The BSS loader should not load the first entry point into an even location. This scheme has two distinct disadvantages:

a. It violates the basic philosophy of symbolic programming - that of position independence of words labeled with symbols. This means that a programmer must decide what his first entry point will be and must never change it thereafter.

b. Secondly the assembly listing does not display the information about even and odd locations in a convenient manner. In order to find out whether a particular instruction will be loaded into an even or odd location, the programmer must examine both the relative address of the first entry point as well as the relative address of the instruction.

To eliminate these objections and also to permit the easy incorporation into Fap of a pseudo-operation "EVEN", the following modifications are proposed:

1. The loader should load all subprograms with an even origin. In addition, common storage should always begin at an odd location. The supervisor should load programs at an even origin.

2. The Fap assembler should guarantee that the first symbolic instruction has an even location. In order to do this, Fap must force the length of the transfer vector plus linkage director to be an even number. If required, the assembler should generate a word of zeroes between the transfer vector and the linkage director, or in the case of main subprograms, between the transfer vector and the first instruction.

3. Since the assembler could then be sure that the first instruction will have an even location, the addition of the following pseudo-operations will be trivial:

   a.) EVEN

   The program counter will be examined; if and only if it is odd, an instruction ANT 0,0 will be generated and both the program counter and the location counter will be advanced. A symbol in the location field will be assigned the even location.
b.) EVENC

If the common counter is odd, it will be reduced by 1. A symbol in the location field will be assigned the even location in common.

Note: For compatibility with this scheme, the library routine (FPT) would have to be adjusted so that it can locate the subprogram name in the linkage director relative to the transfer vector.