

# 6.838: Geometric Computing

Spring 2005

Problem Set 4

Due: Thursday, May 12

## MANDATORY PART

### Problem 1. External queue

Show how to efficiently implement a queue in external memory. That is, show how to support two operations:

- PUSH( $a$ ): pushes the element  $a$  at the *end* of the queue
- POP: removes the element from the *front* of the queue and returns it

Your implementation should support any  $m$  operations on the queue using  $O(m/B + 1)$  block operations.

### Problem 2. Metrics under translation

Consider any metric  $D(A, B)$  defined for sets  $A, B \subset \mathbb{R}^2$ . Assume that  $D(\cdot, \cdot)$  is symmetric and satisfies triangle inequality. Let  $T$  be the set of all translations in  $\mathbb{R}^2$ , and define  $D_T(A, B) = \min_{t \in T} D(t(A), B)$ . Show that  $D_T(A, B)$  is symmetric and satisfies triangle inequality as well.

### Problem 3. Deterministic packing.

In the lecture, we have seen a polynomial-time approximation scheme for the unit disk packing problem. The algorithm was randomized, since it used a randomly shifted grid.

Show that there is a *deterministic* polynomial-time approximation scheme for this problem.

**Hint:** The randomized algorithm uses a grid shifted by a vector chosen uniformly at random from  $S = [0, k]^2$ , where  $k$  is the side length of the grid cell. Thus, a deterministic algorithm could be obtained by trying out “all” vectors from  $S$ , and choosing the one that gives the best packing.

## OPTIONAL THEORETICAL PART

### Problem A. Reference point for EMD

Recall that, for two points sets  $A, B \in \mathbb{R}^2$ ,  $|A| = |B|$

$$EMD(A, B) = \min_{\pi: A \rightarrow B} \sum_{a \in A} \|a - \pi(a)\|$$

where  $\pi$  is 1 – to – 1.

Construct a reference point function  $r(A)$  for EMD (under translations). That is, give a function  $r(\cdot)$  such that for any two sets  $A, B$ , if  $t^* = r(B) - r(A)$  is a translation that moves  $r(A)$  to  $r(B)$ , then

$$EMD(t^*(A), B) \leq c \cdot \min_{t \in T} EMD(t(A), B)$$

for some constant  $c \geq 1$ . As before,  $T$  is the set of all translations in  $\mathfrak{R}^2$ .

#### OPTIONAL PROGRAMMING PART

##### **Problem B. Directed Hausdorff Under Translation**

Implement a Java applet that simulates the algorithm from slide 13 of lecture 20. The algorithm should take as an input two sets of points  $A$  and  $B$ . Then it should do the following: if there is a translation  $t$  such that  $DH(t(A), B) \leq r$ , it should find a translation  $t'$  such that  $D(t'(A), B) \leq (1 + \epsilon)r$ . It is OK if  $\epsilon r$  is set to be equal to the diameter of a screen pixel (i.e., the grid used in the algorithm can coincide with the grid induced by screen pixels).

To implement the above, the algorithm should construct approximations of the sets  $T(a)$  for all  $a \in A$ , and depict them on the screen. Then it should check if the intersections of all those sets is non-empty. If so, it should visualize the resulting translation.