

MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 Department of Electrical Engineering and Computer Science
 6.001—Structure and Interpretation of Computer Programs
 Spring Semester, 1999

Recitation – Friday, February 12

1. Cons Pairs and Lists

Recall the contract for `cons`, `car`, `cdr`, `pair?`, and `null?`.

```
(car (cons a b)) == a           (list a b c) == (cons a
(cdr (cons a b)) == b           (cons b
(pair? (cons a b)) == #t        (cons c nil)))
(null? nil) == #t
```

There are three main methods of representing cons and list structures. You should be able to convert between them.

Scheme Expression	Box & Pointer	Scheme Printout
(cons 1 2)		(1 . 2)
(cons 1 (cons 2 nil))		(1 2)
(cons 1 nil)		(1)
(cons 1 (cons 2 3))		(1 2 . 3)
(cons (cons 1 2) nil)		((1 . 2))
(list 1 2 3 4)		(1 2 3 4)
(list 1 (cons 2 3) (list 4 5))		(1 (2 . 3) (4 5))

2. Other Accessors

In scheme, we often want to access elements deep in a cons structure. Therefore, the following accessors have been defined to help us out:

```
(cadr x) == (car (cdr x))           (cddr x) == (cdr (cdr x))
(caddr x) == (car (cdr (cdr x)))   (cdadar x) == (cdr (car (cdr (car x))))
(cdaar x) == (cdr (car (car x)))   etc, etc...
```

For lists, we also often want to easily access the n'th element of a list. The accessors **first**, **second**, **third**, ..., **tenth** are defined to access the corresponding values of a list. For example,

```
(sixth (list 1 2 3 4 5 6 7 8 9))
;Value: 6
```

How could you define **first**, **second**, **third**, and **fourth** using the `c???r` functions?

```
(first x) == (car x)               (third x) == (caddr x)
(second x) == (cadr x)             (fourth x) == (caddr x)
```

3. Practice

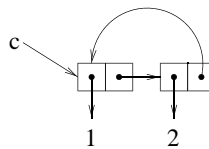
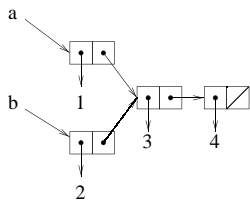
Draw box and pointer diagrams and write what will Scheme print for the following expressions.

```
⇒ (define x (cons 5 2))
⇒ (car x)
;Value: 5
⇒ (cdr x)
;Value: 2
⇒ (car (cdr x))
;Error: 2 passed car
⇒ (define y (cons sqrt x))
⇒ (car (cdr y))
;Value: 5
⇒ (car y)
;Value: #[compiled-procedure 24]
⇒ (define z (cons ((car y) 49) x))
⇒ z
;Value: (7 5 . 2)
```

Write a Scheme expression that will print each of the following. Also draw box and pointer diagrams.

```
⇒ (list 1 2 3)
;Value: (1 2 3)
⇒ (cons 1 (cons 2 3))
;Value: (1 2 . 3)
⇒ (cons (list 1 (list 2)) 3)
;Value: ((1 (2)) . 3)
```

Write Scheme expressions that correspond to the following.



```
(define c (list 3 4))
(define a (cons 1 c))
(define b (cons 2 c))
```

Trick Question! We can't do this yet!

4. Functions on Lists

We saw that we have the primitive function `pair?` to see if an object is a pair. What if we wanted to write the function `list?` to see if an object is a list?

What is the contract for `list?` $\forall x_1, x_2, \dots, x_n \quad (\text{list? } (\text{list } x_1 \ x_2 \ \dots \ x_n)) == \#t$

What's another way to write it? $(\text{list? } \text{nil}) = \#t$
 $(\text{list? } (\text{cons } x \ l)) = \#t \implies (\text{list? } l)$

Now, how can we write `list?` in scheme?

```
(define (list? x)
  (or (null? x)
      (and (pair? x)
            (list? (cdr x)))))
)
```

What is the Order of Growth of `pair?` and `list?` ?

`pair?` is $\Theta(1)$ and `list?` is $\Theta(n)$, where n is the length of the list.