_5$ 

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$

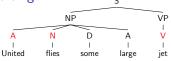


## Penalties u(i, t) = 0 for all i, t

$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,t} u(i,t)y(i,t))$$

#### Viterbi Decoding

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$



#### Penalties

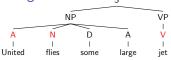
$$u(i,t) = 0$$
 for all  $i,t$ 

$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,t} u(i,t)y(i,t))$$

#### Viterbi Decoding

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$

$$\begin{array}{lllll} f(y) & \Leftarrow & \mathsf{CFG} & g(z) & \Leftarrow & \mathsf{HMM} \\ \mathcal{Y} & \Leftarrow & \mathsf{Parse Trees} & \mathcal{Z} & \Leftarrow & \mathsf{Tagging} \\ y(i,t) = 1 & \mathsf{if} & y \; \mathsf{contains \; tag} \; t \; \mathsf{at \; position} \; i \end{array}$$



$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,t} u(i,t)y(i,t))$$

#### Viterbi Decoding

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$

#### Key

$$\begin{array}{lllll} f(y) & \Leftarrow & \mathsf{CFG} & g(z) & \Leftarrow & \mathsf{HMM} \\ \mathcal{Y} & \Leftarrow & \mathsf{Parse Trees} & \mathcal{Z} & \Leftarrow & \mathsf{Tagging} \\ y(i,t) = 1 & \mathsf{if} & y \; \mathsf{contains \; tag} \; t \; \mathsf{at \; position} \; i \end{array}$$

#### **Penalties**

$$u(i,t) = 0$$
 for all  $i,t$ 

Iteration 1	
u(1,A)	-1
u(1, N)	1
u(2, N)	-1
u(2, V)	1
u(5, V)	-1
u(5, N)	1

#### Penalties

$$u(i, t) = 0$$
 for all  $i, t$ 

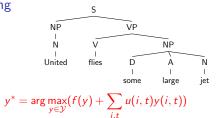
Iteration 1	
u(1,A)	-1
u(1, N)	1
u(2, N)	-1
u(2, V)	1
u(5, V)	-1
u(5, N)	1

$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,t} u(i,t)y(i,t))$$

#### Viterbi Decoding

United<sub>1</sub> flies<sub>2</sub> some<sub>3</sub> large<sub>4</sub> jet<sub>5</sub>

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$



#### Viterbi Decoding

## United<sub>1</sub> flies<sub>2</sub> some<sub>3</sub> large<sub>4</sub> jet<sub>5</sub>

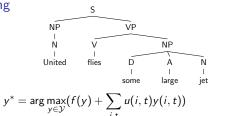
$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$

#### Key

#### **Penalties**

$$u(i,t) = 0$$
 for all  $i,t$ 

Iteration 1	
u(1,A)	-1
u(1, N)	1
u(2, N)	-1
u(2, V)	1
u(5, V)	-1
u(5. N)	1



#### Viterbi Decoding

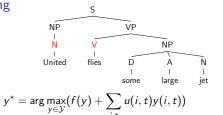
$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$

#### Key

#### **Penalties**

$$u(i,t) = 0$$
 for all  $i,t$ 

Iteration 1	
u(1,A)	-1
u(1, N)	1
u(2, N)	-1
u(2, V)	1
u(5, V)	-1
u(5, N)	1



#### Viterbi Decoding

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$

#### Key

#### **Penalties**

$$u(i,t) = 0$$
 for all  $i,t$ 

Iteration 1	
u(1,A)	-1
u(1, N)	1
u(2, N)	-1
u(2, V)	1
u(5, V)	-1

#### Iteration 2

u(5, N)

u(5, V)	-1
u(5, N)	1

1

## Penalties

#### u(i,t) = 0 for all i,t

Iteration 1	
u(1, A)	-1
u(1, N)	1
u(2, N)	-1
u(2, V)	1
u(5, V)	-1
u(5, N)	1

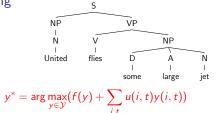
## $y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,t} u(i,t)y(i,t))$

#### Viterbi Decoding

 $\mathsf{United}_1 \ \mathsf{flies}_2 \ \mathsf{some}_3 \ \mathsf{large}_4 \ \mathsf{jet}_5$ 

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$

# Iteration 2 u(5, V) -1 u(5, N) 1



#### Viterbi Decoding

United<sub>1</sub> flies<sub>2</sub> some<sub>3</sub> large<sub>4</sub>

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$

#### Key

#### **Penalties**

$$u(i,t) = 0$$
 for all  $i,t$ 

Iteration 1	
u(1, A)	-1
u(1, N)	1
u(2, N)	-1
u(2, V)	1
u(5, V)	-1

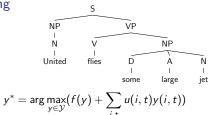
Iteration 2	
u(5, V)	-1
u(5, N)	1

u(5, N)

$$u(5, V)$$
 -1  $u(5, N)$  1

1

#### **CKY Parsing**



$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$

#### Key

#### **Penalties**

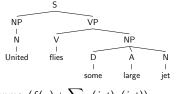
$$u(i,t) = 0$$
 for all  $i,t$ 

Iteration 1	
u(1,A)	-1
u(1, N)	1
u(2, N)	-1
u(2, V)	1
u(5, V)	-1

Iteration 2	
u(5, V)	-1
u(5, N)	1

u(5, N)

#### **CKY Parsing**



$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,t} u(i,t)y(i,t))$$

### Viterbi Decoding

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,t} u(i,t)z(i,t))$$

#### **Penalties**

$$u(i,t) = 0$$
 for all  $i,t$ 

Iteration 1	
u(1,A)	-1
u(1, N)	1
u(2, N)	-1
u(2, V)	1

$$u(5, V)$$
 -1  $u(5, N)$  1

Iteration 2	
u(5, V)	-1
u(5, N)	1

### Converged

# Key

$$\begin{array}{lll} f(y) & \Leftarrow & \mathsf{CFG} \\ \mathcal{Y} & \Leftarrow & \mathsf{Parse Trees} \\ y(i,t) = 1 & \mathsf{if} & y \; \mathsf{contains } \; \mathsf{tag} \; t \; \mathsf{at position} \; i \end{array}$$

$$y^* = \arg\max_{y \in \mathcal{Y}} f(y) + g(y)$$

**HMM** 

**Taggings** 

g(z)

# Main theorem

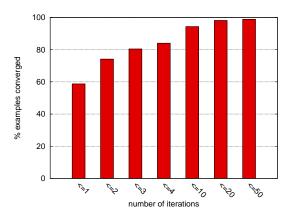
**theorem:** if at any iteration, for all  $i, t \in \mathcal{T}$ 

$$y^{(k)}(i,t) = z^{(k)}(i,t)$$

then  $(y^{(k)}, z^{(k)})$  is the global optimum

proof: focus of the next section

# Convergence



# 2. Formal properties

**aim:** formal derivation of the algorithm given in the previous section

- derive Lagrangian dual
- prove three properties
  - upper bound
  - convergence
  - optimality
- describe subgradient method

# Lagrangian

goal:

$$\arg\max_{y\in\mathcal{V}} f(y) + g(z)$$
 such that  $y(i,t) = z(i,t)$ 

#### Lagrangian:

$$L(u, y, z) = f(y) + g(z) + \sum_{i} u(i, t) (y(i, t) - z(i, t))$$

redistribute terms

$$L(u,y,z) = \left(f(y) + \sum_{i,t} u(i,t)y(i,t)\right) + \left(g(z) - \sum_{i,t} u(i,t)z(i,t)\right)$$

# Lagrangian dual

#### Lagrangian:

$$L(u,y,z) = \left(f(y) + \sum_{i,t} u(i,t)y(i,t)\right) + \left(g(z) - \sum_{i,t} u(i,t)z(i,t)\right)$$

#### Lagrangian dual:

$$L(u) = \max_{y \in \mathcal{Y}, z \in \mathcal{Z}} L(u, y, z)$$

$$= \max_{y \in \mathcal{Y}} \left( f(y) + \sum_{i, t} u(i, t) y(i, t) \right) + \max_{z \in \mathcal{Z}} \left( g(z) - \sum_{i, t} u(i, t) z(i, t) \right)$$

# Theorem 1. Upper bound

#### define:

•  $y^*, z^*$  is the optimal combined parsing and tagging solution with  $y^*(i, t) = z^*(i, t)$  for all i, t

**theorem:** for any value of u

$$L(u) \ge f(y^*) + g(z^*)$$

L(u) provides an upper bound on the score of the optimal solution **note:** upper bound may be useful as input to branch and bound or  $A^*$  search

# Theorem 1. Upper bound (proof)

**theorem:** for any value of u,  $L(u) \ge f(y^*) + g(z^*)$  **proof:** 

$$L(u) = \max_{y \in \mathcal{Y}, z \in \mathcal{Z}} L(u, y, z)$$

$$\geq \max_{y \in \mathcal{Y}, z \in \mathcal{Z}: y = z} L(u, y, z)$$

$$= \max_{y \in \mathcal{Y}, z \in \mathcal{Z}: y = z} f(y) + g(z)$$

$$= f(y^*) + g(z^*)$$

$$(1)$$

$$(2)$$

$$(3)$$

# Formal algorithm (reminder)

Set 
$$u^{(1)}(i,t) = 0$$
 for all  $i, t \in \mathcal{T}$ 

For 
$$k = 1$$
 to  $K$ 

$$y^{(k)} \leftarrow \arg\max_{y \in \mathcal{Y}} f(y) + \sum_{i,t} u^{(k)}(i,t)y(i,t)$$
 [Parsing]

$$z^{(k)} \leftarrow \arg\max_{z \in \mathcal{Z}} g(z) - \sum_{i,t} u^{(k)}(i,t)z(i,t)$$
 [Tagging]

If 
$$y^{(k)}(i,t) = z^{(k)}(i,t)$$
 for all  $i, t$  Return  $(y^{(k)}, z^{(k)})$ 

Else 
$$u^{(k+1)}(i,t) \leftarrow u^{(k)}(i,t) - \alpha_k(y^{(k)}(i,t) - z^{(k)}(i,t))$$

# Theorem 2. Convergence

#### notation:

- $u^{(k+1)}(i,t) \leftarrow u^{(k)}(i,t) + \alpha_k(y^{(k)}(i,t) z^{(k)}(i,t))$  is update
- $u^{(k)}$  is the penalty vector at iteration k
- $\alpha_k > 0$  is the update rate at iteration k

**theorem:** for any sequence  $\alpha^1, \alpha^2, \alpha^3, \dots$  such that

$$\lim_{t\to\infty}\alpha^t=0\quad\text{and}\quad\sum_{t=1}^\infty\alpha^t=\infty,$$

we have

$$\lim_{t\to\infty} L(u^t) = \min_{u} L(u)$$

i.e. the algorithm converges to the tightest possible upper bound **proof:** by subgradient convergence (next section)

### **Dual solutions**

#### define:

• for any value of u

$$y_u = \arg \max_{y \in \mathcal{Y}} \left( f(y) + \sum_{i,t} u(i,t)y(i,t) \right)$$

and

$$z_u = \arg \max_{z \in \mathcal{Z}} \left( g(z) - \sum_{i,t} u(i,t)z(i,t) \right)$$

•  $y_u$  and  $z_u$  are the dual solutions for a given u

# Theorem 3. Optimality

**theorem:** if there exists *u* such that

$$y_u(i,t) = z_u(i,t)$$

for all *i*, *t* then

$$f(y_u) + g(z_u) = f(y^*) + g(z^*)$$

i.e. if the dual solutions agree, we have an optimal solution

$$(y_u, z_u)$$

# Theorem 3. Optimality (proof)

**theorem:** if u such that  $y_u(i, t) = z_u(i, t)$  for all i, t then

$$f(y_u) + g(z_u) = f(y^*) + g(z^*)$$

**proof:** by the definitions of  $y_u$  and  $z_u$ 

$$L(u) = f(y_u) + g(z_u) + \sum_{i,t} u(i,t)(y_u(i,t) - z_u(i,t))$$
  
=  $f(y_u) + g(z_u)$ 

since  $L(u) \ge f(y^*) + g(z^*)$  for all values of u

$$f(y_u) + g(z_u) \ge f(y^*) + g(z^*)$$

but  $y^*$  and  $z^*$  are optimal

$$f(y_u) + g(z_u) \leq f(y^*) + g(z^*)$$

# **Dual optimization**

#### Lagrangian dual:

$$L(u) = \max_{y \in \mathcal{Y}, z \in \mathcal{Z}} L(u, y, z)$$

$$= \max_{y \in \mathcal{Y}} \left( f(y) + \sum_{i,t} u(i, t) y(i, t) \right) + \max_{z \in \mathcal{Z}} \left( g(z) - \sum_{i,t} u(i, t) z(i, t) \right)$$

goal: dual problem is to find the tightest upper bound

$$\min_{u} L(u)$$

# Dual subgradient

$$L(u) = \max_{y \in \mathcal{Y}} \left( f(y) + \sum_{i,t} u(i,t)y(i,t) \right) + \max_{z \in \mathcal{Z}} \left( g(z) - \sum_{i,t} u(i,t)z(i,t) \right)$$

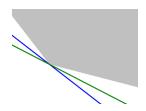
#### properties:

- L(u) is convex in u (no local minima)
- L(u) is not differentiable (because of max operator)

handle non-differentiability by using subgradient descent

**define:** a subgradient of L(u) at u is a vector  $g_u$  such that for all v

$$L(v) \geq L(u) + g_u \cdot (v - u)$$



# Subgradient algorithm

$$L(u) = \max_{y \in \mathcal{Y}} \left( f(y) + \sum_{i,t} u(i,t)y(i,t) \right) + \max_{z \in \mathcal{Z}} \left( g(z) - \sum_{i,j} u(i,t)z(i,t) \right)$$

recall,  $y_u$  and  $z_u$  are the argmax's of the two terms subgradient:

$$g_u(i,t) = y_u(i,t) - z_u(i,t)$$

subgradient descent: move along the subgradient

$$u'(i,t) = u(i,t) - \alpha \left( y_u(i,t) - z_u(i,t) \right)$$

guaranteed to find a minimum with conditions given earlier for  $\alpha$ 

# 3. More examples

**aim:** demonstrate similar algorithms that can be applied to other decoding applications

- · context-free parsing combined with dependency parsing
- combined translation alignment

# Combined constituency and dependency parsing (Rush et al., 2010)

**setup:** assume separate models trained for constituency and dependency parsing

**problem:** find constituency parse that maximizes the sum of the two models

#### example:

combine lexicalized CFG with second-order dependency parser

# Lexicalized constituency parsing

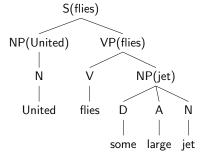
#### notation:

- $m{\cdot}$   $\mathcal{Y}$  is set of lexicalized constituency parses for input
- $y \in \mathcal{Y}$  is a valid parse
- f(y) scores a parse tree

### goal:

$$\arg\max_{y\in\mathcal{Y}}f(y)$$

example: a lexicalized context-free grammar

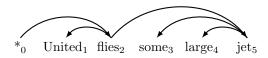


# Dependency parsing

#### define:

- $\mathcal{Z}$  is set of dependency parses for input
- $z \in \mathcal{Z}$  is a valid dependency parse
- g(z) scores a dependency parse

#### example:

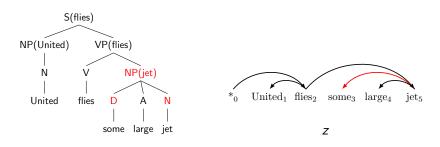


# Identifying dependencies

notation: identify the dependencies selected by each model

- y(i,j) = 1 when word i modifies of word j in constituency parse y
- z(i,j) = 1 when word i modifies of word j in dependency parse z

**example:** a constituency and dependency parse with y(3,5) = 1 and z(3,5) = 1



# Combined optimization

goal:

$$arg \max_{y \in \mathcal{Y}, z \in \mathcal{Z}} f(y) + g(z)$$

such that for all  $i = 1 \dots n$ ,  $j = 0 \dots n$ ,

$$y(i,j) = z(i,j)$$

$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,j} u(i,j)y(i,j))$$

# Dependency Parsing

\*<sub>0</sub> United<sub>1</sub> flies<sub>2</sub> some<sub>3</sub> large<sub>4</sub> jet<sub>5</sub>

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,j} u(i,j)z(i,j))$$

# Key

CKY Parsing S(flies)

NP VP(flies)

N V D NP(jet)

United flies some A N large jet

$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,j} u(i,j)y(i,j))$$

# Penalties

# u(i,j) = 0 for all i,j

# Dependency Parsing

 $*_0$  United<sub>1</sub> flies<sub>2</sub> some<sub>3</sub> large<sub>4</sub> jet<sub>5</sub>

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,j} u(i,j)z(i,j))$$

$$\begin{array}{lllll} f(y) & \Leftarrow & \mathsf{CFG} & g(z) & \Leftarrow & \mathsf{Dependency Model} \\ \mathcal{Y} & \Leftarrow & \mathsf{Parse Trees} & \mathcal{Z} & \Leftarrow & \mathsf{Dependency Trees} \\ y(i,j) = 1 & \mathsf{if} & y \mathsf{ contains dependency } i,j & & & & & & \\ \end{array}$$

CKY Parsing

S(flies)

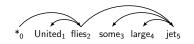
NP
VP(flies)

N
V
D
NP(jet)

United flies some A
N
I
large jet

Penalties 
$$u(i,j) = 0$$
 for all  $i,j$ 

#### Dependency Parsing



 $y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,j} u(i,j)y(i,j))$ 

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,j} u(i,j)z(i,j))$$

# Key

$$\begin{array}{lllll} f(y) & \Leftarrow & \mathsf{CFG} & g(z) & \Leftarrow & \mathsf{Dependency Model} \\ \mathcal{Y} & \Leftarrow & \mathsf{Parse Trees} & \mathcal{Z} & \Leftarrow & \mathsf{Dependency Trees} \\ y(i,j) = 1 & \mathsf{if} & y \mathsf{ contains dependency } i,j & & & & & & \\ \end{array}$$

CKY Parsing S(flies) NP VP(flies) N V D NP(jet) United flies some A N | I arge jet  $y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,j} u(i,j)y(i,j))$ 

# Penalties

u(i,j) = 0 for all i,j

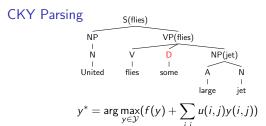
# **Dependency Parsing**



$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,j} u(i,j)z(i,j))$$

# Key

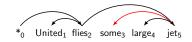
$$\begin{array}{lllll} f(y) & \Leftarrow & \mathsf{CFG} & g(z) & \Leftarrow & \mathsf{Dependency Model} \\ \mathcal{Y} & \Leftarrow & \mathsf{Parse Trees} & \mathcal{Z} & \Leftarrow & \mathsf{Dependency Trees} \\ y(i,j) = 1 & \mathsf{if} & y \mathsf{ contains dependency } i,j & & & & & & \\ \end{array}$$



# Penalties u(i,j) = 0 for all i,jIteration 1 u(2,3) -1

u(5,3)

#### Dependency Parsing



$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,j} u(i,j)z(i,j))$$

# **CKY Parsing**

#### Penalties

$$u(i,j) = 0$$
 for all  $i,j$ 

Iteration 1	
u(2,3)	-1
u(5 3)	1

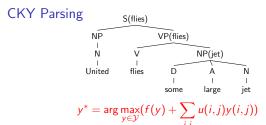
$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,j} u(i,j)y(i,j))$$

#### Dependency Parsing

 $*_0$  United<sub>1</sub> flies<sub>2</sub> some<sub>3</sub> large<sub>4</sub> jet<sub>5</sub>

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,j} u(i,j)z(i,j))$$

# Key



# Penalties u(i,j) = 0 for all i,j $\frac{\text{Iteration 1}}{u(2,3)}$

1

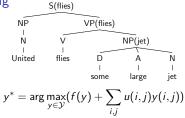
u(5,3)

### Dependency Parsing

\*<sub>0</sub> United<sub>1</sub> flies<sub>2</sub> some<sub>3</sub> large<sub>4</sub> jet<sub>5</sub>

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,j} u(i,j)z(i,j))$$





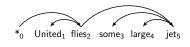
# Penalties

$$u(i,j) = 0$$
 for all  $i,j$   

$$\frac{\text{Iteration 1}}{u(2,3)} - 1$$

$$u(5,3) \qquad 1$$

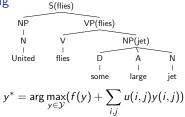
### Dependency Parsing



$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,j} u(i,j)z(i,j))$$

# Key





# Dependency Parsing



$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,i} u(i,j)z(i,j))$$

# Key

$$\begin{array}{lllll} f(y) & \Leftarrow & \mathsf{CFG} & g(z) & \Leftarrow & \mathsf{Dependency Model} \\ \mathcal{Y} & \Leftarrow & \mathsf{Parse Trees} & \mathcal{Z} & \Leftarrow & \mathsf{Dependency Trees} \\ y(i,j) = 1 & \mathsf{if} & y \mathsf{ contains dependency } i,j & & & & & & \\ \end{array}$$

#### Penalties

$$u(i,j) = 0 \text{ for all } i,j$$

$$\frac{\text{Iteration 1}}{u(2,3)}$$

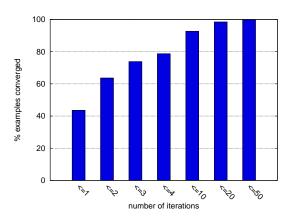
1

#### Converged

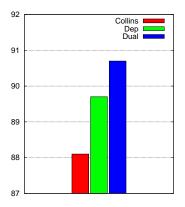
u(5,3)

$$y^* = \arg\max_{y \in \mathcal{Y}} f(y) + g(y)$$

# Convergence



# Integrated Constituency and Dependency Parsing: Accuracy



F<sub>1</sub> Score

- ► Collins (1997) Model 1
- ► Fixed, First-best Dependencies from Koo (2008)
- Dual Decomposition

# Combined alignment (DeNero and Macherey, 2011)

**setup:** assume separate models trained for English-to-French and French-to-English alignment

**problem:** find an alignment that maximizes the score of both models

#### example:

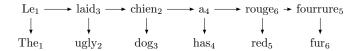
 HMM models for both directional alignments (assume correct alignment is one-to-one for simplicity)

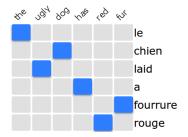
# English-to-French alignment

#### define:

- $oldsymbol{\cdot} \mathcal{Y}$  is set of all possible English-to-French alignments
- $y \in \mathcal{Y}$  is a valid alignment
- f(y) scores of the alignment

#### example: HMM alignment



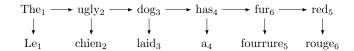


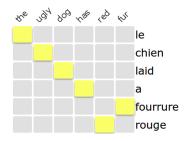
# French-to-English alignment

### define:

- Z is set of all possible French-to-English alignments
- $z \in \mathcal{Z}$  is a valid alignment
- g(z) scores of an alignment

# example: HMM alignment



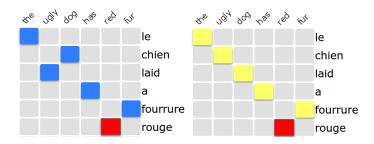


# Identifying word alignments

notation: identify the tag labels selected by each model

- y(i,j) = 1 when e-to-f alignment y selects French word i to align with English word j
- z(i,j) = 1 when f-to-e alignment z selects French word i to align with English word j

**example:** two HMM alignment models with y(6,5) = 1 and z(6,5) = 1



# Combined optimization

goal:

$$arg \max_{y \in \mathcal{Y}, z \in \mathcal{Z}} f(y) + g(z)$$

such that for all  $i = 1 \dots n$ ,  $j = 1 \dots n$ ,

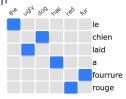
$$y(i,j)=z(i,j)$$

$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,i} u(i,j)y(i,j))$$

French-to-English

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,i} u(i,j)z(i,j))$$

### Key



**Penalties** 

u(i,j) = 0 for all i,j

$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,j} u(i,j)y(i,j))$$

# French-to-English

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,j} u(i,j)z(i,j))$$

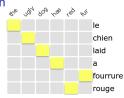
### Key

# 

fourrure rouge

$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,i} u(i,j)y(i,j))$$

# French-to-English



$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,i} u(i,j)z(i,j))$$

# Key

$$f(y) \Leftarrow HMM Alignment$$
 $\mathcal{Y} \Leftarrow English-to-French model$ 
 $y(i,j)=1$  if French word  $i$  aligns to English word  $j$ 

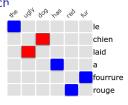
$$g(z) \leftarrow \mathsf{HMM} \; \mathsf{Alignment}$$

**Penalties** 

u(i,j) = 0 for all i,j

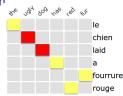
# **Penalties**

$$u(i,j) = 0$$
 for all  $i,j$ 



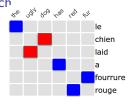
$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,i} u(i,j)y(i,j))$$

French-to-English



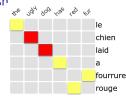
$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,i} u(i,j)z(i,j))$$

# Key



$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,j} u(i,j)y(i,j))$$

# French-to-English



$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,i} u(i,j)z(i,j))$$

### Key

$$f(y) \Leftarrow \mathsf{HMM}$$
 Alignment
 $\mathcal{Y} \Leftarrow \mathsf{English}$ -to-French model
 $y(i,j)=1$  if French word  $i$  aligns to English word  $j$ 

# Penalties

$$u(i,j) = 0 \text{ for all } i,j$$

$$\frac{\text{Iteration 1}}{u(3,2)} -1$$

$$u(2,2) \qquad 1$$

$$u(2,3) \qquad -1$$

$$u(3,3)$$
 1

$$u(3,3)$$
 1

$$g(z) \leftarrow \mathsf{HMM} \; \mathsf{Alignment}$$

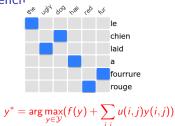
# u(i,j) = 0 for all i,j $\frac{\text{Iteration 1}}{u(3,2)} - 1$ $u(2,2) \qquad 1$ $u(2,3) \qquad -1$ $u(2,3) \qquad -1$ $u(3,3) \qquad 1$

**Penalties** 

French-to-English

$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,j} u(i,j)z(i,j))$$

# Key



# French-to-English

# **Penalties**

$$u(i,j) = 0$$
 for all  $i,j$   

$$\frac{\text{Iteration 1}}{u(3,2)} - 1$$

$$u(2,2) \qquad 1$$

$$u(2,3) \qquad -1$$

$$u(3,3) \qquad 1$$

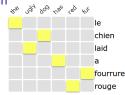
$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,i} u(i,j)z(i,j))$$

# Key



$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,j} u(i,j)y(i,j))$$

# French-to-English



$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,j} u(i,j)z(i,j))$$

# Key

### **Penalties**

$$u(i,j) = 0 \text{ for all } i,j$$

$$\frac{\text{Iteration 1}}{u(3,2)} - 1$$

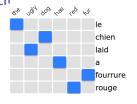
$$u(2,2) \qquad 1$$

$$u(2,3) \qquad -1$$

1

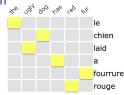
u(3,3)

 $\leftarrow$  French-to-English model



$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,j} u(i,j)y(i,j))$$

French-to-English



$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,i} u(i,j)z(i,j))$$

# Key

$$u(i,j) = 0 \text{ for all } i,j$$

$$\frac{\text{Iteration 1}}{u(3,2)} -1$$

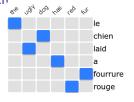
$$u(2,2) \qquad 1$$

$$u(2,3) \qquad -1$$

1

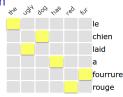
u(3,3)

$$g(z) \Leftarrow \mathsf{HMM} \mathsf{Alignment}$$



$$y^* = \arg\max_{y \in \mathcal{Y}} (f(y) + \sum_{i,i} u(i,j)y(i,j))$$

French-to-English



$$z^* = \arg\max_{z \in \mathcal{Z}} (g(z) - \sum_{i,i} u(i,j)z(i,j))$$

# Key

$$\begin{array}{llll} f(y) & \Leftarrow & \mathsf{HMM} \; \mathsf{Alignment} & g(z) \\ \mathcal{Y} & \Leftarrow & \mathsf{English-to-French} \; \mathsf{model} & \mathcal{Z} \\ y(i,j) = 1 & \mathsf{if} & \mathsf{French} \; \mathsf{word} \; i \; \mathsf{aligns} \; \mathsf{to} \; \mathsf{English} \; \mathsf{word} \; j \\ \end{array}$$

$$u(i,j) = 0 \text{ for all } i,j$$

$$\frac{\text{Iteration } 1}{u(3,2)} - 1$$

$$u(2,2) \qquad 1$$

$$u(2,3) \qquad -1$$

1

u(3,3)

$$g(z) \Leftarrow \mathsf{HMM} \mathsf{Alignment}$$

# 4. Linear programming

**aim:** explore the connections between Lagrangian relaxation and linear programming

- basic optimization over the simplex
- · formal properties of linear programming
- full example with fractional optimal solutions

# **Simplex**

### define:

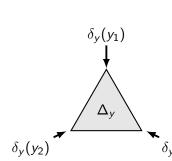
•  $\Delta_y \subset \mathcal{R}^{|\mathcal{Y}|}$  is the simplex over  $\mathcal{Y}$  where  $\alpha \in \Delta_y$  implies

$$lpha_y \geq 0$$
 and  $\sum_y lpha_y = 1$ 

- $\alpha$  is distribution over  $\mathcal{Y}$
- $\Delta_z$  is the simplex over  $\mathcal{Z}$
- $\delta_y: \mathcal{Y} o \Delta_y$  maps elements to the simplex

# example:

$$\mathcal{Y} = \{y_1, y_2, y_3\}$$
  
vertices  
•  $\delta_y(y_1) = (1, 0, 0)$   
•  $\delta_y(y_2) = (0, 1, 0)$   
•  $\delta_y(y_3) = (0, 0, 1)$ 



# Theorem 1. Simplex linear program

optimize over the simplex  $\Delta_y$  instead of the discrete sets  ${\cal Y}$ 

goal: optimize linear program

$$\max_{\alpha \in \Delta_y} \sum_{y} \alpha_y f(y)$$

theorem:

$$\max_{y \in \mathcal{Y}} f(y) = \max_{\alpha \in \Delta_y} \sum_{y} \alpha_y f(y)$$

**proof:** points in  $\mathcal{Y}$  correspond to the exteme points of simplex

$$\{\delta_{y}(y):y\in\mathcal{Y}\}$$

linear program has optimum at extreme point proof shows that best distribution chooses a single parse

# Combined linear program

optimize over the simplices  $\Delta_y$  and  $\Delta_z$  instead of the discrete sets  $\mathcal Y$  and  $\mathcal Z$ 

goal: optimize linear program

$$\max_{\alpha \in \Delta_y, \beta \in \Delta_z} \sum_{y} \alpha_y f(y) + \sum_{z} \beta_z g(z)$$

such that for all i, t

$$\sum_{v} \alpha_{y} y(i, t) = \sum_{z} \beta_{z} z(i, t)$$

**note:** the two distributions must match in expectation of POS tags the best distributions  $\alpha^*,\beta^*$  are possibly no longer a single parse tree or tag sequence

# Lagrangian

# Lagrangian:

$$M(u,\alpha,\beta) = \sum_{y} \alpha_{y} f(y) + \sum_{z} \beta_{z} g(z) + \sum_{i,t} u(i,t) \left( \sum_{y} \alpha_{y} y(i,t) - \sum_{z} \beta_{z} z(i,t) \right)$$

$$= \left( \sum_{y} \alpha_{y} f(y) + \sum_{i,t} u(i,t) \sum_{y} \alpha_{y} y(i,t) \right) + \left( \sum_{z} \beta_{z} g(z) - \sum_{i,t} u(i,t) \sum_{z} \beta_{z} z(i,t) \right)$$

# Lagrangian dual:

$$M(u) = \max_{\alpha \in \Delta_{v}, \beta \in \Delta_{z}} M(u, \alpha, \beta)$$

# Theorem 2. Strong duality

### define:

•  $\alpha^*, \beta^*$  is the optimal assignment to  $\alpha, \beta$  in the linear program

### theorem:

$$\min_{u} M(u) = \sum_{y} \alpha_{y}^{*} f(y) + \sum_{z} \beta_{z}^{*} g(z)$$

proof: by linear programming duality

# Theorem 3. Dual relationship

**theorem:** for any value of u,

$$M(u) = L(u)$$

**note:** solving the original Lagrangian dual also solves dual of the linear program

# Theorem 3. Dual relationship (proof sketch)

focus on  ${\mathcal Y}$  term in Lagrangian

$$L(u) = \max_{y \in \mathcal{Y}} \left( f(y) + \sum_{i,t} u(i,t)y(i,t) \right) + \dots$$

$$M(u) = \max_{\alpha \in \Delta_y} \left( \sum_{y} \alpha_y f(y) + \sum_{i,t} u(i,t) \sum_{y} \alpha_y y(i,t) \right) + \dots$$

by theorem 1. optimization over  $\mathcal Y$  and  $\Delta_y$  have the same max similar argument for  $\mathcal Z$  gives L(u)=M(u)

# Summary

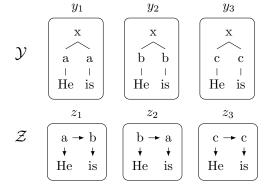
$$f(y)+g(z)$$
 original primal objective  $L(u)$  original dual  $\sum_y \alpha_y f(y) + \sum_z \beta_z g(z)$  LP primal objective LP dual

relationship between LP dual, original dual, and LP primal objective

$$\min_{u} M(u) = \min_{u} L(u) = \sum_{y} \alpha_{y}^{*} f(y) + \sum_{z} \beta_{z}^{*} g(z)$$

# Concrete example

- $\mathcal{Y} = \{y_1, y_2, y_3\}$
- $\mathcal{Z} = \{z_1, z_2, z_3\}$
- $\Delta_{\scriptscriptstyle V}\subset \mathbb{R}^3$ ,  $\Delta_{\scriptscriptstyle Z}\subset \mathbb{R}^3$



# Simple solution

# choose:

- $\alpha^{(1)} = (0,0,1) \in \Delta_V$  is representation of  $y_3$
- $\beta^{(1)} = (0,0,1) \in \Delta_z$  is representation of  $z_3$

# confirm:

$$\sum_{y} \alpha_y^{(1)} y(i,t) = \sum_{z} \beta_z^{(1)} z(i,t)$$

 $\alpha^{(1)}$  and  $\beta^{(1)}$  satisfy agreement constraint

Fractional solution
$$y_1 \qquad y_2 \qquad y_3$$

$$\begin{array}{c|cccc} x & & & & & & & & \\ \hline x & & & & & & & \\ \hline x & & & & & & \\ \hline x & & & & & & \\ \hline x & & & & & & \\ \hline x & & & & & \\ \hline a & a & & & & \\ \hline b & b & & & \\ \hline b & b & & & \\ \hline c & c & \\ \hline d & & & & \\ \hline c & c & \\ \hline d & & & \\ \hline d & & & \\ \hline c & & \\ \hline c & & \\ \hline d & & \\ \hline c & & \\ \hline d & & \\ \hline d$$

### choose:

- $\alpha^{(2)} = (0.5, 0.5, 0) \in \Delta_v$  is combination of  $y_1$  and  $y_2$
- $\beta^{(2)} = (0.5, 0.5, 0) \in \Delta_z$  is combination of  $z_1$  and  $z_2$

# confirm:

$$\sum_{y} \alpha_y^{(2)} y(i,t) = \sum_{z} \beta_z^{(2)} z(i,t)$$

 $\alpha^{(2)}$  and  $\beta^{(2)}$  satisfy agreement constraint, but not integral

# Optimal solution

# weights:

- the choice of f and g determines the optimal solution
- if (f,g) favors  $(\alpha^{(2)},\beta^{(2)})$ , the optimal solution is fractional

**example:** 
$$f = [1 \ 1 \ 2]$$
 and  $g = [1 \ 1 \ -2]$ 

- $f \cdot \alpha^{(1)} + g \cdot \beta^{(1)} = 0$  vs  $f \cdot \alpha^{(2)} + g \cdot \beta^{(2)} = 2$
- $\alpha^{(2)}, \beta^{(2)}$  is optimal, even though it is fractional

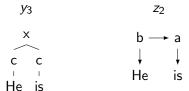
summary: dual and LP primal optimal:

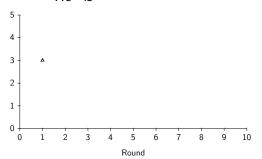
$$\min_{u} M(u) = \min_{u} L(u) = \sum_{y} \alpha_{y}^{(2)} f(y) + \sum_{z} \beta_{z}^{(2)} g(z) = 2$$

original primal optimal:

$$f(y^*) + g(z^*) = 0$$

# dual solutions:





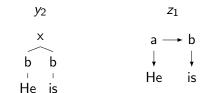
# dual values:

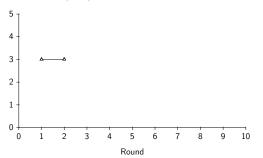
 $y^{(1)}$  2.00  $z^{(1)}$  1.00  $L(u^{(1)})$  3.00

# previous solutions:

 $y_3$   $z_2$ 

# dual solutions:





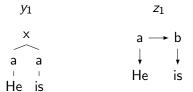
# dual values:

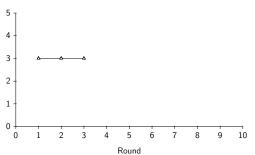
 $y^{(2)}$  2.00  $z^{(2)}$  1.00  $L(u^{(2)})$  3.00

# previous solutions:

 $y_3$   $z_2$   $y_2$   $z_1$ 

# dual solutions:





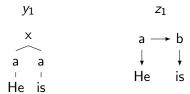
# dual values:

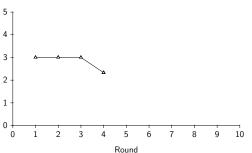
 $y^{(3)}$  2.50  $z^{(3)}$  0.50  $L(u^{(3)})$  3.00

# previous solutions:

 $y_3 Z_2$   $y_2 Z_1$   $y_1 Z_1$ 

### dual solutions:





# dual values:

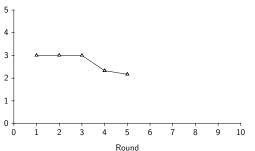
 $y^{(4)}$  2.17  $z^{(4)}$  0.17  $L(u^{(4)})$  2.33

# previous solutions:

 $y_3 Z_2$   $y_2 Z_1$   $y_1 Z_1$   $y_1 Z_1$ 

### dual solutions:





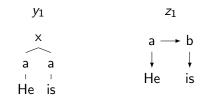
# dual values:

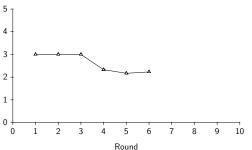
 $y^{(5)}$  2.08  $z^{(5)}$  0.08  $L(u^{(5)})$  2.17

# previous solutions:

 $y_3$   $z_2$   $y_2$   $z_1$   $y_1$   $z_1$   $y_1$   $z_1$  $y_2$   $z_2$ 

### dual solutions:





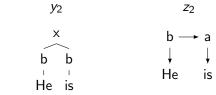
# dual values:

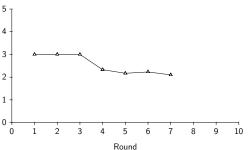
 $y^{(6)}$  2.12  $z^{(6)}$  0.12  $L(u^{(6)})$  2.23

# previous solutions:

y3 Z<sub>2</sub>
y2 Z<sub>1</sub>
y1 Z<sub>1</sub>
y1 Z<sub>1</sub>
y2 Z<sub>2</sub>
y1 Z<sub>1</sub>

# dual solutions:





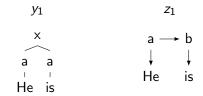
# dual values:

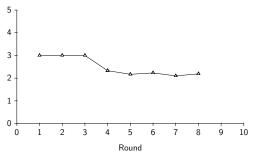
$y^{(7)}$	2.05
$z^{(7)}$	0.05
$L(u^{(7)})$	2.10

# previous solutions:

<i>y</i> <sub>3</sub>	$z_2$
<i>y</i> <sub>2</sub>	$z_1$
<i>y</i> <sub>1</sub>	$z_1$
<i>y</i> <sub>1</sub>	$z_1$
<i>y</i> <sub>2</sub>	$z_2$
$y_1$	$z_1$
<i>y</i> <sub>2</sub>	$z_2$

# dual solutions:





# dual values:

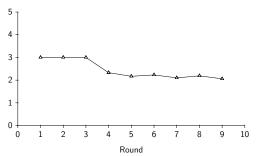
 $y^{(8)}$  2.09  $z^{(8)}$  0.09  $L(u^{(8)})$  2.19

# previous solutions:

*y*<sub>3</sub>  $z_2$ *y*2  $z_1$ *y*<sub>1</sub>  $z_1$ *y*<sub>1</sub>  $z_1$ *y*<sub>2</sub>  $z_2$  $y_1$  $z_1$ *y*<sub>2</sub>  $z_2$ *y*<sub>1</sub>  $z_1$ 

# dual solutions:





# dual values:

$y^{(9)}$	2.03
$z^{(9)}$	0.03
$L(u^{(9)})$	2.06

# previous solutions:

<i>y</i> <sub>3</sub>	$z_2$
<i>y</i> <sub>2</sub>	$z_1$
<i>y</i> <sub>1</sub>	$z_1$
<i>y</i> <sub>1</sub>	$z_1$
<i>y</i> <sub>2</sub>	<i>z</i> <sub>2</sub>
$y_1$	$z_1$
<i>y</i> <sub>2</sub>	$z_2$
<i>y</i> <sub>1</sub>	$z_1$
<i>y</i> <sub>2</sub>	<i>z</i> <sub>2</sub>

# 5. Practical issues

tracking the progress of the algorithm

- know current dual value and (possibly) primal value choice of update rate  $\alpha_{\it k}$
- various strategies; success with rate based on dual progress
   lazy update of dual solutions
- if updates are sparse, can avoid dynamically update soltuions extracting solutions if algorithm does not converge
  - best primal feasible solution; average solutions

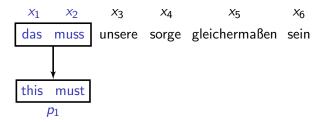
### define:

- ▶ source-language sentence words  $x_1, ..., x_N$
- phrase translation p = (s, e, t)
- ▶ translation derivation  $y = p_1, ..., p_L$

$$x_1$$
  $x_2$   $x_3$   $x_4$   $x_5$   $x_6$  das muss unsere sorge gleichermaßen sein

### define:

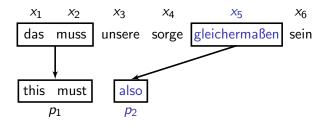
- ▶ source-language sentence words  $x_1, ..., x_N$
- phrase translation p = (s, e, t)
- ▶ translation derivation  $y = p_1, ..., p_L$



$$y = \{(1, 2, \text{this must}),$$

### define:

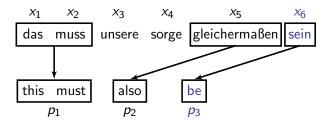
- ▶ source-language sentence words  $x_1, ..., x_N$
- phrase translation p = (s, e, t)
- ▶ translation derivation  $y = p_1, ..., p_L$



$$y = \{(1, 2, \text{this must}), (5, 5, \text{also}), \}$$

### define:

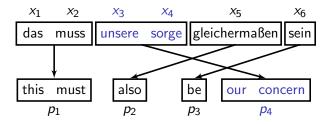
- ▶ source-language sentence words  $x_1, ..., x_N$
- phrase translation p = (s, e, t)
- ▶ translation derivation  $y = p_1, ..., p_L$



$$y = \{(1, 2, \text{this must}), (5, 5, \text{also}), (6, 6, \text{be}), \}$$

### define:

- ▶ source-language sentence words  $x_1, ..., x_N$
- phrase translation p = (s, e, t)
- translation derivation  $y = p_1, \dots, p_L$

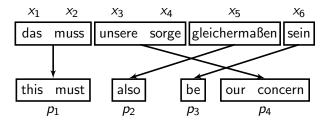


$$y = \{(1, 2, \text{this must}), (5, 5, \text{also}), (6, 6, \text{be}), (3, 4, \text{our concern})\}$$

### define:

- ▶ source-language sentence words  $x_1, ..., x_N$
- phrase translation p = (s, e, t)
- ▶ translation derivation  $y = p_1, ..., p_L$

## example:

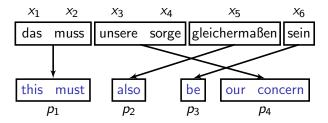


 $y = \{(1, 2, \text{this must}), (5, 5, \text{also}), (6, 6, \text{be}), (3, 4, \text{our concern})\}$ 

### define:

- ▶ source-language sentence words  $x_1, ..., x_N$
- phrase translation p = (s, e, t)
- ▶ translation derivation  $y = p_1, ..., p_L$

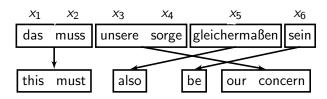
## example:



 $y = \{(1, 2, \text{this must}), (5, 5, \text{also}), (6, 6, \text{be}), (3, 4, \text{our concern})\}$ 

#### derivation:

$$y = \{(1, 2, \text{this must}), (5, 5, \text{also}), (6, 6, \text{be}), (3, 4, \text{our concern})\}$$

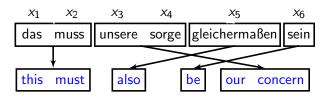


$$f(y) = h(e(y)) + \sum_{k=1}^{L} g(p_k) + \sum_{k=1}^{L-1} \eta |t(p_k) + 1 - s(p_{k+1})|$$

- ► language model score h
- ▶ phrase translation score g
- $\triangleright$  distortion penalty  $\eta$

#### derivation:

$$y = \{(1, 2, \text{this must}), (5, 5, \text{also}), (6, 6, \text{be}), (3, 4, \text{our concern})\}$$

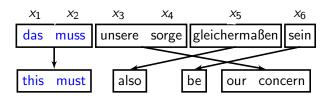


$$f(y) = h(e(y)) + \sum_{k=1}^{L} g(p_k) + \sum_{k=1}^{L-1} \eta |t(p_k) + 1 - s(p_{k+1})|$$

- ► language model score h
- ▶ phrase translation score g
- ightharpoonup distortion penalty  $\eta$

#### derivation:

$$y = \{(1, 2, \text{this must}), (5, 5, \text{also}), (6, 6, \text{be}), (3, 4, \text{our concern})\}$$

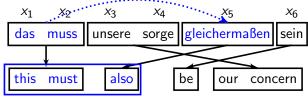


$$f(y) = h(e(y)) + \sum_{k=1}^{L} g(p_k) + \sum_{k=1}^{L-1} \eta |t(p_k) + 1 - s(p_{k+1})|$$

- ► language model score h
- ▶ phrase translation score g
- ightharpoonup distortion penalty  $\eta$

#### derivation:

 $y = \{(1, 2, \text{this must}), (5, 5, 2 \text{lso}), (6, 6, \text{be}), (3, 4, \text{our concern})\}$   $x_1 \dots x_2 \dots x_3 \dots x_4 \dots x_5 \dots x_6$ 



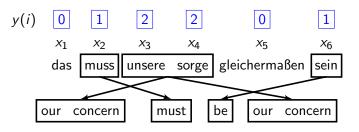
$$f(y) = h(e(y)) + \sum_{k=1}^{L} g(p_k) + \sum_{k=1}^{L-1} \eta |t(p_k) + 1 - s(p_{k+1})|$$

- ▶ language model score h
- ▶ phrase translation score g
- $\blacktriangleright$  distortion penalty  $\eta$

 $\mathcal{Y}'$ : only requires the total number of words translated to be N

$$\mathcal{Y}' = \{y : \sum_{i=1}^{N} y(i) = N \text{ and the distortion limit } d \text{ is satisfied}\}$$

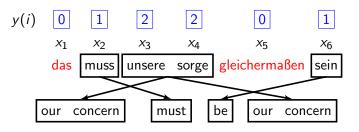
### example:



 $\mathcal{Y}'$ : only requires the total number of words translated to be N

$$\mathcal{Y}' = \{y : \sum_{i=1}^{N} y(i) = N \text{ and the distortion limit } d \text{ is satisfied}\}$$

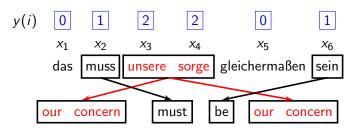
### example:



 $\mathcal{Y}'$ : only requires the total number of words translated to be N

$$\mathcal{Y}' = \{y : \sum_{i=1}^{N} y(i) = N \text{ and the distortion limit } d \text{ is satisfied}\}$$

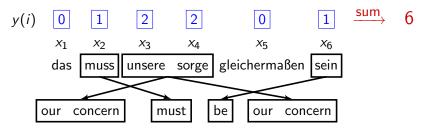
### example:



 $\mathcal{Y}'$ : only requires the total number of words translated to be N

$$\mathcal{Y}' = \{y : \sum_{i=1}^{N} y(i) = N \text{ and the distortion limit } d \text{ is satisfied}\}$$

### example:



## original:

$$\underset{y \in \mathcal{Y}}{\operatorname{arg\,max}\, f(y)}$$

$$\mathcal{Y} = \{ y : y(i) = 1 \ \forall i = 1 \dots N \}$$

$$oxed{1} oxed{1} \ldots oxed{1}$$

$$\arg\max_{y\in\mathcal{Y}'}f(y)$$

$$\arg\max_{y\in\mathcal{Y}'}f(y)$$
 such that  $\underbrace{y(i)=1\ \forall i=1\dots N}$ 

$$\mathcal{Y}' = \{ y : \sum_{i=1}^{N} y(i) = N \}$$

$$2 0 \dots 1$$

## original:

$$\underbrace{\operatorname{arg\,max} f(y)}_{y \in \mathcal{Y}}$$
exact DP is NP-hard

$$\mathcal{Y} = \{ y : y(i) = 1 \ \forall i = 1 \dots N \}$$

$$oxed{1}oxed{1}oxed{1}...oxed{1}$$

$$\underset{y \in \mathcal{Y'}}{\operatorname{arg} \max} f(y) \qquad \text{such that} \qquad \underbrace{y(i) = 1 \ \forall i = 1 \dots N}_{}$$

$$\mathcal{Y}' = \{ y : \sum_{i=1}^{N} y(i) = N \}$$

$$2 0 \dots 1$$

## original:

$$\underbrace{\operatorname{arg\,max} f(y)}_{y \in \mathcal{Y}}$$
exact DP is NP-hard

$$\mathcal{Y} = \{ y : y(i) = 1 \ \forall i = 1 \dots N \}$$

$$oxed{1}oxed{1}oxed{1}...oxed{1}$$

$$\underset{y \in \mathcal{Y'}}{\operatorname{arg} \max} f(y) \qquad \text{such that} \qquad \underbrace{y(i) = 1 \ \forall i = 1 \dots N}_{}$$

$$\mathcal{Y}' = \{ y : \sum_{i=1}^{N} y(i) = N \}$$

$$2 0 \dots 1$$

## original:

$$\underbrace{\operatorname{arg\,max} f(y)}_{y \in \mathcal{Y}}$$
exact DP is NP-hard

$$\mathcal{Y} = \{ y : y(i) = 1 \ \forall i = 1 \dots N \}$$

### rewrite:

$$\arg\max_{y\in\mathcal{Y'}}f(y)\qquad \text{such that}\qquad \underbrace{y(i)=1\ \forall i=1\dots N}$$
 can be solved efficiently by DP

an be solved efficiently by Di

$$\mathcal{Y}' = \{ y : \sum_{i=1}^{N} y(i) = N \}$$

original:

$$\underset{y \in \mathcal{Y}}{\operatorname{arg max}} f(y)$$
exact DP is NP-hard

$$\mathcal{Y} = \{ y : y(i) = 1 \ \forall i = 1 \dots N \}$$

$$oxed{1}oxed{1}\dotsoxed{1}$$

$$\underbrace{\arg\max_{y\in\mathcal{Y}'}f(y)}_{\text{can be solved efficiently by DP}} \text{ such that } \underbrace{y(i)=1 \ \forall i=1\dots N}_{\text{using Lagrangian relaxation}}$$

$$\mathcal{Y}' = \{ y : \sum_{i=1}^{N} y(i) = N \}$$

$$2 0 \dots 1$$
sum to N

### Iteration 1:

▶ update 
$$u(i)$$
:  $u(i) \leftarrow u(i) - \alpha(y(i) - 1)$ 

$$\alpha = 1$$

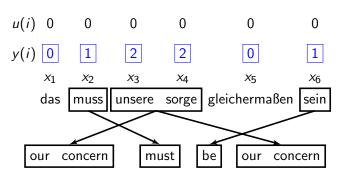
$$u(i) \quad 0 \quad 0 \quad 0 \quad 0 \quad 0$$

$$y(i)$$

$$x_1 \quad x_2 \quad x_3 \quad x_4 \quad x_5 \quad x_6$$
das muss unsere sorge gleichermaßen sein

### Iteration 1:

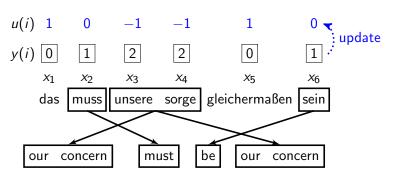
▶ update u(i):  $u(i) \leftarrow u(i) - \alpha(y(i) - 1)$   $\alpha = 1$ 



### Iteration 1:

▶ update 
$$u(i)$$
:  $u(i) \leftarrow u(i) - \alpha(y(i) - 1)$ 

$$\alpha = 1$$



das muss unsere sorge gleichermaßen sein

### Iteration 2:

▶ update 
$$u(i)$$
:  $u(i) \leftarrow u(i) - \alpha(y(i) - 1)$ 

$$\alpha = 0.5$$

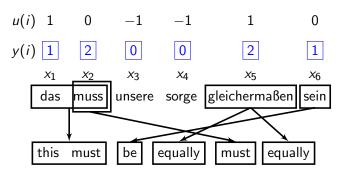
$$u(i) \ 1 \quad 0 \quad -1 \quad -1 \quad 1 \quad 0$$

$$y(i)$$

$$x_1 \quad x_2 \quad x_3 \quad x_4 \quad x_5 \quad x_6$$

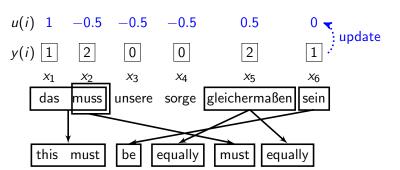
### Iteration 2:

▶ update u(i):  $u(i) \leftarrow u(i) - \alpha(y(i) - 1)$  $\alpha = 0.5$ 



### Iteration 2:

▶ update 
$$u(i)$$
:  $u(i) \leftarrow u(i) - \alpha(y(i) - 1)$   
 $\alpha = 0.5$ 

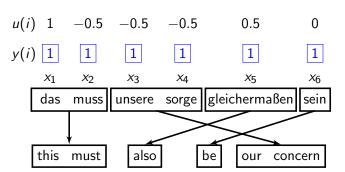


### Iteration 3:

▶ update 
$$u(i)$$
:  $u(i) \leftarrow u(i) - \alpha(y(i) - 1)$ 
 $\alpha = 0.5$ 
 $u(i) \ 1 \ -0.5 \ -0.5 \ -0.5 \ 0.5 \ 0$ 
 $y(i)$ 
 $x_1 \ x_2 \ x_3 \ x_4 \ x_5 \ x_6$ 
das muss unsere sorge gleichermaßen sein

### Iteration 3:

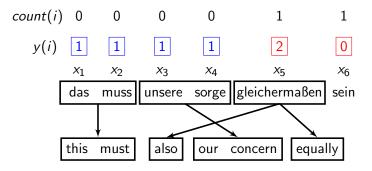
▶ update u(i):  $u(i) \leftarrow u(i) - \alpha(y(i) - 1)$  $\alpha = 0.5$ 



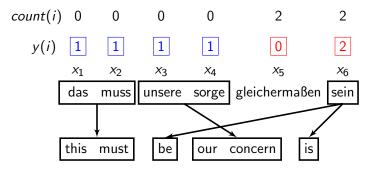
In some cases, we never reach y(i) = 1 for i = 1 ... N

If dual L(u) is not decreasing fast enough run for 10 more iterations count number of times each constraint is violated add 3 most often violated constraints

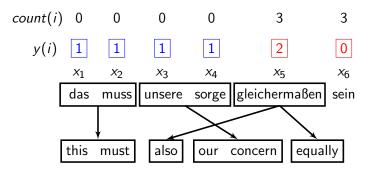
### Iteration 41:



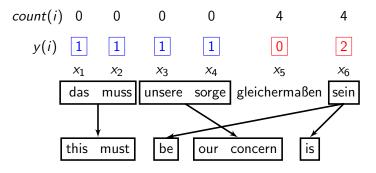
### Iteration 42:



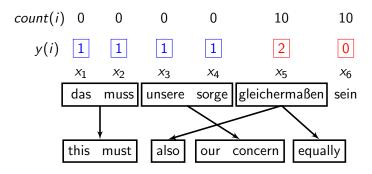
Iteration 43:



### Iteration 44:



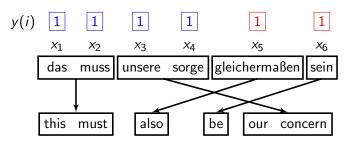
Iteration 50:



Iteration 51:

Add 2 hard constraints  $(x_5, x_6)$  to the dynamic program

Iteration 51:

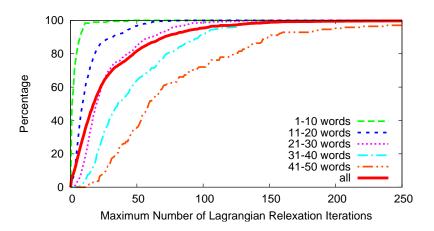


Add 2 hard constraints  $(x_5, x_6)$  to the dynamic program

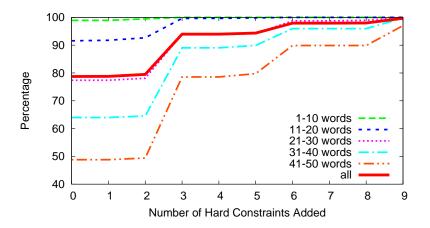
## Experiments: German to English

- ► Europarl data: German to English
- ► Test on 1,824 sentences with length 1-50 words
- Converged: 1,818 sentences (99.67%)

## **Experiments: Number of Iterations**



# Experiments: Number of Hard Constraints Required



# Experiments: Mean Time in Seconds

# words	1-10	11-20	21-30	31-40	41-50	All
mean	0.8	10.9	57.2	203.4	679.9	120.9
median	0.7	8.9	48.3	169.7	484.0	35.2

# Comparison to ILP Decoding

	(sec.)	(sec.)
1-10	275.2	132.9
11-15	2,707.8	1,138.5
16-20	20,583.1	3,692.6

## Summary

presented Lagrangian relaxation as a method for decoding in NLP

#### formal guarantees

- gives certificate or approximate solution
- can improve approximate solutions by tightening relaxation

#### efficient algorithms

- uses fast combinatorial algorithms
- can improve speed with lazy decoding

#### widely applicable

 demonstrated algorithms for a wide range of NLP tasks (parsing, tagging, alignment, mt decoding)

# Higher-order non-projective dependency parsing

**setup:** given a model for higher-order non-projective dependency parsing (sibling features)

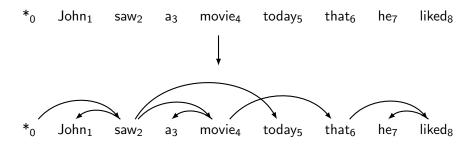
**problem:** find non-projective dependency parse that maximizes the score of this model

#### difficulty:

- model is NP-hard to decode
- complexity of the model comes from enforcing combinatorial constraints

**strategy:** design a decomposition that separates combinatorial constraints from direct implementation of the scoring function

# Non-Projective Dependency Parsing



Important problem in many languages.

Problem is NP-Hard for all but the simplest models.

## **Dual Decomposition**

A classical technique for constructing decoding algorithms.

Solve complicated models

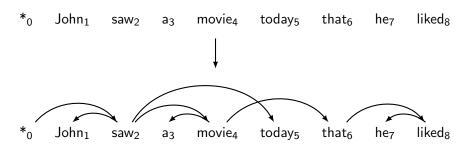
$$y^* = \arg\max_{y} f(y)$$

by decomposing into smaller problems.

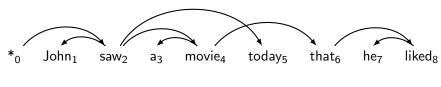
Upshot: Can utilize a toolbox of combinatorial algorithms.

- Dynamic programming
- Minimum spanning tree
- Shortest path
- Min-Cut
- ...

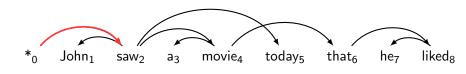
## Non-Projective Dependency Parsing



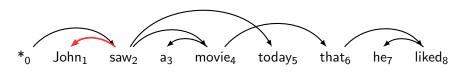
- ► Starts at the root symbol \*
- ▶ Each word has a exactly one parent word
- Produces a tree structure (no cycles)
- ► Dependencies can cross



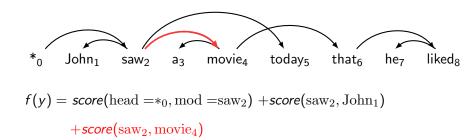
$$f(y) =$$

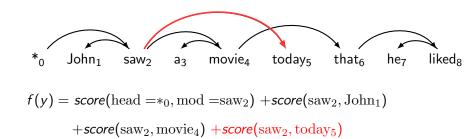


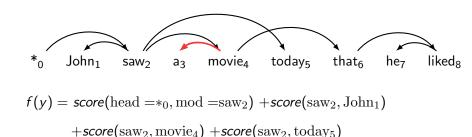
$$f(y) = score(head = *_0, mod = saw_2)$$



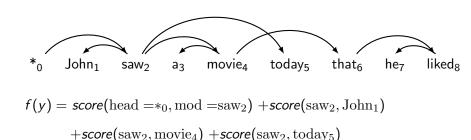
$$f(y) = score(head = *_0, mod = saw_2) + score(saw_2, John_1)$$





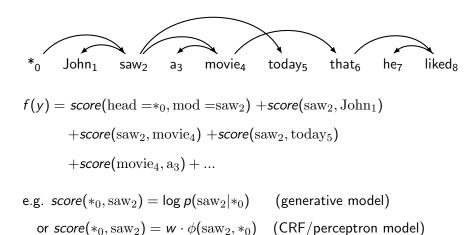


 $+score(movie_4, a_3) + ...$ 



e.g. 
$$score(*_0, saw_2) = log p(saw_2|*_0)$$
 (generative model)

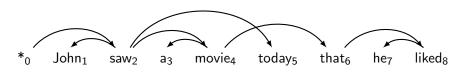
 $+score(movie_4, a_3) + ...$ 



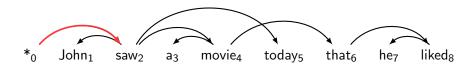
\*0 John<sub>1</sub> saw<sub>2</sub> a<sub>3</sub> movie<sub>4</sub> today<sub>5</sub> that<sub>6</sub> he<sub>7</sub> liked
$$f(y) = score(\text{head} = *_0, \text{mod} = \text{saw}_2) + score(\text{saw}_2, \text{John}_1)$$

$$+ score(\text{saw}_2, \text{movie}_4) + score(\text{saw}_2, \text{today}_5)$$

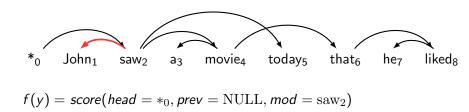
$$+ score(\text{movie}_4, \text{a}_3) + \dots$$
e.g.  $score(*_0, \text{saw}_2) = \log p(\text{saw}_2 | *_0)$  (generative model)
or  $score(*_0, \text{saw}_2) = w \cdot \phi(\text{saw}_2, *_0)$  (CRF/perceptron model)
$$y^* = \arg \max_{v} f(y) \Leftarrow \text{Minimum Spanning Tree Algorithm}$$



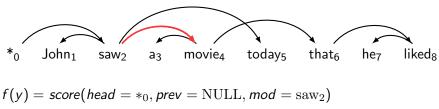
$$f(y) =$$



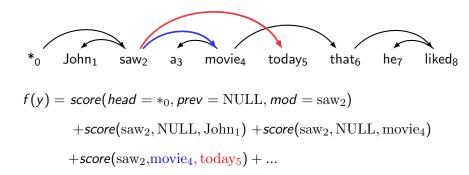
$$f(y) = score(head = *_0, prev = NULL, mod = saw_2)$$

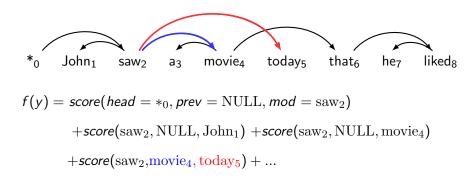


 $+score(saw_2, NULL, John_1)$ 



$$+score(saw2, NULL, John1) +score(saw2, NULL, movie4)$$





e.g.  $score(saw_2, movie_4, today_5) = log p(today_5|saw_2, movie_4)$ 

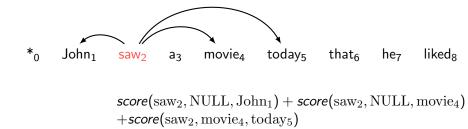
$$*_0$$
 John $_1$  saw $_2$  a $_3$  movie $_4$  today $_5$  that $_6$  he $_7$  liked $_8$   $f(y) = score(head = *_0, prev = NULL, mod = saw $_2$ )$ 

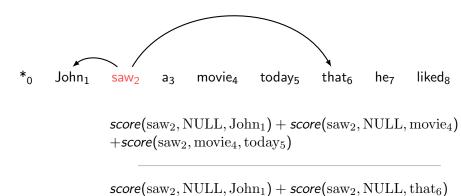
$$+score(saw_2, NULL, John_1) +score(saw_2, NULL, movie_4) +score(saw_2, movie_4, today_5) + ...$$

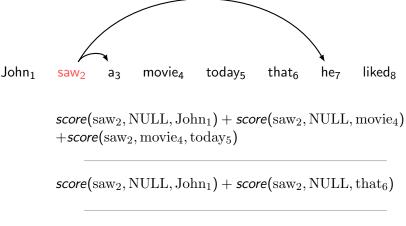
e.g.  $score(saw_2, movie_4, today_5) = log p(today_5|saw_2, movie_4)$ or  $score(saw_2, movie_4, today_5) = w \cdot \phi(saw_2, movie_4, today_5)$ 

\*0 John¹ saw² a³ movie⁴ today⁵ that6 he७ liked8 
$$f(y) = score(head = *_0, prev = \text{NULL}, mod = \text{saw}_2) \\ + score(\text{saw}_2, \text{NULL}, \text{John}_1) + score(\text{saw}_2, \text{NULL}, \text{movie}_4) \\ + score(\text{saw}_2, \text{movie}_4, \text{today}_5) + \dots \\ \text{e.g. } score(\text{saw}_2, \text{movie}_4, \text{today}_5) = \log p(\text{today}_5 | \text{saw}_2, \text{movie}_4) \\ \text{or } score(\text{saw}_2, \text{movie}_4, \text{today}_5) = w \cdot \phi(\text{saw}_2, \text{movie}_4, \text{today}_5) \\ y^* = \arg \max_y f(y) \Leftarrow \text{NP-Hard}$$

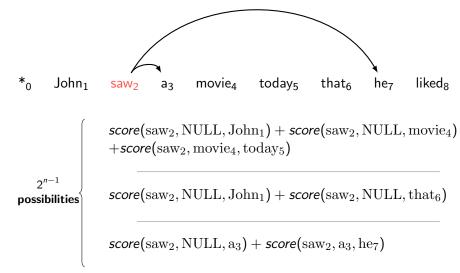
 $*_0$  John<sub>1</sub> saw<sub>2</sub> a<sub>3</sub> movie<sub>4</sub> today<sub>5</sub> that<sub>6</sub> he<sub>7</sub> liked<sub>8</sub>

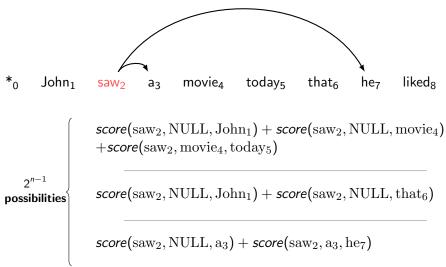






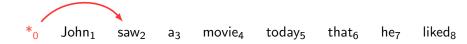
 $score(saw_2, NULL, a_3) + score(saw_2, a_3, he_7)$ 





Under Sibling Model, can solve for each word with Viterbi decoding.

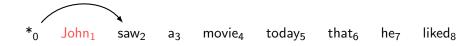
### Thought Experiment Continued



Idea: Do individual decoding for each head word using dynamic programming.

If we're lucky, we'll end up with a valid final tree.

### Thought Experiment Continued



Idea: Do individual decoding for each head word using dynamic programming.

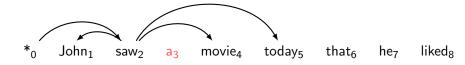
If we're lucky, we'll end up with a valid final tree.

### Thought Experiment Continued



Idea: Do individual decoding for each head word using dynamic programming.

If we're lucky, we'll end up with a valid final tree.



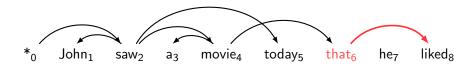
Idea: Do individual decoding for each head word using dynamic programming.



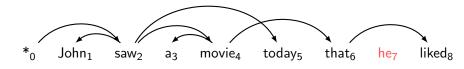
Idea: Do individual decoding for each head word using dynamic programming.



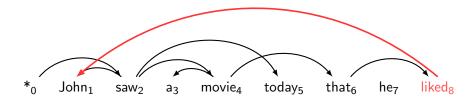
Idea: Do individual decoding for each head word using dynamic programming.



Idea: Do individual decoding for each head word using dynamic programming.



Idea: Do individual decoding for each head word using dynamic programming.



Idea: Do individual decoding for each head word using dynamic programming.

If we're lucky, we'll end up with a valid final tree.

But we might violate some constraints.

Goal 
$$y^* = \arg\max_{y \in \mathcal{Y}} f(y)$$

Goal 
$$y^* = \arg \max_{y \in \mathcal{Y}} f(y)$$

Rewrite as 
$$\underset{z \in \mathcal{Z}, y \in \mathcal{Y}}{\operatorname{argmax}} f(z) + g(y)$$

such that 
$$z = y$$

Goal 
$$y^* = \arg \max_{y \in \mathcal{Y}} f(y)$$

Rewrite as 
$$\underset{z \in \mathcal{Z}, \ y \in \mathcal{Y}}{\operatorname{argmax}} f(z) + g(y)$$
 
$$\underset{\pi}{\underset{\pi}{\underset{\text{All Possible}}{\overbrace{}}}} f(z) + g(y)$$
 such that  $z = y$ 

Goal 
$$y^* = \arg \max_{y \in \mathcal{Y}} f(y)$$

Rewrite as 
$$\underset{z \in \mathcal{Z}, \ y \in \mathcal{Y}}{\operatorname{argmax}} f(z) + g(y)$$

$$\underset{z \in \mathcal{Z}, \ y \in \mathcal{Y}}{\operatorname{All Possible}} \text{ Valid Trees}$$

$$\operatorname{such that } z = y$$

Goal 
$$y^* = \arg\max_{y \in \mathcal{Y}} f(y)$$

Sibling

Rewrite as  $\arg\max_{z \in \mathcal{Z}, y \in \mathcal{Y}} f(z) + g(y)$ 

All Possible Valid Trees

such that  $z = y$ 

Goal 
$$y^* = \arg\max_{y \in \mathcal{Y}} f(y)$$

Sibling Arc-Factored

Rewrite as  $\arg\max_{z \in \mathcal{Z}, \ y \in \mathcal{Y}} f(z) + g(y)$ 

All Possible Valid Trees

such that  $z = y$ 

Goal 
$$y^* = \arg\max_{y \in \mathcal{Y}} f(y)$$

Sibling Arc-Factored

Rewrite as  $\arg\max_{z \in \mathcal{Z}, y \in \mathcal{Y}} f(z) + g(y)$ 

All Possible Valid Trees

Such that  $z = y$ 

Constraint

Set penalty weights equal to 0 for all edges.

For k = 1 to K

Set penalty weights equal to 0 for all edges.

For k = 1 to K

 $z^{(k)} \leftarrow \mathsf{Decode} \; (f(z) + \mathsf{penalty}) \; \mathsf{by} \; \mathsf{Individual} \; \mathsf{Decoding}$ 

Set penalty weights equal to 0 for all edges.

For 
$$k = 1$$
 to  $K$ 

$$z^{(k)} \leftarrow \text{Decode}(f(z) + \text{penalty})$$
 by Individual Decoding

$$y^{(k)} \leftarrow \mathsf{Decode}\; (g(y) - \mathsf{penalty}) \; \mathsf{by} \; \mathsf{Minimum} \; \mathsf{Spanning} \; \mathsf{Tree}$$

Set penalty weights equal to 0 for all edges.

For 
$$k = 1$$
 to  $K$ 

$$z^{(k)} \leftarrow \text{Decode } (f(z) + \text{penalty}) \text{ by Individual Decoding}$$
  $y^{(k)} \leftarrow \text{Decode } (g(y) - \text{penalty}) \text{ by Minimum Spanning Tree}$  If  $y^{(k)}(i,j) = z^{(k)}(i,j)$  for all  $i,j$  Return  $(y^{(k)},z^{(k)})$ 

Set penalty weights equal to 0 for all edges.

For 
$$k = 1$$
 to  $K$ 

$$z^{(k)} \leftarrow \text{Decode } (f(z) + \text{penalty}) \text{ by Individual Decoding}$$
  $y^{(k)} \leftarrow \text{Decode } (g(y) - \text{penalty}) \text{ by Minimum Spanning Tree}$  If  $y^{(k)}(i,j) = z^{(k)}(i,j)$  for all  $i,j$  Return  $(y^{(k)},z^{(k)})$  Else Update penalty weights based on  $y^{(k)}(i,j) - z^{(k)}(i,j)$ 

# Penalties u(i,j) = 0 for all i,j

$$*_0$$
 John $_1$  saw $_2$  a $_3$  movie $_4$  today $_5$  that $_6$  he $_7$  liked $_8$   $z^* = rg \max_{z \in \mathcal{Z}} (f(z) + \sum_{i,j} u(i,j)z(i,j))$ 

$$*_0$$
 John $_1$  saw $_2$  a $_3$  movie $_4$  today $_5$  that $_6$  he $_7$  liked $_8$   $y^* = rg \max_{y \in \mathcal{Y}} (g(y) - \sum_{i,j} u(i,j)y(i,j))$ 

#### **Penalties**

u(i,j) = 0 for all i,j

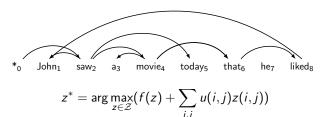
$$*_0$$
 John<sub>1</sub> saw<sub>2</sub>  $a_3$  movie<sub>4</sub> today<sub>5</sub> that<sub>6</sub> he<sub>7</sub> liked<sub>8</sub> 
$$z^* = \arg\max_{z \in \mathcal{Z}} (f(z) + \sum_{i,i} u(i,j)z(i,j))$$

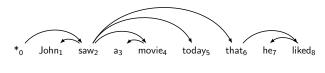
$$*_0$$
 John $_1$  saw $_2$  a $_3$  movie $_4$  today $_5$  that $_6$  he $_7$  liked $_8$   $y^*=rg\max_{y\in\mathcal{Y}}(g(y)-\sum_{i,j}u(i,j)y(i,j))$  Key

$$f(z) \Leftarrow \text{Sibling Model} \qquad g(y) \Leftarrow \text{Arc-Factored Model} \\ \mathcal{Z} \Leftarrow \text{No Constraints} \qquad \mathcal{Y} \Leftarrow \text{Tree Constraints} \\ y(i,j) = 1 \quad \text{if} \qquad y \text{ contains dependency } i,j$$

#### **Penalties**

u(i,j) = 0 for all i,j

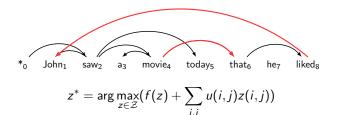


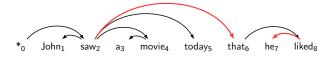


$$y^* = \arg\max_{y \in \mathcal{Y}} (g(y) - \sum_{i,j} u(i,j)y(i,j))$$

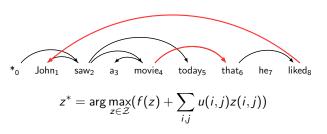
#### **Penalties**

u(i,j) = 0 for all i,j





$$y^* = \arg\max_{y \in \mathcal{Y}} (g(y) - \sum_{i,j} u(i,j)y(i,j))$$

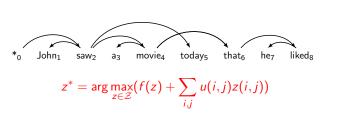


# Penalties u(i,j) = 0 for all i,j $\frac{\text{Iteration 1}}{u(8,1)} -1$ u(4,6) -1 u(2,6) 1

1

u(8,7)

$$y^* = \arg\max_{y \in \mathcal{Y}} (g(y) - \sum_{i,j} u(i,j)y(i,j))$$



#### Penalties

$$u(i,j) = 0 \text{ for all } i,j$$

$$\frac{\text{Iteration } 1}{u(8,1)} - 1$$

$$u(4,6) - 1$$

$$u(2,6) - 1$$

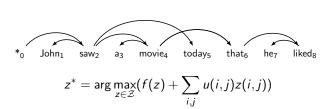
1

u(8,7)

#### Minimum Spanning Tree

$$*_0$$
 John $_1$  saw $_2$  a $_3$  movie $_4$  today $_5$  that $_6$  he $_7$  liked $_8$  
$$y^* = \arg\max_{y \in \mathcal{Y}} (g(y) - \sum_{i,j} u(i,j)y(i,j))$$
 Key

 $f(z) \Leftarrow \text{Sibling Model} \qquad g(y) \Leftarrow \text{Arc-Factored Model} \\ \mathcal{Z} \Leftarrow \text{No Constraints} \qquad \mathcal{Y} \Leftarrow \text{Tree Constraints} \\ y(i,j) = 1 \quad \text{if} \quad y \text{ contains dependency } i,j$ 



#### **Penalties**

$$u(i,j) = 0$$
 for all  $i,j$   

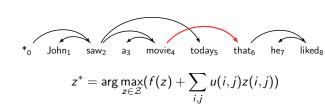
$$\frac{\text{Iteration 1}}{u(8,1)} -1$$

$$u(4,6) -1$$

1

u(2,6)u(8,7)

$$y^* = \arg\max_{y \in \mathcal{Y}} (g(y) - \sum_{i,j} u(i,j)y(i,j))$$



#### Minimum Spanning Tree



$$y^* = \arg\max_{y \in \mathcal{Y}} (g(y) - \sum_{i,j} u(i,j)y(i,j))$$

# Key

#### **Penalties**

$$u(i,j) = 0$$
 for all  $i,j$   
Iteration 1

$$u(8,1)$$
 -1  
 $u(4,6)$  -1  
 $u(2,6)$  1  
 $u(8,7)$  1

# Iteration 2 u(8,1)

$$u(4,6)$$
 -2  $u(2,6)$  2  $u(8,7)$  1

-1

John<sub>1</sub>

John<sub>1</sub>

movie<sub>4</sub>

$$z^* = \arg\max_{z \in \mathcal{Z}} (f(z) + \sum_{i,j} u(i,j)z(i,j))$$

today<sub>5</sub>

# Minimum Spanning Tree

today<sub>5</sub>

that<sub>6</sub> he<sub>7</sub>

he<sub>7</sub>

that<sub>6</sub>

liked<sub>8</sub>

liked<sub>s</sub>

u(4,6)u(2,6)u(8,7)

1

1

-1

-2

**Penalties** u(i,j) = 0 for all i,jIteration 1 u(8,1) -1 u(4,6) -1

u(2,6) 1 u(8,7)

Iteration 2 u(8,1)

\*0

# $y^* = \arg\max_{y \in \mathcal{Y}} (g(y) - \sum_{i:i} u(i,j)y(i,j))$

y(i, j) = 1 if y contains dependency i, j

movie<sub>4</sub>

$$(ey \qquad \qquad f(z) \qquad \Leftarrow \quad Sibling \; Model$$

аз

аз

$$\begin{array}{cccc} \mathsf{Key} & & & \\ f(z) & & \Leftarrow & \mathsf{Sibling\ Model} \\ \mathcal{Z} & & \Leftarrow & \mathsf{No\ Constraints} \end{array}$$

saw<sub>2</sub>

$$\mathsf{del} \qquad \qquad \mathsf{g}(y) \; \leftarrow \;$$

$$g(y) \Leftarrow Arc$$
-Factored Model  $\mathcal{Y} \Leftarrow Tree Constraints$ 

#### that6 liked<sub>8</sub> John₁ movie<sub>4</sub> today<sub>5</sub> $z^* = \arg\max_{z \in \mathcal{Z}} (f(z) + \sum_{i,j} u(i,j)z(i,j))$

# Minimum Spanning Tree

sawa

he<sub>7</sub>

liked<sub>8</sub>

$$y^* = arg \max_{y \in \mathcal{Y}} (g(y) - \sum_{i:i} u(i,j)y(i,j))$$

аз

movie<sub>4</sub>

\*0

John<sub>1</sub>

$$f(z) \Leftarrow \text{Sibling Model} \qquad g(y) \Leftarrow \text{Arc-Factored Model} \\ \mathcal{Z} \Leftarrow \text{No Constraints} \qquad \mathcal{Y} \Leftarrow \text{Tree Constraints} \\ y(i,j) = 1 \quad \text{if} \quad y \text{ contains dependency } i,j$$

today<sub>5</sub>

that<sub>6</sub>

#### **Penalties** u(i,j) = 0 for all i,j

$$\frac{\text{Iteration 1}}{u(8,1)} -1$$

$$u(4,6)$$
 -1  $u(2,6)$  1

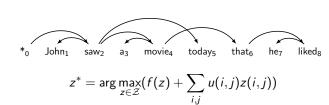
$$u(8,7)$$
 1

Iteration 2 
$$u(8,1)$$

$$u(4,6)$$
 -2  $u(2,6)$  2

-1

$$u(8,7)$$
 1



# Minimum Spanning Tree



$$y^* = \arg\max_{y \in \mathcal{Y}} (g(y) - \sum_{i,j} u(i,j)y(i,j))$$

## **Penalties**

$$u(i,j) = 0$$
 for all  $i,j$   
Iteration 1

-1

$$u(8,1)$$
 -  $u(4,6)$  -

$$u(2,6)$$
 1  $u(8,7)$  1

$$u(8,1)$$
 -1  $u(4,6)$  -2

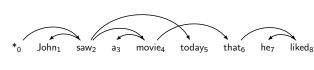
$$u(2,6)$$
 2  $u(8,7)$  1

$$u(8,7)$$
 1

liked<sub>8</sub>

$$z^* = arg \max_{z \in \mathcal{Z}} (f(z) + \sum_{i,j} u(i,j)z(i,j))$$

# Minimum Spanning Tree



$$y^* = \arg\max_{y \in \mathcal{Y}} (g(y) - \sum_{i,i} u(i,j)y(i,j))$$

#### **Penalties** u(i,j) = 0 for all i,j

-1

1

2

$$u(8,1)$$
  
 $u(4,6)$ 

Iteration 2

u(2,6)

$$u(8,1)$$
 -1  $u(4,6)$  -2

$$u(8,7)$$
 1

#### Converged

$$y^* = \arg \max_{y \in \mathcal{V}} f(y) + g(y)$$

f(z)

$$\Leftarrow$$
 Sibling Model  $g(y) \Leftarrow$  Arc-Factored Model

No Constraints

y contains dependency i, j

Tree Constraints

#### Guarantees

#### Theorem

If at any iteration  $y^{(k)}=z^{(k)}$ , then  $(y^{(k)},z^{(k)})$  is the global optimum.

In experiments, we find the global optimum on 98% of examples.

#### Guarantees

#### **Theorem**

If at any iteration  $y^{(k)}=z^{(k)}$ , then  $(y^{(k)},z^{(k)})$  is the global optimum.

In experiments, we find the global optimum on 98% of examples.

If we do not converge to a match, we can still return an approximate solution (more in the paper).

#### **Extensions**

► Grandparent Models



$$f(y) = ... + score(gp = *_0, head = saw_2, prev = movie_4, mod = today_5)$$

► Head Automata (Eisner, 2000)

Generalization of Sibling models

Allow arbitrary automata as local scoring function.

# **Experiments**

#### Properties:

- Exactness
- Parsing Speed
- ► Parsing Accuracy
- Comparison to Individual Decoding
- ► Comparison to LP/ILP

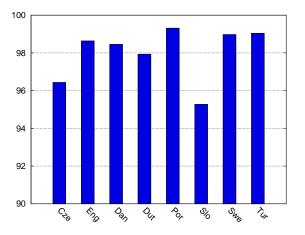
#### Training:

Averaged Perceptron (more details in paper)

#### Experiments on:

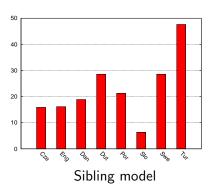
- CoNLL Datasets
- ► English Penn Treebank
- Czech Dependency Treebank

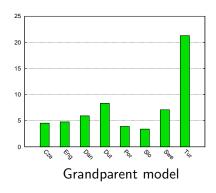
# How often do we exactly solve the problem?



 Percentage of examples where the dual decomposition finds an exact solution.

# Parsing Speed





- ▶ Number of sentences parsed per second
- Comparable to dynamic programming for projective parsing

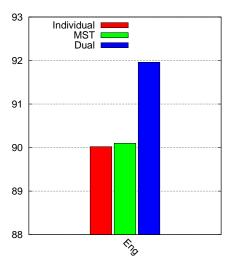
# Accuracy

	Arc-Factored	Prev Best	Grandparent
Dan	89.7	91.5	91.8
Dut	82.3	85.6	85.8
Por	90.7	92.1	93.0
Slo	82.4	85.6	86.2
Swe	88.9	90.6	91.4
Tur	75.7	76.4	77.6
Eng	90.1		92.5
Cze	84.4	_	87.3

Prev Best - Best reported results for CoNLL-X data set, includes

- ► Approximate search (McDonald and Pereira, 2006)
- ▶ Loop belief propagation (Smith and Eisner, 2008)
- ▶ (Integer) Linear Programming (Martins et al., 2009)

# Comparison to Subproblems



F<sub>1</sub> for dependency accuracy

# Comparison to LP/ILP

Martins et al.(2009): Proposes two representations of non-projective dependency parsing as a linear programming relaxation as well as an exact ILP.

- ► LP (1)
- ▶ LP (2)
- ► II P

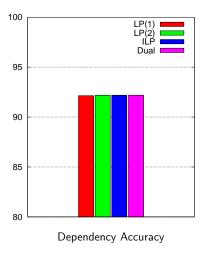
Use an LP/ILP Solver for decoding

#### We compare:

- Accuracy
- Exactness
- Speed

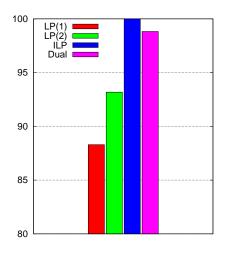
Both LP and dual decomposition methods use the same model, features, and weights w.

# Comparison to LP/ILP: Accuracy



▶ All decoding methods have comparable accuracy

# Comparison to LP/ILP: Exactness and Speed



12 10 6

Percentage with exact solution

Sentences per second

#### References I

- Y. Chang and M. Collins. Exact Decoding of Phrase-based Translation Models through Lagrangian Relaxation. In *To appear proc. of EMNLP*, 2011.
- J. DeNero and K. Macherey. Model-Based Aligner Combination Using Dual Decomposition. In *Proc. ACL*, 2011.
- J. Duchi, D. Tarlow, G. Elidan, and D. Koller. Using Combinatorial Optimization within Max-Product Belief Propagation. In *NIPS*, pages 369–376, 2007.
- D. Klein and C.D. Manning. Factored A\* Search for Models over Sequences and Trees. In *Proc IJCAI*, volume 18, pages 1246–1251. Citeseer, 2003.
- N. Komodakis, N. Paragios, and G. Tziritas. Mrf energy minimization and beyond via dual decomposition. *IEEE Transactions on Pattern Analysis and Machine Intelligence*, 2010. ISSN 0162-8828.

#### References II

- Terry Koo, Alexander M. Rush, Michael Collins, Tommi Jaakkola, and David Sontag. Dual decomposition for parsing with non-projective head automata. In *EMNLP*, 2010. URL http://www.aclweb.org/anthology/D10-1125.
- B.H. Korte and J. Vygen. *Combinatorial Optimization: Theory and Algorithms*. Springer Verlag, 2008.
- A.M. Rush and M. Collins. Exact Decoding of Syntactic Translation Models through Lagrangian Relaxation. In *Proc.* ACL, 2011.
- A.M. Rush, D. Sontag, M. Collins, and T. Jaakkola. On Dual Decomposition and Linear Programming Relaxations for Natural Language Processing. In *Proc. EMNLP*, 2010.
- D.A. Smith and J. Eisner. Dependency Parsing by Belief Propagation. In *Proc. EMNLP*, pages 145–156, 2008. URL http://www.aclweb.org/anthology/D08-1016.