

Moreover, the algorithm works in higher dimensions too (still poly-time)

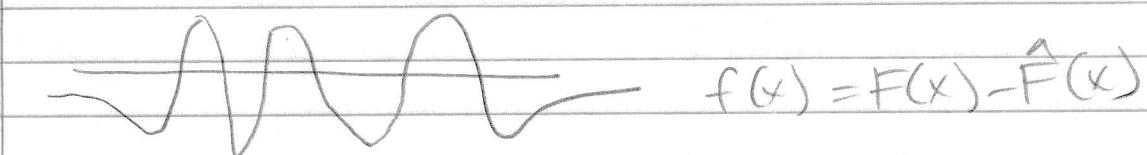
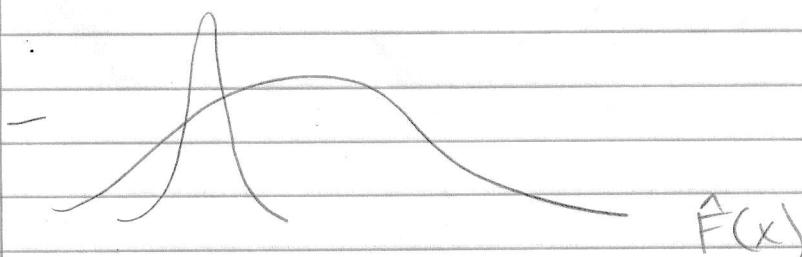
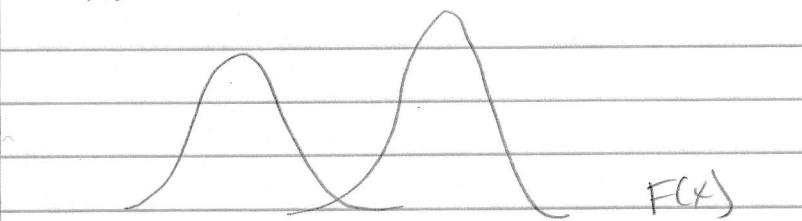
Six Moments S suffice

what if we are given the first six moments exactly? Does this determine the parameters?

Alternatively:

Do any two different mixtures F and \hat{F} necessarily differ on at least one of their first six moments?

Consider the difference between their pdfs



Lemma: If $f(x)$ has at most six zero crossings, then $F(x)$ and $\hat{F}(x)$ differ on one of their first six moments

Proof: We can find a degree ≤ 6 polynomial $p(x)$ that agrees in sign with $f(x)$.

$$\begin{aligned} 0 &< \left| \int p(x) f(x) dx \right| = \left| \int \sum_{r=1}^6 p_r x^r f(x) dx \right| \\ &\leq \sum_{r=1}^6 |p_r| \left| \int x^r f(x) dx \right| \\ &= \sum_{r=1}^6 |p_r| |M_r(F) - M_r(\hat{F})| \end{aligned}$$

So $\exists r \in \{1, \dots, 6\}$ s.t. $|M_r(F) - M_r(\hat{F})| > 0$ \square

What remains is to show:

Proposition If $f(x) = \sum_{i=1}^k \alpha_i N(\mu_i, \sigma_i^2, x)$ is not positive/negative

Identically zero, then $f(x)$ has at most $2k-2$ zero crossings

(Can check with $k=4 \Rightarrow$ six moments suffice)

We'll use the heat equation.

Question: If the initial heat distribution on a 1-d infinite rod (\mathcal{R}) is $f(x) = f(x, 0)$, what is the heat distribution at time t ?

Probabilistic interpretation ($\sigma^2 = 2\kappa t$)

$$f(x, t) = \mathbb{E}_{z \sim N(0, \sigma^2)} [f(x+z, 0)]$$

Alternatively this is a convolution

$$\begin{aligned} f(x, t) &= \int_{-\infty}^{\infty} f(x+z) N(0, \sigma^2, z) dz \\ &\stackrel{\Delta}{=} f(x) * N(0, \sigma^2, x) \end{aligned}$$

Theorem [Hummel, Gidas] Suppose $f(x) : \mathbb{R} \rightarrow \mathbb{R}$ is analytic and has N zeros. Then

$$f(x) * N(0, \sigma^2, x)$$

has at most N zeros (for any $\sigma^2 > 0$)

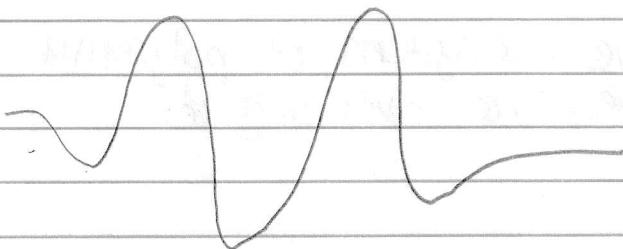
i.e. convolving by a Gaussian / running the heat equation does not increase # of zeros

Last ingredient:

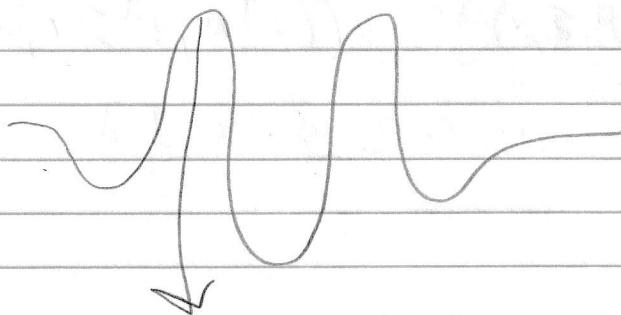
Fact: $N(\mu_1, \sigma_1^2, x) * N(\mu_2, \sigma_2^2, x) =$

$$N(\mu_1 + \mu_2, \sigma_1^2 + \sigma_2^2, x)$$

has at most 4 ZC's. Hence



when we add back in the delta function,
we add at most two new ZC's,



and convolving by $N(0, \sigma_{\min}^2 x)$ gets us
back to the original linear combination,
but does not increase the number of ZC's. TS

Sampling Noise

What if we only have estimates of the
first six moments?

definition Let Θ be the set of valid parameters,
i.e. $w_i \in [0, 1]$ and $\sigma_i^2 \geq 0$

what we just showed is:

$$\{\hat{\theta} \in \Theta \mid M_r(\hat{\theta}) = M_r(\theta) \text{ for } r=1 \text{ to } 6\}$$

the only solutions are $\theta = [w_1, \mu_1, \sigma_1^2, \mu_2, \sigma_2^2]$
and the relabeling $\theta' = [1-w_1, \mu_2, \sigma_2^2, \mu_1, \sigma_1^2]$

Are these equations stable, when given noisy estimates?

Again using deconvolution, can show

Proposition \exists constants c, C s.t. if $\varepsilon < c$, the means and variances $\leq \frac{1}{\varepsilon}$, and mixing weights are in $[\varepsilon, 1-\varepsilon]$ and

$$|M_r(\theta) - M_r(\hat{\theta})| \leq \varepsilon^C \quad r=1 \text{ to } 6$$

then there is a permutation^{II} s.t.

$$\sum_{i=1}^2 |w_i - \hat{w}_{\pi(i)}| + |\mu_i - \hat{\mu}_{\pi(i)}| + |\sigma_i^2 - \hat{\sigma}_{\pi(i)}^2| \leq \varepsilon$$

Hence close enough estimates for the first six moments guarantee close parameters too

This is called polynomial identifiability

reduction
from high d
to 1-d

A View from Algebraic Geometry

Following Belkin and Sinha, we'll give a more general framework

definition: A class of distributions $F(\theta)$ is a polynomial family if $\forall r$

$$M_r(\theta) \stackrel{\Delta}{=} \mathbb{E}[x^r] \\ x \sim F(\theta)$$

is a polynomial in $\theta = (\theta_1, \dots, \theta_k)$

e.g. GMMs, mixtures of uniform, exponential, Poisson, gamma

definition: the moment generating function (mgf) is defined as

$$f(t) = \sum_{n=0}^{\infty} \mathbb{E}[x^n] \frac{t^n}{n!}$$

Fact: If the mgf converges in a neighborhood of zero then

$$\{M_r(\theta) = M_r(\hat{\theta}) \quad \forall r\} \Rightarrow F(\theta) = F(\hat{\theta})$$

i.e. the infinite sequence of moments determines the density function

Now we'll need some notions/tools from algebraic geometry

definition: Given a ring R , an ideal I generated by $g_1, \dots, g_n \in R$ is

$$\langle g_1, \dots, g_n \rangle \stackrel{\Delta}{=} I = \{ \sum r_i g_i \mid \forall i, r_i \in R \}$$

Moreover:

definition: A Noetherian ring is a ring s.t.
for any sequence of ideals

$$I_1 \subseteq I_2 \subseteq \dots$$

$$\exists N \text{ s.t. } I_N = I_{N+1} = I_{N+2} = \dots$$

And our main tool:

Theorem [Hilbert's Basis theorem] If R is
a Noetherian ring, then $R[X]$ is too.

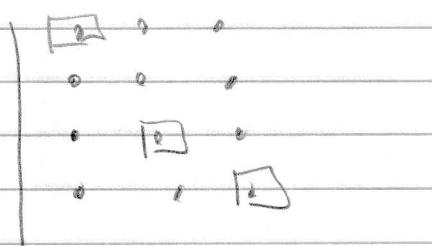
↑
polynomials w/ coeffs in R

Let's pause for some intuition

Lemma: [Dickson] For any subset $K \subseteq \mathbb{N} \times \mathbb{N}$
there are a finite number of minimal elements

i.e. $(a, b) \in K$ s.t. $\nexists (a', b') \in K$
with $a' \leq a$, $b' \leq b$ and one
inequality is strict

We can visualize this



containing only
If I is an ideal generated by polynomials of the form

$$x^a y^b$$

then we know I is Noetherian, and hence has a finite list of generators. This is what Dickson's lemma tells us too.

Warning: An ideal I over the ring $\mathbb{R}[x, y]$ is Noetherian, but there is no effective bound of how many generators we need.

Theorem [Belkin, Sinha] Let $F(\theta)$ be a polynomial family. If the mgf converges in a neighborhood of zero then $\exists N$ s.t.

$$F(\theta) = F(\hat{\theta}) \iff M_r(\theta) = M_r(\hat{\theta}) \text{ for } r=1 \text{ to } N$$

Proof: Let $P_r(\theta, \hat{\theta}) = M_r(\theta) - M_r(\hat{\theta})$

Now let

$$I_1 = \langle \varphi_1 \rangle, I_2 = \langle \varphi_1, \varphi_2 \rangle, \text{ etc.}$$

By Hilbert's basis theorem, $\exists N$ s.t.

$$I_N = I_{N+j} \quad \forall j \geq 0$$

Thus we have

$$\varphi_{N+j}(\theta, \hat{\theta}) = \sum_{i=1}^N P_{ij}(\theta, \hat{\theta}) \varphi_i(\theta, \hat{\theta})$$

for some polynomials $p_{ij} \in \mathbb{R}[\theta, \hat{\theta}]$. thus if

$$M_r(\theta) = M_r(\hat{\theta}) \text{ for } r=1 \dots N$$

we have $Q_r(\theta, \hat{\theta}) = 0$ for $r=1 \dots N$ which implies

$$Q_{N+1}(\theta, \hat{\theta}) = 0 \text{ too!}$$

Now from the fact, we get that $F_r(\theta) = F_r(\hat{\theta})$ \square

so we know finitely many moments suffice,
but don't have an effective bound

Belkin and Sinha also gave a stability analysis
through quantifier elimination

def: A set S is semialgebraic if \exists polynomials s.t.

$$S = \{x_1, \dots, x_d \mid \begin{matrix} p_i(x_1, \dots, x_d) \geq 0 \\ \text{or} \\ p_i(x_1, \dots, x_d) \leq 0 \end{matrix}\}$$

or if S is the finite union or intersection of such sets

① we can define the projection i.e.

$$T = \{x_1, \dots, x_{d-1} \mid \exists x_d \text{ st. } (x_1, \dots, x_d) \in S\}$$

Theorem: [Tarski] the projection of a semialgebraic
set is also semi-algebraic

Notice, if you only use polynomial equations,
you might need polynomial inequalities for the proj.

②

Fact: the complement \bar{S} of a semialgebraic set is also semialgebraic.

Corollary: The set $U = \{x_1, \dots, x_d \mid \forall x_d (x_1, \dots, x_d) \in S\}$ is also semialgebraic.

Proof: We can write

$$\bar{U} = \{x_1, \dots, x_{d-1} \mid \exists x_d (x_1, \dots, x_d) \notin S\}$$

Thus we have

$$S \stackrel{\text{Fact}}{\Rightarrow} \bar{S} \stackrel{\text{Tarski}}{\Rightarrow} \bar{U} \stackrel{\text{Fact}}{\Rightarrow} U$$

are all semialgebraic \square

Now let's prove stability via quantifier elimination

Consider the set $H(\varepsilon, \delta)$ defined to be the set of all ε and δ s.t.

$$\left. \begin{array}{l} \forall \theta, \hat{\theta} \quad |M_r(\theta) - M_r(\hat{\theta})| \leq \delta \text{ for } r=1 \dots 6 \\ d_p(\theta, \hat{\theta}) \leq \varepsilon \end{array} \right\} (P)$$

Claim: $H(\varepsilon, \delta)$ is a semialgebraic set

Proof: Consider the tuple $(\varepsilon, \delta, \theta, \hat{\theta})$ that satisfy the predicate. Equivalently

$$\left. \begin{array}{l} |M_r(\theta) - M_r(\hat{\theta})| \leq \delta \text{ for } r=1 \dots 6 \\ d_p(\theta, \hat{\theta}) \leq \varepsilon \end{array} \right\}$$

Thus it's semialgebraic

Now applying quantifier elimination to get rid of the \forall completes the proof. \square

Theorem [Beltz, Sinha] \exists constants C_1, C_2 s.t.
if $\delta \leq 1/C_1$, then $\varepsilon(\delta) \leq C_2 \delta^{1/s}$
[sketch]

Proof: Since $H(\varepsilon, \delta)$ is semialgebraic,

need ε

show $a(\delta) > 0$

$\forall \delta > 0$

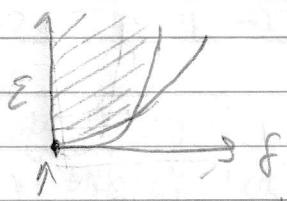
Proof by wiggling

$\forall \delta > 0$

$$\varepsilon^*(\delta) = \text{smallest } \varepsilon \text{ s.t. } (\varepsilon, \delta) \in H(\varepsilon, \delta)$$

$$\varepsilon^*(\delta) \geq \text{poly}(\delta)$$

for sufficiently small δ . \square



Robust Statistics

Basic estimation problem, but we'll go in a new direction:

Given samples from a 1-d Gaussian $N(\mu, \sigma^2)$, can we estimate its parameters?

Of course! Use

$$\hat{\mu} = \frac{1}{N} \sum_{i=1}^N x_i, \quad \hat{\sigma}^2 = \frac{1}{N} \sum_{i=1}^N (x_i - \hat{\mu})^2$$

These are examples of maximum likelihood paradigm (Ronald Fisher 1912-1922)

(1) consistent: converges to true parameters as $N \rightarrow \infty$ under tame conditions

(2) asymptotically normal: has smallest variance among all unbiased estimators

Main Question: But what if the samples are only approximately Gaussian?

definition: In the strong contamination model:

(1) m samples are drawn iid from $P \in \mathcal{D}$

known class of distributions