Normal Factor Graphs, Linear Algebra and Probabilistic Modeling

Yongyi Mao University of Ottawa

Workshop on Counting, Inference, and Optimization on Graphs Princeton University, NJ, USA

November 3, 2011

It started on day one of my PhD program ...

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Normal graph duality



• Duality theorem: Dual normal graphs represent dual codes

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• Kschischang:

"Why do normal graphs have duality, but factor graphs do not?"

- Two weeks later ...
- Me: "We need ... Fourier ... we need convolution..."

• Given f(x, y) and g(y, z)

$$f(\mathbf{x},\mathbf{y}) * g(\mathbf{y},z) = \sum_{w} f(\mathbf{x},\mathbf{y}-w)g(w,z)$$

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Convolutional factor graph

- bipartite graph: function nodes and variables nodes
- encodes a convolutional product



 $f_1(x_1, x_2) * f_2(x_1, x_3) * f_3(x_1, x_3)$

\bullet Dual FGs: $(G,\{f_i\},\times)$ and $(G,\{\widehat{f}_i\},\ast)$

 $\bullet \ \ast \stackrel{\mathcal{F}}{\leftrightarrow} \times$

- \Rightarrow Dual FGs encode a FT pair
- Dual FGs of codes: $(G, \{\delta_{C_i}\}, \times)$ and $(G, \{\delta_{C_i}\}, *)$

• $\delta_C \stackrel{\mathcal{F}}{\leftrightarrow} \delta_{C^{\perp}}$

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- Normal graph duality reinterpreted
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- Forney: "It may take years for people to take it [convolutional FG] and run"

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 - Each node: joint distribution of a collection of (latent) variables
 - All collections are independent
 - ⇒ Convoltional CFG represents the joint distribution of observed variables that are linearly transformed from the latent variables

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• CDN:

- multiplicative FG
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Conditional/marginal independence properties



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- Multiplicative FG: conditional indepdence;
- Convolutional FG: marginal independence;
- CDN: marginal independence

Again, it is Frank Kschischang ...

• Early 2010 ...

- Kschischang: "Look at this!"
 - "Constrained coding as networks of relations" [Schwartz & Bruck 2008]

- "Holographic algorithms" [Valiant 2004]
 - Counting problem
 - Holographic reduction: $\sum \prod \rightarrow \sum \prod$
 - Holant theorem
- Apparent "normal" structure
- Do something with holographic algorithm?
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Give it the student ...

• Ali Al-Bashebsheh



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Two concurrent submissions ...

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- "Normal factor graphs and holographic transformations " [Al-Bashebsheh & Mao 2011]
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Normal factor graphs under new semantics ...



$$\begin{aligned} \mathsf{Z}_{\mathcal{G}}(\mathsf{x}_1,\mathsf{x}_2) &:= \sum_{\substack{\mathsf{x}_3,\ldots,\mathsf{x}_7\\ &:= \langle \mathsf{f}_1,\mathsf{f}_2,\mathsf{f}_3,\mathsf{f}_4 \rangle.\\ &:= \langle \mathsf{f}_1(\mathsf{x}_1,\mathsf{x}_3,\mathsf{x}_4),\mathsf{f}_2(\mathsf{x}_3,\mathsf{x}_5,\mathsf{x}_6),\mathsf{f}_3(\mathsf{x}_2,\mathsf{x}_4,\mathsf{x}_5,\mathsf{x}_7),\mathsf{f}_4(\mathsf{x}_6,\mathsf{x}_7). \rangle \end{aligned}$$

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Functions as multidimensional arrays

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$$|\mathcal{X}_1| = 2, |\mathcal{X}_2| = 3$$
, and $|\mathcal{X}_3| = 5$.















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$$\delta(s,s') := \left\{ \begin{array}{ll} 1, & s = s' \\ 0, & \text{otherwise} \end{array} \right.$$

• $\delta(\cdot, \cdot)$: identity matrix

Example ("Transformer"

- Function Φ : χ × χ → C is called a transformer if it corresponds to an invertible matrix.
- "Dual pair of transformers" Φ and Φ: an inverse pair of matrices

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$\langle \mathbf{f},\mathbf{g} \rangle$:

- $\langle f(s), g(s) \rangle$: vector-vector dot product
- $\langle f(s), g(s,t) \rangle$: vector-matrix product
- $\langle f(s,t), g(t,u) \rangle$: matrix-matrix product
- $\langle f(s), g(t) \rangle$: vector outer product, matrix Kronecker product, tensor product etc

NFGs are linear algebraic expressions written graphically.

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Opening/closing the box [Vontobel, Loeliger, 2002]



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Holographic transformation



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Theorem (Generalized Holant Theorem)

 $Z_{\mathcal{G}^{H}}$ is "externally transformed" version of $Z_{\mathcal{G}}$.

- When external transformers are $\delta(\cdot, \cdot)$
 - reduces to Valiant's Holant Theorem
- When all transformers are Fourier kernels
 - reduces to General NFG Duality Theorem
 - when each function is a subgroup-indicator function
 - reduces to Forney's Normal Graph Duality Theorem

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General NFG Duality Theorem









596

Holographic algorithm: PerfMatch(H)

- H: weighted graph
- PerfMatch(H)

$$\pi(\mathsf{H}) := \sum_{\mathsf{M} \in \mathsf{Q}(\mathsf{H})} \prod_{e \in \mathsf{M}} w(e),$$

where Q(H) the set of all perfect matchings of H.

- If H is planar, $\pi(H)$ is poly. solvable by FKT algorithm
- Matchgate (G, U), where G is a weighted graph, $U \subseteq V(G)$.
- \bullet Signature of matchgate (G,U) is a function $\mu(x_U)$...









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Lemma: $\pi(H) = \sum_{x_{\mathcal{E}}} \prod_{i=1}^{m} \mu_i(x_{\mathcal{E}(i)})$.



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Holographic algorithm



↓ FKT

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- Vontobel pointed to Trace Diagram, Birdtracks ...
- "Unshackling linear algebra from linear notation" [Peterson 2009]
 "Group Theory: Birdtracks, Lie's, and Exceptional Group", [Cvitanovic, 2008]

- Trace diagram is contained in NFG framework
 Trace diagrams are NFGs with degree 1, 2, or |x|
- "NFGs: a diagrammatic approach to linear algebra" [Al-Bashebsheh, Mao & Vontobel ISIT2011]

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• Draw the edges of f counter-clockwise, with first edge ciliated



Ciliated function nodes



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Trace



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Proof:



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Read the graph in two ways: $Z = tr(A \cdot B)$ and $Z = tr(B \cdot A)$.

Levi-Civita symbol

- S_n : the permutation group on $\{1, \ldots, n\}$, $n = |\chi|$.
- Levi-Civita symbol, $\varepsilon : \{1, \ldots, n\}^n \to \mathbb{C}$ s.t.

$$\varepsilon(x_1,\ldots,x_n) = \begin{cases} \operatorname{sgn} \begin{pmatrix} 1 & \cdots & n \\ x_1 & \cdots & x_n \end{pmatrix}, & \begin{pmatrix} 1 & \cdots & n \\ x_1 & \cdots & x_n \end{pmatrix} \in S_n \\ 0, & \text{otherwise} \end{cases}$$

• Contraction (n = 3).

 $\sum \varepsilon(\mathbf{y}_1,\mathbf{y}_2,\mathbf{t})\varepsilon(\mathbf{t},\mathbf{x}_2,\mathbf{x}_1) = \delta(\mathbf{x}_1,\mathbf{y}_2)\delta(\mathbf{x}_2,\mathbf{y}_1) - \delta(\mathbf{x}_1,\mathbf{y}_1)\delta(\mathbf{x}_2,\mathbf{y}_2).$





Levi-Civita symbol

- S_n : the permutation group on $\{1, \ldots, n\}$, $n = |\chi|$.
- Levi-Civita symbol, $\epsilon: \{1, \ldots, n\}^n \to \mathbb{C}$ s.t.

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$$\mathfrak{u}, \nu : \{1, 2, 3\} \to \mathbb{F} \Leftrightarrow \mathfrak{u}, \nu \in \mathbb{F}^3.$$

• Cross Product.

 $Z(x)=\sum_{t_1,t_2}u(t_1)\nu(t_2)\varepsilon(t_1,t_2,x).$

Z(1) = u(2)v(3) - u(3)v(2) Z(2) = u(3)v(1) - u(1)v(3)Z(3) = u(1)v(2) - u(2)v(1)



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Cross product identities

$$(\mathbf{u} \times \mathbf{v}) \cdot (\mathbf{s} \times \mathbf{w}) = ((\mathbf{u} \times \mathbf{v}) \times \mathbf{s}) \cdot \mathbf{w} = (\mathbf{w} \times (\mathbf{u} \times \mathbf{v})) \cdot \mathbf{s}$$
$$= ((\mathbf{s} \times \mathbf{w}) \times \mathbf{u}) \cdot \mathbf{v} = (\mathbf{v} \times (\mathbf{s} \times \mathbf{w})) \cdot \mathbf{u}$$
$$= (\mathbf{u} \cdot \mathbf{s})(\mathbf{v} \cdot \mathbf{w}) - (\mathbf{u} \cdot \mathbf{w})(\mathbf{v} \cdot \mathbf{s}).$$



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$$|A| := \sum_{\sigma \in S_n} \operatorname{sgn}(\sigma) \prod_{j=1}^n A(j, \sigma(j)).$$

Lemma:



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Proof:



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 $|A^{\mathsf{T}}| = |A|$

Proof:



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• A is a $2n \times 2n$ skew-symmetric matrix.

• $Pf(A) = \frac{1}{2^n n!} \sum_{\sigma \in S_{2n}} sgn(\sigma) \prod_{i=1}^n A(\sigma(2i-1), \sigma(2i)).$



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- $Z = n! \cdot 2^n \cdot Pf(A)$
- Affirms a conjecture of Peterson [Peterson, 2009].

Pfaffian and Determinant

 $J := \begin{pmatrix} 0 & I_n \\ -I_n & 0 \end{pmatrix}, A := R^T \cdot J \cdot R \text{ for some } R. \text{ Then } Pf(A)^2 = |A|.$ *Proof:* $Pf(J) = \pm 1$



 $\pm Pf(A) = |R| \Rightarrow Pf(A)^2 = |R|^2 = |A|.$ ▲□▶ ▲□▶ ▲□▶ ▲□▶ □ のへの

"NFGs as probabilistic model"

[Al-Bashebsheh & Mao, 2011, soon available on arxiv]

- Internal edges: latent variabels
- External edges: observed variables
- Exterior function: joint distribution of observed variables (up to scale)



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• Function:

- compatibility/constraints/interaction protential
- source of randomness, i.e., distribution
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Special functions



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"Cumulus" and "difference" transformers



• χ : an ordered set (integers) with "<" well defined.

•
$$A(x, x') = [x' \le x]$$

• E.g. $|\chi| = 2, A = \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix};$
 $|\chi| = 3, A = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 1 \end{bmatrix}$
• $\langle f_X(x'), A(x, x') \rangle = F_X(x).$
• $D := A^{-1}, e.g.D = \begin{bmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \\ 0 & -1 & 1 \end{bmatrix}$ for $|\chi|=3$.

- One partitiion: arbitrary potential functions
- The other partition: equality indicators
- ullet \Rightarrow NFG model reduces to FG model





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Bipartite NFG reduces to convolutional FG

- One side: distributions of indep. collections of latent RVs
- The other side: sum indicators
- No RVs from the same collection connect to the same indicator
- $\bullet \Rightarrow$ NFG model reduces to convolutional FG model



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CDN is a holographic transformed NFG



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CDN is a holographic transformed NFG

- One side: distributions of indep. collections of latent RVs
- The other side: max indicators
- No RVs from the same collection connect to the same indicator



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CDN is a holographic transformed NFG





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The independence coincidence of Convolutional FG and CDN



- Both are generative models
- In fact, the independence property holds when changing indicator to conditional distributions
- Changing the right-side functions gives rise to a family of infinite models ... what applications?

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NFGs are linear algebraic expressions written graphically

Opening/closing the box is a powerful technique

NFGs are potential tools for inference

Outlook: continuous alphabets

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