## 6.001 Recitation 3:

# Substitution Model, Recursion, and Iteration, Oh My!

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### The Substitution Model

Newtonian physics, quantum mechanics, relativity, and Scheme...

#### Rules of the substitution model:

- 1. if **self-evaluating**, (e.g. number, string, #t, #f), just return value.
- 2. if **name**, replace it with value associated with that name.
- 3. if **special form** (e.g. if, define, lambda), follow the special form's rules for evaluation
  - (a) if **lambda**, create a procedure
  - (b) ...
- 4. if combination  $(e_0 \ e_1 \ \dots \ e_n)$ ,
  - (a) Evaluate subexpressions  $e_i$  in any order to produce values ( $v_0 \ v_1 \ \dots \ v_n$ )
  - (b) If  $v_0$  is **primitive procedure** (e.g. +), just apply it to  $v_1 \ldots v_n$
  - (c) if  $v_0$  is **compound procedure** (created by lambda):
    - i. Substitute  $v_1 \ldots v_n$  for corresponding parameters in body of procedure, then repeat on body

### Hints:

- Do one thing at a time
- Take small steps
- Sometimes we have a choice of what to do next

```
; Evaluate the following expression, one
; step at a time using the substitution model:
(if (= (+ 5 2) 7) (* (+ 2 3) 5) (/ 4 (- 7 5)))
(if (= 7 7) (* (+ 2 3) 5) (/ 4 (- 7 5)))
(if #t (* (+ 2 3) 5) (/ 4 (- 7 5)))
(if #t (* 5 5) (/ 4 (- 7 5)))
(if #t 25 (/ 4 (- 7 5)))
25
```

### Recursive versus Iterative

#### From Webster:

- recur: to come up again for consideration
- recursion: something that calls itself
- iterate: repeating, each time getting closer to the desired result

#### What makes a procedure recursive?

```
it calls itself (possibly indirectly)
```

#### What makes a process recursive?

```
deferred operations
```

## Designing Recursive Algorithms

- wishful thinking
- decompose the problem
- identify the simple/smallest parts (modular programming)

```
(define (odd? x)
  (not (even? x)))

(define (even? x)
   (not (odd? x)))

; fix even? ...
(define (even? x)
  (if (= x 0) #t (odd? (- x 1))))
```

- Enumerate assumptions
- Test/predicate for base case (termination of recursive unwinding)
- Recursive call

## **High-Level Questions**

What is "recursive recursion?"

- the width of the expression (substitution model) keeps growing
- deferred operations

What is "iterative recursion?"

- constant space
- no pending/deferred operations

Why do we like recursive procedures?

many problems are easy to think about as loops

Why do some programming guides discourage writing recursive programs?

- limited stack depth
- limited memory
- ullet encourage the use of for loops

## $Recursive \rightarrow Iterative Recipe$

- 1. What can we use as our partial answer
  - can we "accumulate" something?
- 2. How do we keep track of what's left to do

- some sort of counter (count up, count down?)
- 3. How do we update these variables
  - ullet the partial answer & counter
- 4. What's the base case?
- 5. Write out a table
  - fixed size (# of variables) because there are no pending operations
- 6. Almost always we will write a helper procedure that is the recursive part because we have extra variables to keep track of!

## biggie-size Returns!

We encode an order as multi-digit number where each digit represents a combo. For example, the order 327 represents a Triple, Double, and biggie-sized Triple. (biggie-size = regular+4)

Is this recursive or iterative?

```
It generates a recursive process.
```

Rewrite as the other type:

## begin Special Form

(begin expr1 expr2 ...): The expressions are evaluated sequentially from left to right, and the value of the last expression is returned. This expression type is used to sequence side effects such as input and output.

- allows multiple statements
- scoping
- implicit begin inside define, cond, lambda, ...

Consider:

```
(define (foo a)
  (begin
     (define b 5)
     (define (foo-helper x) (+ a x))
     (foo-helper b)))
```

```
To what do the following expressions evaluate?

foo \rightarrow #[procedure] foo-helper \rightarrow ERROR, undefined variable

a \rightarrow ERROR, undefined variable

foo-helper \rightarrow ERROR, undefined variable

foo-helper \rightarrow ERROR, undefined variable

a \rightarrow ERROR, undefined variable

b \rightarrow ERROR, undefined variable
```

### Order of Evaluation

Consider the following code:

```
; What does the following generate?
(count1 5)
=> 543210
```

```
; What does the following generate?
(count2 4)
=> 123455
```

Evaluate using the substitution model:

```
(count1 5)
(begin (our-display 5) (count1 4)) <PRINT 5>
(count1 4)
(begin (our-display 4) (count1 3)) <PRINT 4>
(count1 3)
(begin (our-display 3) (count1 2)) <PRINT 3>
(count1 2)
(begin (our-display 2) (count1 1)) <PRINT 2>
(count1 1)
(begin (our-display 1) (count1 0)) <PRINT 1>
(count1 0)
=> 0
```

```
(count2 5)
(begin (count2 4) (our-display 5))
(begin (begin (count2 3) (our-display 4)) (our-display 5))
(begin (begin (begin (count2 2) (our-display 3)) (our-display 4)) (our-display 5))
(begin (begin (begin (begin (count2 1) (our-display 2)) (our-display 3)) (our-display 4)) (our-display 5))
(begin (begin (begin (begin (begin (count2 0) (our-display 1)) (our-display 2)) (our-display 3)) (our-display 6))
(begin (begin (begin (begin (our-display 1) (our-display 2)) (our-display 3)) (our-display 4)) (our-display 6))
(begin (begin (begin (our-display 2) (our-display 3)) (our-display 4)) (our-display 5)) <PRINT 2>
(begin (begin (our-display 3) (our-display 4)) (our-display 5)) <PRINT 3>
(begin (our-display 4) (our-display 5)) <PRINT 4>
(our-display 5) <PRINT 5>
=> 5
```