

## Lecture 9

### Szemerédi's Regularity Lemma

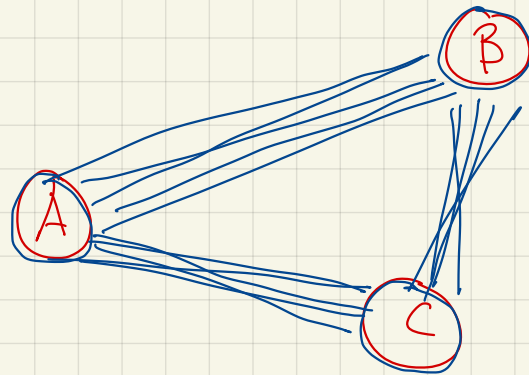
Testing dense graph properties via SRL:

$\Delta$ -freeness

Graphs with "random" properties:

Example question:

How many triangles in a random tripartite graph?



density  $\eta$

$\forall u \in A, v \in B, w \in C:$

$$\Pr[u \sim v \sim w] = \eta^3$$

$$\delta_{u,v,w} = \begin{cases} 1 & \text{if } u \sim v \sim w \\ 0 & \text{o.w.} \end{cases}$$

$$E[\delta_{u,v,w}] = \eta^3$$

$$E[\# \text{ triangles}] = E\left[\sum_{\substack{u \in A \\ v \in B \\ w \in C}} \delta_{u,v,w}\right] = \eta^3 \cdot |A| \cdot |B| \cdot |C|$$

Can we make weaker assumptions + still get  
reasonable bounds?

# Density & Regularity of set pairs:

def. For  $A, B \subseteq V$  s.t.

(1)  $A \cap B = \emptyset$

(2)  $|A|, |B| > 1$

Let  $e(A, B) = \# \text{ edges between } A \text{ \& B}$

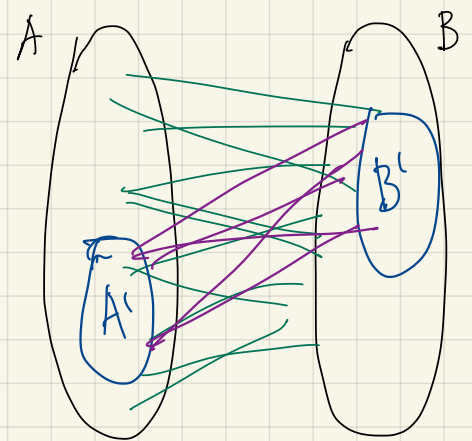
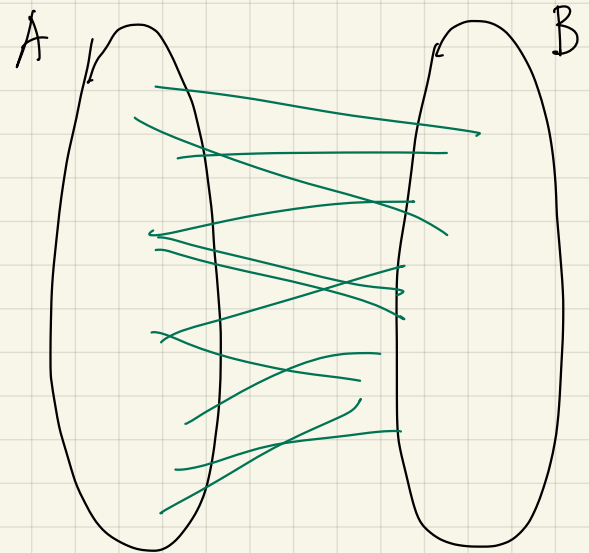
$\&$  density  $d(A, B) = \frac{e(A, B)}{|A| \cdot |B|}$

Say  $A, B$  is  $\gamma$ -regular if  $\forall A' \subseteq A, B' \subseteq B$

s.t.  $|A'| \geq \gamma |A|$

$|B'| \geq \gamma |B|$

$$|d(A', B') - d(A, B)| < \gamma$$



behaves like "a random graph"

Lemma density

$$\forall \eta > 0$$

regularity parameter, depends only on  $\eta$

$$\exists \gamma = \frac{1}{2} \eta \equiv \gamma^\Delta(\eta)$$

$$\delta = (1-\eta) \frac{\eta^3}{8} \geq \frac{\eta^3}{16} \equiv \delta^\Delta(\eta)$$

# triangles, depends only on  $\eta$

if  $\eta < \frac{1}{2}$

$$d(A,B) = \frac{e(A,B)}{|A| \cdot |B|}$$

$A, B$  is  $\gamma$ -regular if  $\forall A' \subseteq A, B' \subseteq B$

$$\text{s.t. } |A'| \geq \gamma |A|$$

$$|B'| \geq \gamma |B|$$

$$|d(A',B') - d(A,B)| < \gamma$$

s.t. if  $A, B, C$  disjoint subsets of  $V$  s.t. each pair

is  $\gamma$ -regular with density  $> \eta$

then  $G$  contains  $\geq \delta \cdot |A| \cdot |B| \cdot |C|$  distinct  $\Delta$ 's

with node in each of  $A, B, C$ .

if  $A, B, C$  disjoint subsets of  $V$  s.t. each pair is  $\gamma$ -regular with density  $> \eta$  then  $G$  contains  $\geq \delta \cdot |A| \cdot |B| \cdot |C|$  distinct  $\Delta$ 's

$$d(A, B) = \frac{e(A, B)}{|A| \cdot |B|}$$

$A, B$  is  $\gamma$ -regular if  $\forall A' \subseteq A, B' \subseteq B$  s.t.  $|A'| \geq \gamma |A|$   $|B'| \geq \gamma |B|$

$$|d(A', B') - d(A, B)| < \gamma$$

Proof:  $A^* \leftarrow$  nodes in  $A$  with  $\geq (\eta - \gamma) \cdot |B|$  nbrs in  $B$   $\geq (\eta - \gamma) \cdot |C|$  nbrs in  $C$

Claim  $|A^*| \geq (1 - 2\gamma) |A|$

Why? (PF of claim)

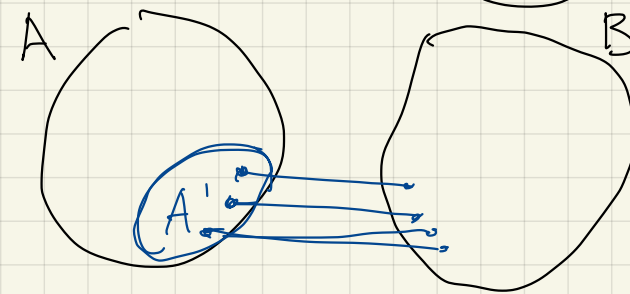
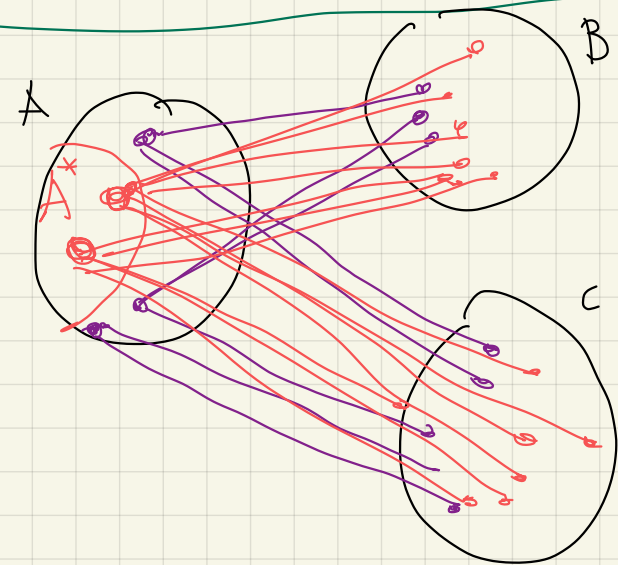
$A' \leftarrow$  "bad" nodes w.r.t.  $B$  ( $< (\eta - \gamma) \cdot |B|$  nbrs in  $B$ )  
 $A'' \leftarrow$  " " " " ( $<$  " " " "  $C$ )

then  $|A'| \leq \gamma |A|$  ( $\& \ |A''| \leq \gamma |A|$ )

Why? consider pair  $A', B$ .  
 $d(A', B) < \frac{|A'| \cdot (\eta - \gamma) \cdot |B|}{|A'| \cdot |B|} = \eta - \gamma$   $\leftarrow$  def of  $A'$

but  $d(A, B) > \eta$   
 so  $|d(A', B) - d(A, B)| > \gamma$

$\& \ |B'| \geq \gamma |B|$   
 so if  $|A'| \geq \gamma |A|$  then  $(A, B)$  is not  $\gamma$ -regular  $\rightarrow \leftarrow$



$$\text{Let } A^* = A \setminus (A' \cup A'') \text{ then } |A^*| \geq |A| - |A'| - |A''| \geq |A| - 2\delta|A| = (1-2\delta) \cdot |A| \quad \square$$

if  $A, B, C$  disjoint subsets of  $V$  s.t. each pair is  $\gamma$ -regular with density  $> \eta$  then  $G$  contains  $\geq \delta \cdot |A| \cdot |B| \cdot |C|$  distinct  $\Delta$ 's

$$d(A, B) = \frac{e(A, B)}{|A| \cdot |B|}$$

$A, B$  is  $\gamma$ -regular if  $\forall A' \subseteq A, B' \subseteq B$  s.t.  $|A'| \geq \gamma |A|$   $|B'| \geq \gamma |B|$

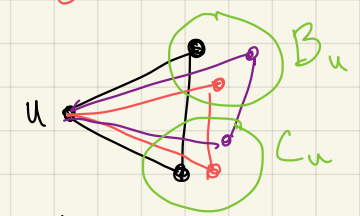
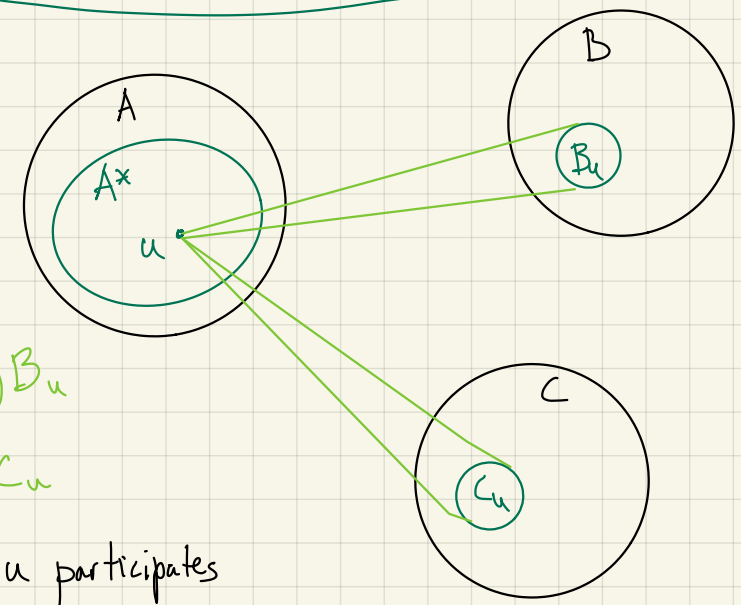
$$|d(A', B') - d(A, B)| < \gamma$$

Proof:  $A^* \leftarrow$  nodes in  $A$  with  $\geq (\eta - \gamma) \cdot |B|$  nbrs in  $B$   
 $\geq (\eta - \gamma) \cdot |C|$  nbrs in  $C$

Claim  $|A^*| \geq (1 - 2\gamma) |A|$

For each  $u \in A^*$ : define  $B_u \equiv$  nbrs of  $u$  in  $B$   
 $C_u \equiv$  nbrs of  $u$  in  $C$

both pretty big by def of  $A^*$



since  $\gamma < \frac{\eta}{2}$ ,  $|B_u| \geq (\eta - \gamma) |B| \geq \gamma |B|$   
 $(\eta - \gamma > \gamma)^2$ ,  $|C_u| \geq (\eta - \gamma) |C| \geq \gamma |C|$

Find lots of distinct  $\Delta$ 's

#Edges between  $B_u + C_u \Rightarrow$  lower bound on # distinct  $\Delta$ 's in which  $u$  participates  
 $d(B, C) \geq \eta \Rightarrow d(B_u, C_u) \geq \eta - \gamma \Rightarrow e(B_u, C_u) \geq (\eta - \gamma) |B_u| |C_u| \geq (\eta - \gamma)^3 \cdot |B| \cdot |C|$   
 +  $B_u, C_u$  big enough +  $(B, C)$  is  $\gamma$  regular

so total #  $\Delta$ 's  $\geq (1 - 2\gamma) |A| \cdot (\eta - \gamma)^3 |B| |C| \geq (1 - \eta) (\frac{\eta}{2})^3 |A| |B| |C|$   
 choose  $\gamma \leq \frac{\eta}{2}$





Do interesting graphs have regularity properties?

Yes in some sense, all graphs do "can be approximated as small collection of random graphs"

## Szemerédi's Regularity Lemma

would like it to say:

"one can equipartition nodes of  $V$  into  $V_1 \dots V_k$  (for const  $k$ ) s.t.

all pairs  $(V_i, V_j)$  are  $\epsilon$ -regular"

only most  
 $\leq \epsilon \binom{k}{2}$   
are not

↑  
Sometimes need  $k > m$   
for some  $m$   
( $k=1, k=n$  trivial)

# Szemerédi's Regularity Lemma: (especially useful version)

$\forall m, \epsilon > 0 \quad \exists T = T(m, \epsilon)$  s.t. given  $G = (V, E)$  s.t.  $|V| > T$

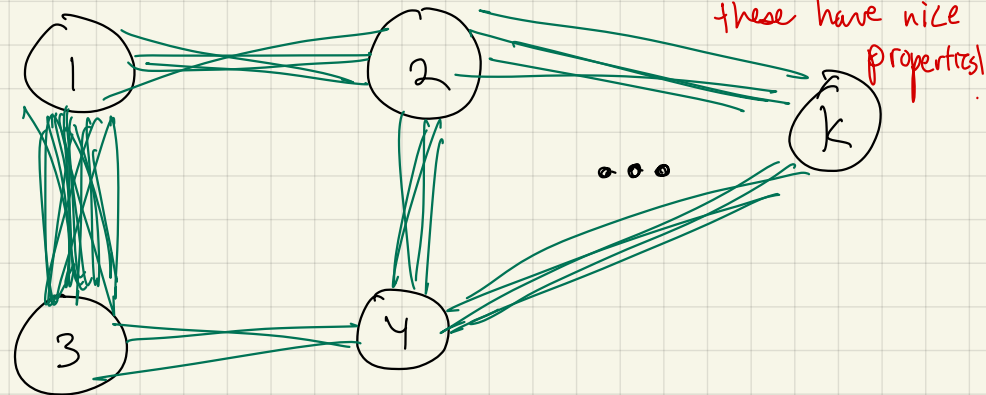
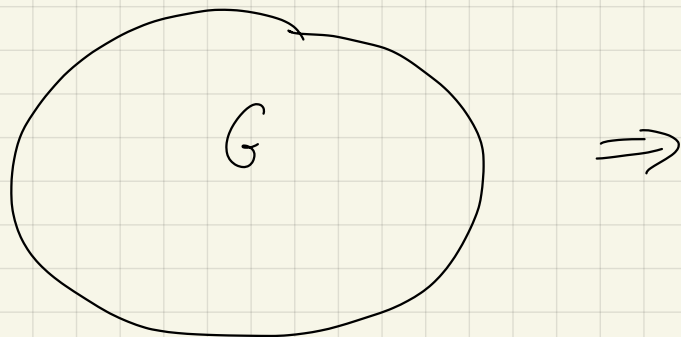
$\downarrow$   $\mathcal{A}$  an equipartition of  $V$  into sets

then exists equipartition  $\mathcal{B}$  into  $k$  sets which refines  $\mathcal{A}$

s.t.  $m \leq k \leq T$

$\downarrow$   $\leq \epsilon \binom{k}{2}$  set pairs not  $\epsilon$ -regular

Note:  $T$  does not depend on  $|V|$



Why was SRL first studied?

to prove conjecture of Erdős + Turán: sequences of ints have long arithmetic progressions

Very rough idea of proof:

"expectation of  $d^2(v_i, v_j)$ "

→

$$\text{ind}(V_1, \dots, V_k) = \frac{1}{k^2} \sum_{i=1}^k \sum_{j=i+1}^k d^2(v_i, v_j) \leq \frac{1}{2}$$

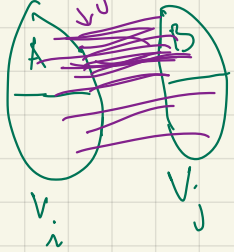
← same densities

"variance of  $d$ "

note:

$$E[d(v_i, v_j)] = \frac{|E|}{|V|^2}$$

high density



if a partition violates, can refine st.

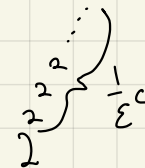
$\text{ind}(V'_1, \dots, V'_k)$  grows significantly (ie. by  $\approx \epsilon^c$ )

so in less than  $\frac{1}{\epsilon^c}$  refinements, have good

partition

} note, if refine, Cauchy Schwartz  $\Rightarrow$  ind can't decrease

How big is  $k$ ? u.b. tower of size  $\frac{1}{\epsilon^c}$   
 l.b. " " "  $\frac{1}{\epsilon^c}$



issue: what if

split  $v_i$  for many  $v_j$ ?

$\Rightarrow$  split into exponential subsets