An Efficient Evolutionary Algorithm for Solving Incrementally Structured Problems

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- The PetaBricks language is a collaboration between:
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	- A evolutionary algorithms research group
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- Our goal is to make programs run faster
- We use evolutionary algorithms to search for faster programs
- The PetaBricks language defines search spaces of algorithmic choices

• How would you write a *fast* sorting algorithm?

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A motivating example

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- Insertion sort
- Quick sort
- **•** Merge sort
- Radix sort

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- Binary tree sort, Bitonic sort, Bubble sort, Bucket sort, Burstsort, Cocktail sort, Comb sort, Counting Sort, Distribution sort, Flashsort, Heapsort, Introsort, Library sort, Odd-even sort, Postman sort, Samplesort, Selection sort, Shell sort, Stooge sort, Strand sort, Timsort?

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- Poly-algorithms

```
/usr/include/c++/4.5.2/bits/stl_algo.h lines 3350-3367
/// This is a helper function for the stable sorting routines.
template<typename RandomAccessIterator>
  void
  inplace stable sort( RandomAccessIterator first,
                       RandomAccessIterator last)
    if ( last first < 15)
      \mathcal{L}std:: insertion sort( first, last);
        return:
    RandomAccessIterator middle = first + ( last - first) / 2;
    std:: inplace stable sort( first, middle);
    std:: inplace stable sort( middle, last);
    std:: merge without buffer( first, middle, last,
                               minedle first,
                               \frac{1}{2} last \frac{1}{2} middle);
  λ
```
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Is 15 the right number?

- The best cutoff (CO) changes
- Depends on competing costs:
	- \bullet Cost of computation (< operator, call overhead, etc)
	- Cost of communication (swaps)
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- Sorting 100000 doubles with std::stable_sort:
	- $CO \approx 200$ optimal on a Phenom 905e (15% speedup over $CO = 15$)
	- $CO \approx 400$ optimal on a Opteron 6168 (15% speedup over $CO = 15$)
	- $CO \approx 500$ optimal on a Xeon E5320 (34% speedup over $CO = 15$)
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- If the best cutoff has changed, perhaps best algorithm has also changed

Language

```
either \{InsertionSort( out, in );\} or \{QuickSort (out, in);
\} or \{MergeSort( out, in);\} or \{R a d i x S ort ( out, in );
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Representation

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⇒ Decision tree synthesized by our evolutionary algorithm

Text notation (will be used later): I 600 Q 1420 M^2

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Text notation: M^{16} 75 M^8 1461 M^4 2400 M^2

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- **•** Decision trees
- Algorithm parameters (integers, floats)
- Parallel scheduling / blocking parameters (integers)
- Synthesized scalar functions (not used in the benchmarks shown)
- The average PetaBricks benchmark's genome has:
	- 1.9 decision trees
	- 10.1 algorithm/parallelism/blocking parameters
	- 0.6 synthesized scalar functions
	- 2³¹⁰⁷ possible configurations

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[INCREA](#page-24-0)

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PetaBricks programs at runtime

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• Evaluating objective function is expensive

- Must run the program (at least once)
- More expensive for unfit solutions
- Scales poorly with larger problem sizes
- **•** Fitness is noisy
	- Randomness from parallel races and system noise
	- Testing each candidate only once often produces an worse algorithm
	- Running many trials is expensive
- Decision tree structures are complex
	- Theoretically infinite size
	- We artificially bound them to 2^{736} bits (23 ints) each

- **GPEA: General Purpose Evolutionary Algorithm**
	- Used as a baseline
- INCREA: Incremental Evolutionary Algorithm
	- Bottom-up approach
	- Noisy fitness evaluation strategy
	- Domain informed mutation operators

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State

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Cost of autotuning front-loaded in initial (unfit) population

We could speed up tuning if we start with a faster initial population

Cost of autotuning front-loaded in initial (unfit) population

We could speed up tuning if we start with a faster initial population

Key insight

Smaller input sizes can be used to form better initial population

Train on input size 64

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- Train on input size 32, to form initial population for:
- Train on input size 64

- Train on input size 16, to form initial population for:
- Train on input size 32, to form initial population for:
- **o** Train on input size 64

- Train on input size 1, to form initial population for:
- Train on input size 2, to form initial population for:
- Train on input size 8, to form initial population for:
- Train on input size 16, to form initial population for:
- Train on input size 32, to form initial population for:
- Train on input size 64
- Naturally exploits optimal substructure of problems

- Both strategies terminate slow tests early
- GPEA uses 1 trial per candidate algorithm
- INCREA adaptively changes the number of trials
- Represents fitness as a probability distribution
- Runs a single tailed t-test to get confidence in differences
- **Runs more trails if confidence is low**

- Mutation operators deal with larger structures in the genome
	- "Add algorithm Y to the top of decision tree X"
	- "Scale cutoff X using a lognormal distribution"
- **•** Generated fully automatically by our compiler

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[INCREA](#page-24-0)

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• Measuring convergence time

- Important to both program users and developers
- Vital in online autotuning
- Three fixed-accuracy PetaBricks programs:
	- Sort 2^{20} (small input size)
	- Sort 2^{23} (large input size)
	- Matrix multiply
	- Eigenvector solve
- Representative runs
- Average of 30 runs, with tests for statistical significance in paper
- Run an 8-core Xeon running Debian 5.0

Sort 2²⁰: training input size

Training Input Size Training Input Size

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Sort 2²⁰: candidates tested

Tests Conducted Tests Conducted

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Sort 2^{23} : performance

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Matrix Multiply (input size 1024x1024)

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Eigenvector Solve (input size 1024x1024)

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Take away

The technique of solving incrementally structured problems by exploiting knowledge from smaller problem instances may be more broadly applicable.

Take away

PetaBricks is a useful framework for comparing techniques for autotuning programs.

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- Questions?
- <http://projects.csail.mit.edu/petabricks/>

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