

MASSACHUSETTS INSTITUTE OF TECHNOLOGY  
 Department of Electrical Engineering and Computer Science  
 6.001—Structure and Interpretation of Computer Programs  
 Spring 2006

**Recitation 4**  
**Orders of Growth**

## Definitions

Theta ( $\Theta$ ) notation:

$$f(n) = \Theta(g(n)) \rightarrow k_1 \cdot g(n) \leq f(n) \leq k_2 \cdot g(n), \text{ for } n > n_0$$

Big-O notation:

$$f(n) = O(g(n)) \rightarrow f(n) \leq k \cdot g(n), \text{ for } n > n_0$$

Adversarial approach: For you to show that  $f(n) = \Theta(g(n))$ , you pick  $k_1$ ,  $k_2$ , and  $n_0$ , then I (the adversary) try to pick an  $n$  which doesn't satisfy  $k_1 \cdot g(n) \leq f(n) \leq k_2 \cdot g(n)$ .

## Implications

Ignore constants. Ignore lower order terms. For a sum, take the larger term. For a product, multiply the two terms. Orders of growth are concerned with how the effort scales up as the size of the problem increases, rather than an exact measure of the cost.

## Typical Orders of Growth

- $\Theta(1)$  - Constant growth. Simple, non-looping, non-decomposable operations have constant growth.
- $\Theta(\log n)$  - Logarithmic growth. At each iteration, the problem size is scaled down by a constant amount: (`call-again (/ n c)`).
- $\Theta(n)$  - Linear growth. At each iteration, the problem size is decremented by a constant amount: (`call-again (- n c)`).
- $\Theta(n \log n)$  - Nifty growth. Nice recursive solution to normally  $\Theta(n^2)$  problem.
- $\Theta(n^2)$  - Quadratic growth. Computing correspondence between a set of  $n$  things, or doing something of cost  $n$  to all  $n$  things both result in quadratic growth.
- $\Theta(2^n)$  - Exponential growth. Really bad. Searching all possibilities usually results in exponential growth.

## What's $n$ ?

Order of growth is *always* in terms of the size of the problem. Without stating what the problem is, and what is considered primitive (what is being counted as a “unit of work” or “unit of space”), the order of growth doesn't have any meaning.

## Problems

1. Give order notation for the following:

- (a)  $5n^2 + n$   
 $\Theta(n^2)$
- (b)  $\sqrt{n} + n$   
 $\Theta(n)$
- (c)  $3^n n^2$   
 $\Theta(3^n n^2)$

2. `(define (fact n)`

```
  (if (= n 0)
      1
      (* n (fact (- n 1)))))
```

Running time?  $\Theta(n)$  Space?  $\Theta(n)$

3. `(define (find-e n)`

```
  (if (= n 0)
      1.
      (+ (/ (fact n)) (find-e (- n 1)))))
```

Running time?  $\Theta(n^2)$  Space?  $\Theta(n)$

4. Assume you have a procedure `(divisible? n x)` which returns `#t` if `n` is divisible by `x`. It runs in  $O(n)$  time and  $O(1)$  space. Write a procedure `prime?` which takes a number and returns `#t` if it's prime and `#f` otherwise. You'll want to use a helper procedure.

```
(define (prime? p)
  (define (helper n)
    (if (> n (sqrt p))
        #t
        (if (divisible? p n)
            #f
            (helper (+ n 1)))))
  (helper 2))
```

```
; iterative process
; assuming sqrt takes  $O(1)$  time,  $\text{sqrt}(n) * n$ 
```

Running time?  $\Theta(n\sqrt{n})$  Space?  $\Theta(1)$

5. Write an iterative version of `find-e`. What is its running time and space?

```
(define (find-e-iter n)
  (define (helper n s)
    (if (= n 0) s
        (helper (- n 1) (+ s (/ (fact n))))))
  (helper n 1.0))
```

Running time?  $\Theta(n^2)$  Space?  $\Theta(n)$

## Micro Quiz

Name:

1. Write a procedure that computes the number of decimal digits in its input. Do not use logs. You may use `quotient` (integer division) if you wish.  
`(num-digits 102) → 3`

```
(define (num-digits n)
  (if (= n 0)
      0
      (+ 1 (num-digits (quotient n 10)))))
```

Running time?  $\Theta(\log n)$  Space?  $\Theta(\log n)$

2. Write a procedure that will multiply two positive integers together, but the only arithmetic operation allowed is addition (ie multiplication through repeated addition). In addition, your procedure should be iterative, not recursive.  
`(slow-mul 3 4) → 12`

```
(define (mul-helper a b total)
  (if (= a 0)
      total
      (mul-helper (- a 1) b (+ total b)))) ; or (+ a -1) if picky
(define (slow-mul a b)
  (mul-helper a b 0))
```

Running time?  $\Theta(n)$  Space?  $\Theta(1)$