

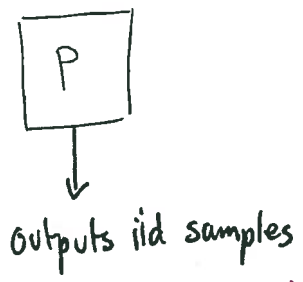
Testing Distributions :

Uniformity

Turning to a new model:

prob dists

Probability distributions - get samples of distribution



Domain D , $|D|=n$ ← known
 $p_i = \Pr[p \text{ outputs } i]$ ← unknown

← this is all we can learn from

Examples:

- Lottery data
- Shopping choices
- experimental outcomes
- ⋮

What do we want to know?

- is it uniform? eg. lottery
- is it high entropy?
- large support? (many distinct elements have >0 probability)
- is it monotone increasing, K-modal, monotone hazard rate...?

how can we do it?

χ^2 test

plug in estimate

learn distribution, Maximum likelihood estimates

Goal: sample complexity **SUBLINEAR** in n

Testing Uniformity

The goal:

Uniform dist on D

- if $P \equiv U_D$ then tester outputs PASS \leftarrow with prob $\geq 3/4$
- if $\underbrace{\text{dist}(P, U_D)} > \epsilon$ then tester outputs FAIL

which measure of distance?

$l_1, l_2, \text{KL-divergence, Earthmover, Jensen-Shannon}$

today's focus \nearrow

good direction for projects!

Distances

l_1 -distance : $\|p-q\|_1 = \sum_{i \in D} |p_i - q_i|$

l_2 -distance : $\|p-q\|_2 = \sqrt{\sum_{i \in D} (p_i - q_i)^2}$

$\|p-q\|_1 = 2 \cdot \text{TVD dist}(p,q)$
 $\equiv \max_{S \subseteq D} \left\{ \sum_{i \in S} |p_i - q_i| \right\}$
 "Total Variation distance"

Fact: $\|p-q\|_2 \leq \|p-q\|_1 \leq n^{1/2} \|p-q\|_2$

examples:

① $p = (1, 0, 0, \dots, 0)$



$q = (\frac{1}{n}, \frac{1}{n}, \dots, \frac{1}{n})$

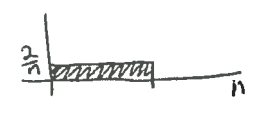


l_1 distance:
 $\|p-q\|_1 = \frac{(n-1)}{n} + (n-1) \cdot \frac{1}{n} \approx 2$

l_2 distance:
 $\|p-q\|_2^2 = (1 - \frac{1}{n})^2 + (n-1) \cdot (\frac{1}{n})^2 \approx 1$

②

$p = (\frac{2}{n}, \frac{2}{n}, \dots, \frac{2}{n}, 0, 0, \dots, 0)$



$q = (0, 0, \dots, 0, \frac{2}{n}, \frac{2}{n}, \dots, \frac{2}{n})$



l_1 distance:
 $\|p-q\|_1 = n \cdot (\frac{2}{n}) = 2$

l_2 distance: $\|p-q\|_2^2 = n \cdot (\frac{2}{n})^2 = \frac{4}{n}$

$\|p-q\|_2 = \frac{2}{\sqrt{n}}$

as far as possible

\Rightarrow so l_2 -distance can be weird (but above fact limits how weird)

↑
pretty close

"Plug-in Estimate"

Algorithm:

- Take m samples from p
- $\forall x$, estimate $p(x)$ by $\hat{p}(x) = \frac{\# \text{ times } x \text{ appears in sample}}{m}$
- if $\sum_x |\hat{p}(x) - \frac{1}{n}| \geq \epsilon$ reject
else accept

Analysis: (better analyses exist, e.g. next page)

Correct
behavior

pick m st. whp $\forall x |p(x) - \hat{p}(x)| < \epsilon/n$ (assume this holds in following)
 $\Rightarrow \|\hat{p} - p\|_1 < \epsilon$

so if $p = u$, test will accept

to show if $\|p - u\|_1 > 2\epsilon$, likely to reject,
will show contrapositive:

If test accepts, good approx test accepts

by $\Delta \#$: if $\|p - \hat{p}\|_1 < \epsilon + \|\hat{p} - u\|_1 < \epsilon$
then $\|p - u\|_1 < 2\epsilon$ ▣

but how big should m be?

$\Omega(m)$? coupon collector? ... (2)

Above algorithm gives good approx in $O(n/\epsilon^2)$ samples:

Thm if $m = \Theta\left(\frac{n}{\epsilon^2}\right)$, $\Pr[\|\hat{p} - p\|_1 \leq \epsilon] \geq 3/4$

⑤
p.d

Better analysis: (not done in lecture)

Claim $E[\|\hat{p} - p\|_1] \leq \sqrt{\frac{n}{m}}$

Pf

$$E[\|\hat{p} - p\|_1] = \sum_x E[|\hat{p}(x) - p(x)|]$$

$$\leq \sum_x \sqrt{E[(\hat{p}(x) - p(x))^2]}$$

$$= \sum_x \sqrt{\text{Var}(\hat{p}(x))}$$

$$\leq \sum_x \sqrt{\frac{p(x)}{m}}$$

$$\leq \frac{1}{\sqrt{m}} \cdot \sqrt{n}$$

def of \hat{p} + lin of expectation

note:

$$E[\hat{p}(x)] = \frac{1}{m} E\left[\sum_{i=1}^m \mathbb{1}_{i^{\text{th}} \text{ sample is } x}\right]$$

$$= \frac{1}{m} \sum_{i=1}^m E[\mathbb{1}_{i^{\text{th}} \text{ sample is } x}]$$

$$= \frac{m \cdot p(x)}{m} = p(x)$$

Jensen's \neq

$$\text{Var}(\hat{p}(x)) = \frac{1}{m^2} m p(x)(1-p(x)) \leq \frac{p(x)}{m}$$

since $\max_{p \in \text{prob dist over domain of size } n} \sum \sqrt{p(x)}$ is \sqrt{n}

So picking $m = \Omega\left(\frac{n}{\epsilon^2}\right)$ gives

$$E[\|\hat{p} - p\|_1] \leq \frac{\epsilon}{2}$$

by Markov's \neq : with prob $1 - \frac{1}{2}$, $\|\hat{p} - p\|_1 \leq \epsilon$

Note, this says we can "learn" (approximate) any dist wr.t. L_1 distance in $\Theta(n/\epsilon^2)$ samples

Lets consider an "easier" problem - L_2 -distance

⑥
pd

L_2 -Distance (squared):

$$\begin{aligned}\|p - u\|_2^2 &= \sum_{i \in [n]} (p_i - \frac{1}{n})^2 \\ &= \sum p_i^2 - \underbrace{\frac{2}{n} \sum p_i}_{=1} + \underbrace{\sum (\frac{1}{n})^2}_{=\frac{1}{n}} \\ &= \sum p_i^2 - \frac{1}{n}\end{aligned}$$

Collision probability of p :

$$\|p\|_2^2 \equiv \Pr_{s, t \in p} [s = t] = \sum p_i^2$$

$$\text{for } p = u, \quad \|p\|_2^2 = \frac{1}{n}$$

$$\text{for } p \neq u, \quad \|p\|_2^2 > \frac{1}{n}$$

$$= \|p\|_2^2 - \|u\|_2^2$$

we can estimate this

we know this since we know n

Algorithm

1. take s samples from p
2. let $\hat{c} \leftarrow$ estimate of $\|p\|_2^2$ from sample
3. if $\hat{c} < \frac{1}{n} + \delta$ pass
else fail

① how many samples?

② how?

③ what should δ be?

Thm if $s = \Theta(\sqrt{n}/\epsilon^4)$, $\Pr[|\hat{c} - \|p\|_2^2| > \frac{\epsilon^2}{2}] \leq 1/4$

† Algorithm is property tester for uniformity under L_2 -dist.

Question

⑦
p.0

② How to estimate $\|p\|_2^2$?

Naive idea: (pair off samples)

take two raw samples:

$$b_x \leftarrow \begin{cases} 1 & \text{if samples are equal} \\ 0 & \text{o.w.} \end{cases}$$

} b_x 's are independent

" gives $\Theta(k)$ samples of collision probability from k samples of p "

Better idea: recycle - use all pairs in sample

" gives $\Theta(k^2)$ samples of collision probability from k samples of p "

} b_{ij} 's are not independent

$$b_{ij} \leftarrow \begin{cases} 1 & \text{if sample } i \text{ \& } j \text{ are equal} \\ 0 & \text{o.w.} \end{cases}$$

Estimate by recycling:

• Take s samples from p : X_1, \dots, X_s

• for each $1 \leq i < j \leq s$

$$b_{ij} \leftarrow \begin{cases} 1 & \text{if } X_i = X_j \\ 0 & \text{if } X_i \neq X_j \end{cases}$$

• Output $\hat{c} \leftarrow \frac{\sum_{i < j} b_{ij}}{\binom{s}{2}}$

} b_{ij} 's not independent so can't use Chernoff

Analysis:
$$E[\hat{c}] = \frac{1}{\binom{s}{2}} \cdot \binom{s}{2} \cdot E[b_{ij}] = \|p\|_2^2$$

Question 3

How well do we need to estimate $\|p\|_2^2$?

(+ how to pick δ ?)

Pick $\delta = \Delta = \frac{\epsilon^2}{2}$

Assumption \star : $|\hat{C} - \|p\|_2^2| < \Delta$

will take enough samples so that this holds with prob $\geq 3/4$

this is our parameter that determines whether our approximation is good. Spoiler: will set $\Delta = \frac{\epsilon^2}{2}$

What happens if \star holds with $\Delta = \frac{\epsilon^2}{2}$?

Correct behavior!

• if $p = U_{[n]}$ then

$\hat{C} \leq \|U_{[n]}\|_2^2 + \Delta = \frac{1}{n} + \frac{\epsilon^2}{2}$

so test will PASS

• if $\|p - U_{[n]}\|_2 > \epsilon$ then

$\|p - U_{[n]}\|_2^2 > \epsilon^2$

but

$\|p\|_2^2 = \|p - U_{[n]}\|_2^2 + \frac{1}{n} > \epsilon^2 + \frac{1}{n}$

← see p.6

+

$\hat{C} > \|p\|_2^2 - \Delta$

$\geq \epsilon^2 + \frac{1}{n} - \Delta = \epsilon^2 + \frac{1}{n} - \frac{\epsilon^2}{2} = \frac{\epsilon^2}{2} + \frac{1}{n}$

so test will FAIL

Remaining Question: (Question 1)

How many samples do we need to estimate \hat{C} to within Δ ?

Question ①:

Analysis: $E[\delta_{ij}] = \Pr[b_{ij} = 1] = \|p\|_2^2$

$$E[\hat{C}] = \frac{1}{\binom{s}{2}} \cdot \binom{s}{2} \cdot E[b_{ij}] = \|p\|_2^2$$

$$\Pr[|\hat{C} - \|p\|_2^2| > \rho] \leq \frac{\text{Var}[\hat{C}]}{\rho^2} \quad \text{Chebyshev}$$

Fact $\text{Var}[aX] = a^2 \text{Var}[X]$

$$\begin{aligned} \text{So } \text{Var}[\hat{C}] &= \text{Var}\left[\frac{1}{\binom{s}{2}} \sum_{i < j} b_{ij}\right] \\ &= \frac{1}{\binom{s}{2}^2} \text{Var}\left[\sum_{i < j} b_{ij}\right] \end{aligned}$$

Lemma $\text{Var}\left[\sum_{i < j} b_{ij}\right] \leq O(s^3 \cdot \|p\|_2^3)$

Corr $\text{Var}[\hat{C}] \leq O(\|p\|_2^3 / s)$

Proof of lemma def $\bar{b}_{ij} = b_{ij} - E[b_{ij}]$ (nice trick)

note $E[\bar{b}_{ij}] = 0$ and $\bar{b}_{ij} < b_{ij}$ since $E[b_{ij}] > 0$

also $E[\bar{b}_{ij} \bar{b}_{kl}] \leq E[b_{ij} b_{kl}]$

$$\text{Var} \left[\sum_{i < j} b_{ij} \right] = E \left[\left(\sum_{i < j} b_{ij} - E \left[\sum_{i < j} b_{ij} \right] \right)^2 \right]$$

$$= E \left[\left(\sum_{i < j} \bar{b}_{ij} \right)^2 \right]$$

$$= E \left[\sum_{\substack{i < j \\ k < l}} \bar{b}_{ij} \bar{b}_{kl} + \sum_{\substack{i < j \\ k < l}} \bar{b}_{ij} \bar{b}_{kl} + \sum_{\substack{i < j \\ k < l}} \bar{b}_{ij} \bar{b}_{kl} \right]$$

$$|\{i, j, k, l\}| = 2$$

$$|\{i, j, k, l\}| = 3$$

$$|\{i, j, k, l\}| = 4$$

①

②

③

Bounding ①:

$$E \left[\sum_{\substack{i < j \\ k < l}} \bar{b}_{ij} \bar{b}_{kl} \right] = E \left[\sum_{i < j} \bar{b}_{ij}^2 \right] \leq E \left[\sum_{i < j} b_{ij}^2 \right] = \binom{s}{2} \|p\|_2^2$$

$$|\{i, j, k, l\}| = 2$$

note $b_{ij}^2 = \bar{b}_{ij}$
since indicator variable

Bounding ③:

$$E \left[\sum_{\substack{i < j \\ k < l}} \bar{b}_{ij} \bar{b}_{kl} \right] = \sum_{\substack{i < j \\ k < l}} E[\bar{b}_{ij}] E[\bar{b}_{kl}] = 0$$

independence
linearity of expectations

note: moving to \bar{b}_{ij}
means all of these
terms drop out!

Bounding ②:

$$E \left[\sum_{\substack{i < j \\ k < l}} \bar{b}_{ij} \bar{b}_{kl} \right] \leq E \left[\sum_{\substack{i < j \\ k < l}} b_{ij} b_{kl} \right] = \sum_{\substack{i < j \\ k < l}} E[b_{ij} b_{kl}] = \sum_{\{a, b, c\}} P[X_a = X_b = X_c]$$

$$|\{i, j, k, l\}| = 3$$

$$|\{i, j, k, l\}| = 3$$

$$|\{i, j, k, l\}| = 3$$

$$|\{a, b, c\}| = 3$$

e.g. $i=k, j=l, i \neq l, j \neq k \dots$

ways to
pick $i < j$
 $k < l$

st. $|\{i, j, k, l\}| = 3$:

(pick 3 indices, pick
one to be repeated twice
leaves 2 options)

$$\leq 6 \cdot \binom{s}{3} \cdot \sum_x p(x)^3$$

$$\leq 6 \cdot \binom{s}{3} \cdot \left(\sum_x p(x)^2 \right)^{3/2}$$

$$\leq O \left(s^3 \cdot (\|p\|_2^2)^{3/2} \right) = O \left(s^3 \|p\|_2^3 \right)$$

$$\text{note } \left(\sum p(x)^3 \right)^{1/3} \leq \left(\sum p(x)^2 \right)^{1/2}$$

$$\text{So: } \text{Var} \left[\sum_{i < j} b_{ij} \right] = O \left(\binom{s}{2} \|p\|_2^2 + 0 + s^3 \cdot \|p\|_2^3 \right) \\ = O \left(s^3 \|p\|_2^3 \right)$$

So how many samples?

for property tester wrt L_2 -distance

need to estimate $\|p\|_2^2$ to within (additive) $\Delta = \frac{\varepsilon^2}{2}$

$$\Pr \left[|\hat{c} - \|p\|_2^2| > \frac{\varepsilon^2}{2} \right] \leq \frac{\text{Var}[\hat{c}]}{\varepsilon^4/4} = \frac{1}{\binom{s}{2}^2} \cdot \frac{s^3 \cdot \|p\|_2^3}{\varepsilon^4} \cdot \text{Const}$$

$$\leq O \left(\frac{1}{s} \cdot \frac{1}{\varepsilon^4} \cdot \|p\|_2^3 \right)$$

want this
to be small

≤ 1

$$\text{Pick } s = \Omega \left(\frac{1}{\varepsilon^4} \right) \quad (\text{can do better})$$

what about L_1 -distance?

Now: Distinguish $\|p - u\|_1 \geq \varepsilon$ from $p = u$
 via L_2 -testing

Thm there is distribution testing algorithm
 which tests uniformity (in L_1) & outputs
 correct answer w/ prob $\geq 1 - \delta$
 using $O(\frac{1}{\varepsilon^4} \sqrt{n} \log(1/\delta))$ samples

Why?

$$\text{if } \|p - u\|_1 = 0 \iff \|p - u\|_2 = 0 \iff \|p\|_2^2 = \frac{1}{n}$$

$$\text{if } \|p - u\|_1 > \varepsilon \implies \|p - u\|_2 > \frac{\varepsilon}{\sqrt{n}} \implies \|p - u\|_2^2 > \frac{\varepsilon^2}{n}$$

$$\implies \|p\|_2^2 > \frac{1 + \varepsilon^2}{n}$$

need to get *multiplicative* estimate of $\|p\|_2^2$ to w/in $(1 \pm \frac{\varepsilon}{4\sqrt{n}})$
 or *additive* " " " " " " $\frac{\varepsilon^2}{2n}$ do this

$$\Pr[|\hat{C} - \|p\|_2^2| > \gamma \cdot \|p\|_2^2] \leq \frac{\text{Var}[\hat{C}]}{\gamma^2 \|p\|_2^4} \leq \frac{\text{Const} \cdot \|p\|_2^3 / s}{\gamma^2 \|p\|_2^4} = \frac{\text{Const}}{s \gamma^2 \cdot \|p\|_2}$$

$$\leq O\left(\frac{\sqrt{n}}{s \gamma^2}\right)$$

$\uparrow \approx \frac{1}{\sqrt{n}}$

so pick $s = \Omega\left(\frac{\sqrt{n}}{\varepsilon^4}\right)$

