



On Abstraction Refinement for Program Analyses in Datalog

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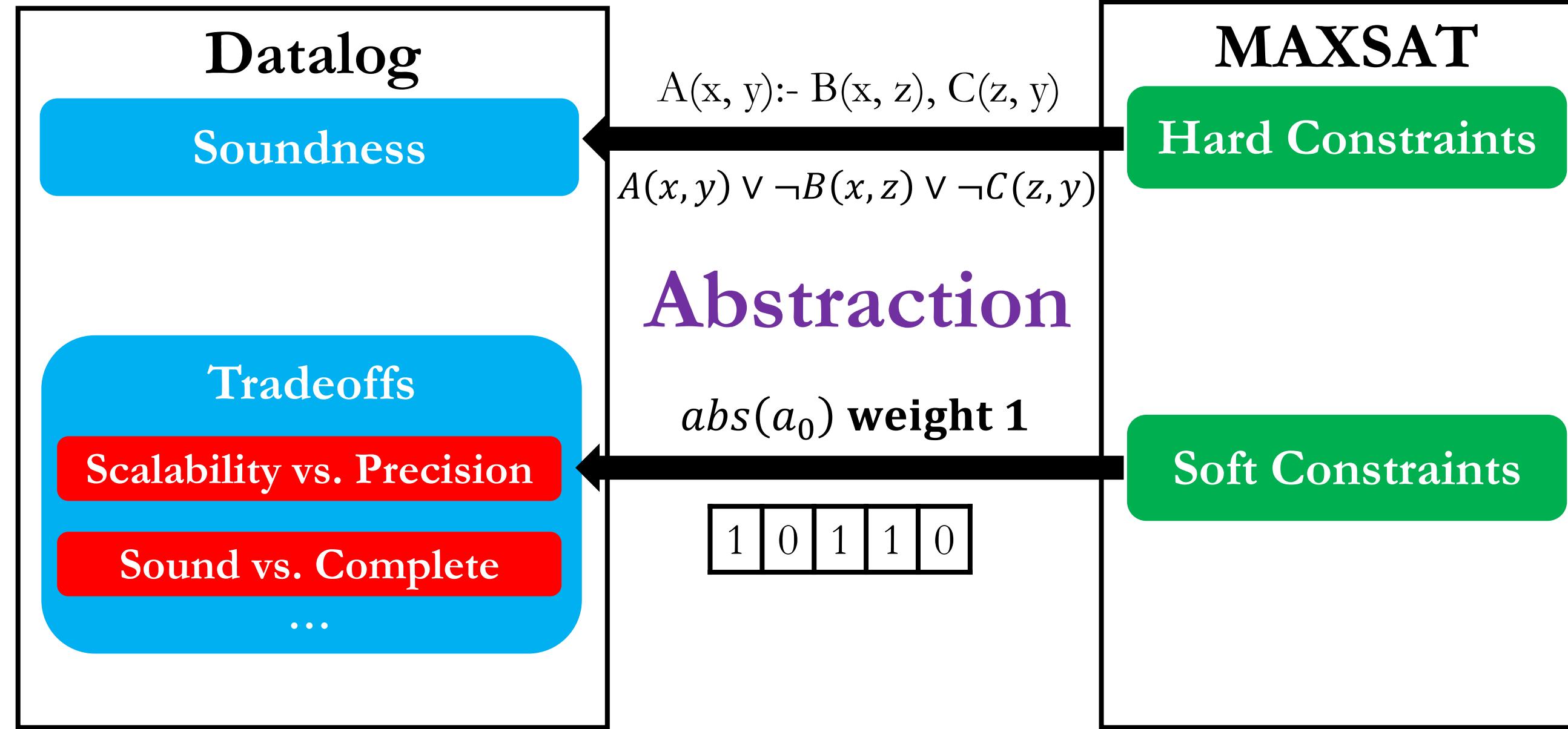
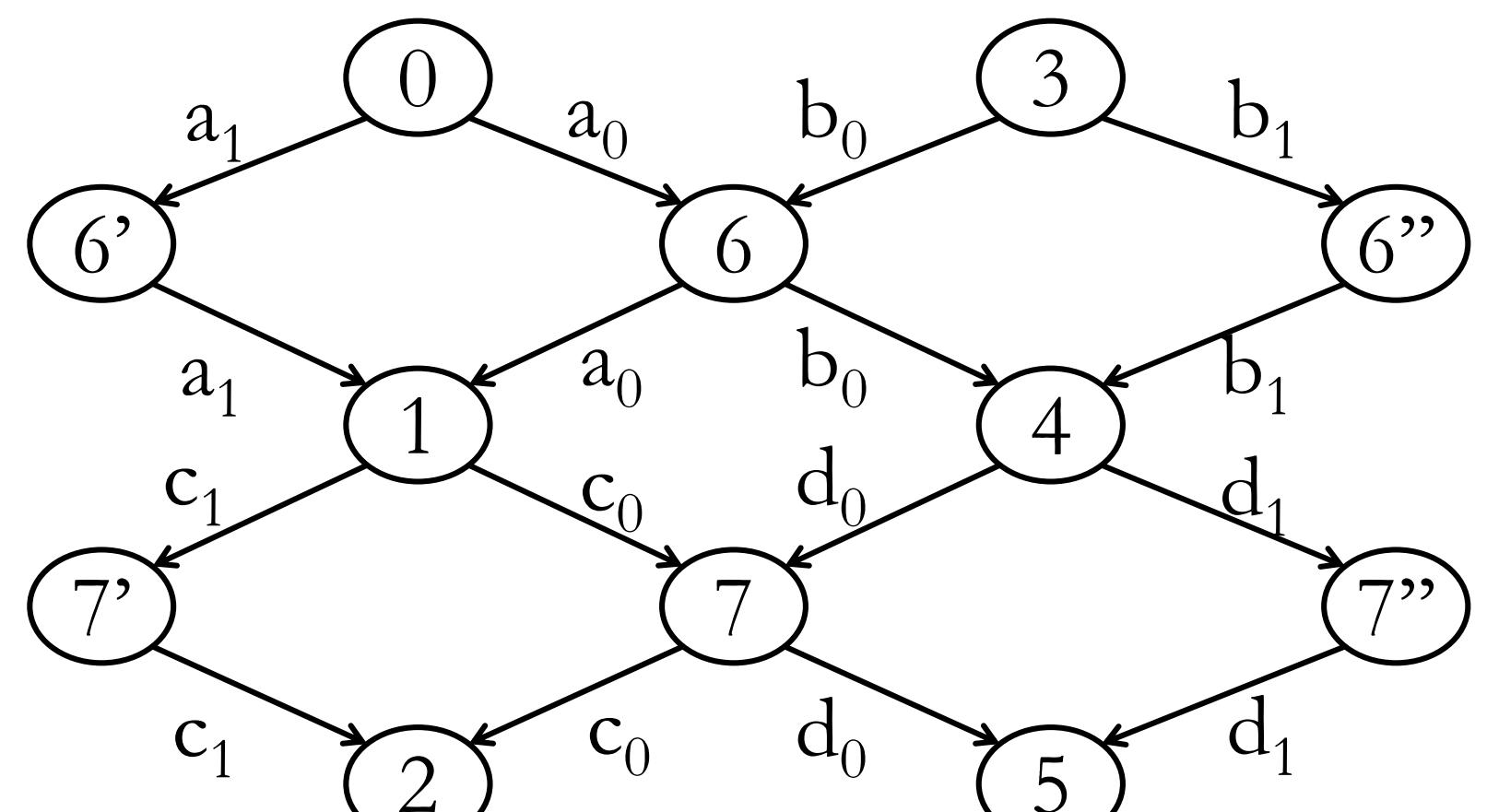
Pointer Analysis As Graph Reachability

```
f () {
    v1 = new ...;
    v2 = id1(v1);
    v3 = id2(v2);
q2:assert(v3 != v1);
}

id1(v) { return v; }
```

```
g () {
    v4 = new ...;
    v5 = id1(v4);
    v6 = id2(v5);
q1:assert(v6 != v1);
}

id2(v) { return v; }
```



Encode As MAXSAT

Hard constraints:

$$\begin{aligned} & \text{path}(0, 0) \wedge \\ & (\text{path}(0, 6) \vee \neg \text{path}(0, 0) \vee \neg \text{abs}(a_0)) \wedge \\ & (\text{path}(0, 1) \vee \neg \text{path}(0, 6) \vee \neg \text{abs}(a_0)) \wedge \\ & (\text{path}(0, 7) \vee \neg \text{path}(0, 1) \vee \neg \text{abs}(c_0)) \wedge \\ & (\text{path}(0, 4) \vee \neg \text{path}(0, 6) \vee \neg \text{abs}(b_0)) \wedge \\ & \dots \end{aligned}$$

Soft constraints:

$$\begin{aligned} & (\text{abs}(a_0) \text{ weight 1}) \wedge (\text{abs}(b_0) \text{ weight 1}) \wedge \\ & (\text{abs}(c_0) \text{ weight 1}) \wedge (\text{abs}(d_0) \text{ weight 1}) \wedge \\ & (\neg \text{path}(0, 2) \text{ weight 5}) \wedge (\neg \text{path}(0, 5) \text{ weight 5}) \end{aligned}$$

Solution:

$\text{path}(0, 0) = \text{true}$, $\text{path}(0, 6) = \text{false}$, $\text{path}(0, 1) = \text{false}$,
 $\text{path}(0, 4) = \text{false}$, $\text{path}(0, 7) = \text{false}$, $\text{path}(0, 2) = \text{false}$,
 $\text{path}(0, 5) = \text{false}$, $\text{abs}(a_0) = \text{false}$, $\text{abs}(b_0) = \text{true}$,
 $\text{abs}(c_0) = \text{true}$, $\text{abs}(d_0) = \text{true}$.

Graph Reachability in Datalog

Input relations:

$\text{edge}(i, j, n)$, $\text{abs}(n)$

Output relations:

$\text{path}(i, j)$

Rules:

$\text{path}(i, i)$.

$\text{path}(i, j) :- \text{path}(i, k), \text{edge}(k, j, n), \text{abs}(n)$.

Input tuples:

$\text{edge}(0, 6, a_0)$, $\text{edge}(0, 6', a_1)$, $\text{edge}(3, 6, b_0)$,
 \dots

$\text{abs}(a_0) \oplus \text{abs}(a_1)$, $\text{abs}(b_0) \oplus \text{abs}(b_1)$,
 $\text{abs}(c_0) \oplus \text{abs}(c_1)$, $\text{abs}(d_0) \oplus \text{abs}(d_1)$.

Query Tuple

Original Query

q1: $\text{path}(0, 5)$

assert ($v_6 \neq v_1$)

q2: $\text{path}(0, 2)$

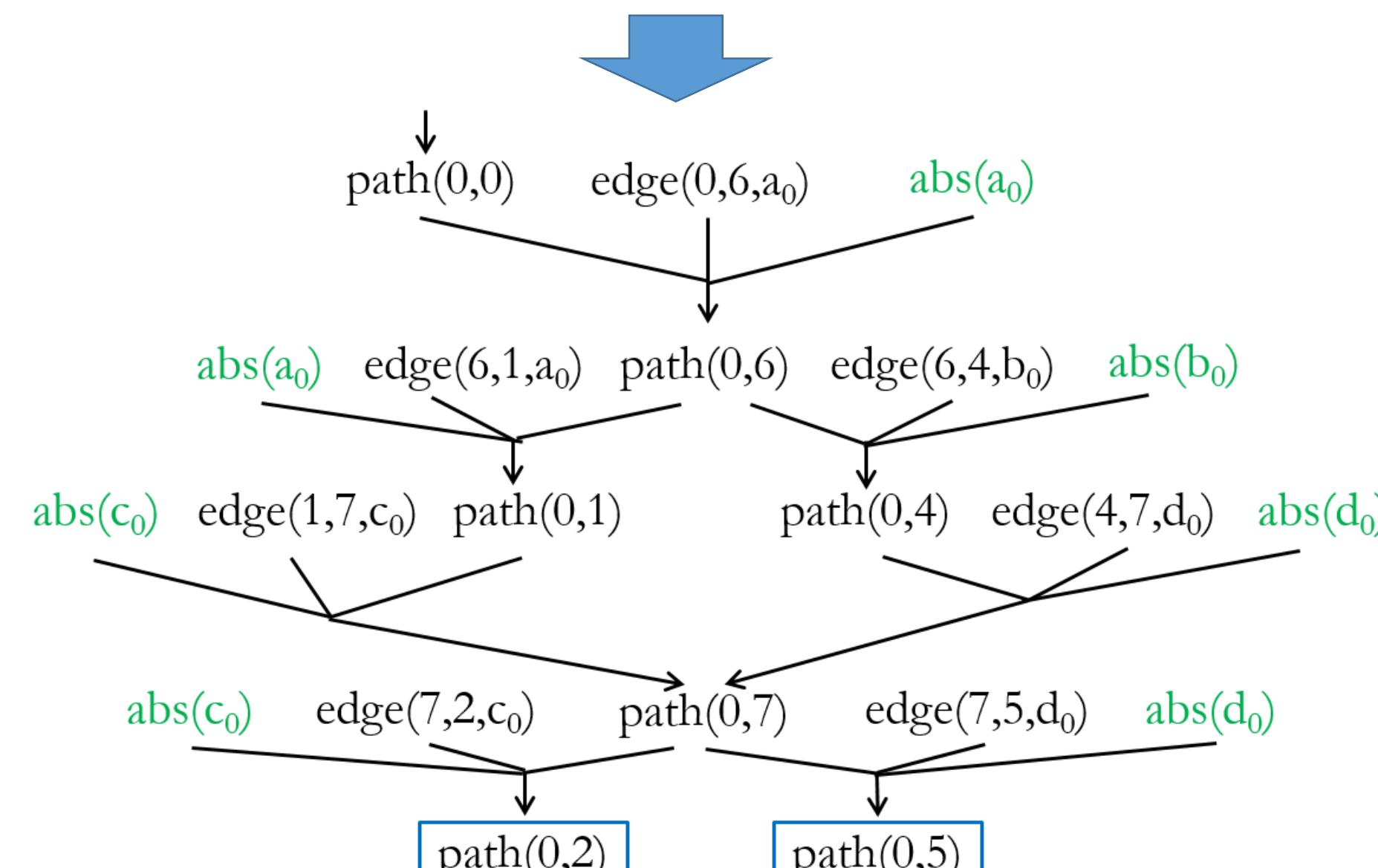
assert ($v_3 \neq v_1$)

Derivation As Counterexample

$\text{abs}(a_0)$
 $\text{abs}(b_0)$
 $\text{abs}(c_0)$
 $\text{abs}(d_0)$



$\text{path}(0, 0)$.
 $\text{path}(0, 6) :- \text{path}(0, 0), \text{edge}(0, 6, a_0), \text{abs}(a_0)$.
 $\text{path}(0, 1) :- \text{path}(0, 6), \text{edge}(6, 1, a_0), \text{abs}(a_0)$.
 $\text{path}(0, 7) :- \text{path}(0, 1), \text{edge}(1, 7, c_0), \text{abs}(c_0)$.
 $\text{path}(0, 2) :- \text{path}(0, 7), \text{edge}(7, 2, c_0), \text{abs}(c_0)$.
 $\text{path}(0, 4) :- \text{path}(0, 6), \text{edge}(6, 4, b_0), \text{abs}(b_0)$.
 $\text{path}(0, 7) :- \text{path}(0, 4), \text{edge}(4, 7, d_0), \text{abs}(d_0)$.
 $\text{path}(0, 5) :- \text{path}(0, 7), \text{edge}(7, 5, d_0), \text{abs}(d_0)$.
 \dots



Experiment Result

	queries	abstraction size			iterations
		total	resolved	final	
toba-s	7	7	0	170	18K
javasrc-p	46	46	0	470	18K
weblech	5	5	2	140	31K
hedc	47	47	6	730	29K
antlr	143	143	5	970	29K
luindex	138	138	67	1K	40K
lusearch	322	322	29	1K	39K
schroeder-m	51	51	25	450	58K
k = 4, 3h28m					
k = 3, 590s					
k = 2, 214s					
k = 1, 153s					

