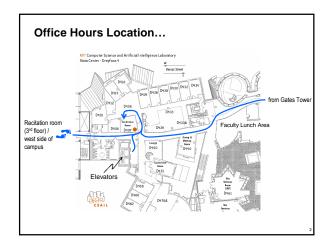
# 6.001 Recitation 7: Data Abstraction II

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## **Announcements**

- Solutions, handouts, etc.:
  - http://people.csail.mit.edu/dalleyg/6.001/SP2007/
  - primes-in-range discussion & orders of growth
- Office Hours
  - Thursdays, 2-3PM, 32-D407



### Overview

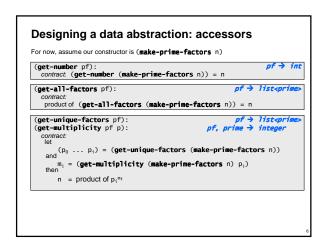
- Today: prime factorization, an extended example
- This is a nice example for several reasons:
  - Interesting design decisions
  - · Practice with writing types
    - $-\mbox{ (using prime and pf as new types)}$
  - Related to primality testing from yesterday's lecture
  - primes are also important to Project 1...which is due next Friday.

# Designing a data abstraction

• Prime factorization: representing an integer as the product of its prime factors

What about 1? What about 0? What about negative integers?

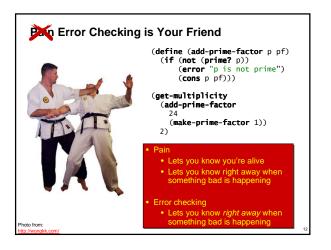
# Designing a data abstraction: constructors New types prime = subset of integers that are prime pf = prime factorization data type (make-prime-factors n): (make-prime-factors 40) → 2\*2\*2\*5 (make-prime-factors factors): (make-prime-factors (1ist 2 2 2 5)) → 2\*2\*2\*5 (make-prime-factors p): (make-prime-factor p pf): (add-prime-factor p pf): (make-prime-factor 5 (make-prime-factors 2)) → 2\*5 prime-factors-of-1: (add-prime-factor p pf): (add-prime-factor ppf): (add-pri



## Designing a data abstraction: operators pf,pf → boolean (**=pf** pf1 pf2): tests whether two factorizations are the same (divides-pf? pf1 pf2): pf,pf → boolean tests whether pf1 divides evenly into pf2 pf,prime → boolean (has-factor? pf p): tests whether p is a prime factor of pf(\***pf** pf1 pf2): $pf, pf \rightarrow pf$ returns factorization of n1\*n2 $pf, pf \rightarrow pf$ (/pf pf1 pf2): returns factorization of n1/n2 if n2 divides n1 pf.pf → pt (gcd-pf pf1 pf2): returns factorization of greatest common divisor of pf1 $\,$ and pf2 $\,$ +pf, -pf Not really appropriate for this data type. The only way to do it is converting to integer and then factorizing again.

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Many different representations are possible

(factorize 40); 40 = 2*2*2*5
\Rightarrow (2 2 2 5) \qquad 2*2*2*5 \text{ (sorted order)}
\Rightarrow (2 5 2 2) \qquad 2*5*2*2 \text{ (order doesn't matter)}
\Rightarrow ((2 3) (5 1)) \qquad 2^3 * 5^1
\Rightarrow (40 (2 5)) \qquad \text{stores n and its unique factors}
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# Respect abstraction boundaries (define (\*pf-clean pf1 pf2) (make-prime-factors (append (get-all-factors pf1) (get-all-factors pf2))) (define (\*pf-dirty pf1 pf2) (append pf1 pf2)) Procedures inside the abstraction boundary "know" that the real representation is (2 5 2 2), and depend on it make-prime-factors \*pf-clean get-all-factors \*pf-dirty Abstraction boundary Procedures outside don't care about the representation

## Summary of data abstraction design

- Choose <u>constructors</u> and <u>accessors</u> that are <u>useful</u> to clients and that make it possible to write the operators you need
  - Constructors and accessors should be <u>complete</u>: you need to be able to construct every possible object in the domain, and you need to be able to get out enough data to reconstruct the object
  - Write down the <u>contract</u> between the constructors and accessors
- 2. Choose <u>representation</u> that is appropriate to the operators you need (that makes the operators <u>readable</u> and <u>efficient</u>)
  - $\bullet \quad \text{Write down the } \underline{\text{assumptions}} \text{ implicit in your representation} \\$
- 3. Respect abstraction boundaries as much as possible
  - Even within your abstraction's own code
  - Another way to say it: Minimize the amount of code that "knows" what the real representation is.

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