Homework guidelines: You may work with other students, as long as (1) they have not yet solved the problem, (2) you write down the names of all other students with which you discussed the problem, and (3) you write up the solution on your own. No points will be deducted, no matter how many people you talk to, as long as you are honest. If you already knew the answer to one of the problems (call these ”famous” problems), then let me know that in your solution writeup – it will not affect your score, but will help me in the future. It’s ok to look up famous sums and inequalities that help you to solve the problem, but don’t look up an entire solution.

The following problems are to be turned in. You should upload your solution to Stellar as a pdf file.

1. Say that $f_1, f_2, f_3$, mapping from group $G$ to $H$, are linear consistent if there exists a linear function $\phi : G \rightarrow H$ (that is $\forall x, y \in G, \phi(x) + \phi(y) = \phi(x + y)$) and $a_1, a_2, a_3 \in H$ such that $a_1 + a_2 = a_3$ and $f_i(x) = \phi(x) + a_i$ for all $x \in G$. A natural choice for a test of linear consistency is to verify that

$$Pr_{x,y \in G}[f_1(x) + f_2(y) \neq f_3(x + y)] \leq \delta$$

for some small enough choice of $\delta$.

(a) Assume $G, H$ are Abelian. Show that $f, g, h$ are linear-consistent iff for every $x, y \in G$

$$f(x) + g(y) = h(x + y).$$

(b) Let $G = \{+1, -1\}^n$ and $H = \{+1, -1\}$. First note that since $a_i \in \{+1, -1\}$,
then linear consistent $f_i$ must be linear functions or “negations” of linear functions. We refer to the union of linear functions and the negations of linear functions as the affine functions. In class we expressed the minimum distance of $f$ to a linear function. Express the minimum distance of a function $f$ to an affine function.

(c) Show that if $f_1, f_2, f_3$ satisfy the above test, then for each $i \in \{1, 2, 3\}$, there is an affine function $g_i$ such that $Pr_{x \in G}[f_i(x) \neq g_i(x)] \leq \delta$.

(d) (Extra credit) Show that there are linear consistent functions $g_1, g_2, g_3$ such that for

$$i \in \{1, 2, 3\}, Pr_{x \in G}[f_i(x) \neq g_i(x)] \leq \frac{1}{2} - \frac{2\gamma}{3}$$

where $\gamma = \frac{1}{2} - \delta$.

2. Dictator functions, also called projection functions, are the functions mapping $\{+1, -1\}^n$ to $\{+1, -1\}$ of the form $f(x) = x_i$ for $i$ in $\{1, 2, 3\}$.

Consider the following test for whether a function $f$ is a dictator: Given parameter $\delta$, the test chooses $x, y, z \in \{1, -1\}^n$ by first choosing $x, y$ uniformly from $\{1, -1\}^n$, next choosing $w$ by setting each bit $w_i$ to $-1$ with probability $\delta$ and $+1$ with probability $1 - \delta$ (independently for each $i$), and finally setting $z$ to be $x \circ y \circ w$, where $\circ$ denotes the bitwise multiply operation. Finally, the test accepts if $f(x)f(y)f(z) = 1$ and rejects otherwise.

(a) Show that the probability that the test accepts is

$$\frac{1}{2} + \frac{1}{2} \sum_{S \subseteq [n]} (1 - 2\delta)^{|S|} f(S)^3.$$

(b) Show that if $f$ is a dictator function, then $f$ passes with probability at least $1 - \delta$. 


(c) Show that if $f$ passes with probability at least $1 - \epsilon$ then there is some $S$ such that 
$$\hat{f}(S) \geq 1 - 2\epsilon$$ and such that $f$ is $\epsilon$-close to $\chi_S$.

(d) Why isn’t this enough to give a dictator test? (i.e., what nondictators might pass?)

Give a simple fix.

3. Show that if there is a PAC learning algorithm for a class $C$ with sample complexity 
$$\text{poly}(\log n, 1/\epsilon, 1/\delta),$$ then there is a PAC learning algorithm for $C$ with sample complexity 
dependence on $\delta$ (the confidence parameter) that is only $\log 1/\delta$ – i.e., the “new” PAC 
algorithm should have sample complexity $\text{poly}(\log n, 1/\epsilon, \log 1/\delta)$. (It is ok to assume that 
the learning algorithm is over the uniform distribution on inputs, although the claim is 
true in general.)

4. Consider the following graph-based linearity test. Let $G = (V, E)$ be a graph on $k = |V|$ 
vertices and let $f : \{\pm 1\}^n \to \{\pm 1\}$ be given.

• Sample $x_1, \ldots, x_k \in \{\pm 1\}^n$
• Query $f(x_i)$ for all $i \in [k]$ and $f(x_i \odot x_j)$ for all $(i, j) \in E$ where $x_i \odot x_j$ denotes the 
coordinate-wise product of $x_i$ and $x_j$.
• Accept if and only if $f(x_i)f(x_j) = f(x_i \odot x_j)$ for all $(i, j) \in E$.

Note that if $f$ is linear, then this graph-test always accepts.

(a) Prove that: For all $S \subseteq E$ such that $S \neq \emptyset$, then 
$$E[\Pi_{(i,j) \in S} f(x_i)f(x_j)f(x_ix_j)] \leq \max_\alpha |\hat{f}(\alpha)|$$

(b) Conclude that the probability that the above graph-test accepts is at most 
$$\frac{1}{2|E|} + \max_\alpha |\hat{f}(\alpha)|$$