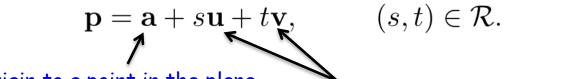
## Support vector machines (SVMs) Lecture 3

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Slides adapted from Luke Zettlemoyer, Vibhav Gogate, and Carlos Guestrin

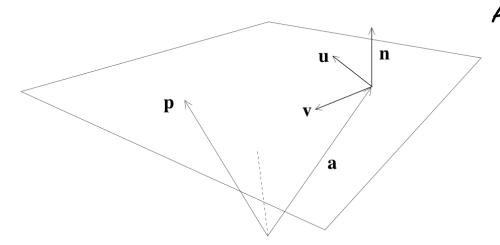
## Geometry of linear separators (see blackboard)

A plane can be specified as the set of all points given by:



Vector from origin to a point in the plane

Two non-parallel directions in the plane



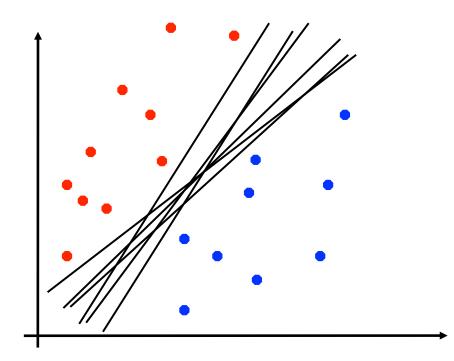
Alternatively, it can be specified as:  $(\mathbf{p} - \mathbf{a}) \cdot \mathbf{n} = 0 \Leftrightarrow \mathbf{p} \cdot \mathbf{n} = \mathbf{a} \cdot \mathbf{n}$ Normal vector (we will call this w)

Only need to specify this dot product, a scalar (we will call this the offset, b)

Barber, Section A.1.1-4

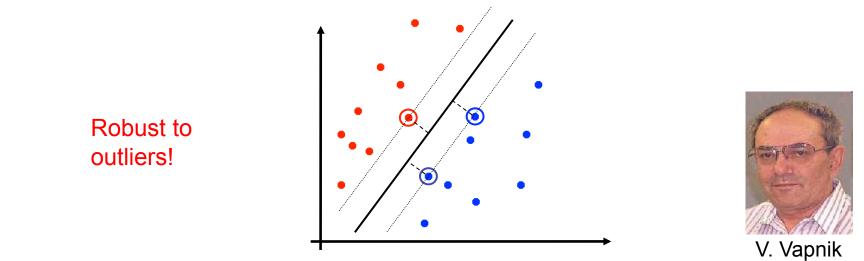
## **Linear Separators**

- If training data is linearly separable, perceptron is guaranteed to find *some* linear separator
- Which of these is **optimal**?



### Support Vector Machine (SVM)

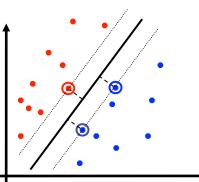
SVMs (Vapnik, 1990's) choose the linear separator with the largest margin



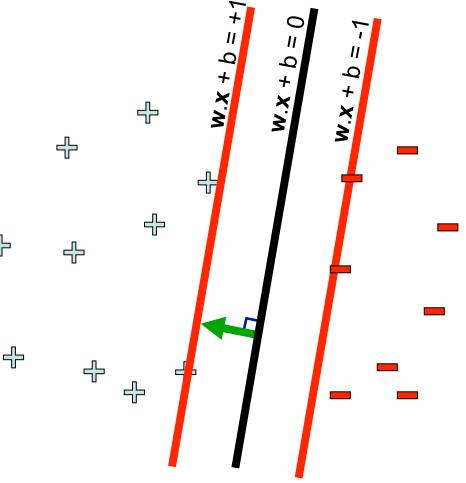
- Good according to intuition, theory, practice
- SVM became famous when, using images as input, it gave accuracy comparable to neural-network with hand-designed features in a handwriting recognition task

#### Support vector machines: 3 key ideas

- 1. Use **optimization** to find solution (i.e. a hyperplane) with few errors
- 2. Seek **large margin** separator to improve generalization
- 3. Use **kernel trick** to make large feature spaces computationally efficient



# Finding a perfect classifier (when one exists) using linear programming



For every data point  $(x_t, y_t)$ , enforce the constraint

for y\_t = +1, 
$$w \cdot x_t + b \ge 1$$
  
and for y\_t = -1,  $w \cdot x_t + b \le -1$ 

Equivalently, we want to satisfy all of the linear constraints

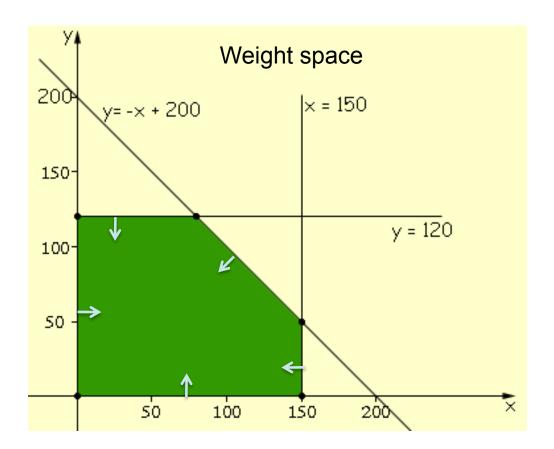
$$y_t \left( w \cdot x_t + b \right) \ge 1 \quad \forall t$$

This *linear program* can be efficiently solved using algorithms such as simplex, interior point, or ellipsoid

# Finding a perfect classifier (when one exists) using linear programming

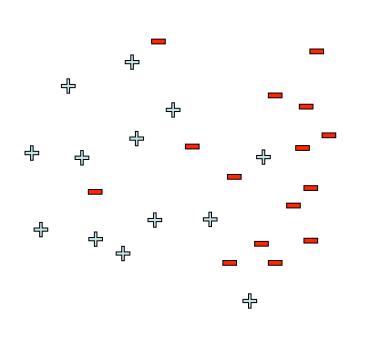
Example of 2-dimensional linear programming (feasibility) problem:

For SVMs, each data point gives one inequality:  $y_t (w \cdot x_t + b) \ge 1$ 



#### What happens if the data set is not linearly separable?

### Minimizing number of errors (0-1 loss)



• Try to find weights that violate as few constraints as possible?

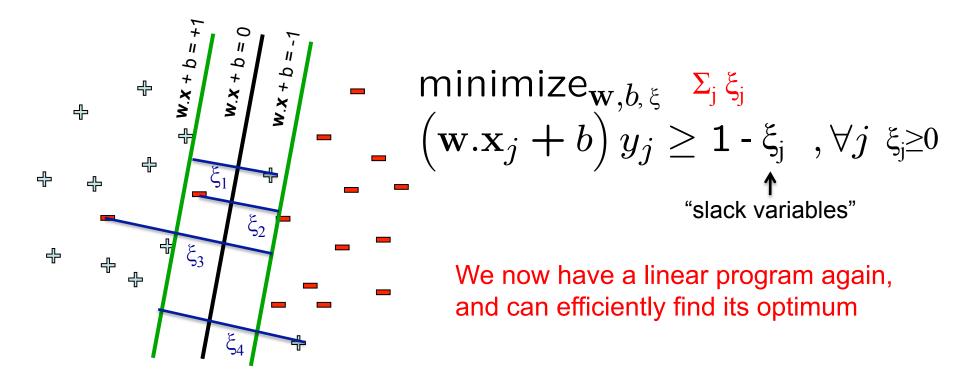
 $\begin{array}{ll} \text{minimize}_{\mathbf{w},b} & \text{#(mistakes)} \\ \left(\mathbf{w}.\mathbf{x}_{j}+b\right)y_{j} \geq 1 & ,\forall j \end{array}$ 

• Formalize this using the 0-1 loss:

$$\begin{split} \min_{\mathbf{w},b} \sum_{j} \ell_{0,1}(y_j, \, w \cdot x_j + b) \\ \text{where } \ell_{0,1}(y, \hat{y}) = \mathbf{1}[y \neq \operatorname{sign}(\hat{y})] \end{split}$$

- Unfortunately, minimizing 0-1 loss is NP-hard in the worst-case
  - Non-starter. We need another approach.

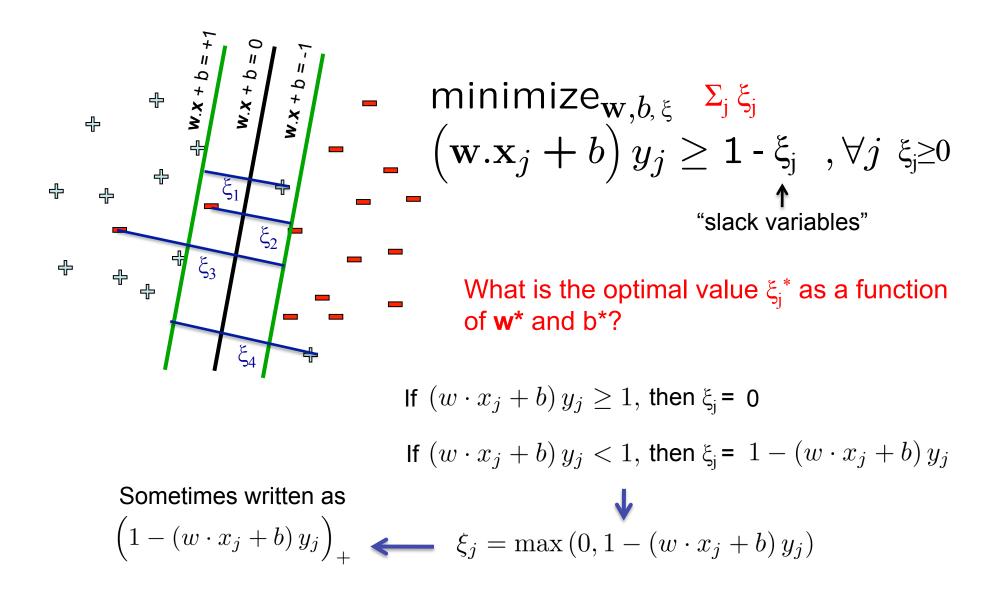
#### Key idea #1: Allow for slack



#### For each data point:

- •If functional margin  $\geq$  1, don't care
- •If functional margin < 1, pay linear penalty

#### Key idea #1: Allow for *slack*



#### Equivalent hinge loss formulation

$$egin{aligned} & ext{minimize}_{\mathbf{w},b,\,\xi} & \Sigma_{\mathrm{j}}\,\xi_{\mathrm{j}} \ & \left(\mathbf{w}.\mathbf{x}_{j}+b
ight)y_{j} \geq \mathbf{1}$$
 -  $\xi_{\mathrm{j}}$  ,  $orall j$   $\xi_{\mathrm{j}} \geq 0$ 

Substituting  $\xi_j = \max(0, 1 - (w \cdot x_j + b) y_j)$  into the objective, we get:

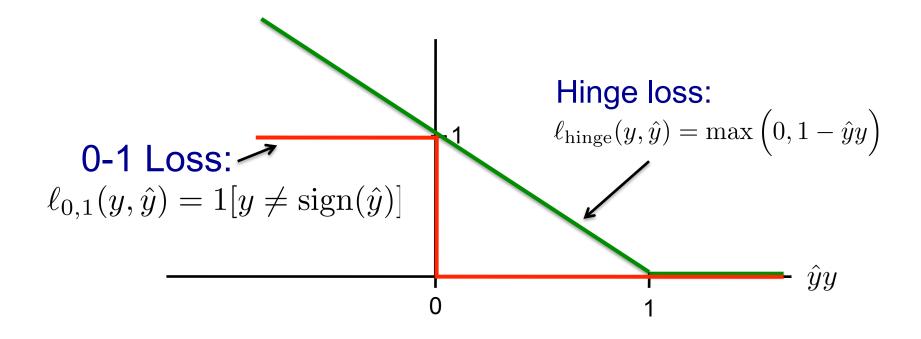
$$\min_{w,b} \sum_{j} \max\left(0, 1 - \left(w \cdot x_{j} + b\right) y_{j}\right)$$

The hinge loss is defined as  $\ell_{\text{hinge}}(y, \hat{y}) = \max\left(0, 1 - \hat{y}y\right)$ 

$$\min_{\mathbf{w},b} \sum_{j} \ell_{\text{hinge}}(y_j, \, w \cdot x_j + b)$$

This is empirical risk minimization, using the hinge loss

#### Hinge loss vs. 0/1 loss



#### Hinge loss upper bounds 0/1 loss!

It is the tightest *convex* upper bound on the 0/1 loss

### Key idea #2: seek large margin

