Rediscovering the Joys of Pebbling

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Theory reading group April 17, 2013

a.k.a. "On the Relative Strength of Pebbling and Resolution"

- Satisfiability algorithms
 - Dramatic developments last 10-15 years
 - ► SAT solvers used on regular basis for large-scale real-world problems
 - ► Best algorithms based on resolution proof system
 - ► Bottlenecks: time and memory consumption

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- This talk
 - ▶ What can proof complexity say about time vs space?
 - ► Connections between resolution and pebble games?

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- ▶ Best algorithms based on resolution proof system
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Pebble games

- Used in 70s-80s to study programming languages, compiler optimization, et cetera
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This talk

- ▶ What can proof complexity say about time vs space?
- Connections between resolution and pebble games?
- ► And after the break: some pebbling results and proofs

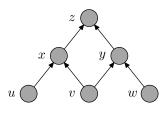
Outline

- 1 The Resolution Proof System
 - Basics
 - Some of What We Know (and What We Don't)
- Pebble Games
 - Black and Black-White Pebbling
 - Pebbling Contradictions
 - Reductions Between Resolution and Pebbling
- Our Results
 - Pebbling Trade-offs
 - Pebbling-to-Resolution Reduction
- Open Problems

Just to Check We're on the Same Page. . .

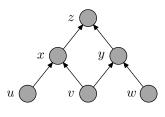
- Literal a: variable x or its negation \overline{x}
- Clause $C = a_1 \lor \cdots \lor a_k$: disjunction of literals
- CNF formula $F = C_1 \wedge \cdots \wedge C_m$: conjunction of clauses
- k-CNF formula: CNF formula with clauses of size $\leq k$ (assume k fixed)
- Refer to clauses of CNF formula as axioms (as opposed to derived clauses)

- $1. \quad u$
- 2. v
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \vee \overline{y} \vee z$
- 7. \overline{z}



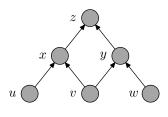
- source vertices true
- truth propagates upwards
- but sink vertex is false

- $1. \quad u$
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- 3. w
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- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}



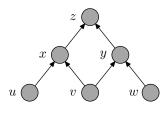
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- 7. \overline{z}



Blackboard bookkeeping		
total # clauses on board	0	
max # lines on board	0	
max # literals on board	0	

Can write down axioms, erase used clauses or infer new clauses by resolution rule $B \vee x \qquad C \vee \overline{x}$

$$\frac{B \vee x \qquad C \vee \overline{x}}{B \vee C}$$

(but only from clauses currently on the board!)

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

u

Blackboard bookkeeping		
total # clauses on board	1	
max # lines on board	1	
max # literals on board	1	

Write down axiom 1: u

- 1. u
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

u		
v		

Blackboard bookkeeping		
total # clauses on board	2	
max # lines on board	2	
max # literals on board	2	

Write down axiom 1: u Write down axiom 2: v

- 1. u
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

Blackboard bookkeeping	g
total # clauses on board	3
max # lines on board	3
max # literals on board	5

Write down axiom 1: u

Write down axiom 2: $\it v$

Write down axiom 4: $\overline{u} \vee \overline{v} \vee x$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

Diackboard bookkeeping		
total # clauses on board	3	
max # lines on board	3	
max # literals on board	5	

Disabbased basklesseine

 $egin{array}{c} u \\ v \\ \overline{u} ee \overline{v} ee x \end{array}$

Write down axiom 1: u Write down axiom 2: v Write down axiom 4: $\overline{u} \lor \overline{v} \lor x$ Infer $\overline{v} \lor x$ from u and $\overline{u} \lor \overline{v} \lor x$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

u	
v	
$\overline{u} \vee \overline{v} \vee x$	
$\overline{v} \lor x$	

Blackboard bookkeeping		
total # clauses on board	4	
max # lines on board	4	
max # literals on board	7	

Write down axiom 1: u Write down axiom 2: v

Write down axiom 4: $\overline{u} \vee \overline{v} \vee x$

 $\mathsf{Infer}\ \overline{v} \vee x\ \mathsf{from}$

 $u \text{ and } \overline{u} \vee \overline{v} \vee x$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

u	
v	
$\overline{u} \vee \overline{v} \vee x$	
$\overline{v} \lor x$	

Blackboard bookkeeping		
total # clauses on board	4	
max # lines on board	4	
max # literals on board	7	

Write down axiom 2: v Write down axiom 4: $\overline{u} \lor \overline{v} \lor x$ Infer $\overline{v} \lor x$ from u and $\overline{u} \lor \overline{v} \lor x$ Erase the line $\overline{u} \lor \overline{v} \lor x$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

u	
v	
$\overline{v} \lor x$	

Blackboard bookkeeping		
total # clauses on board	4	
max # lines on board	4	
max # literals on board	7	

Write down axiom 2: v Write down axiom 4: $\overline{u} \lor \overline{v} \lor x$ Infer $\overline{v} \lor x$ from u and $\overline{u} \lor \overline{v} \lor x$ Erase the line $\overline{u} \lor \overline{v} \lor x$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

u			
v			
\overline{v}	$\vee x$		

Blackboard bookkeeping		
total # clauses on board	4	
max # lines on board	4	
max # literals on board	7	

Write down axiom 4: $\overline{u} \lor \overline{v} \lor x$ Infer $\overline{v} \lor x$ from u and $\overline{u} \lor \overline{v} \lor x$ Erase the line $\overline{u} \lor \overline{v} \lor x$ Erase the line u

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

v		
$\overline{v}\vee x$		

Blackboard bookkeeping		
total # clauses on board	4	
max # lines on board	4	
max # literals on board	7	

Write down axiom 4: $\overline{u} \lor \overline{v} \lor x$ Infer $\overline{v} \lor x$ from u and $\overline{u} \lor \overline{v} \lor x$ Erase the line $\overline{u} \lor \overline{v} \lor x$ Erase the line u

- 1. u
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

v	
$\overline{v}\vee x$	

Blackboard bookkeeping		
total # clauses on board	4	
max # lines on board	4	
max # literals on board	7	

u and $\overline{u} \lor \overline{v} \lor x$ Erase the line $\overline{u} \lor \overline{v} \lor x$ Erase the line uInfer x from v and $\overline{v} \lor x$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

ι	,	
ī	$\overline{y} \vee x$	
а	c	

Blackboard bookkeeping		
total # clauses on board	5	
max # lines on board	4	
max # literals on board	7	

u and $\overline{u} \lor \overline{v} \lor x$ Erase the line $\overline{u} \lor \overline{v} \lor x$ Erase the line uInfer x from v and $\overline{v} \lor x$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

v
$\overline{v} \lor x$
x

Blackboard bookkeeping		
total # clauses on board	5	
max # lines on board	4	
max # literals on board	7	

Erase the line $\overline{u} \lor \overline{v} \lor x$ Erase the line uInfer x from v and $\overline{v} \lor x$ Erase the line $\overline{v} \lor x$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
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- 7. \overline{z}

v		
x		

Blackboard bookkeeping	g
total # clauses on board	5
max # lines on board	4
max # literals on board	7

Erase the line $\overline{u} \lor \overline{v} \lor x$ Erase the line uInfer x from v and $\overline{v} \lor x$ Erase the line $\overline{v} \lor x$

- 1. u
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

v		
x		

Blackboard bookkeeping		
total # clauses on board	5	
max # lines on board	4	
max # literals on board	7	

Erase the line uInfer x from v and $\overline{v} \lor x$ Erase the line $\overline{v} \lor x$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

x		

Blackboard bookkeeping		
total # clauses on board	5	
max # lines on board	4	
max # literals on board	7	

Erase the line u Infer x from v and $\overline{v} \lor x$ Erase the line $\overline{v} \lor x$ Erase the line v

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

x	
$\overline{x} \vee \overline{y} \vee z$	

Blackboard bookkeeping		
total # clauses on board	6	
max # lines on board	4	
max # literals on board	7	

Infer x from $v \text{ and } \overline{v} \vee x$ Erase the line $\overline{v} \vee x$ Erase the line v Write down axiom 6: $\overline{x} \vee \overline{y} \vee z$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

\boldsymbol{x}	
$\overline{x} \vee \overline{y} \vee z$	

Blackboard bookkeeping	g
total # clauses on board	6
max # lines on board	4
max # literals on board	7

Erase the line $\overline{v} \vee x$ Erase the line vWrite down axiom 6: $\overline{x} \vee \overline{y} \vee z$ Infer $\overline{y} \vee z$ from x and $\overline{x} \vee \overline{y} \vee z$

- 1. u
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

x		
$\overline{x} \vee$	$\overline{y} \vee$	z
$\overline{y} \vee$	z	

Blackboard bookkeeping		
total # clauses on board	7	
max # lines on board	4	
max # literals on board	7	

Erase the line $\overline{v} \lor x$ Erase the line vWrite down axiom 6: $\overline{x} \lor \overline{y} \lor z$ Infer $\overline{y} \lor z$ from x and $\overline{x} \lor \overline{y} \lor z$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

x		
$\overline{x} \vee$	$\overline{y} \vee$	z
$\overline{y} \vee$	z	

Blackboard bookkeeping	
total # clauses on board	7
max # lines on board	4
max # literals on board	7

Erase the line v Write down axiom 6: $\overline{x} \lor \overline{y} \lor z$ Infer $\overline{y} \lor z$ from x and $\overline{x} \lor \overline{y} \lor z$ Erase the line $\overline{x} \lor \overline{y} \lor z$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\frac{x}{\overline{x}} \vee z$		
$\overline{y} \lor z$		

Blackboard bookkeeping	
total # clauses on board	7
max # lines on board	4
max # literals on board	7

Erase the line v Write down axiom 6: $\overline{x} \lor \overline{y} \lor z$ Infer $\overline{y} \lor z$ from x and $\overline{x} \lor \overline{y} \lor z$ Erase the line $\overline{x} \lor \overline{y} \lor z$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

x	
$\overline{y}\vee z$	

Blackboard bookkeeping	
total # clauses on board	7
max # lines on board	4
max # literals on board	7

Write down axiom 6: $\overline{x} \lor \overline{y} \lor z$ Infer $\overline{y} \lor z$ from x and $\overline{x} \lor \overline{y} \lor z$ Erase the line $\overline{x} \lor \overline{y} \lor z$ Erase the line x

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{y} \vee z$		

Blackboard bookkeeping	
total # clauses on board	7
max # lines on board	4
max # literals on board	7

Write down axiom 6: $\overline{x} \lor \overline{y} \lor z$ Infer $\overline{y} \lor z$ from x and $\overline{x} \lor \overline{y} \lor z$ Erase the line $\overline{x} \lor \overline{y} \lor z$ Erase the line x

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{y} \vee z$	
$\overline{v} \vee \overline{w} \vee y$	

Blackboard bookkeeping	
total # clauses on board	8
max # lines on board	4
max # literals on board	7

Infer $\overline{y} \lor z$ from x and $\overline{x} \lor \overline{y} \lor z$ Erase the line $\overline{x} \lor \overline{y} \lor z$ Erase the line x Write down axiom 5: $\overline{v} \lor \overline{w} \lor y$

- 1. u
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{y} \lor z$	
$\overline{v} \vee \overline{w} \vee y$	

Blackboard bookkeeping	
total # clauses on board	8
max # lines on board	4
max # literals on board	7

Erase the line $\overline{x} \vee \overline{y} \vee z$ Erase the line xWrite down axiom 5: $\overline{v} \vee \overline{w} \vee y$ Infer $\overline{v} \vee \overline{w} \vee z$ from $\overline{y} \vee z$ and $\overline{v} \vee \overline{w} \vee y$

- 1. u
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{y} \lor z$	
$\overline{v} \vee \overline{w} \vee y$	
$\overline{v} \vee \overline{w} \vee z$	

Blackboard bookkeeping	
total # clauses on board	9
max # lines on board	4
max # literals on board	8

Erase the line $\overline{x} \vee \overline{y} \vee z$ Erase the line xWrite down axiom 5: $\overline{v} \vee \overline{w} \vee y$ Infer $\overline{v} \vee \overline{w} \vee z$ from $\overline{y} \vee z$ and $\overline{v} \vee \overline{w} \vee y$

- $1. \quad u$
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{y} \lor z$	
$\overline{v} \lor \overline{w} \lor y$	
$\overline{v} \vee \overline{w} \vee z$	

Blackboard bookkeeping	
total # clauses on board	9
max # lines on board	4
max # literals on board	8

Erase the line x Write down axiom 5: $\overline{v} \lor \overline{w} \lor y$ Infer $\overline{v} \lor \overline{w} \lor z$ from $\overline{y} \lor z$ and $\overline{v} \lor \overline{w} \lor y$ Erase the line $\overline{v} \lor \overline{w} \lor y$

- 1. u
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{y} \vee z$	
$\overline{v} \vee \overline{w} \vee z$	

Blackboard bookkeeping	
total # clauses on board	9
max # lines on board	4
max # literals on board	8

Erase the line x Write down axiom 5: $\overline{v} \lor \overline{w} \lor y$ Infer $\overline{v} \lor \overline{w} \lor z$ from $\overline{y} \lor z$ and $\overline{v} \lor \overline{w} \lor y$ Erase the line $\overline{v} \lor \overline{w} \lor y$

- 1. u
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{y} \lor z$	
$\overline{v} \vee \overline{w} \vee z$	

Blackboard bookkeeping	
total # clauses on board	9
max # lines on board	4
max # literals on board	8

Write down axiom 5: $\overline{v} \lor \overline{w} \lor y$ Infer $\overline{v} \lor \overline{w} \lor z$ from $\overline{y} \lor z$ and $\overline{v} \lor \overline{w} \lor y$ Erase the line $\overline{v} \lor \overline{w} \lor y$ Erase the line $\overline{y} \lor z$

- 1. u
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{v} \vee \overline{w} \vee z$	

Blackboard bookkeeping	
total # clauses on board	9
max # lines on board	4
max # literals on board	8

Write down axiom 5: $\overline{v} \lor \overline{w} \lor y$ Infer $\overline{v} \lor \overline{w} \lor z$ from $\overline{y} \lor z$ and $\overline{v} \lor \overline{w} \lor y$ Erase the line $\overline{v} \lor \overline{w} \lor y$ Erase the line $\overline{y} \lor z$

- $1. \quad u$
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{v} \vee \overline{w} \vee z$	
v	

Blackboard bookkeeping	
total # clauses on board	10
max # lines on board	4
max # literals on board	8

Infer $\overline{v} \lor \overline{w} \lor z$ from $\overline{y} \lor z$ and $\overline{v} \lor \overline{w} \lor y$ Erase the line $\overline{v} \lor \overline{w} \lor y$ Write down axiom 2: v

- 1. *u*
- 2. *v*
- 3. *w*
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{v} \vee \overline{w} \vee z$
v
w

Blackboard bookkeeping	
total # clauses on board	11
max # lines on board	4
max # literals on board	8

 $\overline{y} \lor z$ and $\overline{v} \lor \overline{w} \lor y$ Erase the line $\overline{v} \lor \overline{w} \lor y$ Erase the line $\overline{y} \lor z$ Write down axiom 2: vWrite down axiom 3: w

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{v} \vee \overline{w} \vee z$
v
w
\overline{z}

Blackboard bookkeeping	
total # clauses on board	12
max # lines on board	4
max # literals on board	8

Erase the line $\overline{v} \lor \overline{w} \lor y$ Erase the line $\overline{y} \lor z$ Write down axiom 2: vWrite down axiom 3: wWrite down axiom 7: \overline{z}

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

Blackboard bookkeeping	
total # clauses on board	12
max # lines on board	4
max # literals on board	8

$\overline{v} \vee \overline{w} \vee z$	
v	
w	
\overline{z}	

Write down axiom 2: v Write down axiom 3: w Write down axiom 7: \overline{z} Infer $\overline{w} \lor z$ from v and $\overline{v} \lor \overline{w} \lor z$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{v} \vee \overline{w} \vee z$	
v	
w	
\overline{z}	
$\overline{w} \lor z$	

Blackboard bookkeeping	
total # clauses on board	13
max # lines on board	5
max # literals on board	8

Write down axiom 2: v Write down axiom 3: w Write down axiom 7: \overline{z} Infer $\overline{w} \lor z$ from v and $\overline{v} \lor \overline{w} \lor z$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

Blackboard bookkeeping	
13	
5	
8	

$\overline{v} \vee \overline{w} \vee z$
v
w
\overline{z}
$\overline{w} \lor z$

Write down axiom 3: w Write down axiom 7: \overline{z} Infer $\overline{w} \lor z$ from v and $\overline{v} \lor \overline{w} \lor z$ Erase the line v

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{v} \vee \overline{w} \vee z$	
w	
\overline{z}	
$\overline{w}\vee z$	

Blackboard bookkeeping	
total # clauses on board	13
max # lines on board	5
max # literals on board	8

Write down axiom 3: w Write down axiom 7: \overline{z} Infer $\overline{w} \lor z$ from v and $\overline{v} \lor \overline{w} \lor z$ Erase the line v

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

$\overline{v} \vee \overline{w} \vee z$	
w	
\overline{z}	
$\overline{w}\vee z$	

Blackboard bookkeeping		
total # clauses on board	13	
max # lines on board	5	
max # literals on board	8	

Write down axiom 7: \overline{z} Infer $\overline{w} \lor z$ from v and $\overline{v} \lor \overline{w} \lor z$ Erase the line v

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

w	
\overline{z}	
$\overline{w}\vee z$	

Blackboard bookkeeping	
total # clauses on board	13
max # lines on board	5
max # literals on board	8

Write down axiom 7: \overline{z} Infer $\overline{w} \lor z$ from v and $\overline{v} \lor \overline{w} \lor z$ Erase the line v

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

Blackboard bookkeeping	
total # clauses on board	13
max # lines on board	5
max # literals on board	8

 $\frac{w}{\overline{z}}$ $\overline{w} \lor z$

v and $\overline{v} \lor \overline{w} \lor z$ Erase the line v Erase the line $\overline{v} \lor \overline{w} \lor z$ Infer z from w and $\overline{w} \lor z$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

w
\overline{z}
$\overline{w}\vee z$
z

Blackboard bookkeeping	
total # clauses on board	14
max # lines on board	5
max # literals on board	8

```
v and \overline{v} \lor \overline{w} \lor z

Erase the line v

Erase the line \overline{v} \lor \overline{w} \lor z

Infer z from w and \overline{w} \lor z
```

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

Blackboard bookkeeping		
total # clauses on board	14	
max # lines on board	5	
max # literals on board	8	

 $\frac{w}{\overline{z}}$ $\overline{w} \lor z$ z

Erase the line vErase the line $\overline{v} \lor \overline{w} \lor z$ Infer z from w and $\overline{w} \lor z$ Erase the line w

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

\overline{z}	
$\overline{w}\vee z$	
z	

Blackboard bookkeeping	
total # clauses on board	14
max # lines on board	5
max # literals on board	8

Erase the line vErase the line $\overline{v} \vee \overline{w} \vee z$ Infer z from w and $\overline{w} \vee z$ Erase the line w

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

\overline{z}	
$\overline{w}\vee z$	
z	

Blackboard bookkeeping		
total # clauses on board	14	
max # lines on board	5	
max # literals on board	8	

Erase the line $\overline{v} \vee \overline{w} \vee z$ Infer z from w and $\overline{w} \vee z$ Erase the line wErase the line $\overline{w} \vee z$

- 1. u
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

\overline{z}	
z	

Blackboard bookkeeping		
total # clauses on board	14	
max # lines on board	5	
max # literals on board	8	

Erase the line $\overline{v} \lor \overline{w} \lor z$ Infer z from w and $\overline{w} \lor z$ Erase the line wErase the line $\overline{w} \lor z$

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

 \overline{z}

Blackboard bookkeeping		
total # clauses on board	14	
max # lines on board	5	
max # literals on board	8	

w and $\overline{w} \lor z$ Erase the line wErase the line $\overline{w} \lor z$ Infer \bot from \overline{z} and z

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \lor \overline{y} \lor z$
- 7. \overline{z}

\overline{z}		
z		
\perp		

Blackboard bookkeeping		
total # clauses on board	15	
max # lines on board	5	
max # literals on board	8	

w and $\overline{w} \lor z$ Erase the line wErase the line $\overline{w} \lor z$ Infer \bot from \overline{z} and z

- Length: Lower bound on time for proof search algorithm
- Space: Lower bound on memory for proof search algorithm

- Length: Lower bound on time for proof search algorithm
- Space: Lower bound on memory for proof search algorithm

Length

clauses written on blackboard counted with repetitions

- Length: Lower bound on time for proof search algorithm
- Space: Lower bound on memory for proof search algorithm

Length

clauses written on blackboard counted with repetitions

Space

Somewhat less straightforward — several ways of measuring

$$\begin{array}{c} x \\ \overline{y} \vee z \\ \overline{v} \vee \overline{w} \vee y \end{array}$$

- Length: Lower bound on time for proof search algorithm
- Space: Lower bound on memory for proof search algorithm

Length

clauses written on blackboard counted with repetitions

Space

Somewhat less straightforward — several ways of measuring

- **1**. *x*
- $2. \ \overline{y} \ \lor \ z$
- 3. $\overline{v} \vee \overline{w} \vee y$

Clause space:

3

- Length: Lower bound on time for proof search algorithm
- Space: Lower bound on memory for proof search algorithm

Length

clauses written on blackboard counted with repetitions

Space

Somewhat less straightforward — several ways of measuring

$$\begin{array}{c} x^1 \\ \overline{y}^2 \lor z^3 \\ \overline{v}^4 \lor \overline{w}^5 \lor y^6 \end{array}$$

3

- Length: Lower bound on time for proof search algorithm
- Space: Lower bound on memory for proof search algorithm

Length

clauses written on blackboard counted with repetitions (in our example resolution refutation 15)

Space

Somewhat less straightforward — several ways of measuring

$$\begin{array}{c} x \\ \overline{y} \ \lor \ z \\ \overline{v} \ \lor \ \overline{w} \ \lor \ y \end{array}$$

```
Let n = \text{size of formula } (\# \text{ symbols})
```

Length: at most 2^n

Lower bound $\exp(\Omega(n))$ [Urquhart '87, Chvátal & Szemerédi '88]

Clause space: at most *n*

Lower bound $\Omega(n)$ [Torán '99, Alekhnovich et al. '00]

Length-Space Trade-offs

Small space \Rightarrow short length

 \exists constant clause space refutation $\Rightarrow \exists$ polynomial length refutation [Atserias & Dalmau '03]

Converse **not** true

 \exists formulas refutable in linear length requiring $n/\log n$ clause space [Ben-Sasson & Nordström '08]

Dramatic length-space trade-offs in worst case

[Ben-Sasson & Nordström '11] and [Beame, Beck & Impagliazzo '12] showed ∃ formulas that

- are refutable in small length
- are refutable in small space
- allow no meaningful simultaneous optimization (minimizing one measure incurs severe penalty in the other)

What We **Don't** Know About Space

Open Question

Total space quadratic in worst case — is this tight? Not even superlinear lower bounds known!

What We **Don't** Know About Space

Open Question

Total space quadratic in worst case — is this tight? Not even superlinear lower bounds known!

Open Question

3-CNF formula refutable in clause space $s \Rightarrow \text{length } \mathcal{O}(n^s)$. Can you do space $\mathcal{O}(s)$ and length $n^{\mathcal{O}(s)}$ simultaneously?

We Really Don't Understand Space That Well...

All lower bounds on space seem to follow (with hindsight) from

- bounds for other measures that we understand better (e.g. width), or
- connections to pebble games

Pebbling and Time-Space Trade-offs

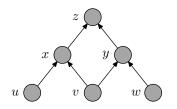
Questions about time-space trade-offs fundamental in TCS

In particular, well-studied (and well-understood) for pebble games modelling calculations described by DAGs ([Cook & Sethi '76] and many others)

- Time needed for calculation: # pebbling moves
- Space needed for calculation: max # pebbles required

The Black-White Pebble Game

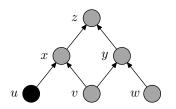
Goal: get single black pebble on sink vertex of G



# moves	0
Current # pebbles	0
Max # pebbles so far	0

The Black-White Pebble Game

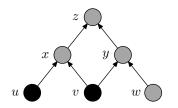
Goal: get single black pebble on sink vertex of G



# moves	1
Current # pebbles	1
Max # pebbles so far	1

• Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them

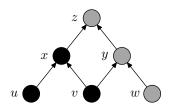
Goal: get single black pebble on sink vertex of G



# moves	2
Current # pebbles	2
Max # pebbles so far	2

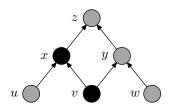
• Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them

Goal: get single black pebble on sink vertex of G



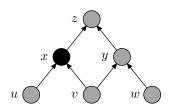
# moves	3
Current # pebbles	3
Max # pebbles so far	3

• Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them



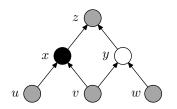
# moves	4
Current # pebbles	2
Max # pebbles so far	3

- Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them
- Can always remove black pebble from vertex



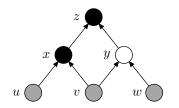
# moves	5
Current # pebbles	1
Max # pebbles so far	3

- Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them
- Can always remove black pebble from vertex



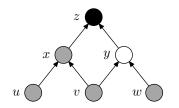
# moves	6
Current # pebbles	2
Max # pebbles so far	3

- Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them
- 2 Can always remove black pebble from vertex
- Can always place white pebble on (empty) vertex



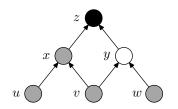
# moves	7
Current # pebbles	3
Max # pebbles so far	3

- Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them
- 2 Can always remove black pebble from vertex
- Can always place white pebble on (empty) vertex



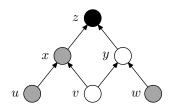
# moves	8
Current # pebbles	2
Max # pebbles so far	3

- Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them
- 2 Can always remove black pebble from vertex
- Can always place white pebble on (empty) vertex



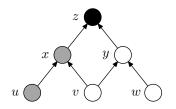
# moves	8
Current # pebbles	2
Max # pebbles so far	3

- Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them
- 2 Can always remove black pebble from vertex
- Can always place white pebble on (empty) vertex
- Can remove white pebble if all immediate predecessors have pebbles



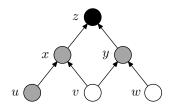
# moves	9
Current # pebbles	3
Max # pebbles so far	3

- Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them
- 2 Can always remove black pebble from vertex
- Can always place white pebble on (empty) vertex
- Can remove white pebble if all immediate predecessors have pebbles



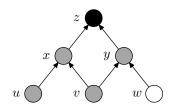
# moves	10
Current # pebbles	4
Max # pebbles so far	4

- Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them
- 2 Can always remove black pebble from vertex
- Can always place white pebble on (empty) vertex
- Can remove white pebble if all immediate predecessors have pebbles



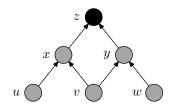
# moves	11
Current # pebbles	3
Max # pebbles so far	4

- Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them
- 2 Can always remove black pebble from vertex
- Can always place white pebble on (empty) vertex
- Can remove white pebble if all immediate predecessors have pebbles



# moves	12
Current # pebbles	2
Max # pebbles so far	4

- Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them
- Can always remove black pebble from vertex
- Can always place white pebble on (empty) vertex
- Can remove white pebble if all immediate predecessors have pebbles



# moves	13
Current # pebbles	1
Max # pebbles so far	4

- Can place black pebble on (empty) vertex if all immediate predecessors have pebbles on them
- Can always remove black pebble from vertex
- Can always place white pebble on (empty) vertex
- Can remove white pebble if all immediate predecessors have pebbles

More About Pebbling

- Black pebbling: Same game but black pebbles only
- Rich literature on both black and black-white pebbling
- Black-white pebbling can save square root over black pebbling space [Wilber '85, Kalyanasundaram & Schnitger '88]*
- But never more [Meyer auf der Heide '81]
- However, transformation from black-white to black-only pebbling incurs exponential time blow-up

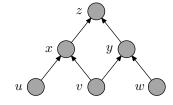
More About Pebbling

- Black pebbling: Same game but black pebbles only
- Rich literature on both black and black-white pebbling
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- But never more [Meyer auf der Heide '81]
- However, transformation from black-white to black-only pebbling incurs exponential time blow-up
- (*) Last two papers in the pebbling literature?

Pebbling Contradictions

CNF formulas encoding pebble game on DAGs

- 1. *u*
- 2. *v*
- 3. w
- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \vee \overline{y} \vee z$
- 7. \overline{z}

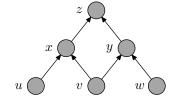


- sources are true
- truth propagates upwards
- but sink is false

Pebbling Contradictions

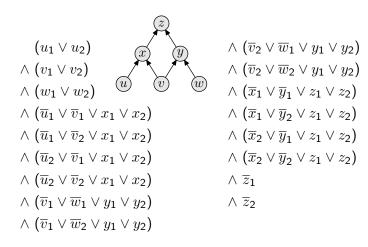
CNF formulas encoding pebble game on DAGs

- 1. *u*
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- 4. $\overline{u} \vee \overline{v} \vee x$
- 5. $\overline{v} \vee \overline{w} \vee y$
- 6. $\overline{x} \vee \overline{y} \vee z$
- 7. \overline{z}

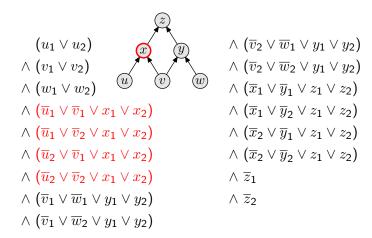


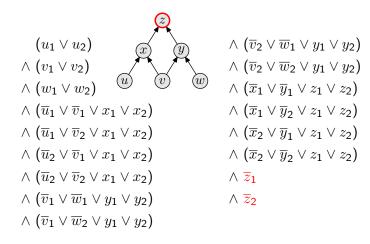
- sources are true
- truth propagates upwards
- but sink is false

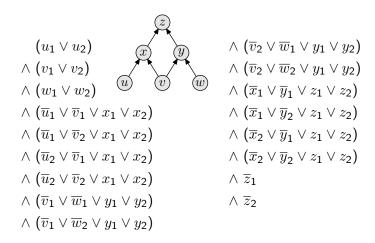
Studied by [Bonet et al. '98, Raz & McKenzie '99, Ben-Sasson & Wigderson '99] and others



$$(u_1 \vee u_2) \qquad \qquad \wedge (\overline{v}_2 \vee \overline{w}_1 \vee y_1 \vee y_2) \\ \wedge (v_1 \vee v_2) \qquad \qquad \wedge (\overline{v}_2 \vee \overline{w}_2 \vee y_1 \vee y_2) \\ \wedge (w_1 \vee w_2) \qquad \qquad \wedge (\overline{x}_1 \vee \overline{y}_1 \vee z_1 \vee z_2) \\ \wedge (\overline{u}_1 \vee \overline{v}_1 \vee x_1 \vee x_2) \qquad \qquad \wedge (\overline{x}_1 \vee \overline{y}_2 \vee z_1 \vee z_2) \\ \wedge (\overline{u}_1 \vee \overline{v}_2 \vee x_1 \vee x_2) \qquad \qquad \wedge (\overline{x}_2 \vee \overline{y}_1 \vee z_1 \vee z_2) \\ \wedge (\overline{u}_2 \vee \overline{v}_1 \vee x_1 \vee x_2) \qquad \qquad \wedge (\overline{x}_2 \vee \overline{y}_2 \vee z_1 \vee z_2) \\ \wedge (\overline{u}_2 \vee \overline{v}_2 \vee x_1 \vee x_2) \qquad \qquad \wedge \overline{z}_1 \\ \wedge (\overline{v}_1 \vee \overline{w}_1 \vee y_1 \vee y_2) \qquad \qquad \wedge \overline{z}_2$$







(*) In fact, they are a bit more involved, but let's stick with this for the purposes of this talk

From Resolution to Black-White Pebbling

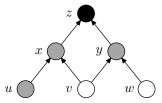
Black-white pebbling models non-deterministic computation

- black pebbles ⇔ computed results
- white pebbles ⇔ guesses needing to be verified

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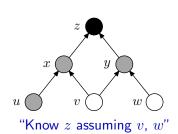


"Know z assuming v, w"

From Resolution to Black-White Pebbling

Black-white pebbling models non-deterministic computation

- black pebbles ⇔ computed results
- white pebbles ⇔ guesses needing to be verified



Corresponds to $(v \wedge w) \rightarrow z$, i.e., blackboard clauses

$$\overline{v}_1 \vee \overline{w}_1 \vee z_1 \vee z_2
\overline{v}_1 \vee \overline{w}_2 \vee z_1 \vee z_2
\overline{v}_2 \vee \overline{w}_1 \vee z_1 \vee z_2
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Formal Refutation-Pebbling Correspondence

Theorem (Ben-Sasson & Nordström '11)

Any refutation translates into black-white pebbling with

- # moves < refutation length
- # pebbles ≤ clause space

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- # pebbles ≤ clause space

Observation (Ben-Sasson et al. '00)

Any black-pebbles-only pebbling translates into refutation with

- refutation length ≤ # moves
- total space ≤ # pebbles

Proof: Just derive $v_1 \lor v_2$ inductively when vertex v is pebbled.

A Fatal Gap and How to Close It

There is a gap in the reductions!

- From resolution to black-white pebbling
- From pebbling to resolution only for black pebbling
- Why worry lose only square root? No, everything! (Due to exponential time blow-up)

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What to do?

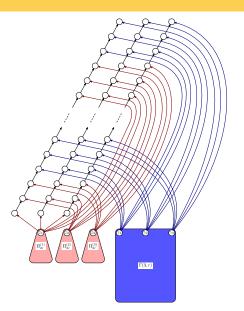
- Find graphs with (essentially) same trade-off properties for black-white and black-only pebbling
- Improve reductions between resolution and pebbling

We contribute in both directions

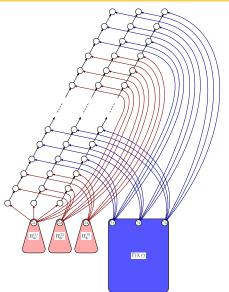
A Picture Says More Than a Thousand Words...

A couple of words about the pebbling result nevertheless:

- Take graph family from [Carlson & Savage '80]
- Black pebbling bounds known (upper and lower)
- Tweak graphs slightly...
- And prove matching black-white lower bounds

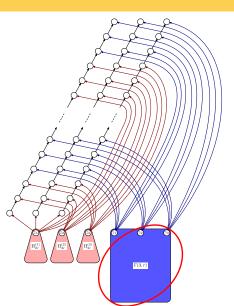


Slightly more detailed proof sketch:



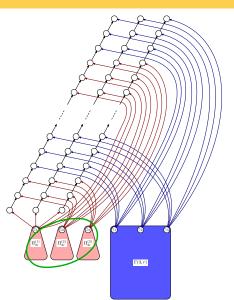
Slightly more detailed proof sketch:

 By induction, small-space pebbling of recursively constructed subgraph requires lots of time



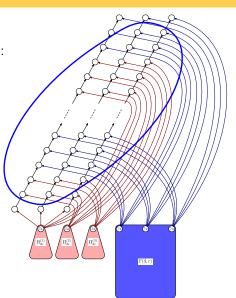
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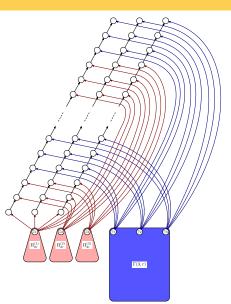
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Return to this after the break...

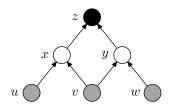


A Naive Idea for Simulating Black-White Pebbling

Run the intuition from [Ben-Sasson & Nordström '11] in reverse

A Naive Idea for Simulating Black-White Pebbling

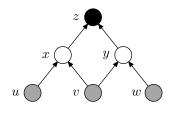
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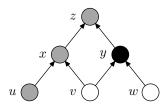


$$\overline{x}_1 \vee \overline{y}_1 \vee z_1 \vee z_2
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\overline{x}_2 \vee \overline{y}_1 \vee z_1 \vee z_2
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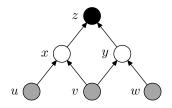


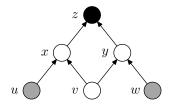
$$\overline{v}_1 \vee \overline{w}_1 \vee y_1 \vee y_2$$

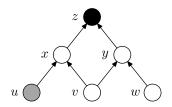
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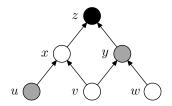
$$\overline{v}_2 \vee \overline{w}_1 \vee y_1 \vee y_2$$

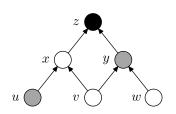
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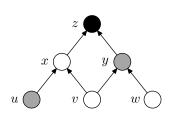








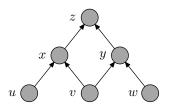
What happens when we try to simulate a pebbling that "combines" these two configurations?



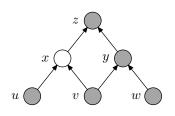
Went only from 2 to 3 white pebbles, but # clauses doubled

Exponential blow-up for naive simulation in worst case

Keep track of for each black pebble which white pebbles it depends on

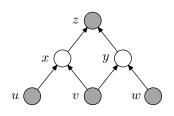


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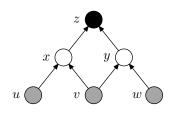
No black pebbles, so no dependencies

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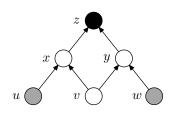
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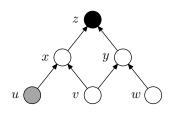
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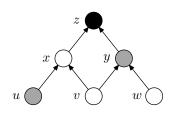
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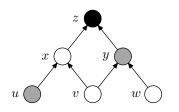
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No black pebbles, so no dependencies Black on z dependent on whites on $\{x, y\}$ Update dependence for z to $\{x, v, w\}$

Require that each black pebble depend on at most $\mathcal{O}(1)$ white pebbles

Black-white pebbling with "limited nondeterminism"

• Pebbling with limited nondeterminism easy to simulate for resolution

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- Turns out all known pebbling separation results for black-white vs. black pebbling can be matched by pebblings with limited nondeterminism
- Yields tight space bounds and time-space trade-offs for pebbling formulas over such graphs
- So, in particular, not possible to reduce from resolution to black-only pebbling

Resolution and Pebbling

Can we reduce from general black-white pebbling to resolution?

Open Question 1

Can resolution on pebbling formulas always simulate black-white pebbling?

Might or might not be true...

Pebbling with Limited Nondeterminism

Open Question 2

Can pebbling with limited nondeterminism always simulate black-white pebbling?

Pebbling with Limited Nondeterminism

Open Question 2

Can pebbling with limited nondeterminism always simulate black-white pebbling?

Affirmative answer to Question 2 would immediately answer Question ${\bf 1}$ as well

Would seem a bit surprising, though...

Candidate for refuting Question 2: Graphs in [Wilber '85]

Space in Resolution

Open Question 3

Total space quadratic in worst case — is this tight? Not even superlinear lower bounds known!

Open Question 4

3-CNF formula refutable in clause space $s\Rightarrow$ length $\mathcal{O}(n^s)$. Can you do space $\mathcal{O}(s)$ and length $n^{\mathcal{O}(s)}$ simultaneously?

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Open Question 4

3-CNF formula refutable in clause space $s \Rightarrow \text{length } \mathcal{O}(n^s)$. Can you do space $\mathcal{O}(s)$ and length $n^{\mathcal{O}(s)}$ simultaneously?

For space s=3 (minimum): No, space-3 refutation can always be carried out in linear length

But space s = 4 already wide open...

Take-Home Message

- There are strong (and surprising!) connections between resolution and pebble games
- But still not fully clarified how tight reductions can we get?
- Also proof space not well-understood many (simple) remaining open questions
- See survey Pebble Games, Proof Complexity, and Time-Space Trade-offs on my webpage for details

Thank you for your attention!