Marius: Learning Massive Graph Embeddings on a Single Machine

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Figure 2. Starting from a knowledge graph, embedding methods generate representations of the elements of the knowledge graph that are embedded in a vector space. For example, these representations could be vectors... Expand
Problem and Motivation: Learning Graph Embeddings

- Learning graph embeddings is widely used in industry for a variety of downstream tasks.
- This paper focuses on link prediction.

Set up:

- Parameters: 1 high-dimensional vector in $\mathbb{R}^{400}$ per node and per edge type.
- Training: minimize contrastive loss:
  - Example $f(\theta_s, \theta_r, \theta_d) = \theta_s \text{diag}(\theta_r) \theta_d$; makes $f(e) \sim 1$ if $e \in E$, $\sim 0$ otherwise
  - Second term is approximated.

\[
\mathcal{L} = - \sum_{s,r,d \in E} (f(e_\theta) - \log(\sum_{s',r',d' \notin E} e^{f(e_{\theta'})})) \quad (1)
\]

where $e_\theta = (\theta_s, \theta_r, \theta_d)$ and $e_{\theta'} = (\theta'_s, \theta'_r, \theta'_d)$.
Scaling issues

- Very memory intensive!
  - For a graph of 100M nodes, and 400 floats per embedding vector, 4 bytes per float:
  - \( \sim 160 \text{GB} \) Just for the node embeddings.
  - Highest GPU memory: 80 GB available (NVIDIA A100)
  - Modest GPU: 12GB (NVIDIA TITAN V)

- Key challenge:
  - How do you facilitate data movement to and from GPU?
  - CPU RAM \( \leftrightarrow \) GPU? (Amazon’s DGL-KE, 2020)
  - DISK \( \leftrightarrow \) GPU? (PyTorch BigGraph, 2019)
  - Distributed GPU Compute? (both ^)
  - How to overcome memory bottleneck and IO bottleneck?
Solution!

- Pipelining for hiding IO overhead
- Disk $\rightarrow$ CPU RAM (used as cache) $\rightarrow$ GPU
  - Partition nodes
  - Load and store partitions with smart replacement schedule

Figure 4: Marius training pipeline.
Partitions

- Nodes are partitioned into $p$ partitions
- To calculate loss, need to cycle through $p^2$ pairs of partitions.

Figure 3: Partitions and edge-buckets with $p = 4$. All edges in edge-bucket $(0, 2)$ have a source node in node-partition 0 and a destination node in node-partition 2.
Figure 5: Example BETA ordering for \( p = 6 \) and \( c = 3 \). The sequence of partition buffers corresponds to first fixing \( \{0, 1\} \), then replacing \( \{0, 1\} \) with \( \{2, 3\} \), fixing \( \{2, 3\} \), and finally replacing \( \{2, 3\} \) with \( \{5\} \). Each successive buffer differs by one swap. A corresponding edge bucket ordering is shown above the buffers. For each partition buffer in the sequence, all previously unprocessed edge buckets which have their source and destination node partitions in the buffer are added to the ordering (red edge buckets). For each buffer, these edge buckets can be added in any order.
Results

See paper.
Discussion

- Exposing a general interface for implementing similar algorithms?
- Machine-aware self-configuring implementation?