

TO: MSPM Distribution  
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This revision corrects the calling sequences for read  
and write.

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### Identification

User I/O Procedures to Read and Write  
read\_in, write\_out  
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### Purpose

Two procedures, read\_in and write\_out, are provided to read from the Multics standard input stream, user\_input, and write into the Multics standard output stream, user\_output. They are provided to simplify calls for the user and to interpret I/O status returns using the Multics standard error handling mechanism described in BY.11.00 - BY.11.04.

### Usage

```
call read_in (workspace, n, nin);  
call write_out (workspace, n);  
call write_out$nl (workspace, n);
```

where

workspace is a character string (either varying or non-varying) at least n characters in length.

n is a fixed decimal integer indicating how many characters are to be read or written. If n is zero or negative, n is taken to be the length of workspace.

nin is the number of characters actually read.

In read\_in, input (up to n characters) is read from the stream user\_input into the character string, workspace. Input is read up to and including the first break character found in the stream, user\_input. That is, read\_in reads the smaller of 1) n characters, or 2) the number of characters up to and including the break character; then returns in nin, the number of characters read. In normal command usage, the only break character is the new line, <NL>. However, the user may define his break characters using the I/O call, breaks, described in BF.1.12.

Write\_out writes from the character string, workspace, into the stream, user\_output. Write\_out writes n characters. The \$nl entry adds a new line character to the end of the character string contained in workspace before writing.

Implementation

The procedures `read_in` and `write_out` are much alike except for their call to the I/O system and the value returned by `read_in`. Hence they will be considered together.

The procedures examine `n` to determine whether it is negative. If so, `n = length (workspace)`;

where `length` is a PL/I built-in function. If `write_out` was called, `n` is incremented by 1 and a new line character is placed in `workspace`. Then one of the following calls (described in BF.1.12) is made:

```
call read ("user_input", addr (workspace), n, nin,
          status);
```

```
call write ("user_output", addr (workspace), n, nout,
           status);
```

where `addr` is a PL/I built-in function which returns a pointer to the character string, `workspace`; `nout` is a return argument which indicates the number of characters written. Declarations are:

```
dcl workspace char (*),
      (n,nin, nout) fixed binary (17),
      status bit (144);
```

The bit string, `status`, is returned by the I/O system containing status information about the transaction. BF.1.21 contains a complete description of status information. When control returns, `read_in` and `write_out` call `check_io_status` (see BY.4.03) and if an error is indicated, call `seterr` (see BY.11.01) with the error information. They then signal `read_in_err` or `write_out_err` respectively.

Asynchronous I/O

The streams `user_input` and `user_output` are asynchronous by default. This means that at any time, more information may have been read by the device than the user has yet requested in read calls. Subsequent read calls collect that information. On writing, some of the information which the user thinks he has written may not actually have appeared yet on the device. Normally the user is not affected by these asynchronous characteristics.

Occasionally, however, the user may want to start all I/O over fresh. He can wipe out read-ahead or write-behind data by calling the procedures `reset_user_in` and `reset_user_out` described in BY.4.04.